Noname manuscript No.

(will be inserted by the editor)

New subspace method for unconstrained black box optimization

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Received: date / Accepted: date

Abstract This paper defines an efficient subspace method, called SSBBO, for unconstrained black box optimization problems where the objective function has Lipschitz continuous gradients, but only function values contaminated by rounding errors are available. SSBBO employs line searches along directions constructed on the basis of quadratic models. These approximate the objective function in a subspace spanned by some of the previous search directions. A worst case complexity bound on the number of iterations and function evaluations is derived for a basic algorithm using this technique. Numerical results for a practical variant with additional heuristic features show that SSBBO has superior performance compared to the best solvers from the literature.

Keywords Unconstrained black box optimization \cdot Subspace technique \cdot Line search approach \cdot Complexity results

Mathematics Subject Classification (2000) primary 90C56

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^{1.} The first author acknowledges the financial support of the Doctoral Program $Vienna\ Graduate\ School\ on\ Computational\ Optimization\ (VGSCO)$ funded by the Austrian Science Foundation under Project No W1260-N35

1 Introduction

In this paper, we discuss a new subspace technique for solving the unconstrained black box optimization problem

$$\min_{x \in \mathbb{R}^n, \\ \sin x \in \mathbb{R}^n, } (1)$$

where $f: \mathbb{R}^n \to \mathbb{R}$ is continuously differentiable with Lipschitz continuous gradients, but accessible only through function values contaminated by rounding errors. In particular, exact function values, gradients, Lipschitz constants, or other information about the structure of f are not available. However, a numerical gradient approximation can be estimated by a finite difference method. We denote by $\widetilde{g} := \widetilde{g}(x)$ the estimated gradient at x and by $\widetilde{f} := \widetilde{f}(x)$ the inexact function value available at x.

In general, the goal of local optimization is to find at least one point x that is an ε -approximate stationary point if the norm of the exact gradient g(x) of f is $\leq \varepsilon$. This criterion cannot be verified under our assumptions. Instead, our algorithm tries to find a point whose exact function value (although unknown) is at least as good as the exact function value of an ε -approximate stationary point whose approximated function value has been evaluated, for a threshold $\varepsilon = O(\sqrt{\omega})$, where ω is a bound on the error in the function values. If f has only one stationary point and ω is sufficiently small, this guarantees that the point found is arbitrarily close to the minimizer (if there are only finitely many stationary points, a point with small gradient could be guaranteed by a variant with restarts).

Our algorithm uses line searches based on the inexact function values evaluated along the search directions based on quadratic models. This algorithm estimates the gradient of the objective function by a finite difference method and the Hessian matrix of the objective function by a non-traditional limited memory quasi-Newton method. We construct quadratic models in a subspace of the previous search directions and then solve these models approximately such that, if possible, reductions of the model function value and gradient norm can be enforced. If this is not possible, we solve these quadratic models in a full subspace. Assuming that the exact gradient norm is not small, we prove that the (unknown) exact gradient and the function value found by the line search satisfy the Wolfe conditions. Therefore, our algorithm finds (at an unknown time) a point whose unknown gradient norm is below a computable threshold related to the level of rounding errors. The point with the best inexact function value $\tilde{f}(x)$ among the evaluated values then has the required properties.

We say that a solver is robust if it has the highest number of solved problems, and efficient if it has the lowest relative cost of function evaluations. We say a solver is competitive if it is both robust and efficient. In this paper, we design an algorithm to be competitive in low, medium and high dimensions compared to other well-known black box optimization solvers.

We improve the state of the art in several respects:

• We ensure that our search directions satisfy the approximate angle condition

$$\frac{\widetilde{g}(x)^T p}{\|\widetilde{g}(x)\| \|p\|} \le -\Delta_{\mathbf{a}} < 0, \tag{2}$$

i.e., the angle between the direction p and the approximate gradient $\tilde{g}(x)$ at the point x obtained from a finite difference method remains bounded away from 90°. Here $0 < \Delta_a < 1$ is a given threshold for the approximate angle condition (2).

• We show numerically that the new subspace method is more efficient than the traditional limited memory BFGS quasi-Newton method in low and high dimensions, and that it is more efficient than the best solvers from the literature.

1.1 Related work

There are many different algorithms for solving the problem (1). They can be classified into deterministic and randomized. An excellent reference for the deterministic methods is Conn et al. [6]. More recently, LARSON et al. [22] have given a comprehensive discussion of both deterministic and stochastic methods. A comprehensive treatment of black box optimization solvers can be found in [18,20,31]. As shown in [18,20], **UOBYQA** by POWELL [30] and **BCDFO** by GRATTON et al. [14] were the two competitive model-based solvers for low-dimensional problems $(n \leq 20)$ and two Nelder-Mead algorithms, NELDER by Kelley [17] and NMSMAX by Higham [16], were also competitive for very small n. A more recent version of the BFO solver was described by PORCELLI & TOINT [29], which showed average behaviour in the noiseless case in low to high dimensions, but was robust only for large noise. Two recent randomized line search solvers are VRBBO by KIMIAEI & Neumaier [20] and VRBBON by Kimiaei [18]. In [18,20], it was shown that VRBBO and VRBBON are robust and moderately efficient compared to other competing local and global solvers for low-dimensional problems, but are competitive for problems in medium and high dimensions for both noiseless and noisy cases. Recently, researchers have shown increased interest in applying quasi-Newton methods that use the finite difference method to estimate the gradient to solve unconstrained black box optimization problems, e.g. [3,4,10,20,23-25,27]. Although this method is very useful for small scale problems, full inverse matrices require a lot of memory. A common suggestion is to use a quasi-Newton method with limited memory instead of such a standard method, which has a low memory requirement, see Liu & Nocedal [24]. FMINUNC from the Matlab optimization toolbox uses Wolfe line search along standard quasi-Newton directions. This solver was efficient and moderately robust for low and medium dimensional problems (cf. [20, Section 7]) in the noiseless case, but not in the presence of high noise (cf. [18, Section 7]), since a finite difference method leads to misleading information and gives a bad quasi-Newton direction.

To estimate the gradient, BERAHAS et al. [2, Table 2] discussed three bounds on the number of function evaluations, step sizes, and the exact gradient norm for all existing randomized and deterministic methods, such that the reduction in the error of the gradient norm is asymptotically bounded by the reduction in the exact gradient norm.

Line search methods proceed iteratively by generating a sequence of estimated or exact step sizes in the hope of finding a good decrease in the function value. Their global convergence was proved by AL-BAALI & FLETCHER [1], FLETCHER [10], and NOCEDAL & WRIGHT [28]. BERAHAS et al. [2] discuss complexity bounds for nonconvex, convex, and strongly convex cases for the accuracy of two different gradient estimates. Other useful references for first and second order complexity bounds are CURTIS et al. [7,8] and GRATTON et al. [13] for gradient-based trust region algorithms and GRAPIGLIA et al. [12] for gradient-based nonlinear step size control algorithms. CARTIS et al. [5] introduced a non-monotone gradient-related algorithm for nonconvex smooth unconstrained optimization problems that finds at most $\mathcal{O}(\varepsilon^{-2})$ function and gradient evaluations where the 2-norm gradient of a point x is below ε in the worst case.

BERAHAS et al. [3] estimated noise in the same way as HAMMING [15] did. This estimate was added to the Armijo line search method and used to estimate the gradient by the finite difference technique. However, the effects of noise in BFGS updating have not been studied. Recently, XIE et al. [32] showed that the Wolfe conditions can be satisfied with the exact function value and gradient when these conditions are satisfied with the inexact function and gradient and the exact gradient is not small. They modified the BFGS method with the Wolfe line search without estimating the noise and showed that this method converges to a neighborhood of the solution, while the effects of the noise are taken into account in BFGS updating.

Kimiaei et al. [21] have proposed a limited memory method with minimal memory requirements for bound constrained optimization using the exact gradient. In this paper, we propose an extension of such a limited memory method for unconstrained black box optimization. Our algorithm uses Wolfe line searches based on the inexact function values evaluated along the search directions based on quadratic models in the subspace of the previous search directions. These line searches are called approximate Wolfe line searches, where in each iteration, instead of estimating the gradient of the objective function, the directional derivative is estimated by a new variant of finite difference methods only with one function evaluation, and at termination, the gradient is estimated by a finite difference method with n function evaluations. Hence these approximate line searches are not costly for problems in medium and high dimensions. The Hessian matrix of the objective function is approximated by a non-traditional limited memory quasi-Newton method. We improve these line searches in a heuristic way when they fail in the presence of rounding errors. Hoping to find reductions of the model function value and gradient norm, we construct quadratic models in a subspace of the previous search directions and then solve these models approximately. When these reductions are not possible, we solve these quadratic models in a full subspace, resulting in non-traditional limited memory quasi-Newton directions. In order to obtain the complexity bound, it is important that we can estimate the extent of the reduction of the function value, which can be achieved by a variant of the Wolfe line search method [1] with inexact function values and approximated gradients, as discussed later. Although our algorithm stops when the estimated gradient norm is below a computable threshold related to the level of rounding errors, it finds a point whose exact gradient norm is below a computable threshold related to the level of rounding errors. In this case, the approximate Wolfe conditions are numerically satisfied for both the inexact function and the approximate gradient. Assuming that the discretization and function evaluation errors occur in the estimation of the gradient and that the exact gradient norm is not small, we show in the same way as XIE et al. [32] that the Wolfe conditions are satisfied with the exact function value and gradient, but with the difference that these conditions are independent of the choice of search directions. In fact, we take into account the effects of noise in an estimated Hessian update. In contrast to XIE et al. [32], however, we enforce that our non-traditional limited memory quasi-Newton direction enforces the approximate angle condition (2). Therefore, we need not worry about whether the approximate Wolfe conditions can produce a quasi-Newton direction with both the inexact function and the gradient whose angle is bounded away from 90° by the approximate gradient. Ignoring some parts of the Hessian information reduces the efficiency of the traditional limited memory technique. We attempt to show numerically that the improved subspace method is more efficient and robust than the standard BFGS quasi-Newton method, even for problems in low and medium dimensions, and than the traditional limited memory BFGS quasi-Newton method for low and high dimensions.

1.2 The Wolfe line search in exact arithmetic

Our new algorithm uses an inexact version of the Wolfe line search method of AL-BAALI & FLETCHER [1], which is proposed there assuming exact arithmetic; we therefore call it **WLS-exact**. This method finds a step size that not only leads to a sufficient reduction in the function value, but is also close to a minimizer of the objective function with respect to the step size. Each iteration of **WLS-exact** is the scheme (S2) in [1, Section 2].

AL-BAALI & FLETCHER [1] assume that the objective function is bounded below, i.e., $f(\alpha) \geq \overline{f}$ for $\alpha \geq 0$ and a given finite threshold \overline{f} . When function values and gradients are known exactly, **WLS-exact** enforces the **Wolfe conditions**

$$f(x + \alpha p) \le f(x) + \rho \alpha g(x)^T p, \tag{3}$$

$$|g(x + \alpha p)^T p| \le -\sigma g(x)^T p,\tag{4}$$

where $\sigma > \rho > 0$. Let $\phi(\alpha) := f(x + \alpha p)$. Then the ρ -line intersects the line $f = \overline{f}$ at point

$$\mu := \frac{\overline{f} - f(x)}{\rho g(x)^T p} = \frac{\overline{f} - \phi(0)}{\rho \phi'(0)^T p}.$$
 (5)

WLS-exact restricts the step sizes to $(0, \mu]$. It finds an interval with suitable step sizes that satisfy the sufficient descent condition (3), and then finds a desired step size within that interval that satisfies the curvature condition (4). Interpolation and extrapolation are used to generate a new step size. Let a_j (for $j=1,2,\cdots,\infty$) be a sequence of the current best points of the function ϕ satisfying the condition (3) but not the condition (4), let b_j (for $j=1,2,\cdots,\infty$) be a sequence of points violating either (3) or $\phi(b_j) \geq \phi(a_j)$ or both, and let α_j (for $j=1,2,\cdots,\infty$) be a sequence of the current trial points of the function ϕ . Then the interval (a_j,b_j) found by **WLS-exact** contains the acceptable points or the points that $f(\alpha_j) \leq \overline{f}$. For j=1, $a_1=0$, $b_1=\infty$, and $0<\alpha_1\leq\mu$. Let $0<\tau_1\leq\tau_2\leq\frac{1}{2}$, $1<\tau_3\leq\tau_4$, $0<\tau_5\leq\tau_6\leq\frac{1}{2}$, and $a_j,b_j>0$ for $j\geq 1$. Whenever the condition $\phi(\alpha_j)>\phi(0)+\alpha_j\rho\phi'(0)$ or $\phi(\alpha_j)\geq\phi(a_j)$ is satisfied, the new step size

$$\alpha_{j+1} \in T(a_j, \alpha_j) := \begin{cases} [a_j + \tau_1(\alpha_j - a_j), \alpha_j - \tau_2(\alpha_j - a_j)] & \text{if } a_j < \alpha_j, \\ [\alpha_j + \tau_2(a_j - \alpha_j), a_j - \tau_1(a_j - \alpha_j)] & \text{if } \alpha_j < a_j \end{cases}$$
(6)

is obtained and the interval is updated as $a_{j+1} = a_j$ and $b_{j+1} = \alpha_j$. If the condition (4) is not satisfied the new step size

$$\alpha_{j+1} \in E(a_j, \alpha_j, b_j) := \begin{cases} [\min(\tau_3 \alpha_j, \mu), \min(\tau_4 \alpha_j, \mu)] & \text{if } a_j < \alpha_j < b_j = \infty, \\ [\alpha_j + \tau_5(b_j - \alpha_j), b_j - \tau_6(b_j - \alpha_j)] & \text{if } a_j < \alpha_j < b_j \le \mu, \\ [b_j + \tau_6(\alpha_j - b_j), \alpha_j - \tau_5(\alpha_j - b_j)] & \text{if } b_j < \alpha_j < a_j \le \mu \end{cases}$$
(7)

is computed and the interval is updated as $a_{j+1} = \alpha_j$ and $b_{j+1} = b_j$. In (6) and (7), α_{j+1} can be found within the defined intervals by minimizing quadratic/cubic polynomial interpolations; for more details see [1, Section 3].

Proposition 1 [1, Theorem 2.2]

Given $\sigma \geq \rho > 0$ and the finite threshold value \overline{f} , one of the following statements is valid:

- (i) WLS-exact ends with $f(x + \alpha_j p) < \overline{f}$ for some $\alpha_j \in (0, \mu]$. Here μ comes from (5).
- (ii) WLS-exact ends with α_j satisfying (3) and (4).
- (iii) **WLS-exact** fails, each interval $I_j := (a_j, b_j)$, $j = 1, 2, \dots, \infty$, contains an interval of acceptable points and there exists a limit point $\widehat{\alpha} \in I_j$, for all j, such that, for sufficiently large j, a_j are monotonically (not strictly) increasing, $\lim_{j \to \infty} a_j = \widehat{\alpha}$ and b_j are monotonically (not strictly) decreasing,

 $\lim_{j\to\infty}b_j=\widehat{\alpha},\ and\ \widehat{\alpha}\ is\ an\ acceptable\ point\ as\ in\ (ii).$

In particular, if $\sigma > \rho > 0$, (iii) cannot occur and WLS-exact ends. Since we do not know in I_j whether $a_j < b_j$ or not, we replace a_j by b_j whenever $b_j < a_j$ for some $j = 1, 2, \dots, \infty$.

