ICPC Team Reference

University of Brasilia

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```
set ts=4 sw=4 sta nu rnu sc stl+=%F cindent
set bg=dark ruler clipboard=unnamed,unnamedplus
  timeoutlen=100
imap {<CR> {<CR>}}<Esc>0
nmap <F2> 0V$%d
nmap <C-down> :m+1<CR>
nmap <C-up> :m-2<CR>
vmap < C-c > "+y
nmap <C-a> ggVG
syntax on
alias cmp='g++ -Wall -Wformat=2 -Wshadow -Wconversion -
  fsanitize=address -fsanitize=undefined -fno-sanitize-
  recover -std=c++14'
```

Data Structures

Merge Sort Tree

```
struct MergeTree{
   int n:
   vector<vector<int>> st;
   void build(int p, int L, int R, const int v[]){
          st[p].push_back(v[L]);
          return;
       }
      int mid = (L+R)/2;
      build(2*p, L, mid, v);
      build(2*p+1, mid+1, R, v);
       st[p].resize(R-L+1);
      merge(st[2*p].begin(), st[2*p].end(),
              st[2*p+1].begin(), st[2*p+1].end(),
              st[p].begin());
   }
   int query(int p, int L, int R, int i, int j, int x)
     const{
       if(L > j || R < i) return 0;
       if(L >= i \&\& R <= j){
          int id = lower_bound(st[p].begin(), st[p].end
            (), x) - st[p].begin();
          return int(st[p].size()) - id;
       int mid = (L+R)/2;
       return query(2*p, L, mid, i, j, x) +
          query(2*p+1, mid+1, R, i, j, x);
   }
public:
   MergeTree(int sz, const int v[]): n(sz), st(4*sz){
      build(1, 1, n, v);
   //number of elements >= x on segment [i, j]
   int query(int i, int j, int x) const{
       if(i > j) swap(i, j);
       return query(1, 1, n, i, j, x);
   }
};
Wavelet Tree
template<typename T>
```

```
class wavelet{
   T L, R;
   vector<int> 1;
    vector<T> sum; // <<</pre>
```

```
wavelet *lef, *rig;
   int r(int i) const{ return i - l[i]; }
public:
   template<typename ITER>
   wavelet(ITER bg, ITER en){
       lef = rig = nullptr;
       L = *bg, R = *bg;
       for(auto it = bg; it != en; it++)
          L = min(L, *it), R = max(R, *it);
       if(L == R) return;
       T mid = L + (R - L)/2;
       1.reserve(std::distance(bg, en) + 1);
       sum.reserve(std::distance(bg, en) + 1);
       1.push_back(0), sum.push_back(0);
       for(auto it = bg; it != en; it++)
          1.push_back(1.back() + (*it <= mid)),</pre>
          sum.push_back(sum.back() + *it);
       auto tmp = stable_partition(bg, en, [mid](T x){
          return x <= mid;</pre>
       });
       if(bg != tmp) lef = new wavelet(bg, tmp);
       if(tmp != en) rig = new wavelet(tmp, en);
   }
    ~wavelet(){
       delete lef;
       delete rig;
   }
   // 1 index, first is 1st
   T kth(int i, int j, int k) const{
       if(L >= R) return L;
       int c = l[j] - l[i-1];
       if(c \ge k) return lef \ge kth(l[i-1]+1, l[j], k);
       else return rig->kth(r(i-1)+1, r(j), k - c);
   }
   // # elements > x on [i, j]
   int cnt(int i, int j, T x) const{
       if(L > x) return j - i + 1;
       if(R <= x || L == R) return 0;</pre>
       int ans = 0;
       if(lef) ans += lef->cnt(l[i-1]+1, l[j], x);
       if(rig) ans += rig->cnt(r(i-1)+1, r(j), x);
       return ans;
   }
   // sum of elements <= k on [i, j]</pre>
   T sumk(int i, int j, T k){
       if(L == R) return R <= k ? L * (j - i + 1) : 0;
       if(R <= k) return sum[j] - sum[i-1];</pre>
       int ans = 0;
       if(lef) ans += lef->sumk(l[i-1]+1, l[j], k);
       if(rig) ans += rig->sumk(r(i-1)+1, r(j), k);
       return ans;
   // swap (i, i+1) just need to update "array" l[i]
};
```

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```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
                                                            11 binary_search(const vector<Line> &cht, 11 x){
                                                               int L = 0, R = (int)cht.size()-2;
#include <ext/pb_ds/detail/standard_policies.hpp>
                                                               int bans = (int)cht.size()-1;
                                                               while(L <= R){</pre>
using namespace __gnu_pbds; // or pb_ds;
                                                                   int mid = (L+R)/2;
                                                                   if(cht[mid].eval(x) >= cht[mid+1].eval(x)) // <<<</pre>
template<typename T, typename B = null_type>
                                                                      L = mid + 1;
using oset = tree<T, B, less<T>, rb_tree_tag,
                                                                   else bans = mid, R = mid - 1;
  tree_order_statistics_node_update>;
// find_by_order / order_of_key
                                                               return cht[bans].eval(x);
Hash table
                                                            Convex Hull Trick
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
                                                            * Author: Simon Lindholm
struct custom_hash {
                                                            * source: https://github.com/kth-competitive-
   static uint64_t splitmix64(uint64_t x) {
                                                               programming/kactl/blob/master/content/data-structures
       // http://xorshift.di.unimi.it/splitmix64.c
                                                               /LineContainer.h
                                                             * License: CC0
      x += 0x9e3779b97f4a7c15;
      x = (x \hat{ } (x >> 30)) * 0xbf58476d1ce4e5b9;
      x = (x \hat{ } (x >> 27)) * 0x94d049bb133111eb;
      return x \hat{ } (x >> 31);
                                                            struct Line {
                                                               mutable 11 m, b, p;
   }
                                                               bool operator<(const Line& o) const { return m < o.m</pre>
   size_t operator()(uint64_t x) const {
                                                                 ; }
       static const uint64_t FIXED_RANDOM = chrono::
                                                               bool operator<(ll x) const { return p < x; }</pre>
         steady_clock::now().time_since_epoch().count();
                                                            };
       return splitmix64(x + FIXED_RANDOM);
   }
                                                            struct LineContainer : multiset<Line, less<>> { // CPP14
};
                                                               // (for doubles, use inf = 1/.0, div(a,b) = a/b)
gp_hash_table<long long, int, custom_hash> table;
                                                               const 11 inf = LLONG_MAX;
unordered_map<long long, int, custom_hash> uhash;
                                                               ll div(ll a, ll b) { // floored division
                                                                   return a / b - ((a ^ b) < 0 && a % b); }
uhash.reserve(1 << 15);</pre>
uhash.max_load_factor(0.25);
                                                               bool isect(iterator x, iterator y) {
                                                                   if (y == end()) { x->p = inf; return false; }
Convex Hull Trick Simple
                                                                   if (x->m == y->m) x->p = x->b > y->b? inf: -inf
struct Line{
                                                                   else x->p = div(y->b - x->b, x->m - y->m);
   11 m, b;
                                                                   return x->p >= y->p;
   inline ll eval(ll x) const{
                                                               }
      return x * m + b;
                                                               void add(ll m, ll b) {
   }
                                                                   auto z = insert(\{m, b, 0\}), y = z++, x = y;
};
                                                                   while (isect(y, z)) z = erase(z);
                                                                   if (x != begin() \&\& isect(--x, y)) isect(x, y =
// min => cht.back().m >= L.m
                                                                     erase(y));
// max => cht.back().m <= L.m
                                                                   while ((y = x) != begin() && (--x)->p >= y->p)
void push_line(vector<Line> &cht, Line L){
                                                                      isect(x, erase(y));
 while((int)cht.size() >= 2){
   int sz = (int)cht.size();
                                                               11 query(ll x) {
   if((long double)(L.b-cht[sz-1].b)*(cht[sz-2].m-L.m)
                                                                   assert(!empty());
  <= (long double)(L.b-cht[sz-2].b)*(cht[sz-1].m-L.m)){</pre>
                                                                   auto 1 = *lower_bound(x);
     cht.pop_back();
                                                                   return 1.m * x + 1.b;
   }
                                                               }
   else break;
                                                            };
 cht.push_back(L);
                                                            Min queue
                                                            template<typename T>
// x increasing; pos = 0 in first call
                                                            class minQ{
11 linear_search(const vector<Line> &cht,ll x,int &pos){
                                                               deque<tuple<T, int, int> > p;
   while(pos+1 < (int)cht.size()){</pre>
                                                               T delta:
 *>>*/ if(cht[pos].eval(x) >= cht[pos+1].eval(x)) pos++;
                                                               int sz:
                                                            public:
      else break;
   }
                                                               minQ() : delta(0), sz(0) {}
   return cht[pos].eval(x);
                                                               inline int size() const{ return sz; }
                                                               inline void add(T x){ delta += x; }
}
```

