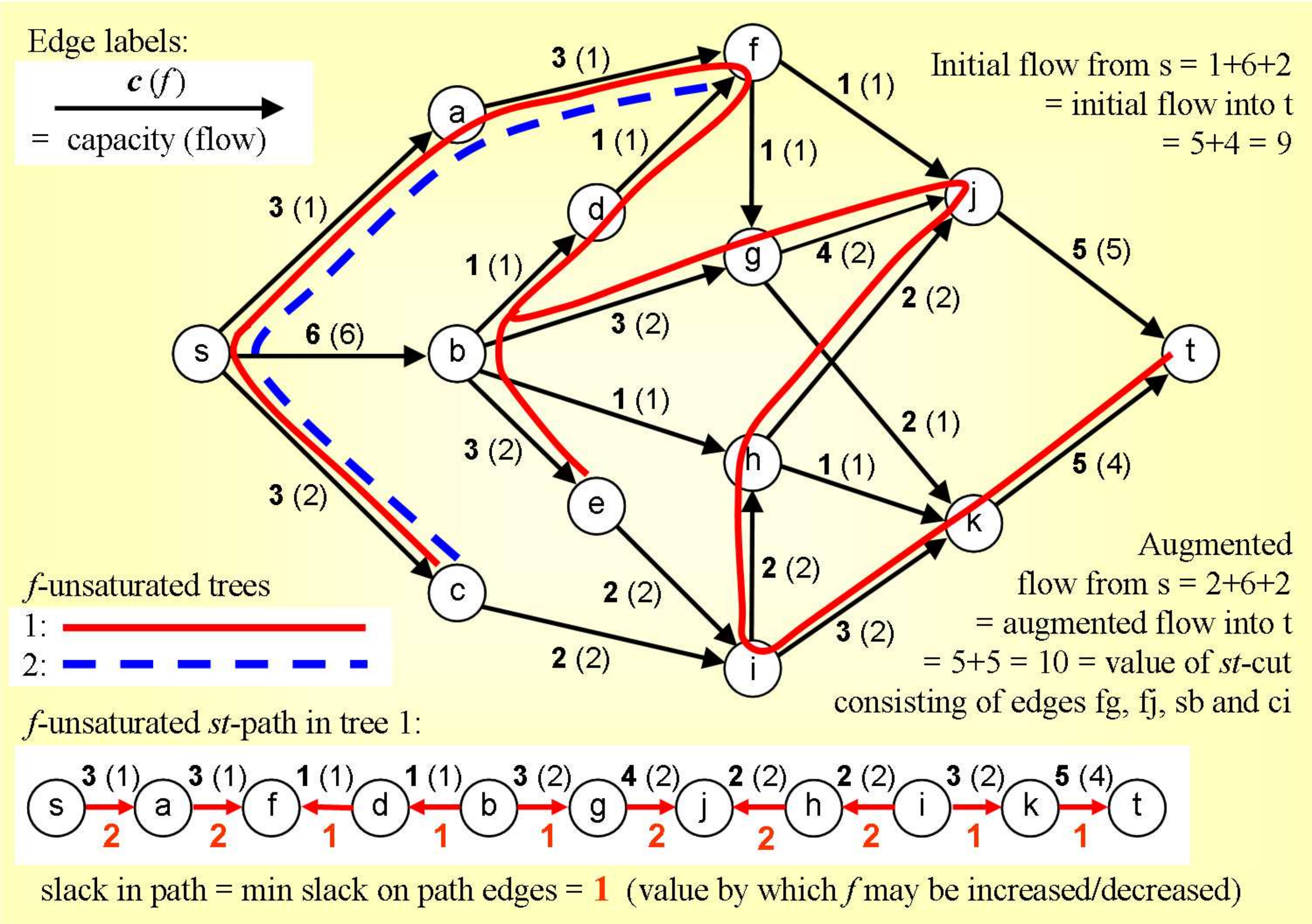
THEOREM OF THE DAY

The Max-Flow Min-Cut Theorem *Let $N = (V, E, s, t)$ be an st-network with vertex set V and edge set E, and with distinguished vertices s and t. Then for any capacity function $c : E \to \mathbb{R}^{\geq 0}$ on the edges of N, the maximum value of an st-flow is equal to the minimum value of an st-cut.*



In the slightly simplified version illustrated here, an *st*-network is a directed graph in which no edges enter s nor exit t. An *st*-flow is a function which, like the capacity function, maps each edge to a nonnegative real number. Additionally it must satisfy:

1. $f(e) \leq c(e)$ for all edges e;
2. the total flow into any vertex $v \neq s, t$ must equal the total flow leaving it.

Under condition (2), the total flow into t will equal the total flow leaving s; this total is called the *value* of the flow. Condition (1) bounds the flow value by the total capacity of any *st*-cut: a set of edges whose removal separates s from t.

On the left, flow value is augmented from 9 to 10 using the Ford-Fulkerson Algorithm: this searches, breadth-first, for an undirected path from s to t in which forward edges have $f(e) < c(e)$ and backward edges have $f(e) > 0$. Along such a so-called f-unsaturated path, f may be increased (decreased) on forward (backward) edges. If the search terminates without reaching t then the set of vertices reached identifies a minimum *st*-cut: and by the theorem, flow has been maximised.

This theorem characterising optimal transportation in capacity constrained networks was published independently in 1956 by: L.R. Ford Jr and D.R. Fulkerson; and by P. Elias, A. Feinstein and C.E. Shannon.

Web link: cse.buffalo.edu/ hungngo/classes/2004/594/notes.html: Week 5; and see homepages.cwi.nl/~lex/files/histtrpclean.pdf for some fascinating prehistory.

Further reading: *Combinatorial Optimization: Algorithms and Complexity* by C. Papadimitriou and K. Steiglitz, Dover Publications, 2000, chapter 6.