1.3 Overview of the new method

This paper constructs a new deterministic solver for unconstrained black box optimization problems, called **SSBBO** for "subspace black box optimization". The basic version of **SSBBO** is a quasi-Newton

algorithm using a finite difference technique for the approximation of the gradient, the limited memory direction by Kimiaei et al. [21], and the Wolfe line search by Al-Baali & Fletcher [1] for finding step sizes but with the inexact function values and the estimated gradients. We call this line search method the **approximate Wolfe line search** method. This method is not costly even for problems in high dimensions, since one function evaluation is required in each iteration to estimate the directional derivative by a **new variant of finite difference methods** and n function evaluations are required to estimate the gradient by the forward finite difference method only in the last iteration.

Assuming that discretization and rounding errors occur when the gradient is estimated by a finite difference method and the exact gradient norm is not small, we prove that the Wolfe line search conditions can be satisfied with both the exact gradient and the function value, while these conditions are numerically satisfied with both the inexact function and the gradient. In this case, we prove the complexity result for both the approximate Wolfe line search algorithm and the basic version of **SSBBO** for the general case regardless of the choice of search directions. The order of our complexity bound is consistent with that of Berahas et al. [2].

The enhancements that make **SSBBO** competitive are:

- The new subspace direction is computed by finding the solution to a linear problem in the subspace and then solving a simple optimization problem to find, if possible, a decrease in both the model function value and its gradient norm.
- In the presence of rounding errors, the subspace direction computed need not satisfy the approximate angle condition (2) and must be modified. We find a robust modification described in Subsection 3.1.
- The Wolfe line search algorithm [1] may fail in finite precision arithmetic. In this case, the step size is generated by a **heuristic formula**. Although the algorithm moves somewhat away from a minimizer in this case, it may approach a minimizer in the next few attempts. At the very least, this idea helps us avoid failure.

SSBBO is available at https://www.mat.univie.ac.at/~neum/software/SSBBO

In Section 2, we introduce a basic version of our algorithm, called **SSBBO-basic**, which uses the approximated Wolfe line search algorithm and then obtain the complexity results of **SSBBO-basic**.

Section 3 discusses new enhancements. In Subsection 3.1 the subspace information and its update are described. Also, the subspace directions are computed while trying to reduce the function value and gradient norm of the quadratic model. In Subsection 3.2 an improved version of the Wolfe line search algorithm [1] is described, but with the inexact function values and estimated gradients, and in Subsection 3.3 an improved version of **SSBBO-basic** is introduced.

Section 4 contains extensive numerical results. They show that **SSBBO** of the present authors is effective and can compete with the best solvers from the literature when all 568 unconstrained problems from the CUTEst collection of GOULD et al. [11] are tested, with dimensions ranging from 1 to 9000. Our conclusions are summarized in Section 5.

2 Complexity of a basic version of SSBBO

For the theoretical analysis we first describe a simplified algorithm **SSBBO-basic**, an algorithm for unconstrained black box optimization problems that uses line searches along directions whose angle with an approximate gradient estimated by the forward finite differences is bounded away from 90°. The line search algorithm, called **AWLS-basic**, is an adaptation of the Wolfe line search algorithm described in Section 1.2 to account for inexact functions values and gradients. Given $\tilde{\sigma} > \tilde{\rho} > 0$, this

algorithm enforces the approximate Wolfe conditions

$$\widetilde{f}(x + \alpha p) \le \widetilde{f}(x) + \widetilde{\rho}\alpha \widetilde{g}(x)^T p,$$
(8)

$$|\widetilde{g}(x+\alpha p)^T p| \le -\widetilde{\sigma}\widetilde{g}(x)^T p,$$
 (9)

with the inexact function and the estimated gradient. The **sufficient descent** (8) guarantees $\widetilde{f}(x + \alpha p) < \widetilde{f}(x)$ while the **curvature condition** (9) guarantees that α is not too far from a minimizer of the function f. **AWLS-basic** is like **WLS-exact**, but with the approximate Wolfe conditions (8) and (9). In fact the gradient is estimated by the forward finite difference method, denoted by **FFD**.

FFD takes the current point x and its inexact function value f(x) as input. It uses the tuning parameter $0 < \varepsilon_h < 1$ for adjusting the finite difference step size and returns the estimated gradient g(x) at x.

Algorithm 1 FFD, estimation of the gradient by the forward finite difference method

1: **for** $i = 1, 2, 3, \cdots$ **do**

Compute the step size

2: Compute $h_i := \sqrt{\varepsilon_h} \operatorname{sign}(x_i) \max\{|x_i|, 1\}$.

Approximation of the *i*th component of the gradient

3: Compute
$$\widetilde{g}_i := (\widetilde{g}(x))_i := \frac{\widetilde{f}(x + h_i e_i) - \widetilde{f}(x)}{h_i}$$
. $\triangleright e_i$ is the i th coordinate direction h_i .

To simplify our notation in the next algorithm, we define $\widetilde{\phi}(\alpha) = \widetilde{f}(x+\alpha p)$ and $\widetilde{\phi}'(\alpha) = \widetilde{g}(x+\alpha p)^T p$. We know that estimating the gradient $\widetilde{g}(x+\alpha p)$ in (9) by **FFD** requires n additional function evaluations, which can be costly in high dimensions. To avoid this drawback, we compute $\widetilde{\phi}'(\alpha_j)$ directly using a new variant of the forward finite difference method with one function evaluation, called **FFDD**, except in the last iteration. Therefore, **AWLS-basic** produces step sizes that are not too far from a minimizer and speeds up the process of reaching a minimizer.

FFDD takes the search direction p, the trial point $x + \alpha p$, and the current inexact function value $\widetilde{\phi}(\alpha)$ as input and uses the tuning parameter $0 < \varepsilon_h < 1$ for adjusting the step size. It returns the directional derivative $\widetilde{\phi}'(\alpha)$ as output.

Algorithm 2 FFDD, estimation of the directional derivative by a new variant of FFD

Compute the step size

1: Compute $h := \varepsilon_h \frac{\|x + \alpha p\|}{\|p\|}$.

Approximation of the directional derivative $\widetilde{\phi}'(\alpha)$ directly

2: Compute
$$\widetilde{\phi}'(\alpha) = \widetilde{g}(x + \alpha p)^T p = \frac{\widetilde{\phi}(\alpha + h) - \widetilde{\phi}(\alpha)}{h}$$

AWLS-basic takes the current point x as input and returns a better point $x + \alpha p$, its inexact function value $\widetilde{\phi}(\alpha) = \widetilde{f}(x + \alpha p)$, and its estimated gradient $\widetilde{g}(x + \alpha p)$ as output. It uses the tuning parameters

Algorithm 3 AWLS-basic, basic approximate Wolfe conditions

```
1: Compute \widetilde{\mu} = (\overline{f} - \widetilde{\phi}(0))/(\widetilde{\rho}\widetilde{\phi}'(0)).
 2: Initialize the initial interval (a_0, b_0) = (0, \widetilde{\mu}).
 3: Choose \alpha_0 \in (0, \infty).
4: for j = 1, 2, 3, \cdots do
           Trying to enforce the sufficient descent condition (8)
 5:
           Compute \phi(\alpha_j).
           First stopping test
                                                                                                                  \triangleright desired decrease in \widetilde{\phi} has been found
           if \widetilde{\phi}(\alpha_i) \leq \overline{f} then
 6:
                Compute the gradient \widetilde{g}(x + \alpha_j p) by FFD and terminate AWLS-basic.
 7:
 8:
           if \widetilde{\phi}(\alpha_j) > \widetilde{\phi}(0) + \alpha_j \widetilde{\rho} \widetilde{\phi}'(0) or \widetilde{\phi}(\alpha_j) \geq \widetilde{\phi}(a_j) then
9:
10:
                Choose \alpha_{j+1} \in T(a_j, \alpha_j).
Update a_{j+1} = a_j and b_{j+1} = \alpha_j.
11:
                                                                                                                                             \triangleright update the interval
12:
                    Approximate the directional derivative \widetilde{\phi}'(\alpha_i) by FFDD
                Approximate \widetilde{\phi}'(\alpha_i) by FFDD.
13:
                Second stopping test: Is the curvature condition (9) satisfied?
                if |\widetilde{\phi}'(\alpha_i)| \leq -\widetilde{\sigma}\widetilde{\phi}'(0) then
14:
15:
                      Compute the gradient \widetilde{g}(x + \alpha_j p) by FFD and terminate AWLS-basic.
16:
                Extrapolation or interpolation to get the new step size?
17:
                Set a_{j+1} = \alpha_j.
18:
                if (b_j - a_j)\widetilde{\phi}'(\alpha_j) < 0 then
                     Choose \alpha_{j+1} \in E(a_j, \alpha_j, b_j) and set b_{j+1} = b_j.
19:
                                                                                                                               \triangleright extrapolation has been done
20:
21:
                     Choose \alpha_{j+1} \in T(a_j, \alpha_j) and set b_{j+1} = a_j.
                                                                                                                                \triangleright interpolation has been done
22:
                end if
23:
           end if
24: end for
```

SSBBO-basic takes the initial point x_0 as input and returns an optimum point x_{best} and its inexact function value $\widetilde{f}_{\text{best}}$ as output. It uses the following tuning parameters:

 $\Delta_{\rm a} \in (0,1)$ (tiny parameter for the approximate angle condition (2)),

 $\widetilde{\sigma} > \widetilde{\rho} > 0$ (parameters for line search),

 $0 < \varepsilon_h < 1$ (parameter for adjusting the finite difference step size h),

 $0<\varepsilon_g<1$ (minimum threshold for the stopping test).

Algorithm 4 SSBBO-basic, basic subspace method for black box optimization

Initialization

- 1: Compute the initial inexact function value $\widetilde{f}_0 := \widetilde{f}(x_0)$.
- 2: Compute $\overline{f} = \tilde{f}_0 10^8 (1 + |\tilde{f}_0|)$.

- ▶ this choice has been suggested by FMINUNC
- 3: Estimate the initial gradient $\widetilde{g}_0 = \widetilde{g}(x_0)$ by **FFD**.
- 4: **for** $\ell = 1, 2, 3, \cdots$ **do**

Stopping test

- 5: if $\|\widetilde{g}_{\ell}\| \leq \varepsilon_g$, then set $x_{\mathrm{best}} := x_{\ell}$, $\widetilde{f}_{\mathrm{best}} := \widetilde{f}_{\ell}$, and stop; end if
 - Computing the search direction
- 6: Choose p_{ℓ} such that the approximate angle condition (2) holds.

Use **AWLS-basic** to find α_{ℓ} satisfying (8) and (9)

- 7: Perform **AWLS-basic** for computing $x_{\ell+1} = x_{\ell} + \alpha_{\ell} p_{\ell}$, $\widetilde{f}_{\ell+1} = \widetilde{f}(x_{\ell+1})$, and $\widetilde{g}_{\ell+1} = \widetilde{g}(x_{\ell+1})$.
- 8: end for

We prove complexity bounds for **AWLS-basic** and **SSBBO-basic** under the following assumptions, using the Euclidean norm:

- (A1) The level set $L(x_0) := \{x \in \mathbb{R}^n \mid f(x) \leq f(x_0), x_0 \in \mathbb{R}^n\}$ of f at the starting point x_0 is compact.
- (A2) There is a convex neighbourhood Ω of $L(x_0)$ such that in Ω , the objective function f is continuously differentiable and its gradient g is Lipschitz continuous, i.e., there exists a constant L > 0 such that

$$||g(x) - g(y)|| \le L ||x - y||, \text{ for all } x, y \in \Omega.$$
 (10)

(A3) The uncertainty of the function value $\widetilde{f}(x)$ obtained by a noisy oracle is globally bounded by a small threshold $\omega > 0$

$$|\widetilde{f}(x) - f(x)| \le \omega$$
, for all $x \in \mathbb{R}^n$. (11)

It is well-known that (A1) and (A2) imply that some global minimizer \hat{x} exists, and

$$\widehat{f} := \inf_{\ell > 0} f(x_{\ell}) \tag{12}$$

is attained there, hence finite. (A2) implies

$$||f(x+s) - f(x) - g(x)^T s|| \le \frac{L}{2} ||s||^2, \quad \text{for } x, x+s \in \Omega.$$
 (13)

Theorem 1 Assume that (A1)-(A3) hold and let p be the search direction. If $\varepsilon_h := \Theta(\omega)$ and

$$c_1 \sqrt{\omega} \le h \|p\| \le c_2 \sqrt{\omega} \tag{14}$$

then the directional derivative approximation defined by

$$\widetilde{g}_p(x) := h^{-1}(\widetilde{f}(x+hp) - \widetilde{f}(x))$$

satisfies

$$|\widetilde{g}_p(x) - g(x)^T p| = \mathcal{O}(\sqrt{\omega}||p||). \tag{15}$$

In particular, the approximate gradient estimated by FFD satisfies (15).

Proof By (11) and (14),

$$h|\widetilde{g}_p(x) - g(x)^T p| = |\widetilde{f}(x + hp) - \widetilde{f}(x) - hg(x)^T p| \le 2\omega + |f(x + hp) - f(x) - hg(x)^T p| \le 2\omega + \frac{L}{2}h^2 ||p||^2,$$

hence by (14)

$$|\widetilde{g}_p(x) - g(x)^T p| \le \frac{2\omega}{h} + \frac{L}{2} h ||p||^2 = O(\sqrt{\omega} ||p||).$$

By choosing $p = e_i$ (coordinate direction) for $i = 1, \dots, n$, we have $|\widetilde{g}_p(x) - g_i(x)| = \mathcal{O}(\sqrt{\omega})$ for $i = 1, \dots, n$, resulting in $\|\widetilde{g}(x) - g(x)\| = \mathcal{O}(\sqrt{\omega})$.

This result is a variant of results discussed in Berahas et al. [2, Section 2.1] for forward finite difference methods for estimating the gradient.

It is well known that optimization methods can get stuck in almost flat regions, and that the gradient for a local minimizer in finite precision arithmetic does not vanish. For a given threshold $\varepsilon > 0$, let $N(\varepsilon)$ be the number of function evaluations to find an ε -approximate stationary point, i.e., approximate with $||g(x)|| \le \varepsilon$. Then, under the assumptions (A1)–(A3), **SSBBO-basic** uses at most $N(\varepsilon)$ function evaluations to find a point x_{best} whose function value $f(x_{\text{best}})$ is at most

$$\sup\{f(x) \mid x \in \mathbb{R}^n, \quad f(x) \le f(x_0), \quad \text{and} \quad ||g(x)|| \le \varepsilon\}.$$

In fact, x_{best} is the final best point found by **SSBBO-basic**. As in [20], x_{best} is unknown to the algorithm because the gradients and Lipschitz constants are unknown. **SSBBO-basic** finds x_{best} whose function value $f(x_{\text{best}})$ is equal to or better than a point whose gradient was small. If the gradients are small only near a global optimizer, a point near the local optimizer is found. If some iterate passes close to a non-global local optimizer or a saddle point, **SSBBO-basic** may escape from its neighbourhood. In this case, only a variant with restarts would lead to convergence to a point with a small gradient.