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```
inline void push(T x, int id){
                                                               if(Y[1] > Y[r]) { R[1] = merge(R[1], r); u = 1; }
      x \rightarrow delta, sz++;
                                                                else { L[r] = merge(1, L[r]); u = r; }
       int t = 1:
                                                                calc(u); return u;
      while(p.size() > 0 \& get<0>(p.back()) >= x)
          t += get<1>(p.back()), p.pop_back();
                                                            void init(int n=N-1) { // XXX call before using other
      p.emplace_back(x, t, id);
                                                                for(int i = en = 1; i <= n; i++) { Y[i] = i; sz[i] =</pre>
   }
   inline void pop(){
                                                                   1; L[i] = R[i] = 0; }
                                                                random_shuffle(Y + 1, Y + n + 1);
       get<1>(p.front())--, sz--;
       if(!get<1>(p.front())) p.pop_front();
                                                            void insert(int &u, int it){
   T getmin() const{ return get<0>(p.front())+delta; }
                                                               unlaze(u);
   int getid() const{ return get<2>(p.front()); }
                                                                if(!u) u = it;
                                                                else if(Y[it] > Y[u]) split_val(u, X[it], L[it], R[
                                                                  it]), u = it;
Sparse Table
                                                                else insert(X[it] < X[u] ? L[u] : R[u], it);</pre>
                                                                calc(u);
const int N = 100005;
                                                            void erase(int &u, num key){
int v[N], n;
                                                               unlaze(u);
int dn[N][20];
                                                                if(!u) return;
int fn(int i, int j){
                                                                if(X[u] == key) u = merge(L[u], R[u]);
   if(j == 0) return v[i];
                                                                else erase(key < X[u] ? L[u] : R[u], key);</pre>
   if(~dn[i][j]) return dn[i][j];
                                                                calc(u):
   return dn[i][j] = min(fn(i, j-1), fn(i + (1 << (j-1)))
     ), j-1));
                                                            int create_node(num key){
}
                                                               X[en] = key;
                                                                sz[en] = 1;
int lg(int x){ return 31 - __builtin_clz(x); }
                                                               L[en] = R[en] = 0;
                                                                return en++;
int getmn(int 1, int r) { // [1, r]
   int lz = lg(r - l + 1);
                                                            int query(int u, int 1, int r){//0 index
   return min(fn(1, lz), fn(r - (1 << lz) + 1, lz));
                                                                unlaze(u);
                                                                if(u! or r < 0 or 1 >= sz[u]) return
Treap
                                                                  identity_element;
                                                                if(1 <= 0 and r >= sz[u] - 1) return subt_data[u];
// source: https://github.com/victorsenam/caderno/blob/
                                                                int ans = query(L[u], 1, r);
 master/code/treap.cpp
                                                                if(1 \le sz[L[u]] and sz[L[u]] \le r)
//const int N = ; typedef int num;
                                                                   ans = max(ans, st[u]);
num X[N]; int en = 1, Y[N], sz[N], L[N], R[N];
                                                                ans = max(ans, query(R[u], l-sz[L[u]]-1, r-sz[L[u]])
void calc (int u) { // update node given children info
                                                                  ]]-1));
   if(!u) return;
                                                                return ans;
   sz[u] = sz[L[u]] + 1 + sz[R[u]];
   // code here, no recursion
                                                            ColorUpdate
void unlaze (int u) {
   if(!u) return;
                                                            // source: https://github.com/tfg50/Competitive-
   // code here, no recursion
                                                              Programming/tree/master/Biblioteca/Data%20Structures
void split_val(int u, num x, int &l, int &r) { // l gets
                                                            #include <set>
                                                            #include <vector>
   <= x, r gets > x
   unlaze(u); if(!u) return (void) (1 = r = 0);
                                                            template <class Info = int>
   if(X[u] <= x) { split_val(R[u], x, 1, r); R[u] = 1;</pre>
     1 = u; }
                                                            class ColorUpdate {
   else { split_val(L[u], x, 1, r); L[u] = r; r = u; }
                                                            public:
   calc(u);
                                                                struct Range {
                                                                   Range(int l = 0) { this->l = 1; }
void split_sz(int u, int s, int &l, int &r) { // l gets
                                                                   Range(int 1, int r, Info v) {
  first s, r gets remaining
                                                                      this \rightarrow l = 1;
   unlaze(u); if(!u) return (void) (1 = r = 0);
                                                                       this->r = r;
   if(sz[L[u]] < s)  { split_sz(R[u], s - sz[L[u]] - 1,
                                                                       this -> v = v;
     1, r); R[u] = 1; 1 = u; }
                                                                   }
   else { split_sz(L[u], s, l, r); L[u] = r; r = u; }
                                                                   int 1, r;
   calc(u);
                                                                   Info v;
int merge(int 1, int r) { // els on l <= els on r</pre>
                                                                   bool operator < (const Range &b) const { return l</pre>
                                                                      < b.1; }
   unlaze(1); unlaze(r); if(!1 || !r) return 1 + r; int
                                                                };
```

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```
std::vector<Range> upd(int 1, int r, Info v) {
                                                           // subtree of u => [ in[u], out[u] )
       std::vector<Range> ans;
                                                           // path from nxt[u] to u => [ in[ nxt[u] ], in[u] ]
       if(1 >= r) return ans;
                                                           Iterative Segtree
       auto it = ranges.lower_bound(1);
       if(it != ranges.begin()) {
                                                           T query(int 1, int r, int &pos){ // [1, r]
          it--;
                                                              T rl, rr;
          if(it->r>1) {
                                                               for (1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1)
              auto cur = *it;
                                                                  if(l & 1) rl = merge(rl, st[l++]);
              ranges.erase(it);
                                                                  if(r & 1) rr = merge(st[--r], rr);
              ranges.insert(Range(cur.1, 1, cur.v));
              ranges.insert(Range(1, cur.r, cur.v));
                                                              return merge(rl, rr);
      it = ranges.lower_bound(r);
                                                           // initially save v[i] in st[n+i] for all i in [0, n)
       if(it != ranges.begin()) {
                                                           void build(){
          it--:
                                                               for(int p = n-1; p > 0; p--)
          if(it->r > r) {
                                                                  st[p] = merge(st[2*p], st[2*p+1]);
              auto cur = *it;
                                                           }
              ranges.erase(it);
              ranges.insert(Range(cur.1, r, cur.v));
                                                           void update(int p, T val){
              ranges.insert(Range(r, cur.r, cur.v));
                                                               st[p += n] = val;
                                                               while(p >>= 1) st[p] = merge(st[2*p], st[2*p+1]);
       for(it = ranges.lower_bound(1); it != ranges.end
                                                           LiChao's Segtree
         () && it->l < r; it++) {
          ans.push_back(*it);
                                                           void add_line(line nw, int v = 1, int l = 0, int r =
                                                             maxn) { // [1, r)}
      ranges.erase(ranges.lower_bound(1), ranges.
                                                               int m = (1 + r) / 2;
         lower_bound(r));
                                                               bool lef = nw.eval(1) < st[v].eval(1);</pre>
       ranges.insert(Range(1, r, v));
                                                               bool mid = nw.eval(m) < st[v].eval(m);</pre>
       return ans;
                                                               if(mid) swap(st[v], nw);
   }
                                                               if(r - 1 == 1) {
private:
                                                                  return;
   std::set<Range> ranges;
                                                               } else if(lef != mid) {
                                                                  add_line(nw, 2 * v, 1, m);
Heavy Light Decomposition
                                                               } else {
                                                                  add_line(nw, 2 * v + 1, m, r);
void dfs_sz(int u){
                                                           }
   sz[u] = 1;
                                                           int get(int x, int v = 1, int l = 0, int r = maxn) {
   for(auto &v : g[u]) if(v == p[u]){
                                                               int m = (1 + r) / 2;
       swap(v, g[u].back());
                                                               if(r - 1 == 1) {
       g[u].pop_back();
                                                                  return st[v].eval(x);
       break;
                                                               else if(x < m) {
                                                                  return min(st[v].eval(x), get(x, 2*v, 1, m));
   for(auto &v : g[u]){
                                                                  return min(st[v].eval(x), get(x, 2*v+1, m, r));
      p[v] = u;
      dfs_sz(v);
       sz[u] += sz[v];
      if(sz[v] > sz[ g[u][0] ])
                                                           Palindromic tree
          swap(v, g[u][0]);
                                                           #include <bits/stdc++.h>
   }
}
                                                           using namespace std;
// nxt[u] = start of path with u
// set nxt[root] = root beforehand
                                                           const int maxn = 3e5 + 1, sigma = 26;
void dfs_hld(int u){
                                                           int len[maxn], link[maxn], to[maxn][sigma];
                                                           int slink[maxn], diff[maxn], series_ans[maxn];
   in[u] = t++;
   rin[in[u]] = u;
                                                           int sz, last, n;
                                                           char s[maxn];
   for(auto v : g[u]){
      nxt[v] = (v == g[u][0] ? nxt[u] : v);
      dfs_hld(v);
                                                           void init()
                                                           {
   out[u] = t;
                                                               s[n++] = -1;
                                                               link[0] = 1;
}
```

