Design and Analysis of Algorithms

L11: MergeSort

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Resources

- Text book 1: Levitin (Mergesort)
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MergeSort

- Problem: Given a set of N elements, sort the elements in ascending (or descending) order
  – Assume that these elements are in an array of size N
- Approaches
  – Divide and Conquer approach

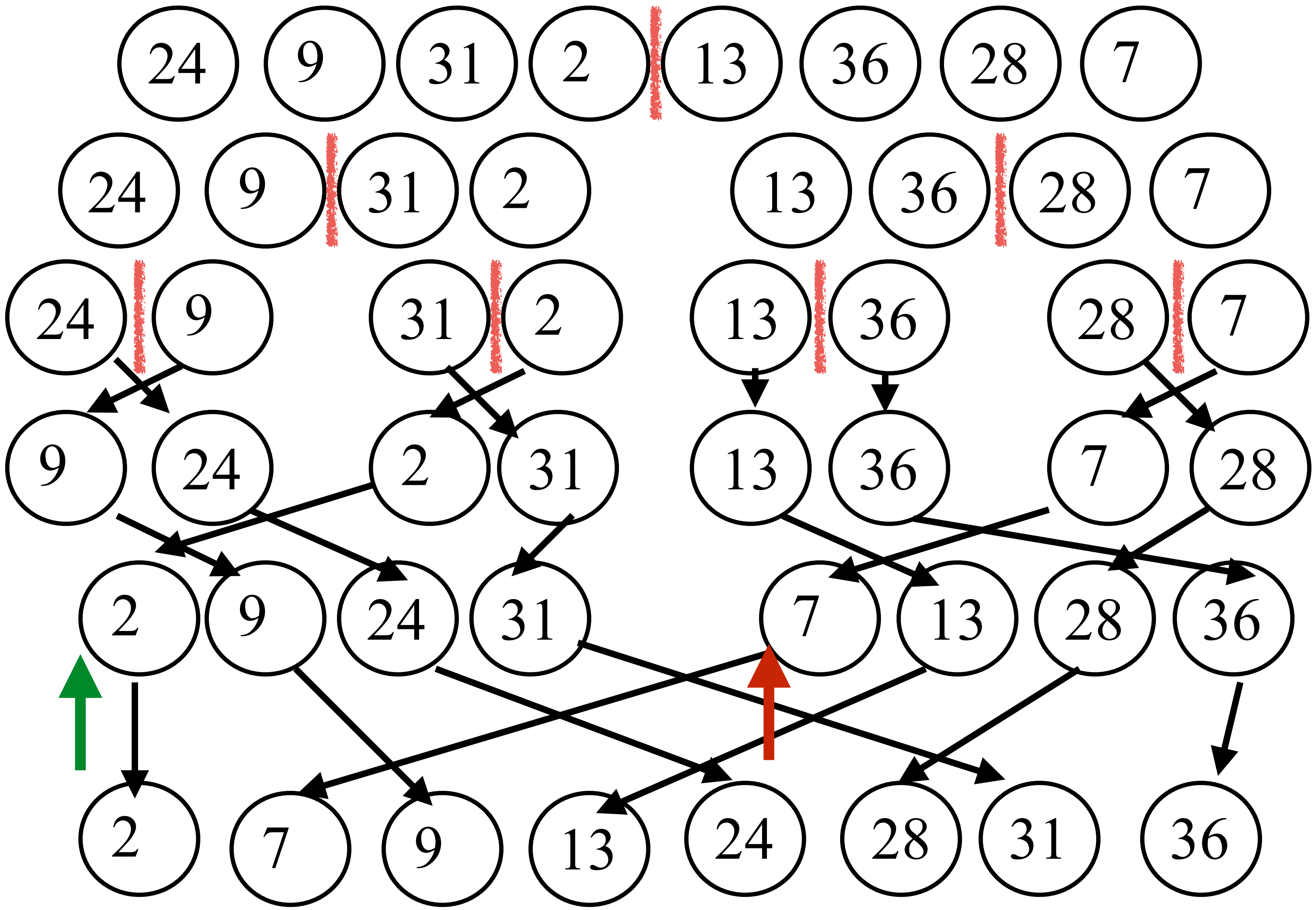
Sort Algorithms

- Bubble sort
- Selection sort
- Insertion sort
- **Mergesort**
- Quicksort
- Shell sort
- Heap sort
- Radix sort

MergeSort

- Basic idea
  - Take two sorted list and merge them into a single sorted list.
- Approach
  - Keep dividing the elements into (almost) equal half size (recursively) till sublist becomes of size 1
  - List of size `1` is sorted by default
  - Merge the sorted lists and keep repeating (recursively back)
  - When all the lists are merged, all elements are sorted.

MergeSort Example



MergeSort

- Split array A[1:n] into about equal halves
  - Make copies of each half in arrays B and C
- Sort arrays B and C recursively
- Merge sorted arrays B and C into A as follows:
  - Repeat until one of the arrays becomes empty
    - Compare the first elements of the remaining unprocessed portions of the arrays
    - Copy the smaller of the two into A,
      - Increment the index of the array (smaller)
  - Once all elements in one of the arrays are copied
    - Copy the remaining unprocessed elements from the other array into A.

Algo: MergeSort

- Algo `MergeSort(1,n,A[])`

```
#Sort array A recursive by merging
#i/p: unsorted array A[1:n]
#o/p: sorted array A[1:n]
if n>1, then
  copy A[1:n/2] to B[1:n/2]
  copy A[n/2+1:n] to C[1:n/2]
  Mergesort(1,n/2,B) #recursive
  Mergesort(1,n/2,C) #recursive
  Merge(B,C,A) # merge two arrays
```

Algo: MergeSort

- Algo Merge(B[1:p],C[1:q],A[1:p+q])

```
#maintain one index for each array
i←1; j←1; k←1;
while (i<p+1) and (j<q+1) do
    if (B[i]≤C[j]), then
        A[k] ← B[i]
        i ← i+1
    else
        A[k] ← C[j]
        j ← j+1
    k ← k+1
if (i > p) then #B has been fully copied to A
    copy C[j:q] to A[k:p+q]
else
    copy B[i:p] to A[k:p+q]
```

MergeSort: Analysis

- Each step of Mergesort
  - Two recursive invocations of size `n/2`: `2T(n/2)`
  - Merging of two n/2 array into one array of size n
    - Time complexity: n
- Recurrence relation for time complexity becomes

$$\begin{aligned} T(n) &= 2T(n/2) + n \\ &= 2(2T(n/4)+n/2)+n = 2^2T(n/2^2)+n+n \\ &= \ldots \\ &= 2^kT(n/2^k)+n+\ldots(\log_2 n \text{ times}) \\ &= n*T(1)+n\log_2 n = n + n\log_2 n \\ &= \Theta(n\log_2 n) \end{aligned}$$

- Space complexity $= \Theta(n)$

Mergesort Shortcomings

- Creates a new array i.e. requires additional O(n) space
  - No obvious way to merge in place in linear time.
- It is inherently recursive.
  - Recursive implemenation requires function invocation and return, a costly operation.
- Thus, Generally, not used in pratice.
- Alternative approaches
  - Can we ensure that left part is always less than the rigth part.
    - Thus, no need to merge the two.
    - Approach taken by **QuickSort.**

Summary

- Mergesort
  - Not in place sort
  - Stable sort
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