We now discuss how to get the complexity result for **AWLS-basic**. It plays a key role in proving the complexity bound for **SSBBO-basic**.

We deduce from Theorem 1 that a complexity bound for AWLS-basic can be obtained .

Proposition 2 Assume that (A1)–(A3) hold, the gradient of f is estimated by **FFD**, the direction p_{ℓ} satisfies the approximate angle condition (2), and the step size α satisfies the strong Wolfe conditions (8) and (9). Denote by $\widetilde{\theta}$ the angle between $\widetilde{q}(x)$ and p and by θ the angle between q(x) and p. Then

if $\varepsilon_h := \Theta(\omega)$ and h satisfies (14):

(i) For all
$$x, y \in \mathbb{R}^n$$
 we have

$$\|\widetilde{g}(x) - \widetilde{g}(y)\| = \mathcal{O}(\sqrt{\omega}) + L\|x - y\|.$$

(ii) p satisfies the angle condition for the exact gradient

$$\frac{g(x)^T p}{\|g(x)\| \|p\|} < -\frac{\Delta_{\mathbf{a}}}{2}.$$
 (16)

(iii) Under the assumptions

$$\|\widetilde{g}(x)\| > \frac{2\sqrt{\omega}}{\Delta_{\mathbf{a}}(1-\sigma)},$$
 (17)

$$||g(x)|| \ge \max\left\{\frac{4\widetilde{\rho}\sqrt{\omega}}{\rho\Delta_{\mathbf{a}}}, \frac{(1+\widetilde{\sigma})\sqrt{\omega}}{\sigma\Delta_{\mathbf{a}}}\right\},$$
 (18)

and

$$\|\widetilde{g}(x)\|\|g(x)\| \ge \frac{16L\omega}{\rho(1-\widetilde{\sigma})\Delta_a^2},\tag{19}$$

 α satisfies the strong Wolfe conditions

$$f(x + \alpha p) \le f(x) + \rho_1 \alpha g(x)^T p, \tag{20}$$

$$|g(x+\alpha p)^T p| \le -\sigma_1 g(x)^T p,\tag{21}$$

where $0 < \rho_1 := \widetilde{\rho} - \rho < \widetilde{\rho} < \widetilde{\sigma} < \sigma_1 := \widetilde{\sigma} + \sigma$.

(iv) Given $0 < \rho_1 < \sigma_1$ and the minimal threshold $\alpha_{\min} \in (0,1)$ for the step size, we define

$$\overline{\mu} := \overline{\mu}(x) = \frac{\overline{f} - f(x)}{\rho_1 g(x)^T p}.$$
(22)

and choose the step size $\alpha \in (0,\overline{\mu}]$, for $\ell \geq 1$, and $\tau \in (0,1)$. Then AWLS-basic needs at most

$$K := K(x) = \left\lceil \frac{\log \alpha_{\min}/\overline{\mu}(x)}{\log \tau} \right\rceil \tag{23}$$

iterations and at most

$$N^f := N^f(x) = 2K(x) + n (24)$$

function evaluations to satisfy (20) and (21) (The factor n comes from the approximated gradient evaluations needed in the last iteration of Wolfe line search).

Proof (i) From (A2) and Theorem 1, we have

$$\|\widetilde{g}(x) - \widetilde{g}(y)\| = \|\widetilde{g}(x) - g(x) + g(x) - g(y) + g(y) - \widetilde{g}(y)\| = \mathcal{O}(\sqrt{\omega}) + L\|x - y\|.$$

- (ii) By [32, Lemma 3.11], since $\cos \tilde{\theta} \ge \Delta_a > 0$ by (2) and (15) holds, $\cos \theta \ge \Delta_a/2$ and so (16) is obtained.
- (iii) By (i), Cauchy–Schwarz inequality, and (9), we obtain

$$(\mathcal{O}(\sqrt{\omega}) + \alpha L \|p\|) \|p\| = \|\widetilde{g}(x + \alpha p) - \widetilde{g}(x)\| \|p\|$$

$$\geq (\widetilde{g}(x + \alpha p) - \widetilde{g}(x))^T p \geq (\widetilde{\sigma} - 1)\widetilde{g}(x)^T p,$$

resulting in

$$\alpha \geq \frac{1-\widetilde{\sigma}}{L} \Big(\frac{-\widetilde{g}(x)^T p}{\|p\|^2}\Big) - \frac{\mathcal{O}(\sqrt{\omega})}{L\|p\|} \geq \frac{(1-\widetilde{\sigma})\Delta_{\mathbf{a}} \|\widetilde{g}(x)\| - \sqrt{\omega}}{L\|p\|} \geq \frac{(1-\widetilde{\sigma})\Delta_{\mathbf{a}} \|\widetilde{g}(x)\|}{2L\|p\|}.$$

(iii) Under the condition (18), the statements (ii) and (iii) result in

$$-\rho \alpha g(x)^{T} p \geq \frac{\rho(1-\widetilde{\sigma})\Delta_{\mathbf{a}} \|\widetilde{g}(x)\|}{2L\|p\|} g(x)^{T} p \geq \frac{\rho(1-\widetilde{\sigma})\Delta_{\mathbf{a}} \|\widetilde{g}(x)\|}{2L\|p\|} \|g(x)\| \|p\| \cos \theta$$

$$\geq \frac{\rho(1-\widetilde{\sigma})\Delta_{\mathbf{a}}^{2}}{4L} \|\widetilde{g}(x)\| \|g(x)\| \geq 4\omega. \tag{25}$$

On the other hand, (18) gives

$$-\rho \alpha g(x)^{T} p \ge \frac{\rho \Delta_{\mathbf{a}}}{2} \alpha \|p\| \|g(x)\| \ge \frac{\rho \Delta_{\mathbf{a}}}{2} \alpha \|p\| \left(\frac{4\widetilde{\rho}\sqrt{\omega}}{\rho \Delta_{\mathbf{a}}}\right) = 2\widetilde{\rho} \alpha \|p\| \sqrt{\omega}. \tag{26}$$

By summing both sides of (25) with (26), it results in

$$-\rho \alpha g(x)^T p \ge 2\omega + \widetilde{\rho} \alpha \|p\| \sqrt{\omega}.$$

From (8) and (A3), we get

$$f(x + \alpha p) \leq f(x) + 2\omega + \widetilde{\rho}\alpha \widetilde{g}(x)^T p \leq f(x) + 2\omega + \widetilde{\rho}\alpha \left(g(x)^T p + ||p||\sqrt{\omega}\right)$$
$$= f(x) + 2\omega + \widetilde{\rho}\alpha ||p||\sqrt{\omega} + \widetilde{\rho}\alpha g(x)^T p$$
$$\leq f(x) + (\widetilde{\rho} - \rho)\alpha g(x)^T p = f(x) + \rho_1 \alpha g(x)^T p;$$

hence (20) holds. From (9) we conclude that

$$g(x + \alpha p)^T p + ||p|| \sqrt{\omega} \ge \widetilde{g}(x + \alpha p)^T p \ge \widetilde{\sigma} \widetilde{g}(x)^T p \ge \widetilde{\sigma} \Big(g(x)^T p - ||p|| \sqrt{\omega}\Big),$$

so that

$$g(x + \alpha p)^T p \ge \widetilde{\sigma}g(x)^T p - (1 + \widetilde{\sigma})\|p\|\sqrt{\omega} \ge (\widetilde{\sigma} + \sigma)g(x)^T p = \sigma_1 g(x)^T p$$

by (18); hence (21) is obtained.

(iv) In the worst case, neither Theorem 1(i) nor Theorem 1(ii) is satisfied. In this case, Theorem 1(iii) holds such that the limit point $\hat{\alpha}$ of $I_j := (a_j, b_j)$ enforces the Wolfe condition (3) and (4) for sufficiently large j. To obtain a bound on the iterations, we use the condition given by interpolation

$$|b_j - a_j| \le (1 - \tau_1)|b_{j-1} - a_{j-1}| \text{ with } \tau_1 \in (0, 1)$$
 (27)

and the condition given by the extrapolation

$$|b_i - a_i| \le (1 - \tau_2)|b_{i-1} - a_{i-1}| \text{ with } \tau_2 \in (0, 1).$$
 (28)

The condition (27) is [1, (16)] obtained from (6) and the condition (28) is [1, (17)] obtained from (7). Denoting $\tau := \min\{1 - \tau_1, 1 - \tau_2\} \in (0, 1)$, we have

$$|b_j - a_j| \le \tau |b_{j-1} - a_{j-1}|. \tag{29}$$

We assume that **AWLS-basic** terminates after K iterations. Applying the inequality (29) recursively, for sufficiently large j we get

$$\alpha_{\min} \le \widehat{\alpha} = |b_K - a_K| \le \tau^{K-1} |b_1 - a_1| \le \tau^{K-1} |\overline{\mu} - a_1| \le \tau^{K-1} |\overline{\mu}|$$

so that the number is bounded by (23) and hence the number of function evaluations is bounded by (24), since 2K + n function evaluations are required. One of the stopping tests (in Lines 6-8 and 14-16 of **AWLS-basic**) requires n function evaluations in the last iteration to estimate the gradient with **FFD**, and two function evaluations in Lines 5 and 13 to compute the function value and directional derivative with **FFDD** in each iteration. Note that since it is assumed that Theorem 1(i) does not hold, $b_1 \leq \overline{\mu} < \infty$. Otherwise, if $b_1 = \infty$, no bracket is found and extrapolation is performed until α_j reaches $\overline{\mu}$. Then $f(\overline{\mu}) \leq \overline{f}$ and **AWLS-basic** ends.

This leads to the complexity bound for SSBBO-basic.

Theorem 2 Suppose that assumptions (A1)-(A3) hold and let f_0 be the initial function value of f. Given $0 < \mu_{\min} < \infty$, define

$$\overline{\mu}_{\min} := \min_{\ell \ge 0} (\mu_{\min}, \overline{\mu}_{\ell}) \quad with \ \overline{\mu}_{\ell} := \overline{\mu}(x_{\ell}) = \frac{\overline{f} - f(x_{\ell})}{\rho_1 g(x_{\ell})^T p_{\ell}}.$$

Here ρ_1 comes from Proposition 2(iii) and \overline{f} comes from Line 2 of Algorithm 4. Then **SSBBO-basic** needs at most $\mathcal{O}(\omega^{-1})$ iterations and $\mathcal{O}((2\overline{K}+n)\omega^{-1})$ function evaluations to find a point x with $||g(x)|| = \mathcal{O}(\sqrt{\omega})$. Here

$$K_{\ell} := K(x_{\ell}) \le \overline{K} := \left\lceil \frac{\log \alpha_{\min} / \overline{\mu}_{\min}}{\log \tau} \right\rceil,$$

 α_{\min} and τ come from Proposition 2, and the factor n comes from the approximate gradient evaluations needed in the last iteration of AWLS-basic.

Proof Let S be the set of iterations generated by **AWLS-basic** such that

$$||g(x_{\ell})|| > \max\left\{\frac{4\widetilde{\rho}}{\rho \Delta_{a}}, \frac{(1+\widetilde{\sigma})}{\sigma \Delta_{a}}\right\} \sqrt{\omega} = \mathcal{O}(\sqrt{\omega}) \text{ for all } \ell \in \mathcal{S};$$

S is not empty by Proposition 1. In other words, all iterations of **SSBBO-basic** decrease the function value; hence, they are successful. Based on Proposition 1, we consider the two following cases: Case 1. Suppose that the statement (i) in Proposition 1 holds. Then, we get

$$f_{\ell+1} := f(x_{\ell} + \alpha_{\ell} p_{\ell}) \leq \overline{f}, \text{ for some } \alpha_{\ell} \in (0, \overline{\mu}_{\ell}];$$

hence

$$f_{\ell} - f_{\ell+1} \ge f_{\ell} - \overline{f} \ge 0. \tag{30}$$

CASE 2. Suppose that the statement (ii) in Proposition 1 holds. Then, by the angle condition (2), we get from Proposition 2

$$f_{\ell} - f_{\ell+1} \ge -\rho_1 \alpha_{\ell} g_{\ell}^T p_{\ell} \ge \frac{4\rho_1 \omega}{\rho}. \tag{31}$$

Then (30) and (31) result in $f_{\ell} - f_{\ell+1} \ge \frac{4\rho_1 \omega}{\rho}$. Summing both sides and using (12) gives

$$f_0 - \widehat{f} \ge \sum_{\ell \in \mathcal{S}} (f_\ell - f_{\ell+1}) \ge |\mathcal{S}| \frac{4\rho_1 \omega}{\rho},$$

leading to $|\mathcal{S}| \leq \frac{\rho(f_0 - \widehat{f})}{4\rho_1\omega}$. Denote by $N_\ell^f = N^f(x_\ell)$ the number of function evaluations in each call to **AWLS-basic** for $\ell \geq 1$. By Proposition 2(iii), the total number of function evaluations by **SSBBO-basic** is

$$N(\omega) = \sum_{\ell \in \mathcal{S}} N_{\ell}^f = \sum_{\ell \in \mathcal{S}} (2K_{\ell} + n) \le (2\overline{K} + n)|\mathcal{S}| = \mathcal{O}\Big((2\overline{K} + n)\omega^{-1}\Big).$$

The order of our bound is the same as that found by Berahas et al. [2].

3 Enhancements

In this section we discuss how new improvements make **SSBBO-basic** very competitive. A limited memory quasi-Newton direction and an update to relevant subspace information are given, leading to an algorithm for computing a subspace step that reduces, if possible, both the value of the model function and its gradient norm. The basic formulation of the subspace direction comes from [21]. We modify the raw subspace direction to enforce that the angle away from 90° is bounded by an approximate gradient estimated by forward finite differences. We discuss an improved Wolfe line search algorithm enriched with a heuristic formula that helps the algorithm to avoid failure. Finally, the **SSBBO** algorithm is explained.

3.1 Decreasing model function value and its gradient norm

In this subsection, we give some details about subspace information, the limited memory quasi-Newton direction, and how to update the subspace information. Then, using a subspace technique, we try to find decreases in the function value and gradient norm of quadratic models. This can be done by solving a linear problem and a one-dimensional bound constrained problem in a cheap way. Our subspace direction is computed by an algorithm called **subspaceDir**. The basic version of this algorithm has already been proposed in [21, Section 2.2], but with the exact gradient, while here the gradient is estimated by **FFD**. Indeed, the new ingredient of **subspaceDir** is to find reductions of the function value and the gradient norm of the quadratic models by new subspace steps.