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```
len[1] = -1;
   sz = 2;
}
int get_link(int v)
   while(s[n - len[v] - 2] != s[n - 1]) v = link[v];
   return v;
void add_letter(char c)
   s[n++] = c -= 'a';
   last = get_link(last);
   if(!to[last][c])
       len[sz] = len[last] + 2;
       link[sz] = to[get_link(link[last])][c];
       diff[sz] = len[sz] - len[link[sz]];
       if(diff[sz] == diff[link[sz]])
          slink[sz] = slink[link[sz]];
          slink[sz] = link[sz];
       to[last][c] = sz++;
   }
   last = to[last][c];
}
int main()
{
   ios::sync_with_stdio(0);
   cin.tie(0);
   init();
   string s;
   cin >> s;
   int n = s.size();
   int ans[n + 1];
   memset(ans, 63, sizeof(ans));
   ans [0] = 0;
   for(int i = 1; i <= n; i++)</pre>
   {
       add_letter(s[i - 1]);
       for(int v = last; len[v] > 0; v = slink[v])
          series_ans[v] = ans[i - (len[slink[v]] + diff
            [v])];
          if(diff[v] == diff[link[v]])
              series_ans[v] = min(series_ans[v],
                series_ans[link[v]]);
          ans[i] = min(ans[i], series_ans[v] + 1);
       cout << ans[i] << "\n";</pre>
   return 0;
```

Math

Extended Euclidean Algorithm

```
// a*x + b*y = gcd(a, b), < gcd, x, y>
tuple<int, int, int> gcd(int a, int b) {
   if(b == 0) return make_tuple(a, 1, 0);
   int q, w, e;
   tie(q, w, e) = gcd(b, a % b);
   return make_tuple(q, e, w - e * (a / b));
}
```

Chinese Remainder Theorem

```
// x = vet[i].first (mod vet[i].second)
11 crt(vector<pair<11, 11>> vet){
   11 ans = vet[0].first, lcm = vet[0].second;
   11 a, b, g, x, y;
   for(int i = 1; i < (int)vet.size(); i++){</pre>
       tie(a, b) = vet[i];
       tie(g, x, y) = gcd(lcm, b);
       ans = ans + x * (a - ans) / g % (b / g) * lcm;
       lcm = lcm * b / g;
       ans = (ans \% lcm + lcm) \% lcm;
   return ans;
Preffix inverse
inv[1] = 1;
for(int i = 2; i < p; i++)</pre>
   inv[i] = (p - (p/i) * inv[p%i] % p) % p;
Pollard Rho
11 rho(11 n){
   if(n \% 2 == 0) return 2;
   ll d, c, x, y;
   do{
       c = 11rand() \% n, x = 11rand() \% n, y = x;
          x = add(mul(x, x, n), c, n);
          y = add(mul(y, y, n), c, n);
          y = add(mul(y, y, n), c, n);
           d = \_gcd(abs(x - y), n);
       }while(d == 1);
   }while(d == n);
   return d;
}
11 pollard_rho(11 n){
   ll x, c, y, d, k;
   int i;
   do{
       i = 1;
       x = 11rand() % n, c = 11rand() % n;
       y = x, k = 4;
       do{
          if(++i == k) y = x, k *= 2;
          x = add(mul(x, x, n), c, n);
          d = \_gcd(abs(x - y), n);
       }while(d == 1);
   }while(d == n);
   return d;
void factorize(ll val, map<ll, int> &fac){
   if(rabin(val)) fac[ val ]++;
   else{
       11 d = pollard_rho(val);
       factorize(d, fac);
       factorize(val / d, fac);
   }
}
map<ll, int> factor(ll val){
```

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```
map<ll, int> fac;
   if(val > 1) factorize(val, fac);
                                                               for(11 m = 2; ; m++) if(\_gcd(m, n) == 1)
   return fac;
                                                                   if(check(m)) return m;
                                                           }
Miller Rabin
                                                            // Let's denote R(n) as the set of primitive roots
                                                              modulo n, p is prime
bool rabin(ll n){
                                                            // g \ln R(p) => (pow(g, p-1, p * p) == 1 ? g+p : g) \ln r
   if(n <= 1) return 0;</pre>
                                                               R(pow(p, k)), for all k > 1
   if(n <= 3) return 1;
                                                            // g \text{ in } R(pow(p, k)) \Rightarrow (g \% 2 == 1 ? g : g + pow(p, k))
   11 s = 0, d = n - 1;
                                                               while(d % 2 == 0) d /= 2, s++;
   for(int k = 0; k < 64; k++){
                                                            Mobius Function
      11 a = (11rand() \% (n - 3)) + 2;
      11 x = fexp(a, d, n);
                                                            memset(mu, 0, sizeof mu);
       if(x != 1 \&\& x != n-1){
                                                            mu[1] = 1;
          for(int r = 1; r < s; r++){
                                                            for(int i = 1; i < N; i++)</pre>
              x = mul(x, x, n);
                                                               for(int j = i + i; j < N; j += i)
              if(x == 1) return 0;
                                                                  mu[j] -= mu[i];
              if(x == n-1) break;
                                                            // g(n) = sum{f(d)} => f(n) = sum{mu(d)*g(n/d)}
                                                            Mulmod TOP
          if(x != n-1) return 0;
                                                            constexpr uint64_t mod = (1ull<<61) - 1;</pre>
   }
                                                            uint64_t modmul(uint64_t a, uint64_t b){
   return 1;
                                                               uint64_t 11 = (uint32_t)a, h1 = a>>32, 12 = (
                                                                 uint32_t)b, h2 = b>>32;
Totiente
                                                               uint64_t 1 = 11*12, m = 11*h2 + 12*h1, h = h1*h2;
                                                               uint64_t ret = (1&mod) + (1>>61) + (h << 3) + (m >>
11 totiente(11 n){
                                                                 29) + (m \ll 35 \gg 3) + 1;
   11 \text{ ans} = n;
                                                               ret = (ret & mod) + (ret>>61);
   for(ll i = 2; i*i <= n; i++){</pre>
                                                               ret = (ret & mod) + (ret>>61);
       if(n \% i == 0){
                                                               return ret-1;
          ans = ans / i * (i - 1);
          while(n \% i == 0) n /= i;
                                                            Matrix Determinant
   }
                                                            int n:
                                                            long double a[n][n];
   if(n > 1) ans = ans / n * (n - 1);
   return ans:
                                                            long double gauss(){
                                                               long double det = 1;
Primitive root
                                                               for(int i = 0; i < n; i++){
                                                                   int q = i;
// a primitive root modulo n is any number g such that
                                                                   for(int j = i+1; j < n; j++){
  any c coprime to n is congruent to a power of g modulo
                                                                      if(abs(a[j][i]) > abs(a[q][i]))
  n.
                                                                          q = j;
bool exists_root(ll n){
                                                                   if(abs(a[q][i]) < EPS){
   if(n == 1 || n == 2 || n == 4) return true;
                                                                      det = 0;
   if(n \% 2 == 0) n /= 2;
                                                                      break;
   if(n % 2 == 0) return false;
                                                                   if(i != q){
   // test if n is a power of only one prime
                                                                      for(int w = 0; w < n; w++)
   for(11 i = 3; i * i <= n; i += 2) if(n % i == 0){
                                                                          swap(a[i][w], a[q][w]);
      while(n \% i == 0) n /= i;
                                                                      det = -det;
      return n == 1;
                                                                   det *= a[i][i];
   return true;
                                                                   for(int j = i+1; j < n; j++) a[i][j] /= a[i][i];</pre>
                                                                   for(int j = 0; j < n; j++) if(j != i){
11 primitive_root(ll n){
                                                                      if(abs(a[j][i]) > EPS)
   if(n == 1 || n == 2 || n == 4) return n - 1;
                                                                          for(int k = i+1; k < n; k++)
   if(not exists_root(n)) return -1;
                                                                              a[j][k] -= a[i][k] * a[j][i];
   11 x = phi(n);
                                                                   }
   auto pr = factorize(x);
                                                               }
   auto check = [x, n, pr](ll m){
       for(ll p : pr) if(fexp(m, x / p, n) == 1)
                                                               return det;
          return false;
                                                           }
       return true;
```