Let S_{ℓ} be the $n \times m$ matrix whose columns are (in the actual implementation a permutation of) the previous m search directions,

$$S_{\ell} := \{ s_{\ell-m+1}, \cdots, s_{\ell} \} = \{ x_{\ell-m+1} - x_{\ell-m}, \cdots, x_{\ell} - x_{\ell-1} \}, \tag{32}$$

and $Y_{\ell} \in \mathbb{R}^{n \times m}$ be the corresponding gradient differences,

$$Y_{\ell} := \{ y_{\ell-m+1}, \cdots, y_{\ell} \} = \{ \widetilde{g}_{\ell-m+1} - \widetilde{g}_{\ell-m}, \cdots, \widetilde{g}_{\ell} - \widetilde{g}_{\ell-1} \}.$$

$$(33)$$

We first describe the raw subspace direction corresponding to the limited memory quasi-Newton direction described in [21, Section 2.2] but with the gradient estimated by **FFD**. We denote by YY the componentwise squares of Y_{ℓ} , by SS the componentwise squares of S_{ℓ} , by // the componentwise division, and denote $\overline{d} := \sum_{i=1}^{m} \operatorname{YY}_{:i}$ and $\underline{d} := \sum_{i=1}^{m} \operatorname{SS}_{:i}$. Moreover we define the matrices $U_{\ell} := Y_{\ell} - D_{\ell}S_{\ell}$, $D_{\ell} := \operatorname{diag}(d_{\ell})$ with $d_{\ell} := \sqrt{\overline{d}//\underline{d}}$, and $\Sigma_{\ell} := U_{\ell}^{T}S_{\ell}$ (which is symmetric). As in [21], $B_{\ell} := D_{\ell} + U_{\ell}\Sigma_{\ell}^{-1}U_{\ell}^{T}$ is an approximate Hessian matrix, determined by unique S_{ℓ} and Y_{ℓ} and the **quasi-Newton condition**

$$B_{\ell}S_{\ell} = Y_{\ell}. \tag{34}$$

If we define the matrix $H_{\ell} = S_{\ell}^T Y_{\ell}$, we obtain from (34) a further approximation by

$$H_{\ell} = S_{\ell}^T Y_{\ell} = S_{\ell}^T B_{\ell} S_{\ell}. \tag{35}$$

We define the matrix $M_{\ell} := Y_{\ell}^T D_{\ell}^{-1} Y_{\ell} - H_{\ell}$ and obtain the solution of the linear system $M_{\ell} z_{\ell} = U_{\ell}^T D_{\ell}^{-1} \widetilde{g}_{\ell}$. Here \widetilde{g}_{ℓ} is estimated by **FFD**, but is exact in [21]. Then [21, Theorem 1] shows that p_{ℓ} computed by

$$p_{\ell} := D_{\ell}^{-1}(U_{\ell}z_{\ell} - \widetilde{g}_{\ell}), \tag{36}$$

is the solution of $B_{\ell}p_{\ell} = -\widetilde{g}_{\ell}$.

We then update the ingredients of the subspace by **updateSY**. **updateSY** takes the matrices $S_{\ell-1}$, $Y_{\ell-1}$, and $H_{\ell-1}$ and the subspace dimension m as input and uses the tuning parameter m_{\max} as an upper bound for m. It returns the updated matrices S_{ℓ} , Y_{ℓ} , H_{ℓ} and the subspace dimension m.

Algorithm 5 updateSY, update the subspace

```
1: if m < m_{\text{max}} then
```

Update m = m + 1.

3: **else**

Set m=1.

5: end if

6: Replace the mth column of $S_{\ell-1}$ by s_ℓ , resulting it S_ℓ . 7: Replace the mth column of $Y_{\ell-1}$ by y_ℓ , resulting in Y_ℓ .

8: Replace the mth row of H_{ℓ} by $y_{\ell}^T S_{\ell-1}$ and the mth column of H_{ℓ} by $S_{\ell-1}^T y_{\ell}$, resulting H_{ℓ} .

In practice, we don't compute B_{ℓ} , only S_{ℓ} , Y_{ℓ} , and H_{ℓ} are used to compute the subspace directions. **updateSY** enforces that H_{ℓ} becomes symmetric.

Let us describe what is new in subspaceDir. We show how to minimize the maximal norm of the gradient model in the subspace spanned by the columns of the matrix $S_{\ell} \in \mathbb{R}^{n \times m}$, where typically $m \ll n$. We define with $\widetilde{f}_{\ell} := \widetilde{f}(x_{\ell})$ the inexact function value at the current point x_{ℓ} , with $\widetilde{g}_{\ell} := \widetilde{g}(x_{\ell})$ the estimated gradient function at x_{ℓ} , and with

$$c_{\ell} := S_{\ell}^{T} \widetilde{g}_{\ell} \tag{37}$$

the estimated gradient function restricted to S_{ℓ} at x_{ℓ} . Then the quadratic model

$$\widetilde{f}(x_{\ell} + S_{\ell}z) - \widetilde{f}_{\ell} \approx q(z) := c_{\ell}^T z + \frac{1}{2} z^T H_{\ell}z$$

is constructed whose gradient is

$$q'(z) := \partial q(z)/\partial z := c_{\ell} + H_{\ell}z.$$

Here H_{ℓ} has already been defined by (35) and c_{ℓ} by (37). The stationary point is at

$$\widehat{z} := -H_{\ell}^{-1} c_{\ell}. \tag{38}$$

To find a point that reduces not only the model function value q but also, if possible, its gradient norm $\|q\|$, we perform an exact line search on the model function along \hat{z} by solving the simpler univariate problem

$$\min_{\beta} \|q'(\beta \hat{z})\|_{2} = |1 - \beta| \|c_{\ell}\|_{2}
\text{s.t.} \quad q(\beta) = \gamma_{1}\beta + \gamma_{2}\beta^{2} \le 0,$$
(39)

where

$$\gamma_1 := c_\ell^T \widehat{z}, \quad \gamma_2 := \frac{1}{2} \widehat{z}^T H_\ell \widehat{z}. \tag{40}$$

In three situations, we do not want to improve the model gradient norm:

- If γ_1 or γ_2 is contaminated by NaN or $\pm \infty$, the constraint is not satisfied.
- If $\gamma_2 \leq 0$, $q(\beta)$ is flat or unbounded below.

• If $\gamma_2 > 0$ but $\gamma_1 \ge 0$, the minimal q is attained at $\beta = 0$. In these cases, the problem (39) cannot be solved and p_ℓ is computed by (36). In the remaining case, $\gamma_1 < 0 < \gamma_2$ and $\beta^* := -\gamma_1/(2\gamma_2) > 0$. Since $q(\beta) \le 0$ iff $0 \le \beta \le 2\beta^*$, the constraint in (39) is equivalent to $0 \le \beta \le 2\beta^*$. This case is acceptable only if β^* is a real value (otherwise p_ℓ is computed

by (36)). Consequently, (39) reduces to

$$\min |1 - \beta|$$

s.t. $0 \le \beta \le 2\beta^*$

with the solution $\widehat{\beta} := \min(1, 2\beta^*)$. We denote by $\widehat{q} := q(\widehat{\beta})$ the optimal model function value. Indeed, $\widehat{\beta}$ does not necessarily reduce both the function value q and the gradient norm q' of the model in finite precision arithmetic. To check this, we define the computational measure \mathtt{df}_{ℓ} , which is initialized by $\mathtt{df}_0 := \varepsilon_1 |\widetilde{f}_0|$ with $\varepsilon_1 > 0$ (if $\widetilde{f}_0 = 0$ set $\mathtt{df}_0 := 1$) and updated based on the information of the function value

$$\mathrm{df}_{\ell} := \begin{cases} \gamma_f^1(\widetilde{f}_{\ell} - \widetilde{f}_{\ell-1}) & \text{if } \widetilde{f}_{\ell} < \widetilde{f}_{\ell-1} - \mathrm{df}_{\ell-1}, \\ \max\{\gamma_f^2\mathrm{df}_{\ell-1}, \gamma_f^3(|\widetilde{f}_{\ell}| + |\widetilde{f}_{\ell-1}|)\} & \text{otherwise} \end{cases}$$

Here, the parameters $\gamma_f^1 \in (0,1)$, $\gamma_f^2 > 1$, and $\gamma_f^3 \in (0,1)$ are tuning parameters. At the ℓ th iteration, when $\widehat{q} \leq -\mathtt{df}_{\ell}$ we have a good decrease in both the model function value and its gradient norm in the current subspace. Then we can compute the **new subspace direction** by

$$p_{\ell} := \widehat{\beta} S_{\ell} \widehat{z}. \tag{41}$$

To be efficient and robust, **AWLS-basic** must bound the angle between the approximate gradient computed by **FFD** and the search direction computed by **subspaceDir** away from 90°. Inspired by [21, Proposition 1] but using an approximate gradient, we modify the search direction $\bar{p}_{\ell} = p_{\ell}$, which can be obtained from either (36) or (41), if this direction does not satisfy the approximate angle condition (2). In this case, **subspaceDir** computes $\eta_1 = \tilde{g}_{\ell}^T \tilde{g}_{\ell}$, $\eta_2 = \bar{p}_{\ell}^T \bar{p}_{\ell}$, and $\eta = \tilde{g}_{\ell}^T \bar{p}_{\ell}$. Then it finds t satisfying

$$\frac{\eta - t\eta_1}{\sqrt{\eta_1(\eta_2 - 2t\eta + t^2\eta_1)}} \le -\Delta_{\mathbf{a}} \in (0, 1),\tag{42}$$

and recomputes the search direction by $p_{\ell} = \overline{p}_{\ell} - t\widetilde{g}_{\ell}$ such that the approximate angle condition (2) holds; but sometimes it is not possible due to rounding errors. In this case, **subspaceDir** recomputes the search direction by a **diagonal preconditioned steepest descent direction** $p_{\ell} = -D_{\ell}^{-1}\widetilde{g}_{\ell}$.

In the ℓ th iteration, **subspaceDir** takes the matrices S_{ℓ} , Y_{ℓ} , H_{ℓ} and the estimated gradient \widetilde{g}_{ℓ} as input and returns the subspace direction p_{ℓ} as output. It uses the tuning parameter $0 < \Delta_{\rm a} < 1$ for the approximate angle condition.

Algorithm 6 subspaceDir, a subspace direction

```
Obtain the solution of quadratic model in the subspace
```

1: Compute $\widehat{z} = -H_{\ell}^{-1}c_{\ell}$.

Minimize $|1 - \beta|$ subject to $0 \le \beta \le 2\beta^*$ to obtain $\widehat{\beta}$

2: Compute $\gamma_1 = c_\ell^T \widehat{z}, \ \gamma_2 = \frac{1}{2} \widehat{z}^T H_\ell \widehat{z}, \ \beta^* = -\gamma_1/(2\gamma_2) > 0$ and $\widehat{\beta} = \min(1, 2\beta^*)$.

Is a reduction in the function value and gradient norm of the quadratic model?

```
3: if \widehat{q} \leq -df_{\ell} then
                                                                                                                                                                                                                                            \triangleright a reduction is found
             Compute p_{\ell} = \widehat{\beta} S_{\ell} \widehat{z}.
4:
5: else
                Compute the diagonal matrix D_{\ell} := \text{diag}(d_{\ell}) with d_{\ell} = \sqrt{\overline{d}//\underline{d}}, \ \overline{d} = \sum_{i=1}^{m} \mathtt{YY}_{:i}, \ \underline{d} = \sum_{i=1}^{m} \mathtt{SS}_{:i}.
6:
               Compute the matrix U_{\ell} = Y_{\ell} - D_{\ell} S_{\ell}. Compute the matrix M_{\ell} = Y_{\ell}^T D_{\ell}^{-1} Y_{\ell} - H_{\ell}. Solve the linear system M_{\ell} z_{\ell} = U_{\ell}^T D_{\ell}^{-1} \widetilde{g}_{\ell}.
7:
```

8:

9:

Compute the subspace direction by $p_{\ell} = D_{\ell}^{-1}(U_{\ell}z_{\ell} - \widetilde{g}_{\ell})$. 10:

11: **end if**

Enforcing the angle condition

```
12: if \frac{\widetilde{g_\ell^T} p_\ell}{\|\widetilde{g}_\ell\| \|p_\ell\|} \le -\Delta_{\mathbf{a}} < 0 then

13: Set \eta_1 = \widetilde{g}_\ell^T \widetilde{g}_\ell, \eta_2 = p_\ell^T p_\ell, and \eta = \widetilde{g}_\ell^T p_\ell. Then find t satisfying (42).

14: if t is numerically computable then \triangleright but sometimes this is
                                                                                                                                                                \triangleright \Delta_a \in (0,1) is a tuning parameter
                                                                                                              ▷ but sometimes this is not possible due to rounding errors
                       Calculate p_{\ell} = p_{\ell} - t\widetilde{g}_{\ell}.
15:
               else Calculate p_{\ell} = -D_{\ell}^{-1} \widetilde{g}_{\ell}.
16:
17:
                                                                                                            b diagonal preconditioned steepest descent direction
18:
               end if
19: end if
```

3.2 An improved Wolfe line search

The statements (i) and (ii) in Proposition 1 may not be numerically satisfied due to rounding errors. In this case, **AWLS-basic** must be improved. To prevent the failure of **AWLS-basic**, it is proposed to use the heuristic formulas proposed in [21, Section 3.1], but the gradient is estimated by **FFD**. Although the algorithm moves somewhat away from a minimizer in this case, it may approach a minimizer in the next few attempts. At the very least, this idea helps us avoid failure.

An improved version of **AWLS-basic** by Al-Baali & Fletcher [1] is called **AWLS**. It uses the following tuning parameters:

```
\widetilde{\sigma} > \widetilde{\rho} > 0 (parameters for line search), 0 < \Delta_{\alpha} < 1 (tiny parameter for adjusting the heuristic step size), 0 < \varepsilon_h < 1 (parameter for adjusting the finite difference step size h).
```

Algorithm 7 AWLS, an improved approximate Wolfe line search

Perform AWLS-basic in finite precision arithmetic

```
1: Call AWLS-basic to find a step size satisfying the approximate Wolfe line search conditions (8) and (9).
    Find the step size by a heuristic way when AWLS-basic fails in finite precision arithmetic
 2: if AWLS-basic does not enforce the approximate Wolfe conditions then
        \text{Find ind} := \{i \mid p_i \neq 0\}.
        if both x and p are not zero then
 4:
             Compute \alpha := \Delta_{\alpha} |f/\widetilde{g}^T p|.
 5:
 6:
        else if only p is not zero then
             Compute \alpha := \Delta_{\alpha} \max(|\widetilde{f}/\widetilde{g}^T p|, \min\{|x_i/p_i| \mid i \in \mathtt{ind}\}).
 7:
 8:
 9:
10:
         end if
         Compute the new point x + \alpha p and its inexact function value \widetilde{f}(x + \alpha p).
11:
12:
         Estimate the gradient \widetilde{g}(x + \alpha p).
13: end if
```

Our new algorithm is called subspace method for unconstrained black box optimization (SS-BBO). This algorithm takes advantage of the new subspace technique, which attempts to significantly reduce not only the value of the model function, but also its gradient norm. As long as SSBBO cannot reach an ε -approximate stationary point, subspaceDir is used to compute the direction along which AWLS is tried. Then some necessary information about the subspace is updated.

SSBBO takes the initial point x_0 and maximal function evaluations (nfmax) as input and returns the best point x_{best} and its inexact function value $\widetilde{f}_{\text{best}}$ as output. It uses the following tuning parameters: $\varepsilon_1 > 0$ (parameter for adjusting the initial df), $\gamma_f^2 > 1$ (parameter for expanding df), $\gamma_f^1, \gamma_f^3 \in (0,1)$ (parameters for reducing df), m>0 (memory for the subspace), $\Delta_a \in (0,1)$ (tiny parameter for the approximate angle condition (2)), $0 < \Delta_\alpha < 1$ (tiny parameter for adjusting the heuristic step size), $\widetilde{\sigma} > \widetilde{\rho} > 0$ (parameters for line search), $0 < \varepsilon_h < 1$ (parameter for adjusting the finite difference step size h).

Algorithm 8 SSBBO, subspace method for black box optimization

```
Initialization
 1: Compute the initial inexact function value \tilde{f}_0.
 2: Compute \overline{f} = f_0 - 10^8 (1 + |f_0|).
                                                                                                  \triangleright this choice has been suggested by FMINUNC
 3: Estimate the gradient \widetilde{g}_0 by FFD.
 4: Initialize df_0 = \varepsilon_1 |\widetilde{f_0}|.
 5: Set x_{\text{best}} := x_0 and f_{\text{best}} := f_0.
 6: for \ell = 1, 2, 3, \cdots do
            Compute the subspace direction p_{\ell} by subspaceDir.
 7:
 8:
            Find step size \alpha_{\ell}
           Perform AWLS to find \alpha_{\ell} resulting in x_{\ell+1}, \widetilde{f}_{\ell+1}, and \widetilde{g}_{\ell+1}.
 9:
            Stopping test
10:
11:
            if nfmax is reached then
                 Set x_{\text{best}} := x_{\ell+1} and \widetilde{f}_{\text{best}} := \widetilde{f}_{\ell+1}; stop.
12:
13:
            Update the predicated decrease \mathtt{df}_{\ell+1} in the function value
14:
            if \widetilde{f}_{\ell+1} < \widetilde{f}_{\ell} - \mathsf{df}_{\ell} then
15:
                \mathrm{df}_{\ell+1} = \gamma_f^1 (\widetilde{f}_{\ell+1} - \widetilde{f}_{\ell}).
16:
17:
                \mathrm{df}_{\ell+1} = \max\{\gamma_f^2 \mathrm{df}_\ell, \gamma_f^3(|\widetilde{f}_{\ell+1}| + |\widetilde{f}_\ell|).
18:
19:
            Update the ingredients of the subspace
20:
           Set s_{\ell} := x_{\ell+1} - x_{\ell} and y_{\ell} := \widetilde{g}_{\ell+1} - \widetilde{g}_{\ell} and update the subspace by updateSYH.
21:
22: end for
```

4 Numerical results

In this section, we give a comprehensive comparison of **SSBBO** with several robust and fast solvers for all 568 unconstrained problems from the CUTEst [11] collection of test problems for optimization with up to 9000 variables. For all problems, we use the standard initial points given by CUTEst collection.

The number of function evaluations and time in milliseconds (excluding the setup time for the objective function) are denoted by nf and msec, respectively. Let nfmax and secmax be the limit for nf and sec (the time in seconds excluding the setup time for the objective function). We choose

$$\label{eq:nfmax} \begin{split} \text{nfmax} \in \left\{ \begin{aligned} &\{100n, 500n, 1000n\} & \text{ if } 1 \leq n \leq 30, \\ &\{100n, 500n, 1000n\} & \text{ if } 31 \leq n \leq 1000, \\ &\{100n, 500n\} & \text{ if } 1001 \leq n \leq 9000, \end{aligned} \right. \end{split}$$

and

$$\mathtt{secmax} = \begin{cases} 180 \text{ sec} & \text{if } 1 \le n \le 1000, \\ 480 \text{ sec} & \text{if } 1001 \le n \le 9000. \end{cases}$$

Denote by S the list of compared solvers and by P the list of problems. We use the data profile of Moré & Wild [26] and the performance profile of Dolan & Moré [9] to determine the robustness and efficiency of all compared solvers. The data profile

$$\delta_s(\kappa) := \frac{1}{|\mathcal{P}|} \left| \left\{ p \in \mathcal{P} \mid cr_{p,s} := \frac{c_{p,s}}{n_p + 1} \le \kappa \right\} \right| \tag{43}$$

of the solver s is the fraction of problems that the solver s can solve with κ groups of $n_p + 1$ function evaluations. Here n_p is the dimension of the problem p, $c_{p,s}$ is the **cost measure** of the solver s to solve the problem p and $cr_{p,s}$ is the **cost ratio** of the solver s to solve the problem p. The performance profile

$$\rho_s(\tau) := \frac{1}{|\mathcal{P}|} \left| \left\{ p \in \mathcal{P} \mid pr_{p,s} := \frac{c_{p,s}}{\min(c_{p,\overline{s}} \mid \overline{s} \in S)} \le \tau \right\} \right| \tag{44}$$

of the solver s is the fraction of problems that the performance ratio $pr_{p,s}$ is at most τ . Note that $\rho_s(1)$ is the fraction of problems that the solver s wins compared to the other solvers, while $\rho_s(\tau)$ ($\delta_s(\kappa)$) is the fraction of problems for sufficiently large τ (κ) that the solver s can solve.

To compare **SSBBO** with known solvers, two cost measures \mathbf{nf} and \mathbf{msec} are used. The **efficiency** $e_{p,s}$ of the solver s to solve the problem p is the inverse of the performance ratio $pr_{p,s}$. Efficiency measures the ability of a solver $s \in \mathcal{S}$ compared to an ideal solver. In this paper, the number of function evaluations is taken as a suitable cost measure, and the efficiency with respect to this measure is called the \mathbf{nf} efficiency.

To determine the convergence speed of the solver s to reach a minimum of the smooth exact function f, we define the quotients

$$q_s := (f_s - f_{\text{opt}})/(f_0 - f_{\text{opt}})$$
 for $s \in \mathcal{S}$.

Such quotients are not available in real applications. Here, f_s denotes the best function value found by the solver s, f_0 denotes the function value at the starting point (common to all solvers), and $f_{\rm opt}$ denotes the function value at the best point known to the algorithm (in most cases, a global minimizer or at least a better local minimizer) found by performing a sequence of gradient-based and local/global gradient-free solvers; see [20].

We consider a problem solved by the solver s if $q_s \leq 10^{-4}$. Otherwise, the problem is unsolved because either nfmax or secmax is exceeded.

We compared two versions of **SSBBO** with the best solvers for unconstrained optimization problems from the literature, **UOBYQA** of POWELL [30], **BCDFO** of GRATTON et al. [14], **BFO** of PORCELLI & TOINT [29], **NMSMAX** of HIGHAM [16], **NELDER** of KELLEY [17], **VRBBO** of KIMIAEI & NEUMAIER [20], and **VRBBON** of KIMIAEI [18]. They were selected based on the numerical results in [18,20].

SSBBO1 and **SSBBO2** denote **SSBBO** with memory dimensiosn $m = \min(10, n)$ and $m = \min(20, n)$, respectively. The remaining tuning parameters for are chosen for both versions as:

$$\varepsilon_h = \varepsilon_m, \quad \Delta_a = 10^{-8}, \quad \widetilde{\sigma} = 10^{-4}, \quad \widetilde{\rho} = 0.9, \quad \Delta_b = \varepsilon_m, \quad \Delta_\alpha = \varepsilon_m, \quad \gamma_f^1 = 0.5, \\ \gamma_f^2 = 2, \quad \gamma_f^3 = 10^{-12}, \quad \varepsilon_1 = 10^{-8}.$$

• UOBYQA, available at

http://mat.uc.pt/~zhang/software.html,

is an algorithm that forms quadratic models by interpolation by POWELL [30]. Tuning parameters are default.

• BFO, available at

https://github.com/m01marpor/BFO,

is a trainable stochastic derivative-free solver for mixed integer bound-constrained optimization by PORCELLI & TOINT [29]. Tuning parameters are default.

- **BCDFO**, obtained from ANKE TROELTZSCH (personal communication), is a deterministic model-based trust-region algorithm for derivative-free bound-constrained minimization by GRATTON et al. [14]. Tuning parameters are default.
- FMINUNC, available at the Matlab Optimization Toolbox at

https://ch.mathworks.com/help/optim/ug/fminunc.html,

is a deterministic quasi-Newton or trust-region algorithm. $\mathbf{FMINUNC}$ is used with the following options set by $\mathbf{optimoptions}$:

```
opts = optimoptions(@fminunc),'Algorithm','quasi-newton'
'Display', 'Iter','MaxIter',Inf,'MaxFunEvals', limits.nfmax
'TolX', 0,'TolFun',0,'ObjectiveLimit',-1e-50).
```

- FMINUNC1 is FMINUNC using the limited memory BFGS suggested by Liu & NOCEDAL [24] with the memory m=10. Other tuning parameters are default.
- FMINUNC2 is FMINUNC using the limited memory BFGS suggested by Liu & NOCEDAL [24] with the memory m = 20. Other tuning parameters are default.

• VRBBO is an efficient stochastic algorithm by KIMIAEI & NEUMAIER [20]; it can be downloaded from

https://www.mat.univie.ac.at/~neum/software/VRBBO/.

• VRBBON is an efficient stochastic algorithm by Kimiaei [18]; it can be downloaded from

https://www.mat.univie.ac.at/~kimiaei/software/VRBBON.

 \bullet Two versions of Nelder–Mead simplex method for direct search optimization algorithms are **NELDER** by Kelley [17]

https://ctk.math.ncsu.edu/matlab_darts.html

and NMSMAX by HIGHAM, obtained from

http://www.ma.man.ac.uk/~higham/mctoolbox/

For both of them, the default parameters are used.

As described in the introduction, a solver is said to be robust if it has the highest number of solved problems, and efficient if it has the lowest relative cost for function evaluations. A solver is competitive if it is both robust and efficient.

We summarize the numerical results in detail in Subsection 6. **SSBBO** is more robust and efficient than solvers using standard and limited memory quasi-Newton methods not only in low dimensions, but also in medium and high dimensions. In most cases, **SSBBO** is comparable to or even better than well-known model-based solvers in low dimensions in terms of robustness and even efficiency. As a result, **SSBBO** maintains its robustness and efficiency in low to high dimensions compared to the other solvers.

5 Conclusion

An efficient and robust subspace algorithm (SSBBO) for unconstrained black box optimization problems in low to high dimensions was introduced. It was able to achieve a significant reduction in the function value by constructing an efficient direction with limited memory along which a successful prediction of a reduction in the model function and its gradient norm could be achieved.

In the presence of the inexact function value and gradient, a worst-case complexity bound on the number of iterations and function evaluations was found for our algorithm, independent of the choice of search directions, which is consistent with that of Berahas et al. [2].

The numerical results confirm that the new subspace method is more efficient and robust than the standard quasi-Newton method and the traditional limited memory, and even comparable to the model-based methods in terms of the efficiency and robustness.

6 Additional material for SSBBO

6.1 Small scale problems

#100 denotes the total number of test problems in which the solver required the least number of nf, and !100 denotes the total number of test problems in which the solver was the only one to require so many nf. T_{mean} is the mean of the time in seconds taken by a solver to solve the test problems selected from the list of test problems, excluding the time taken for unsolved problems. In all tables,

- the efficiencies are given as percentages, so larger efficiencies indicate better average behaviour and zero efficiency indicates failure.
- all values (against zero) are rounded to whole numbers.

'n' and 't' indicate that $nf \ge nfmax$ and $sec \ge secmax$ have been reached, respectively. If the algorithm failed for other reasons, the sign 'f' is used.

Table 1 compares the summary statistics for $1 < n \le 30$ with the small budget nfmax = 100n:

- For $1 < n \le 30$, **SSBBO1** and **SSBBO2** are more robust than others, solve 163 and 162 out of 192 problems, respectively, and have the second highest **nf** efficiency (48%) while **UOBYQA** has the highest **nf** efficiency (49%).
- For $1 < n \le 2$, BCDFO, SSBBO1, SSBBO2, and NMSMAX are more robust than others, while UOBYQA has the highest nf efficiency (72%).
- For $3 \le n \le 5$, **SSBBO1**, **SSBBO2**, and **UOBYQA** are more robust than others and solve 35 problems out of 39 problems. **UOBYQA**, **BCDFO**, and **SSBBO1** have the highest, second highest, and third highest **nf** efficiency, respectively.
- For $6 \le n \le 10$, **UOBYQA** and **SSBBO1** (**SSBBO2**) are more robust than others, solving 58 out of 76 problems, while **SSBBO1** (**SSBBO2**) has the highest **nf** efficiency (48%).
- For $11 \le n \le 30$, **SSBBO1** and (**SSBBO2**, **FMINUNC**, and **VRBBO**) are more robust than others, solving 40 and 39 out of 43 problems, while **VRBBON**, **SSBBO1**, and **SSBBO2** have the highest, second highest, and third highest **nf** efficiency (54%, 52%, and 51%), respectively.