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Simplex Method

```
typedef long double dbl;
const dbl eps = 1e-6;
const int N = , M = ;
mt19937 rng(chrono::steady_clock::now().time_since_epoch
  ().count()):
struct simplex {
   int X[N], Y[M];
   dbl A[M][N], b[M], c[N];
   dbl ans;
   int n, m;
   dbl sol[N];
   void pivot(int x, int y){
       swap(X[y], Y[x]);
      b[x] /= A[x][y];
       for(int i = 0; i < n; i++)
          if(i != y)
              A[x][i] /= A[x][y];
      A[x][y] = 1. / A[x][y];
       for(int i = 0; i < m; i++)</pre>
          if(i != x && abs(A[i][y]) > eps) {
              b[i] -= A[i][y] * b[x];
              for(int j = 0; j < n; j++) if(j != y)
                 A[i][j] -= A[i][y] * A[x][j];
              A[i][y] = -A[i][y] * A[x][y];
          }
       ans += c[y] * b[x];
       for(int i = 0; i < n; i++)</pre>
          if(i != y)
              c[i] -= c[y] * A[x][i];
      c[y] = -c[y] * A[x][y];
   }
   // maximiza sum(x[i] * c[i])
   // sujeito a
   // sum(a[i][j] * x[j]) <= b[i] para 0 <= i < m (Ax)
   // x[i] >= 0 para 0 <= i < n (x >= 0)
   // (n variaveis, m restricoes)
   // guarda a resposta em ans e retorna o valor otimo
   dbl solve(int _n, int _m) {
       this->n = _n; this->m = _m;
       for(int i = 1; i < m; i++){
          int id = uniform_int_distribution<int>(0, i)(
            rng);
          swap(b[i], b[id]);
          for(int j = 0; j < n; j++)
              swap(A[i][j], A[id][j]);
      }
       ans = 0.;
       for(int i = 0; i < n; i++) X[i] = i;
       for(int i = 0; i < m; i++) Y[i] = i + n;
          int x = min_element(b, b + m) - b;
          if(b[x] >= -eps)
             break:
          int y = find_if(A[x], A[x] + n, [](dbl d) {
            return d < -eps; }) - A[x];</pre>
          if(y == n) throw 1; // no solution
          pivot(x, y);
      while(true) {
          int y = max_element(c, c + n) - c;
```

```
if(c[y] <= eps) break;</pre>
          int x = -1;
          dbl mn = 1. / 0.;
          for(int i = 0; i < m; i++)
              if(A[i][y] > eps && b[i] / A[i][y] < mn)</pre>
                 mn = b[i] / A[i][y], x = i;
          if(x == -1) throw 2; // unbounded
          pivot(x, y);
       memset(sol, 0, sizeof(dbl) * n);
       for(int i = 0; i < m; i++)
          if(Y[i] < n)
              sol[Y[i]] = b[i];
       return ans;
};
FFT
struct base{
   double r, i;
   base(double _r = 0, double _i = 0) : r(_r), i(_i) {}
   base operator*(base &o) const{
       return {r*o.r - i*o.i, r*o.i + o.r*i};
   double real() const{ return r; }
   void operator*=(const base &o){
       (*this) = \{r*o.r-i*o.i, r*o.i+o.r*i\};
   }
   void operator+=(const base &o){r += o.r, i += o.i; }
   void operator/=(const double &o){ r /= o, i /= o; }
   void operator==(const base &o){r == o.r, i == o.i; }
   base operator+(const base &o){return {r+o.r,i+o.i};}
   base operator-(const base &o){return {r-o.r,i-o.i};}
};
double PI = acos(-1);
void fft(vector<base> &a, bool inv){
   int n = (int)a.size();
   for(int i = 1, j = 0; i < n; i++){
       int bit = n \gg 1;
       for(; j >= bit; bit >>= 1) j -= bit;
       j += bit;
       if(i < j) swap(a[i], a[j]);
   for(int sz = 2; sz <= n; sz <<= 1) {
       double ang = 2*PI/sz * (inv ? -1 : 1);
       base wlen(cos(ang), sin(ang));
       for(int i = 0; i < n; i += sz){
          base w(1);
          for(int j = 0; j < sz/2; j++){
              base u = a[i+j], v = a[i+j+sz/2] * w;
              a[i+j] = u + v;
              a[i+j+sz/2] = u - v;
              w *= wlen;
          }
       }
   if(inv) for(int i = 0; i < n; i++) a[i] /= 1.0 * n;
}
void multiply(const vector<int> &a, const vector<int> &b
  , vector<int> &res){
   vector<base> fa(a.begin(), a.end());
   vector<base> fb(b.begin(), b.end());
```