Table 1: Results for $1 \le n \le 30$ with $\mathtt{nfmax} = 100n$

stopping test: 178 of 192 problem		$q_f \le 0.00$	0.1		≤ 180,	m f	· ≤ 100	140		
		$q_f \leq 0.00$	01,	Sec :	<u> </u>	1	<u> > 100</u>)*II	l mea	n efficiency in %
dim∈(1,30	is sorved					# of	anom	alies	mea	for cost measure
solver		solved	#100	!100	$T_{ m mean}$	#n	#t	#f	nf	msec
SSBB01	ssbb1	163	30	0	79	29	0	0	48	40
SSBBO2	ssbb2	162	33	3	75	30	ŏ	ŏ	48	39
UOBYQA	uob	159	32	22	386	30	ŏ	3	49	47
FMINUNC	func	145	12	8	69	16	Ó	31	39	41
NMSMAX	nmsm	144	4	2	95	47	0	1	26	39
FMINUNC1	func1	140	14	0	51	47	0	5	39	42
FMINUNC2	func2	140	14	0	51	47	0	5	40	42
BCDFO	bcd	135	41	36	2961	0	10	47	40	12
VRBBO	vrbb	134	18	9	150	58	0	0	28	22
VRBBON	vrbbn	131	37	33	268	61	0	0	34	15
BFO	bfo	112	6	4	141	0	0	80	18	11
NELDER	neld	101	4	4	131	82	0	9	15	19
33 of 34 problems dim∈(1,2	without be	unds solve	ea			-# of	onom	olioa		n efficiency in % for cost measure
solver		solved	1 // 100	!100	$T_{ m mean}$		anom		nf	msec
BCDFO	bed	31	#100 10	1100	1 mean 63	#n	#t 0	#f	67	msec 24
SSBBO1	ssbb1	30	10	ő	44	4	0	0	41	19
SSBB01	ssbb1	30	1 1	ő	43	4	ő	ő	41	21
NMSMAX	nmsm	30	3	2	17	4	ő	ő	43	62
UOBYQA	uob	29	12	9	16	5	ő	ŏ	72	69
NELDER	neld	28	2	2	21	3	ŏ	3	34	39
FMINUNC1	func1	28	2	ō	32	6	ŏ	ŏ	43	25
FMINUNC2	func2	28	2	0	33	6	0	0	43	27
VRBBON	vrbbn	27	6	4	71	7	0	0	39	11
BFO	bfo	26	1	0	87	0	0	8	29	6
VRBBO	vrbb	26	3	1	40	8	0	0	36	20
FMINUNC	func	25	1	0	27	5	0	4	39	25
36 of 39 problems	without be	ounds solve	ed			L			mea	n efficiency in %
dim∈[3,5]							anom			for cost measure
solver		solved	#100	!100	Tmean	#n	#t	#f	nf	msec
SSBB01	ssbb1	35 35	4	0	70	4	0	0	51 51	35
SSBBO2	ssbb2	35	4 7	ō	69 49	4	0	0	63	35 65
UOBYQA NMSMAX	nmsm	33	1	5 0	37	3 5	0	1 1	42	66
NELDER	neld	33	1	1	65	5	ő	1	29	37
BCDFO	bed	32	14	12	109	ő	ŏ	7	60	23
FMINUNC	func	32	3	2	51	2	ŏ	5	48	47
FMINUNC1	func1	29	1	0	53	10	o	o	36	33
FMINUNC2	func2	29	1	0	52	10	0	0	36	34
BFO	bfo	26	3	2	113	0	0	13	30	14
VRBBON	vrbbn	26	6	4	158	13	0	0	29	18
VRBBO	vrbb	23	2	1	79	16	0	0		
67 of 76 problems	67 of 76 problems without be				13		0	- 0	22	16
		unds solve	_		13	<u> </u>				n efficiency in %
$\dim \in [6,10]$	without be		ed			# of	anom	alies	mea	16 n efficiency in % for cost measure
$\dim \in [6,10]$ solver		solved	#100	!100	$T_{ m mean}$	# of #n	anom #t	alies #f	mea nf	16 n efficiency in % for cost measure msec
dim∈[6,10] solver SSBBO1	ssbb1	solved 58	#100 19	0	T _{mean}	# of #n 18	anom #t	alies #f	mea nf 48	16 in efficiency in % for cost measure msec 41
dim∈[6,10] solver SSBBO1 SSBBO2	ssbb1 ssbb2	solved 58 58	#100 19 19	0	Tmean 82 84	# of #n 18 18	anom #t 0	alies #f 0	nf 48 48	16 m efficiency in % for cost measure msec 41 40
dim∈[6,10] solver SSBBO1 SSBBO2 UOBYQA	ssbb1 ssbb2 uob	58 58 58 58	#100 19 19 19	0 0 6	Tmean 82 84 147	# of #n 18 18 16	anom #t 0 0	alies #f 0 0	nf 48 48 43	16 m efficiency in % for cost measure msec 41 40 45
dim∈[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO	ssbb1 ssbb2 uob bcd	58 58 58 58 49	#100 19 19 9 15	0 0 6 14	Tmean 82 84 147 2399	# of #n 18 18 16 0	#t 0 0 0 0	alies #f 0 0 2 27	nf 48 48 43 34	16 n efficiency in % for cost measure msec 41 40 45 7
dim∈[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC	ssbb1 ssbb2 uob bcd func	58 58 58 58	#100 19 19 19 15 5	0 0 6 14 4	Tmean 82 84 147 2399 73	# of #n 18 18 16	#t 0 0 0 0 0	alies #f 0 0 2 27 21	nf 48 48 43	16 m efficiency in % for cost measure msec 41 40 45
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX	ssbb1 ssbb2 uob bcd func nmsm	58 58 58 58 49 49	#100 19 19 19 15 5	0 0 6 14	Tmean 82 84 147 2399 73 102	# of #n 18 18 16 0 6	#t 0 0 0 0	alies #f 0 0 2 27	nf 48 48 43 34 34	16 n efficiency in % for cost measure msec 41 40 45 7 42
dim∈[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC	ssbb1 ssbb2 uob bcd func	58 58 58 58 49 49 49	#100 19 19 19 15 5	0 0 6 14 4 0	Tmean 82 84 147 2399 73	# of #n 18 18 16 0 6 28	#t 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0	mea nf 48 48 43 34 34 17	16 n efficiency in % for cost measure msec 41 40 45 7 42 26
dime[6,10] solver SSBB01 SSBB02 UOBYQA BCDF0 FMINUNC NMSMAX FMINUNC1	ssbb1 ssbb2 uob bcd func nmsm func1	58 58 58 49 49 48	#100 19 19 19 15 5 0 4	0 0 6 14 4 0	Tmean 82 84 147 2399 73 102 52	# of #n 18 18 16 0 6 28 25	#t 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4	mea nf 48 48 43 34 34 17 35	16 n efficiency in % for cost measure msec 41 40 45 7 42 26 42
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON	ssbb1 ssbb2 uob bcd func nmsm func1 func2	solved 58 58 58 49 49 49 48 47	#100 19 19 19 9 15 5 0 4 4 9 9	0 0 6 14 4 0 0 0 5	Tmean 82 84 147 2399 73 102 52 53	# of #n 18 18 16 0 6 28 25 25	#t 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4	mea nf 48 48 43 34 34 17 35 35	16 n efficiency in % for cost measure msec 41 40 45 7 7 42 26 42 40
dime[6,10] solver SSBB01 SSBB02 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBO VRBBON BFO	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo	solved 58 58 58 49 49 48 47 47 46 43 38	#100 19 19 19 15 5 0 4 4 9 9	0 0 6 14 4 0 0 0 5 9	Tmean 82 84 147 2399 73 102 52 53 166 354 147	# of #n 18 16 0 6 28 25 25 30 33 0	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 0 38	mea nf 48 48 43 34 17 35 35 23 24 13	16 n efficiency in % for cost measure msec 41 40 45 7 42 26 42 40 22 12
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NELDER	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo neld	solved 58 58 58 49 49 48 47 46 43 38 34	#100 19 19 19 5 0 4 4 9 9	0 0 6 14 4 0 0 0 5	Tmean 82 84 147 2399 73 102 52 53 166 354	# of #n 18 18 16 0 6 28 25 25 30 33	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 0	mea nf 48 48 43 34 17 35 35 23 24 13 7	16 n efficiency in % for cost measure msec 41 40 45 57 42 26 42 40 22 12 14 9
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NBLDER 42 of 43 problems	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo neld	solved 58 58 58 49 49 48 47 46 43 38 34	#100 19 19 19 5 0 4 4 9 9	0 0 6 14 4 0 0 0 5 9	Tmean 82 84 147 2399 73 102 52 53 166 354 147	# of #n 18 18 16 0 6 28 25 30 33 0 37	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 0 38 5	mea nf 48 48 43 34 17 35 35 23 24 13 7	16 n efficiency in % for cost measure msec 41 40 45 7 42 26 40 22 14 14 9 n efficiency in %
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NELDER 42 of 43 problems dime[11,30]	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo neld	solved 58 58 58 58 49 49 48 47 46 43 38 34 sunds solve	#100 19 19 19 15 5 0 4 4 9 9 2 1	0 0 6 14 4 0 0 0 5 9 2	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213	# of #n 18 16 0 6 28 25 25 30 33 0 37	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 0 38 5	mea mf 48 48 48 43 34 37 35 23 24 13 7 mea	16 n efficiency in % for cost measure 411 440 445 7 7 422 26 440 222 12 12 14 9 n efficiency in % for cost measure for cost measure 45 for cost measure 46 for cost measure 47 for cost measure 47 for cost measure 47 for cost measure 47 for cost measure 48 for cost me
dim∈[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NELDER 42 of 43 problems dim∈[11,30] solver	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo neld	solved 58 58 58 49 49 49 47 47 47 46 43 38 34 bunds solved	#100 19 19 19 15 5 0 4 4 9 9 2 1 ed	0 0 6 14 4 0 0 0 5 9 2 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213	# of #n 18 18 16 0 6 28 25 30 33 7 # of #n	anom #t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 38 5 alies #f	mea nf 48 48 43 34 17 35 35 23 13 7 mea	16 n efficiency in % for cost measure msec 41 40 45 7 42 26 42 40 22 11 41 9 n efficiency in % for cost measure msec
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBO VRBBON BFO NELDER 42 of 43 problems dime[11,30] solver SSBBO1	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo neld without bc	solved 58 58 58 49 49 48 47 46 43 38 34 bunds solved	#100 #100 19 19 15 5 0 4 4 9 9 2 1	0 0 6 14 4 0 0 0 5 9 2 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213 Tmean 108	# of #n 18 18 16 0 6 28 25 25 33 0 37 # of #n 3	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 #t anom	alies #f 0 0 2 27 21 0 4 0 0 38 5 alies #f	mea nf 48 48 43 34 17 35 23 24 13 7 mea	16 n efficiency in % for cost measure msec 41 40 45 7 42 26 40 22 12 14 9 n efficiency in % for cost measure msec
dime [6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NELDER 42 of 43 problems dime [11,30] solver SSBBO1 SSBBO1 SSBBO2	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo neld without bc	solved 58 58 58 49 49 48 47 47 46 43 38 34 bunds solved 40 39	#100 19 19 19 15 5 0 4 4 9 9 2 1 ed	0 0 6 14 4 0 0 0 5 9 2 1	Tmean 82 84 147 2399 73 102 53 166 354 147 213 Tmean 108 92	# of #n 18 18 16 0 6 28 25 25 30 37 # of #n 3 4	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 0 38 5 alies #f 0 0	mea nf 48 48 43 34 17 35 35 23 24 13 7 mea	16 n efficiency in % for cost measure when the formula is seen as th
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBO NELOER 42 of 43 problems dime[11,30] solver SSBBO1 SSBBO1 SSBBO2 FMINUNC	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn neld without bc	solved 58 58 58 58 49 49 48 47 46 43 38 34 bunds solved solved 40 39 39	#100 #100 19 19 9 15 5 0 4 4 9 9 2 1 ed	0 0 6 14 4 0 0 0 5 5 9 2 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213 Tmean 108 92 106	# of #n 18 18 18 16 0 6 28 25 25 30 33 7 # of #n 3 4 3	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 0 38 5 alies #f 0 0 1	mea nf 48 48 43 34 17 35 23 24 13 7 mea nf 52 51 37	16 n efficiency in % for cost measure msec 41 40 45 7 42 26 40 22 11 49 n efficiency in % for cost measure msec
dime[6,10] solver SSBB01 SSBB02 UOBYQA BCDF0 FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NELDER 42 of 43 problems dime[11,30] solver SSBB01 SSBB01 SSBB02 FMINUNC VRBBO	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo neld without bc	solved 58 58 58 49 49 48 47 47 46 43 38 34 unds solved 50lved 40 39 39 39	#100 19 19 9 15 0 4 4 4 9 9 2 1 3 d #100 6 9 3 4	0 0 6 14 4 0 0 0 5 9 2 1 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213 Tmean 108 92 106 245	# of #n 18 18 16 0 6 28 25 30 33 7 # of #n 3 4 3 4 3 4	### ### ### ### ### ### ### ### ### ##	alies #f 0 0 2 27 21 0 4 4 0 0 38 5 alies #f 0 0 1	mea nf 48 48 43 34 34 35 35 22 4 13 7 mea nf 52 51 37 35	16 n efficiency in % for cost measure 411 440 445 45 7 422 26 420 420 222 112 9 n efficiency in % for cost measure msec 588 566 49 28
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBO VRBON BFO NELDER 42 of 43 problems dime[11,30] solver SSBBO1 SSBBO1 SSBBO1 SSBBO2 FMINUNC VRBBO UOBYQA	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn sbo neld without bc	solved 58 58 58 58 49 49 48 47 46 43 38 34 bunds solved solved 40 39 39	#100 #100 19 19 9 15 5 0 0 4 4 9 9 2 1 ad #100 6 9 3 4 4 4	0 0 6 14 4 0 0 0 5 9 2 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213 Tmean 108 92 106 245 1371	# of #n 18 18 16 0 6 28 25 25 30 37 # of #n 3 4 4 6	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 38 5 alies #f 0 0 0 0 0	mea nf 48 48 43 34 17 35 23 24 13 7 mea nf 52 51 37	16 n efficiency in % for cost measure msec 41 40 45 7 42 26 40 22 14 9 n efficiency in % for cost measure msec 58 56 49 28 19
dime[6,10] solver SSBB01 SSBB02 UOBYQA BCDF0 FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NELDER 42 of 43 problems dime[11,30] solver SSBB01 SSBB01 SSBB02 FMINUNC VRBBO UOBYQA FMINUNC1	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo neld without bc ssbb1 ssbb2 func vrbb uob func1	solved 58 58 58 49 49 48 47 47 46 43 38 34 bunds solved 40 39 39 37 36	#100 19 19 19 19 5 0 4 4 4 9 9 2 1 #100 6 9 3 4 4 7	0 0 6 14 4 4 0 0 0 5 9 9 2 1 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213 Tmean 108 92 106 245 1371 63	# of	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 0 38 5 alies #f 0 0 1	mea nf 48 48 48 43 34 17 35 23 24 13 7 mea nf 52 51 37 35 24 46	16 n efficiency in % for cost measure when the formula is seen as th
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBO VRBON BFO NELDER 42 of 43 problems dime[11,30] solver SSBBO1 SSBBO1 SSBBO1 SSBBO2 FMINUNC VRBBO UOBYQA	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn sbo neld without bc	58 58 58 49 49 48 47 47 46 43 38 34 50 10 10 10 10 10 10 10 10 10 10 10 10 10	#100 #100 19 19 9 15 5 0 0 4 4 9 9 2 1 ad #100 6 9 3 4 4 4	0 0 6 14 4 0 0 0 5 9 2 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213 Tmean 108 92 106 245 1371	# of #n 18 18 16 0 6 28 25 25 30 37 # of #n 3 4 4 6	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 38 5 alies #f 0 0 0 0 0	mea nf 48 48 43 34 37 35 35 24 13 7 mea nf 52 37 37 35 27	16 n efficiency in % for cost measure 41 40 45 57 42 26 42 40 42 12 12 14 9 n efficiency in % for cost measure msec 58 6 49 28 19 63 64 42 11
dime [6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NELDER 42 of 43 problems dime [11,30] solver SSBBO1 SSBBO1 SSBBO1 SSBBO1 SSBBO0 FMINUNC VRBBO UOBYQA FMINUNC1 FMINUNC1 VRBBO NMSMAX	ssbb1 ssbb2 uob bed func nmsm func1 func2 vrbb vrbbn sbc neld without bc ssbb1 ssbb2 func vrbb uob func1 func2 func	solved 58 58 58 49 49 48 47 47 46 43 38 34 bunds solved solved 39 39 37 36 36 35 33	#100 19 19 9 15 5 0 4 4 4 9 9 2 1 1 ad #100 6 9 3 4 4 7 7 16 0	1100 0 0 6 14 4 0 0 0 0 0 5 5 5 5 5 5 5 5 2 1 1 1 1 1 1 1 1 1 1 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213 Timean 108 92 106 245 1371 63 61 395 212	# of ## of ## of ## of ## of ## of ## of ## of ## of	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 38 5 alies #f 0 0 1 1 0 0 0	mea nf 48 48 43 34 34 37 35 23 24 13 7 mea nf 52 51 37 46 47 54	16 nefficiency in % for cost measure msec 41 40 45 5 42 46 42 46 42 47 49 9 nefficiency in % for cost measure msec 58 56 49 28 19 63 64 42 12 20
dime[6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NELDER 42 of 43 problems dime[11,30] solver SSBBO1 SSBBO2 FMINUNC VRBBO UOBYQA FMINUNC1 FMINUNC1 VRBBO UOBYQA FMINUNC1 FMINUNC1 FMINUNC1 FMINUNC1 VRBBON NMSMAX BCDFO	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb vrbbn bfo neld without bc ssbb1 ssbb2 func vrbb uob func1 func2 vrbb	solved 58 58 58 49 49 48 47 47 46 43 38 34 solved solved solved solved 40 39 39 37 36 36 36 36 35 33 23	#100 19 19 19 9 15 5 0 4 4 4 9 9 2 1 #100 6 6 6 6 6 7 7 7 16 0 2	0 0 6 14 4 4 0 0 0 5 5 9 2 1 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213 Tmean 108 92 106 245 1371 63 61 395 212	# of #m 18 18 16 0 6 288 255 25 30 37 # of #m 3 4 4 6 6 6 8 10 0	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 0 0 38 5 0 0 1 0 0 1 1 0 0 10	mea nf 48 48 48 43 34 17 35 35 24 13 7 mea nf 52 51 37 35 27 46 47 54 15	16 n efficiency in % for cost measure 41 40 45 47 42 26 42 40 22 114 9 n efficiency in % for cost measure 586 499 28 19 63 64 21 20 0
dime [6,10] solver SSBBO1 SSBBO2 UOBYQA BCDFO FMINUNC NMSMAX FMINUNC1 FMINUNC2 VRBBO VRBBON BFO NELDER 42 of 43 problems dime [11,30] solver SSBBO1 SSBBO1 SSBBO1 SSBBO2 FMINUNC VRBBO UOBYQA FMINUNC1 FMINUNC1 VRBBO NMSMAX	ssbb1 ssbb2 uob bcd func nmsm func1 func2 vrbb bfo neld without be ssbb1 ssbb2 func vrbb uob func1 func2 vrbb neld	solved 58 58 58 49 49 48 47 47 46 43 38 34 bunds solved solved 39 39 37 36 36 35 33	#100 19 19 9 15 5 0 4 4 4 9 9 2 1 1 ad #100 6 9 3 4 4 7 7 16 0	1100 0 0 6 14 4 0 0 0 0 0 5 5 5 5 5 5 5 5 2 1 1 1 1 1 1 1 1 1 1 1	Tmean 82 84 147 2399 73 102 52 53 166 354 147 213 Timean 108 92 106 245 1371 63 61 395 212	# of # of #n 18 18 16 0 6 28 25 25 30 37 # of #n 3 4 6 6 6 8 10	#t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	alies #f 0 0 2 27 21 0 4 4 0 38 5 alies #f 0 0 1 1 0 0 0	mea nf 48 48 43 34 34 37 35 23 24 13 7 mea nf 52 51 37 46 47 54	16 n efficiency in % for cost measure when the following state of th

The results obtained for $1 < n \le 30$ with the medium budget $\mathtt{nfmax} = 500n$ are summarized in Table 2:

- For $1 < n \le 30$, **SSBBO1**, **SSBBO2** and **UOBYQA** are more robust than others, solving 175 out of 192 problems, while **UOBYQA**, **SSBBO1** (**SSBBO2**) have the highest and second highest **nf** efficiency (52% and 49%), respectively.
- For $1 < n \le 2$, **UOBYQA** (**BCDFO**) and **SSBBO1** (**SSBBO2**) are more robust than others, solving 34 and 33 out of 34 problems, respectively, and **UOBYQA**, **BCDFO** and **NMSMAX** have the highest, second highest, and third highest **nf** efficiency (77%, 70%, and 46%), respectively.
- For $3 \le n \le 5$, SSBBO1 (SSBBO2) are more robust than others and solve 36 out of 39 problems,

and **UOBYQA**, **BCDFO**, and **SSBBO1** (**SSBBO2**) have the highest, second highest, and third highest **nf** (63%, 60%, and 50%), respectively.

- For $6 \le n \le 10$, **UOBYQA** and **SSBBO1** (**SSBBO2**) are more robust than others, solving 65 and 64 out of 76 problems, respectively, while **SSBBO1** (**SSBBO2**) and **UOBYQA** have the highest and second highest **nf** efficiency (50% and 48%), respectively.
- For $11 \le n \le 30$, **SSBBO1** (**SSBBO2**, **FMINUNC**, and **VRBBO**) is more robust than others and solves 42 out of 43 problems, while **SSBBO1** (**SSBBO2**) and **VRBBON** have the highest and second highest **nf** efficiency (56% and 52%), respectively.

Table 2: Results for $1 < n \le 30$ with nfmax = 500n

Section Sec	Table 2: Results for $1 < n \le 30$ with $nfmax = 500n$										
Section	stopping test:		$q_f \le 0.00$	01,	sec	≤ 180,	nf	€ 500)*n		
SSBBO1	183 of 192 problem	s solved								mea	n efficiency in %
SSBBO1							# of	anom			for cost measure
SSBBO2											
UGBYQA											
VRBBO											
FMINUNC1											
FMINUNC2											
VRBBON											
EMINUNC											
BCDFO							27				
NELDER meld 148											
34 of 34 problems without bounds solved 34 of 34 problems without bounds solved \$\sqrt{\pi 0}\$ \$\sqr											
Solver					4	231	<u> </u>	U	49		
SOLOPTO Solved #100 1100 Tinean #n #t #f nf msec		without be	unds sorve	30			# 01	anom	alies	mea	for cost measure
BCDFO			solved	#100	!100	Tmean				nf	
UOBYQA		bed									
SSBBO1											
SSBBO2		ssbb1									
FMINUNC1	SSBBO2	ssbb2	33	1	0	59	1	0	0	44	21
FMINUNC2											
BFO NELDER neld 31 2 2 38 0 0 0 3 37 38				2							
NELDER				2							
VRBBON											
FMINUNCC											
VRBBO								ŏ			
Solver		vrbb	30	3	1	64	4	0	0	40	19
Solver		without be	unds solve	ed						mea	n efficiency in %
SSBBO1											for cost measure
SSBBO2											
NELDER											
BCDFO											
UOBYQA											
VRBBON											
FMINUNC1								ŏ			17
FMINUNC2	NMSMAX	nmsm	34	1	0	43	3	0		42	68
VRBBO											
FMINUNC Func Solved #100 1100 Timean #n #t #f fn mean efficiency in %											
BFO											
To f 76 problems without bounds solved # 100 1100 Tmean											
						114	 				
Solver		without be	dinds sorve	54		# of	anom	alies			
UOBYQA			solved	#100	!100	$T_{ m mean}$					
SSBBO2	UOBYQA	uob	65		8					48	
VRBBO	SSBBO1					109					
NMSMAX											
BCDFO											
NELDER											
FMINUNC1											
VRBBON vrbbn fame 58 bit 9 solved 9 solved 1100 solved 7 mean 11 mean 11 mean 12 mean 12 mean 12 mean 12 mean 13 mean 14 mean 15 mean 14 mean 15 mean 14 mean 15 mean 15 mean 15 mean 15 mean 15 mean 16 mean 17 mean 17 mean 17 mean 17 mean 17 mean 18 mean							10				
FMINUNC func 53 6 5 149 0 0 23 36 42											
BFO bfo 47 2 2 242 0 0 29 14 17 42 of 43 problems without bounds solved #100 1100 Timen #n #t #f for cost measure solver solved #100 1100 Timen #n #t #f for cost measure SSBBO1 ssbb1 42 6 0 154 1 0 0 52 59 SSBBO2 ssbb2 42 9 3 150 1 0 0 52 57 FMINUNC fune 42 3 2 215 0 0 1 38 49 VRBBO vrbb 42 4 2 399 1 0 0 36 24 UOBYQA uob 41 4 2 2402 0 0 2 28 19 FMINUNC1 func1 40 7 0 153 2 0 1 48 63 FMINUNC2 func2 40 7 0 153 2 0 1 49 66 VRBBON vrbb 40 16 16 1190 3 0 0 56 21 BFO bfo 35 0 0 511 0 0 8 9 11 NMSMAX nmsm 35 0 0 332 8 0 0 15 21 BCDFO bcd 28 2 2 23698 0 12 3 14 0											
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$											
					2	242	1 0	U	29		
Solver		without be	unus solve	eu .			# 01	anom	alies	inea	for cost measure
SSBBO1 ssbb1 42 6 0 154 1 0 0 52 59 SSBBO2 ssbb2 42 9 3 150 1 0 0 52 57 FMINUNC func 42 3 2 215 0 0 1 38 49 VRBBO vrbb 42 4 2 399 1 0 0 36 24 UOBYQA uob 41 4 2 2402 0 0 2 28 19 FMINUNC1 func1 40 7 0 159 2 0 1 48 63 FMINUNC2 func2 40 7 0 153 2 0 1 48 63 VRBBON vrbbn 40 16 16 1190 3 0 0 56 21 BFO bfo 35 0 0<			solved	#100	!100	Tmear				nf	
SSBO2 ssbb2 func 42 func 9 state 3 state 150 state 1 state 0 state 0 state 52 state 49 state 62 state 49 state 63 state 63 state 63 state 63 state 64 state 65 state 6		ssbb1				154					
FMINUNC func 42 3 2 215 0 0 1 38 49 VRBBO vrbb 42 4 2 399 1 0 0 36 24 UOBYQA uob 41 4 2 2402 0 0 2 28 19 FMINUNC1 func1 40 7 0 159 2 0 1 48 63 FMINUNC2 func2 40 7 0 153 2 0 1 49 66 VRBBON vrbbn 40 16 16 1190 3 0 0 56 21 BFO bfo 35 0 0 511 0 8 9 11 NMSMAX nmsm 35 0 0 332 8 0 0 15 21 BCDFO bcd 28 2 2 236988<							1	0	Ō		
UOBYQA uob 41 4 2 2402 0 0 2 28 19 FMINUNC1 func1 40 7 0 159 2 0 1 48 63 FMINUNC2 func2 40 7 0 153 2 0 1 49 66 VRBBON vrbbn 40 16 16 1190 3 0 0 56 21 BFO bfo 35 0 0 511 0 0 8 9 11 NMSMAX nmsm 35 0 0 332 8 0 0 15 21 BCDFO bcd 28 2 2 23698 0 12 3 14 0	FMINUNC		42	3	2	215		0	1	38	
FMINUNC1 func1 40 7 0 159 2 0 1 48 63 FMINUNC2 func2 40 7 0 159 2 0 1 49 66 VRBBON vrbbn 40 16 16 1190 3 0 0 56 21 BFO bfo 35 0 0 511 0 0 8 9 11 NMSMAX nmsm 35 0 0 332 8 0 0 15 21 BCDFO bcd 28 2 2 286 0 12 3 14 0											
FMINUNC2 VRBBON func2 vrbbn 40 7 0 153 2 0 1 49 66 VRBBON vrbbn 40 16 16 1190 3 0 0 56 21 BFO bfo 35 0 0 511 0 0 8 9 11 NMSMAX nmsm 35 0 0 332 8 0 0 15 21 BCDFO bcd 28 2 2 23698 0 12 3 14 0											
VRBBON vrbbn 40 16 16 1190 3 0 0 56 21 BFO bfo 35 0 0 511 0 0 8 9 11 NMSMAX nmsm 35 0 0 332 8 0 0 15 21 BCDFO bcd 28 2 2 23698 0 12 3 14 0											
BFO bfo 35 0 0 511 0 0 8 9 11 NMSMAX nmsm 35 0 0 332 8 0 0 15 21 BCDFO bcd 28 2 2 23698 0 12 3 14 0							3				
NMSMAX nmsm 35 0 0 332 8 0 0 15 21 BCDFO bcd 28 2 2 23698 0 12 3 14 0											
BCDFO bcd 28 2 2 23698 0 12 3 14 0											
NELDER neld 24 0 0 2464 17 0 2 4 4	BCDFO	bcd	28	2	2	23698	0		3	14	0
	NELDER	neld	24	0	0	2464	17	0	2	4	4

Table 3 provides the results for $1 < n \le 30$ with the large budget nfmax = 1000n:

- For $1 < n \le 30$, SSBBO1 (SSBBO2) and UOBYQA are more robust than others, solving 178 and 176 out of 192 problems, respectively, while UOBYQA and SSBBO1 (SSBBO2) have the highest and second highest nf efficiency (52% and 50%), respectively.
- For $1 < n \le 2$, **SSBBO1** (**SSBBO2**, **UOBYQA**, and **BCDFO**) is more robust than others and solves 34 out of 34 problems, while **UOBYQA**, **BCDFO**, and **NMSMAX** have the highest, second highest, and third highest **nf** efficiency (77%, 71%, and 46%), respectively.
- For $3 \le n \le 5$, **SSBBO1** (**SSBBO2**, **FMINUNC1**, and **FMINUNC2**) is more robust than others and solves 36 out of 39 problems, while **UOBYQA**, **BCDFO**, and **SSBBO1** (**SSBBO2**) have the highest, second highest, and third highest nf (63%, 61%, and 51%), respectively.
- For $6 \le n \le 10$, **SSBBO1** (**SSBBO2**), **UOBYQA**, and **VRBBO** are more robust than others, solving 66 out of 76 problems, while **SSBBO1** (**SSBBO2**) and **UOBYQA** have the highest and second highest **nf** efficiency (51% and 48%), respectively.
- For $11 \le n \le 30$, SSBBO1 (SSBBO2, FMINUNC1, VRBBO, and FMINUNC2) is more robust than others and solves 42 out of 43 problems, while SSBBO1 (SSBBO2) and VRBBON have the highest and second highest nf efficiency (56% and 52%), respectively.