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```
size_t n = 1;
                                                               void run(double mat[][MAXC], int R, int C) {
   while(n < a.size()) n <<= 1;</pre>
                                                                 REP(i, C) row[i] = -1;
   while(n < b.size()) n <<= 1;</pre>
   n <<= 1:
                                                                 int r = 0;
   fa.resize(n), fb.resize(n);
                                                                 REP(c, C) {
                                                                   int k = r:
   fft(fa, false), fft(fb, false);
                                                                   FOR(i, r, R) if(fabs(mat[i][c]) > fabs(mat[k][c]))
   for(size_t i = 0; i < n; i++)</pre>
                                                                     k = i;
       fa[i] *= fb[i];
                                                                   if(fabs(mat[k][c]) < eps) continue;</pre>
   fft(fa, true);
                                                                   REP(j, C+1) swap(mat[r][j], mat[k][j]);
   res.resize (n);
                                                                   REP(i, R) if (i != r) {
   for(size_t i = 0; i < n; ++i)</pre>
                                                                     double w = mat[i][c] / mat[r][c];
                                                                    REP(j, C+1) mat[i][j] -= mat[r][j] * w;
       res[i] = int(fa[i].real() + 0.5);
                                                                   row[c] = r++;
NTT
                                                                 }
const int mod = 7340033;
                                                                 REP(i, C) {
const int root = 5;
                                                                   int r = row[i];
const int root_1 = 4404020;
                                                                   ans[i] = r == -1 ? 0 : mat[r][C] / mat[r][i];
const int root_pw = 1<<20;</pre>
                                                               }
void fft (vector<int> & a, bool invert) {
                                                             }
   int n = (int) a.size();
                                                              Gauss Xor
   for (int i=1, j=0; i<n; ++i) {</pre>
       int bit = n \gg 1;
                                                              const ll MAX = 1e9;
       for (; j>=bit; bit>>=1)
                                                              const 11 LOG_MAX = 64 - __builtin_clzll((11)MAX);
          j -= bit;
       j += bit;
                                                              struct gauss{
       if (i < j)
                                                                 vector<11> vet;
           swap (a[i], a[j]);
                                                                 gauss(){}
   }
                                                                 gauss(ll val){
                                                                     if(val) vet.push_back(val);
   for (int len=2; len<=n; len<<=1) {</pre>
       int wlen = invert ? root_1 : root;
                                                                 void add(ll val){
       for (int i=len; i<root_pw; i<<=1)</pre>
                                                                     int sig = LOG_MAX;
          wlen = int (wlen * 111 * wlen % mod);
                                                                     for(int i = 0; i < (int)vet.size(); i++){</pre>
       for (int i=0; i<n; i+=len) {</pre>
                                                                        while(!(vet[i] & (1LL << sig))) sig--;</pre>
          int w = 1;
                                                                         if(val & (1LL << sig)) val ^= vet[i];</pre>
           for (int j=0; j<len/2; ++j) {</pre>
                                                                     if(!val) return;
              int u = a[i+j], v = int (a[i+j+len/2] * 1
                11 * w % mod);
                                                                     sig = LOG_MAX;
              a[i+j] = u+v < mod ? u+v : u+v-mod;
                                                                     while(!(val & (1 << sig))) sig--;</pre>
              a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
                                                                     for(auto &x : vet) if(x & (1LL << sig)) x ^= val;</pre>
              w = int (w * 111 * wlen % mod);
                                                                     vet.push_back(val);
          }
                                                                     for(int i = (int)vet.size() - 2; i >= 0 && vet[i]
       }
                                                                        < vet[i + 1]; i--){
   }
                                                                         swap(vet[i], vet[i + 1]);
   if (invert) {
                                                                     }
       int nrev = reverse (n, mod);
                                                                 }
       for (int i=0; i<n; ++i)</pre>
                                                             };
          a[i] = int (a[i] * 111 * nrev % mod);
                                                              Simpson
   }
}
                                                              inline double simpson(double fl,double fr,double fmid,
Gauss
                                                                double 1,double r) { return (fl+fr+4.0*fmid)*(r-1)
// Solves systems of linear equations.
                                                              double rsimpson(double slr,double fl,double fr,double
// To use, build a matrix of coefficients and call run(
                                                                fmid,double 1,double r)
  mat, R, C). If the i-th variable is free, row[i] will
  be -1, otherwise it's value will be ans[i].
                                                                 double mid = (1+r)*0.5;
                                                                 double fml = f((1+mid)*0.5);
namespace Gauss {
                                                                 double fmr = f((mid+r)*0.5);
 const int MAXC = 1001;
                                                                 double slm = simpson(fl, fmid, fml, l, mid);
 int row[MAXC];
                                                                 double smr = simpson(fmid, fr, fmr, mid, r);
 double ans[MAXC];
                                                                 if(fabs(slr-slm-smr) < eps) return slm+smr;</pre>
                                                                 return rsimpson(slm,fl,fmid,fml,l,mid)+rsimpson(smr,
```

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```
fmid, fr, fmr, mid, r);
}
double integrate(double 1, double r) {
    double mid = (l+r)*0.5;
    double fl = f(l);
    double fr = f(r);
    double fmid = f(mid);
    return rsimpson(simpson(fl, fr, fmid, l, r), fl, fr, fmid, l, r);
}
Graphs
Dinic
```

```
const int N = 100005;
const int E = 2000006;
vector<int> g[N];
int ne;
struct Edge{
   int from, to;
   11 flow, cap;
} edge[E];
int lvl[N], vis[N], pass, start = N-2, target = N-1;
int qu[N], qt, px[N];
11 run(int s, int sink, ll minE){
   if(s == sink) return minE;
   11 ans = 0:
   for(; px[s] < (int)g[s].size(); px[s]++){</pre>
       int e = g[s][ px[s] ];
       auto &v = edge[e], &rev = edge[e^1];
       if(lvl[v.to] != lvl[s]+1 || v.flow >= v.cap)
          continue; // v.cap - v.flow < lim</pre>
       11 tmp = run(v.to, sink,min(minE, v.cap-v.flow));
       v.flow += tmp, rev.flow -= tmp;
       ans += tmp, minE -= tmp;
       if(minE == 0) break;
   }
   return ans;
}
bool bfs(int source, int sink){
   at = 0:
   qu[qt++] = source;
   lvl[source] = 1;
   vis[source] = ++pass;
   for(int i = 0; i < qt; i++){</pre>
       int u = qu[i];
       px[u] = 0;
       if(u == sink) return true;
       for(int e : g[u]){
          auto v = edge[e];
          if(v.flow >= v.cap || vis[v.to] == pass)
              continue; // v.cap - v.flow < lim</pre>
          vis[v.to] = pass;
          lvl[v.to] = lvl[u]+1;
          qu[qt++] = v.to;
       }
   }
   return false;
}
```

```
11 flow(int source = start, int sink = target){
   11 \text{ ans} = 0;
   //for(lim = (1LL << 62); lim >= 1; lim /= 2)
   while(bfs(source, sink))
       ans += run(source, sink, oo);
   return ans;
void addEdge(int u, int v, ll c = 1, ll rc = 0){
   edge[ne] = \{u, v, 0, c\};
   g[u].push_back(ne++);
   edge[ne] = \{v, u, 0, rc\};
   g[v].push_back(ne++);
}
void reset_flow(){
   for(int i = 0; i < ne; i++)</pre>
       edge[i].flow = 0;
Push relabel
// Push relabel in O(V^2 E^0.5) with gap heuristic
// It's quite fast
template<typename flow_t = long long>
struct PushRelabel {
   struct Edge { int to, rev; flow_t f, c; };
   vector<vector<Edge> > g;
   vector<flow_t> ec;
   vector<Edge*> cur;
   vector<vector<int> > hs;
   vector<int> H;
   PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H
   void add_edge(int s, int t, flow_t cap, flow_t rcap
     =0) {
       if (s == t) return;
       Edge a = \{t, (int)g[t].size(), 0, cap\};
       Edge b = \{s, (int)g[s].size(), 0, rcap\};
       g[s].push_back(a);
       g[t].push_back(b);
   void add_flow(Edge& e, flow_t f) {
       Edge &back = g[e.to][e.rev];
       if (!ec[e.to] && f)
          hs[H[e.to]].push_back(e.to);
       e.f += f, ec[e.to] += f;
       back.f -= f, ec[back.to] -= f;
   flow_t max_flow(int s, int t) {
       int v = g.size();
       H[s] = v; ec[t] = 1;
       vector<int> co(2 * v);
       co[0] = v-1;
       for(int i = 0; i < v; ++i) cur[i] = g[i].data();
       for(auto &e : g[s]) add_flow(e, e.c);
       if(hs[0].size())
       for (int hi = 0; hi >= 0;) {
          int u = hs[hi].back();
          hs[hi].pop_back();
          while (ec[u] > 0) // discharge u
              if (cur[u] == g[u].data() + g[u].size()) {
                 H[u] = 1e9;
                  for(auto &e:g[u])
                     if (e.c - e.f \&\& H[u] > H[e.to]+1)
                        H[u] = H[e.to]+1, cur[u] = &e;
```

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```
if (++co[H[u]], !--co[hi] && hi < v)</pre>
                     for(int i = 0; i < v; ++i)
                                                                    for(int u = sink; u != source; u = edge[ p[u] ].
                         if (hi < H[i] \&\& H[i] < v){
                                                                        f = min(f, edge[ p[u] ].cap);
                             --co[H[i]];
                            H[i] = v + 1;
                         }
                                                                    for(int u = sink; u != source; u = edge[ p[u] ].
                  hi = H[u];
              } else if (cur[u]->c - cur[u]->f \&\& H[u]
                                                                        edge[ p[u] ].cap -= f;
                == H[cur[u]->to]+1)
                                                                        edge[p[u] ^ 1].cap += f;
                  add_flow(*cur[u], min(ec[u], cur[u]->c
                    - cur[u]->f));
              else ++cur[u];
                                                                    mf += f;
                                                                    ans += f * d[sink];
          while (hi >= 0 && hs[hi].empty()) --hi;
                                                                return {mf, ans};
       return -ec[s];
   }
                                                             }
}:
                                                             void addEdge(int u, int v, ll c, ll cost){
Min Cost Max Flow
                                                                 edge[ne] = \{u, v, c, cost\};
                                                                g[u].push_back(ne++);
const 11 oo = 1e18;
                                                                 edge[ne] = \{v, u, 0, -cost\};
const int N = 222, E = 2 * 1000006;
                                                                 g[v].push_back(ne++);
vector<int> g[N];
                                                             Small to Large
int ne;
struct Edge{
                                                             void cnt_sz(int u, int p = -1){
   int from, to;
                                                                sz[u] = 1;
   11 cap, cost;
} edge[E];
                                                                 for(int v : g[u]) if(v != p)
                                                                    cnt_sz(v, u), sz[u] += sz[v];
int start = N-1, target = N-2, p[N];
                                                             }
int inqueue[N];
11 d[N];
                                                             void add(int u, int p, int big = -1){
                                                                // Update info about this vx in global answer
bool spfa(int source, int sink){
   for(int i = 0; i < N; i++) d[i] = oo, inqueue[i] =</pre>
                                                                for(int v : g[u]) if(v != p && v != big)
                                                                    add(v, u);
                                                             }
   d[source] = 0;
   queue<int> q;
                                                             void dfs(int u, int p, int keep){
   q.push(source);
   inqueue[source] = 1;
                                                                int big = -1, mmx = -1;
   while(!q.empty()){
                                                                 for(int v : g[u]) if(v != p \&\& sz[v] > mmx)
       int u = q.front(); q.pop();
                                                                    mmx = sz[v], big = v;
       inqueue[u] = 0;
                                                                 for(int v : g[u]) if(v != p && v != big)
       for(int e : g[u]){
                                                                    dfs(v, u, 0);
           auto v = edge[e];
           if(v.cap > 0 \text{ and } d[u] + v.cost < d[v.to]){
                                                                if(big != -1) dfs(big, u, 1);
              d[v.to] = d[u] + v.cost;
              p[v.to] = e;
                                                                add(u, p, big);
              if(!inqueue[v.to]){
                  q.push(v.to);
                                                                for(auto x : q[u]){
                  inqueue[v.to] = 1;
                                                                    \ensuremath{//} answer all queries for this \ensuremath{\text{vx}}
              }
          }
       }
                                                                if(!keep){
   }
                                                                    // Remove data from this subtree
   return d[sink] != oo;
}
                                                             Centroid Decomposition
// <max flow, min cost>
pair<11, 11> mincost(int source = start, int sink =
  target){
                                                             vector<pair<int,int>> G[500005];
   11 ans = 0, mf = 0;
                                                             int subtree[500005], k;
                                                             bool erased[500005];
   while(spfa(source, sink)){
       11 f = oo;
                                                             ll dist[500005], ans;
```

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```
for(int u : S) if(!cor[u])
int calc_sz(int v, int p){
                                                                   dfst(u, ++e);
                                                            }
   subtree[v] = 1;
   for(auto x : G[v]) if(!erased[x.ff] && x.ff != p){
                                                            Tarjan
       subtree[v] += calc_sz(x.ff,v);
   return subtree[v];
                                                            int cnt = 0, root;
}
                                                            void dfs(int u, int p = -1){
                                                                low[u] = num[u] = ++t;
int centroid(int v, int p, int treesize){
                                                                for(int v : g[u]){
   for(auto x : G[v]) if(!erased[x.ff] && x.ff != p){
                                                                   if(!num[v]){
       if(subtree[x.ff] * 2 > treesize)
                                                                       dfs(v, u);
          return centroid(x.ff, v, treesize);
                                                                          if(u == root) cnt++;
                                                                       if(low[v] >= num[u]) u PONTO DE ARTICULAÇÃO;
   return v;
                                                                       if(low[v] > num[u]) ARESTA u->v PONTE;
}
                                                                       low[u] = min(low[u], low[v]);
                                                                   }
void procurar_ans(int v, int p, int d_atual, ll custo){
                                                                   else if(v != p) low[u] = min(low[u], num[v]);
   ans = min(ans, dist[k - d_atual] + custo);
                                                               }
   if(d_atual == k) return;
   for(auto x : G[v]) if(!erased[x.ff] && x.ff != p)
       procurar_ans(x.ff, v, d_atual+1, custo+x.ss);
                                                            root PONTO DE ARTICULAÇÃO <=> cnt > 1
}
                                                            void tarjanSCC(int u){
void atualiza_dist(int v, int p, int d_atual, ll custo){
                                                                low[u] = num[u] = ++cnt;
   dist[d_atual] = min(dist[d_atual], custo);
                                                                vis[u] = 1;
   if(d_atual == k) return;
                                                                S.push_back(u);
   for(auto x : G[v]) if(!erased[x.ff] && x.ff != p)
                                                                for(int v : g[u]){
       atualiza_dist(x.ff,v,d_atual+1,custo+x.ss);
                                                                   if(!num[v]) tarjanSCC(v);
}
                                                                   if(vis[v]) low[u] = min(low[u], low[v]);
void decomp(int v, int p){
                                                                if(low[u] == num[u]){
   int treesize = calc_sz(v, v);
                                                                   ssc[u] = ++ssc_cnt; int v;
   if(treesize < k) return;</pre>
                                                                   do{
   int cent = centroid(v, v, treesize);
                                                                       v = S.back(); S.pop_back(); vis[v] = 0;
   erased[cent] = 1;
                                                                       ssc[v] = ssc_cnt;
                                                                   }while(u != v);
   for(int i = 1; i <= treesize; i++) dist[i] = 1e18;</pre>
                                                               }
   for(pair<int,int> x : G[cent]) if(!erased[x.ff]){
                                                            Max Clique
       procurar_ans(x.ff, cent, 1, x.ss);
       atualiza_dist(x.ff, cent, 1, x.ss);
                                                            long long adj[N], dp[N];
   }
                                                            for(int i = 0; i < n; i++){
   for(pair<int,int> x : G[cent]) if(!erased[x.ff])
                                                                for(int j = 0; j < n; j++){
       decomp(x.ff, cent);
                                                                   int x;
                                                                   scanf("%d",&x);
Kosaraju
                                                                   if(x \mid \mid i == j)
                                                                       adj[i] |= 1LL << j;
vector<int> g[N], gt[N], S;
int vis[N], cor[N];
                                                            }
void dfs(int u){
                                                            int resto = n - n/2;
   vis[u] = 1;
                                                            int C = n/2;
   for(int v : g[u]) if(!vis[v]) dfs(v);
                                                            for(int i = 1; i < (1 << resto); i++){</pre>
   S.push_back(u);
                                                                int x = i;
                                                                for(int j = 0; j < resto; j++)
void dfst(int u, int e){
                                                                   if(i & (1 << j))
   cor[u] = e;
                                                                       x \&= adj[j + C] >> C;
   for(int v : gt[u]) if(!cor[v]) dfst(v, e);
                                                                if(x == i){
}
                                                                   dp[i] = __builtin_popcount(i);
void kosaraju(){
                                                            }
   for(int i = 1; i <= n; i++) if(!vis[i]) dfs(i);</pre>
   for(int i = 1; i <= n; i++) for(int j : g[i])</pre>
                                                            for(int i = 1; i < (1 << resto); i++)</pre>
       gt[j].push_back(i);
                                                                for(int j = 0; j < resto; j++)
   int e = 0;
                                                                   if(i & (1 << j))
   reverse(S.begin(), S.end());
                                                                       dp[i] = max(dp[i], dp[i ^ (1 << j)]);
```