Table 3: Results for $1 < n \le 30$ with $\mathtt{nfmax} = 1000n$

)Un									
stopping test: 184 of 192 problem		$q_f \leq 0.000$,,	sec >	≤ 180,	II ni	≤ 100	U*n	I mea	n efficiency in %
dim∈(1,30	ib borred					# 01	anom	alies	11100	for cost measure
solver		solved	#100	!100	$T_{ m mean}$	#n	#t	#f	nf	msec
SSBBO1	ssbb1	178	29	0	127	14	0	0	50	44
SSBBO2	ssbb2	178	32	3	125	14	0	0	50	45
UOBYQA VRBBO	uob vrbb	$\frac{176}{172}$	35 19	26 12	762 299	$\frac{7}{20}$	0	9	52 32	51 23
FMINUNC1	func1	171	14	0	221	20	0	19	41	44
FMINUNC2	func2	170	14	ŏ	207	3	ŏ	19	41	47
VRBBON	vrbbn	166	37	33	676	26	0	0	37	16
NMSMAX	nmsm	162	4	2	214	24	0	6	28	41
FMINUNC BCDFO	func bcd	158 157	13 41	9 36	133 6907	1 0	0 16	33 19	40 43	45 13
NELDER	neld	156	41	4	850	8	0	28	18	22
BFO	bfo	147	6	4	386	ŏ	ŏ	45	20	21
34 of 34 problems	without be	ounds solve	ed						mea	n efficiency in %
dim∈(1,2]							anom			for cost measure
solver SSBBO1	ssbb1	solved	#100	!100	Tmean	#n	#t	#f	nf	msec
SSBBO1	ssbb1	34 34	1 1	0	77 76	0 0	0	0	45 45	27 27
BCDFO	bcd	34	10	8	96	ll ŏ	ŏ	ŏ	71	25
UOBYQA	uob	34	13	10	46	0	0	0	77	77
NMSMAX	nmsm	32	3	2	27	1	0	1	46	68
FMINUNC1 FMINUNC2	func1 func2	32	2 2	0	56 54	1	0	1	45 45	26 33
FMINUNC2 BFO	bfo	32 31	1	0	54 90	0	0	1 3	32 32	33 27
NELDER	neld	31	2	2	35	ő	ŏ	3	37	49
VRBBO	vrbb	31	3	1	96	3	0	0	37	17
VRBBON	vrbbn	31	6	4	102	3 0	0	0	42	16
FMINUNC	func	30	1	0	54	0	0	4	44	37
36 of 39 problems of dim∈[3,5]	without be	ounds solve	ea			# 01	anom	alies	mea	n efficiency in % for cost measure
solver		solved	#100	!100	Tmean	#n	#t	#f	nf	msec
SSBBO1	ssbb1	36	4	0	74	<u>"</u> 3	″o	″o	51	38
SSBBO2	ssbb2	36	4	0	71	3	0	0	51	41
FMINUNC1	func1	36	1	0	135	0	0	3	38	37
FMINUNC2 NELDER	func2 neld	36 35	1 1	0	138 78	$0 \\ 2$	0	3 2	38 30	39 40
BCDFO	bcd	35	14	12	158	l õ	ŏ	4	61	26
UOBYQA	uob	35	7	5	53	2	0	2	63	69
VRBBON	vrbbn	35	6	4	267	4	0	0	32	16
NMSMAX FMINUNC	nmsm func	34 33	1 3	0 2	45 60	3	0	2 5	43 48	62 48
VRBBO	vrbb	33	2	1	148	6	0	0	26	20
BFO	bfo	30	3	2	108	ŏ	ŏ	9	31	31
72 of 76 problems	without be	ounds solve	ed			ii			mea	n efficiency in %
dim∈[6,10]				anom		for cost measure				
solver		solved	#100	!100	Tmean	#n	#t	#f	nf	msec
SSBBO1 SSBBO2	ssbb1 ssbb2	66 66	18 18	0	160 161	10 10	0	0	51 51	45 47
UOBYQA	uob	66	11	9	497	5	ő	5	48	48
VRBBO	vrbb	66	10	8	418	10	0	0	30	25
NMSMAX	nmsm	61	0	0	338	12	0	3	19	29
FMINUNC1 FMINUNC2	func1 func2	61 61	4	0	295 296	1 1	0	14 14	37 37	43 46
BCDFO	bcd	60	15	14	6321	0	2	14	37	8
VRBBON	vrbbn	60	9	9	875	16	0	0	27	14
NELDER	neld	58	1	1	437	0	0	18	10	11
FMINUNC BFO	func bfo	53 48	6 2	5 2	156 287	0	0	23 28	36 14	43 17
42 of 43 problems					201	0	U	20		n efficiency in %
dim∈ 11,30	Dut DC	Janus SOIV	4			# 01	anom	alies	mea	for cost measure
solver		solved	#100	!100	$T_{ m mean}$	#n	#t	#f	nf	msec
SSBBO1	ssbb1	42	6	0	159	1	0	0	52	59
SSBBO2	ssbb2	42	9	3	155	1	0	0	52	59
FMINUNC FMINUNC1	func	42 42	3 7	2 0	218 315	0 0	0	1 1	38 49	52 66
VRBBO	func1 vrbb	42	4	2	315	1	ő	0	36	28
UOBYQA	uob	41	4	2	2387	ō	ő	2	28	19
FMINUNC2	func2	41	7	0	254	1	0	1	49	68
VRBBON BFO	vrbbn	40 38	16	16	1180 971	3	0	0	56	21 13
NMSMAX	bfo nmsm	38	0	0	971 332	0 8	0	5 0	15	13 21
NELDER	neld	32	0	0	3231	6	0	5	4	3
BCDFO	bcd	28	2	2	24869	ő	14	1	14	1
·	•		•	•						•

From performance and data profiles of Figures 1, we conclude that **SSBBO1** (**SSBBO2**) and **UOBYQA** are more robust than others, respectively, while **UOBYQA** and **SSBBO1** (**SSBBO2**) have more efficient than others, respectively.

What is interesting about Tables 1-3 and Figures 1 is that **SSBBO1** and **SSBBO2** are more efficient and robust than **FMINUNC** using the standard BFGS method and **FMINUNC1** and **FMINUNC2** using the limited memory BFGS method in terms of the number of problems solved and the **nf** efficiency, especially at low and medium budgets.

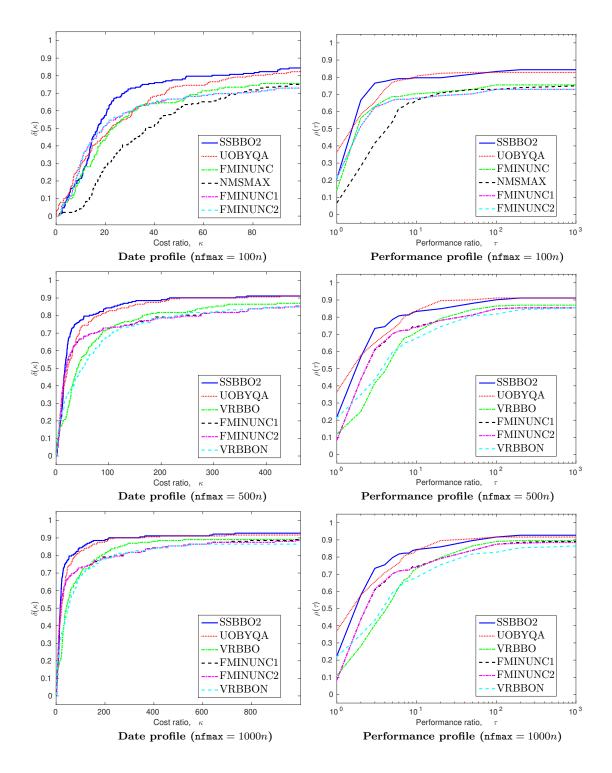


Fig. 1: Results for small scale problems ($n \leq 30$). Left: Data profile $\delta(\kappa)$ in dependence of a bound κ on the cost ratio, see (43). Right: Performance profile $\rho(\tau)$ in dependence of a bound τ on the performance ratio, see (44). Problems solved by no solver are ignored.

6.2 Medium scale problems

Table 4 shows the results for $30 < n \le 1000$ with the small budget nfmax = 100n. **SSBBO2** and **SSBBO1** are more robust than others, solving 212 and 210 out of 267 problems, respectively, while **SSBBO2** and **SSBBO1** have the highest and second highest nf efficiency (58% and 56%), respectively.

Table 4: Results for $30 < n \le 1000$ with $\mathtt{nfmax} = 100n$

stopping test:		$q_f \le 0.00$	01,	sec	≤ 180,	ni	₹ ≤ 10	0*n					
227 of 267 problen	227 of 267 problems without bounds solved								mean efficiency in %				
$\dim \in (30,1000]$	# 01	anom	alies		for cost measure								
solver		solved	#100	!100	Tmean	#n	#t	#f	nf	msec			
SSBBO2	ssbb2	212	82	14	840	54	1	0	58	62			
SSBBO1	ssbb1	209	76	9	809	57	1	0	56	61			
FMINUNC	func	191	37	20	1670	33	0	43	43	33			
VRBBO	vrbb	190	17	15	5080	75	2	0	31	20			
FMINUNC1	func1	188	42	1	1013	68	1	10	46	53			
FMINUNC2	func2	188	43	1	870	69	1	9	46	53			
VRBBON	vrbbn	174	78	71	7737	89	4	0 1	43	14			

Our results, for $30 < n \le 1000$ with the medium budget nfmax = 500n, are shown in Table 5. **SSBBO2** and **SSBBO1** are more robust than others, solving 226 and 223 out of 267 problems, respectively, while **SSBBO2** and **SSBBO1** have the highest and second highest nf efficiency (60% and 59%), respectively.

Table 5: Results for $30 < n \le 1000$ with $\mathtt{nfmax} = 500n$

stopping test:										
244 of 267 problem	ns without	bounds so	lved				n efficiency in %			
$\dim \in (30,1000]$	# 01	anom	alies		for cost measure					
solver	solved	#100	!100	$T_{ m mean}$	#n	#t	#f	nf	msec	
SSBBO2	ssbb2	226	83	16	1679	40	1	0	60	66
SSBBO1	ssbb1	223	77	11	1042	43	1	0	59	66
VRBBO	vrbb	216	22	21	9902	32	19	0	34	19
FMINUNC2	func2	204	37	1	2307	50	1	12	47	56
FMINUNC1	func1	203	36	1	2420	51	1	12	47	55
FMINUNC	func	201	41	24	1997	14	1	51	45	37
VRBBON	vrbbn	192	86	81	7926	45	30	0	45	16

The results, shown in Table 6 for $30 < n \le 1000$ with the large budget nfmax = 1000n, show that **SSBBO2** and **SSBBO1** are more robust than others, solving 234 and 229 out of 267 problems, respectively, while **SSBBO2** and **SSBBO1** have the highest and second highest nf efficiency (62% and 60%), respectively.

Table 6: Results for $30 < n \leq 1000$ with ${\tt nfmax} = 1000n$

stopping test:	g	$q_f \leq 0.000$)1,	sec ≤	180 ,	nf	≤ 100	0*n				
246 of 267 problem	bounds so	mean efficiency in %										
$\dim \in (30,1000]$	# 01	anom	alies		for cost measure							
solver	solved	#100	!100	$T_{ m mean}$	#n	#t	#f	nf	msec			
SSBBO2	ssbb2	234	87	19	3143	25	8	0	62	64		
SSBBO1	ssbb1	229	80	13	2842	32	6	0	60	65		
VRBBO	vrbb	223	21	20	10123	23	21	0	35	20		
FMINUNC1	func1	210	42	1	3758	27	1	29	47	56		
FMINUNC2	func2	208	43	1	3547	38	2	19	48	56		
FMINUNC	func	201	40	23	1953	7	5	54	45	37		
VRBBON	vrbbn	191	80	74	8460	38	38	0	44	15		

From performance and data profiles of Figures 2, we conclude that **SSBBO1** (**SSBBO2**) is more robust and efficient than others, e.g., solvers using standard and limited memory quasi-Newton methods.

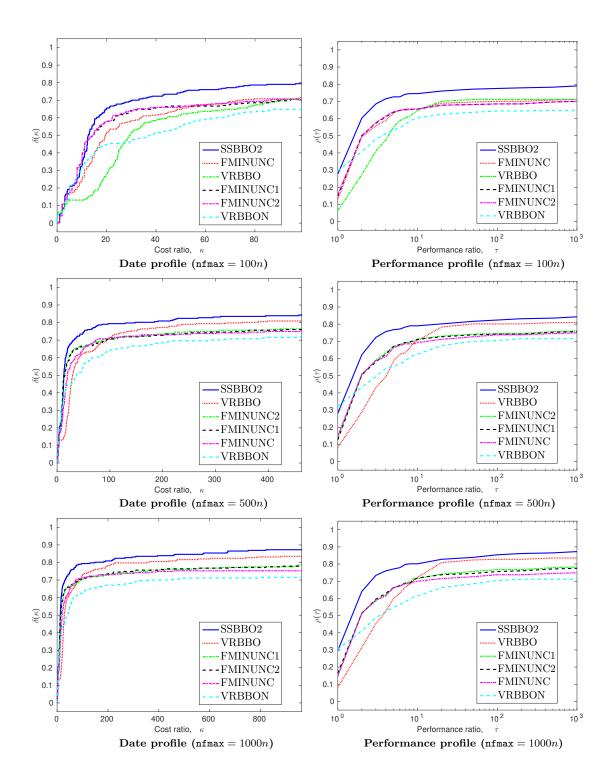


Fig. 2: Results for medium scale problems (30 < $n \le 1000$). Details as in Figure 1.

6.3 Large scale problems

The results, for $1000 < n \le 9000$ with the small budget nfmax = 100n, are summarized in Table 7 with the goal of identifying the competitive solver. **SSBBO2** is more robust than others, solving 87 out of 109 problems, while **VRBBON** has the highest highest nf efficiency.

Table 7: Results for $1000 < n \le 9000$ with $\mathtt{nfmax} = 100n$

99 of 109 proble	99 of 109 problems without bounds solved									mean efficiency in %		
dim∈(1000,9000		# of anomalies			for cost measure							
solver	solved	#100	!100	$T_{ m mean}$	#n	#t	#f	nf	msec			
SSBBO2	ssbb2	87	18	2	33685	17	5	0	48	58		
VRBBO	vrbb	86	5	5	116710	6	17	О	31	21		
SSBBO1	ssbb1	85	17	1	34749	20	4	0	48	58		
FMINUNC1	func1	83	17	1	28944	21	5	0	53	67		
FMINUNC	func	83	15	7	61249	12	5	9	44	35		
FMINUNC2	func2	82	16	О	25069	22	5	0	52	67		
VRBBON	vrbbn	80	59	58	85530	7	22	0	64	24		

The results, for $1000 < n \le 9000$ with the medium budget nfmax = 500n, are summarized in Table 7 with the goal of identifying the competitive solver. **SSBBO2** (**STBBO1**) is more robust than others, solving 90 out of 109 problems, while **VRBBON** has the highest highest nf efficiency.

Table 8: Results for $1000 < n \le 9000$ with nfmax = 500n

					_					
stopping test:		$q_f \leq 0.00$	sec	≤ 180,	ní	² ≤ 500)*n			
99 of 109 problems	ounds solv	mean efficiency in %				n efficiency in %				
$\dim \in (1000,9000]$	# of	anom	alies		for cost measure					
solver	solved	#100	!100	$T_{ m mean}$	#n	#t	#f	nf	msec	
SSBBO1	ssbb1	90	17	1	34796	1	18	0	50	61
SSBBO2	ssbb2	90	18	2	33479	1	18	0	50	60
VRBBO	vrbb	88	4	4	110178	0	21	0	33	22
FMINUNC1	func1	86	19	1	30037	3	20	0	55	67
FMINUNC2	func2	86	18	0	29237	3	20	0	55	68
FMINUNC	func	85	16	7	60847	2	13	9	46	36
VRBBON	vrbbn	80	59	56	81908	0	29	0	64	23

From performance and data profiles of Figures 3, we conclude that **SSBBO2** is more robust than others, while **VRBBON** is more efficient than others. Moreover, **SSBBO2** is robust and efficient than solvers using standard and limited memory quasi-Newton methods.

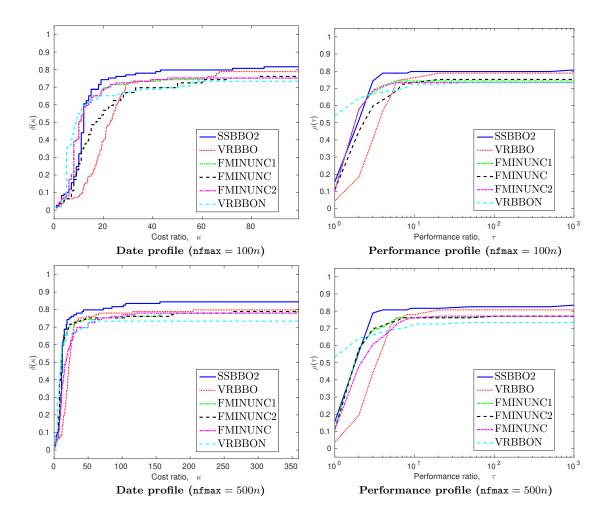


Fig. 3: Results for large scale problems (1000 $< n \leq 9000$). Details as in Figure 1.

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