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```
int maxCliq = 0;
for(int i = 0; i < (1 << C); i++){
    int x = i, y = (1 << resto) - 1;
    for(int j = 0; j < C; j++)
        if(i & (1 << j))
            x &= adj[j] & ((1 << C) - 1), y &= adj[j] >>
            C;
    if(x != i) continue;
    maxCliq = max(maxCliq, __builtin_popcount(i) + dp[y ]);
}

Dominator Tree

vector<int> g[N], gt[N], T[N];
vector<int> S;
int dsu[N] label[N];
```

```
int dsu[N], label[N];
int sdom[N], idom[N], dfs_time, id[N];
vector<int> bucket[N];
vector<int> down[N];
void prep(int u){
   S.push_back(u);
   id[u] = ++dfs_time;
   label[u] = sdom[u] = dsu[u] = u;
   for(int v : g[u]){
       if(!id[v])
          prep(v), down[u].push_back(v);
       gt[v].push_back(u);
   }
}
int fnd(int u, int flag = 0){
   if(u == dsu[u]) return u;
   int v = fnd(dsu[u], 1), b = label[ dsu[u] ];
   if(id[ sdom[b] ] < id[ sdom[ label[u] ] ])</pre>
      label[u] = b;
   dsu[u] = v;
   return flag ? v : label[u];
}
void build_dominator_tree(int root, int sz){
   // memset(id, 0, sizeof(int) * (sz + 1));
   // for(int i = 0; i <= sz; i++) T[i].clear();
   prep(root);
   reverse(S.begin(), S.end());
   int w;
   for(int u : S){
       for(int v : gt[u]){
          w = fnd(v);
          if(id[ sdom[w] ] < id[ sdom[u] ])</pre>
              sdom[u] = sdom[w];
       gt[u].clear();
      if(u != root) bucket[ sdom[u] ].push_back(u);
       for(int v : bucket[u]){
          w = fnd(v);
          if(sdom[w] == sdom[v]) idom[v] = sdom[v];
```

else idom[v] = w;

bucket[u].clear();

```
for(int v : down[u]) dsu[v] = u;
       down[u].clear();
   }
   reverse(S.begin(), S.end());
   for(int u : S) if(u != root){
       if(idom[u] != sdom[u]) idom[u] = idom[ idom[u] ];
       T[ idom[u] ].push_back(u);
   S.clear();
Min Cost Matching
// Min cost matching
// O(n^2 * m)
// n == nro de linhas
// m == nro de colunas
// n <= m | flow == n
// a[i][j] = custo pra conectar i a j
vector < int > u(n + 1), v(m + 1), p(m + 1), way(m + 1);
for(int i = 1; i \le n; ++i){
   p[0] = i;
   int j0 = 0;
   vector<int> minv(m + 1 , oo);
   vector<char> used(m + 1 , false);
   do{
       used[j0] = true;
       int i0 = p[j0] , delta = oo, j1;
       for(int j = 1; j \le m; ++j)
          if(! used[j]){
              int cur = a[i0][j] - u[i0] - v[j];
              if(cur < minv[j])</pre>
                 minv[j] = cur, way[j] = j0;
              if(minv[j] < delta)</pre>
                 delta = minv[j], j1 = j;
          }
       for(int j = 0; j \le m; ++j)
          if(used[j])
             u[p[j]] += delta, v[j] -= delta;
              minv[j] -= delta;
       j0 = j1;
   }while(p[j0] != 0);
   do{
       int j1 = way[j0];
      p[j0] = p[j1];
       j0 = j1;
   }while(j0);
}
// match[i] = coluna escolhida para linha i
vector<int> match(n + 1);
for(int j = 1; j \le m; ++j)
   match[p[j]] = j;
int cost = -v[0];
Strings
Aho Corasick
map<char, int> to[N];
int ne = 1, term[N], fail[N];
void add_string(char *str){
   int p = 0;
```

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```
for(int i = 0; str[i]; i++){
       if(!to[p][ str[i] ]) to[p][ str[i] ] = ne++;
                                                                   memcpy(a, a1, sizeof a1);
      p = to[p][ str[i] ];
                                                                   memcpy(c, c1, sizeof c1);
   term[p] = 1;
                                                                   if(cc == n-1) break;
}
                                                               }
                                                            }
int go(int s, char c){
   while(s && !to[s].count(c)) s = fail[s];
                                                            void build_lcp(char s[], int n, int a[]){ // lcp[i] =
   if(to[s].count(c)) return to[s][c];
                                                              lcp(s[:i], s[:i+1])
   return s;
                                                               int k = 0;
                                                               //memset(lcp, 0, sizeof lcp);
                                                               for(int i = 0; i < n; i++){
void init(){
                                                                   if(c[i] == n-1) continue;
   queue<int> q;
                                                                   int j = a[c[i]+1];
   q.push(0);
                                                                   while(i+k < n \&\& j+k < n \&\& s[i+k] == s[j+k]) k
                                                                     ++:
   int u, v; char c;
                                                                   lcp[c[i]] = k;
   while(!q.empty()){
                                                                   if(k) k--;
      u = q.front(); q.pop();
                                                            }
       for(auto w : to[u]){
                                                            int comp_lcp(int i, int j){
          tie(c, v) = w;
          q.push(v);
                                                               if(i == j) return n - i;
          if(u){
                                                               if(c[i] > c[j]) swap(i, j);
              fail[v] = go(fail[u], c);
                                                               return min(lcp[k] for k in [c[i], c[j]-1]);
              term[v] |= term[ fail[v] ];
                                                            Z Algorithm
      }
   }
                                                            vector<int> z_algo(const string &s) {
                                                               int n = s.size(), L = 0, R = 0;
                                                               vector<int> z(n, 0);
Suffix Array
                                                               for(int i = 1; i < n; i++){
                                                                   if(i \le R) z[i] = min(z[i-L], R - i + 1);
int lcp[N], c[N];
                                                                   while(z[i]+i < n \&\& s[z[i]+i] == s[z[i]])
// Caractere final da string '\0' esta sendo considerado
                                                                   if(i+z[i]-1 > R) L = i, R = i + z[i] - 1;
  parte da string s
void build_sa(char s[], int n, int a[]){
                                                               return z;
   const int A = 300; // Tamanho do alfabeto
   int c1[n], a1[n], h[n + A];
                                                            Prefix function/KMP
   memset(h, 0, sizeof h);
   for(int i = 0; i < n; i++) {
                                                            vector<int> preffix_function(const string &s){
       c[i] = s[i];
                                                               int n = s.size();
      h[c[i] + 1]++;
                                                               vector<int> b(n+1);
                                                               b[0] = -1;
                                                               int i = 0, j = -1;
   partial_sum(h, h + A, h);
                                                               while(i < n){
   for(int i = 0; i < n; i++)
                                                                   while(j >= 0 \&\& s[i] != s[j]) j = b[j];
       a[h[c[i]]++] = i;
                                                                  b[++i] = ++j;
   for(int i = 0; i < n; i++)</pre>
                                                               }
      h[c[i]]--;
                                                               return b:
                                                            }
   for(int L = 1; L < n; L <<= 1) {
       for(int i = 0; i < n; i++) {</pre>
                                                            void kmp(const string &t, const string &p){
          int j = (a[i] - L + n) \% n;
                                                               vector<int> b = preffix_function(p);
          a1[h[c[j]]++] = j;
                                                               int n = t.size(), m = p.size();
                                                               int j = 0;
                                                               for(int i = 0; i < n; i++){
      int cc = -1;
                                                                   while(j \ge 0 \& t[i] != p[j]) j = b[j];
       for(int i = 0; i < n; i++) {</pre>
                                                                   j++;
          if(i == 0 || c[a1[i]] != c[a1[i-1]] || c[(a1[
                                                                   if(j == m){
            i] + L) % n] != c[(a1[i-1] + L) % n])
                                                                      //patern of p found on t
             h[++cc] = i;
                                                                      i = b[i];
          c1[a1[i]] = cc;
                                                                   }
       }
                                                               }
```

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```
Min rotation
int min_rotation(int *s, int N) {
 REP(i, N) s[N+i] = s[i];
 int a = 0;
 REP(b, N) REP(i, N) {
   if (a+i == b \mid \mid s[a+i] < s[b+i]) \{ b += max(0, i-1);
   if (s[a+i] > s[b+i]) { a = b; break; }
 }
 return a;
}
All palindrome
void manacher(char *s, int N, int *rad) {
 static char t[2*MAX];
 int m = 2*N - 1;
 REP(i, m) t[i] = -1;
 REP(i, N) t[2*i] = s[i];
 int x = 0;
 FOR(i, 1, m) {
   int &r = rad[i] = 0;
   if (i <= x+rad[x]) r = min(rad[x+x-i], x+rad[x]-i);</pre>
   while (i-r-1 >= 0 \& i+r+1 < m \& t[i-r-1] == t[i+r]
     +1]) ++r;
   if (i+r >= x+rad[x]) x = i;
 REP(i, m) if (i-rad[i] == 0 || i+rad[i] == m-1) ++rad[
 REP(i, m) rad[i] /= 2;
Suffix Automaton
map<char, int> to[2*N];
int link[2*N], len[2*N], last = 0, sz = 1;
void add_letter(char c){
   int p = last;
   last = sz++;
   len[last] = len[p] + 1;
   for(; !to[p][c]; p = link[p]) to[p][c] = last;
   if(to[p][c] == last){
       link[last] = 0;
      return;
   int u = to[p][c];
   if(len[u] == len[p]+1){
      link[last] = u;
      return;
   int c1 = sz++;
   to[c1] = to[u];
   link[c1] = link[u];
   len[c1] = len[p]+1;
   link[last] = link[u] = c1;
   for(; to[p][c] == u; p = link[p]) to[p][c] = c1;
Suffix Tree
namespace sf {
// const int NS = ; const int N = * 2;
int cn, cd, ns, en = 1, lst;
```

```
string S[NS]; int si = -1;
vector<int> sufn[N]; // sufn[si][i] no do sufixo S[si][i
struct node {
   int 1, r, si, p, suf;
   map<char, int> adj;
   node() : 1(0), r(-1), suf(0), p(0) {}
   node(int L, int R, int S, int P) : 1(L), r(R), si(S)
      . p(P) {}
   inline int len() { return r - l + 1; }
   inline int operator[](int i) { return S[si][l + i];
   inline int& operator()(char c) { return adj[c]; }
} t[N];
inline int new_node(int L, int R, int S, int P) { t[en]
  = node(L, R, S, P); return en++; }
void add_string(string s) {
   s += '; S[++si] = s; sufn[si].resize(s.size() + 1)
     ; cn = cd = 0;
   int i = 0; const int n = s.size();
   for(int j = 0; j < n; j++)
       for(; i <= j; i++) {</pre>
          if(cd == t[cn].len() \&\& t[cn](s[j])) { cn = t}
             [cn](s[j]); cd = 0; 
          if(cd < t[cn].len() \& t[cn][cd] == s[j]) {
              cd++;
              if(j < s.size() - 1) break;</pre>
              else {
                 if(i) t[lst].suf = cn;
                 for(; i <= j; i++) { sufn[si][i] = cn;</pre>
                   cn = t[cn].suf; }
              }
          } else if(cd == t[cn].len()) {
              sufn[si][i] = en;
              if(i) t[lst].suf = en; lst = en;
              t[cn](s[j]) = new_node(j, n - 1, si, cn);
              cn = t[cn].suf; cd = t[cn].len();
          } else {
              int mid = new_node(t[cn].1, t[cn].1 + cd -
                 1, t[cn].si, t[cn].p);
              t[t[cn].p](t[cn][0]) = mid;
              if(ns) t[ns].suf = mid;
              if(i) t[lst].suf = en; lst = en;
              sufn[si][i] = en;
              t[mid](s[j]) = new_node(j, n - 1, si, mid)
              t[mid](t[cn][cd]) = cn;
              t[cn].p = mid; t[cn].l += cd; cn = t[mid].
              int g = cn? j - cd : i + 1; cn = t[cn].suf
              while (g < j \&\& g + t[t[cn](S[si][g])].len
                () <= j) {
                 cn = t[cn](S[si][g]); g += t[cn].len();
              if(g == j) \{ ns = 0; t[mid].suf = cn; cd =
                 t[cn].len(); }
              else { ns = mid; cn = t[cn](S[si][g]); cd
                = j - g;  }
          }
      }
   }
};
Geometry
```

2D basics

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```
typedef double cod;
double eps = 1e-7;
bool eq(cod a, cod b){ return abs(a - b) <= eps; }</pre>
struct vec{
   cod x, y; int id;
   vec(cod a = 0, cod b = 0) : x(a), y(b) {}
   vec operator+(const vec &o) const{
       return \{x + o.x, y + o.y\};
   vec operator-(const vec &o) const{
       return {x - o.x, y - o.y};
   vec operator*(cod t) const{
       return {x * t, y * t};
   }
   vec operator/(cod t) const{
       return {x / t, y / t};
   }
   cod operator*(const vec &o) const{ // cos
       return x * o.x + y * o.y;
   cod operator^(const vec &o) const{ // sin
       return x * o.y - y * o.x;
   }
   bool operator==(const vec &o) const{
       return eq(x, o.x) && eq(y, o.y);
   }
   bool operator<(const vec &o) const{</pre>
       if(!eq(x, o.x)) return x < o.x;
       return y < o.y;</pre>
   cod cross(const vec &a, const vec &b) const{
       return (a-(*this)) ^ (b-(*this));
   int ccw(const vec &a, const vec &b) const{
       cod tmp = cross(a, b);
       return (tmp > eps) - (tmp < -eps);</pre>
   }
   cod dot(const vec &a, const vec &b) const{
       return (a-(*this)) * (b-(*this));
   }
   cod len() const{
       return sqrt(x * x + y * y); // <</pre>
   }
   double angle(const vec &a, const vec &b) const{
       return atan2(cross(a, b), dot(a, b));
   }
   double tan(const vec &a, const vec &b) const{
       return cross(a, b) / dot(a, b);
   }
   vec unit() const{
       return operator/(len());
   int quad() const{
       if(x > 0 \&\& y >= 0) return 0;
       if(x \le 0 \& y > 0) return 1;
       if(x < 0 \&\& y <=0) return 2;
       return 3;
   bool comp(const vec &a, const vec &b) const{
       return (a - *this).comp(b - *this);
   bool comp(vec b){
       if(quad() != b.quad()) return quad() < b.quad();</pre>
       if(!eq(operator^(b), 0)) return operator^(b) > 0;
       return (*this) * (*this) < b * b;</pre>
```

```
template<class T>
   void sort_by_angle(T first, T last) const{
       std::sort(first, last, [=](const vec &a, const
         vec &b){
          return comp(a, b);
       });
   vec rot90() const{ return {-y, x}; }
   vec rot(double a) const{
       return \{\cos(a)*x - \sin(a)*y, \sin(a)*x + \cos(a)*y\};
   vec proj(const vec &b) const{ // proj of *this onto
      cod k = operator*(b) / (b * b);
      return b * k;
   // proj of (*this) onto the plane orthogonal to b
   vec rejection(vec b) const{
       return (*this) - proj(b);
};
struct line{
   cod a, b, c; vec n;
   line(vec q, vec w){ // q.cross(w, (x, y)) = 0
       a = -(w.y-q.y);
      b = w.x-q.x;
      c = -(a * q.x + b * q.y);
      n = \{a, b\};
   cod dist(const vec &o) const{
       return abs(eval(o)) / n.len();
   bool contains(const vec &o) const{
       return eq(a * o.x + b * o.y + c, 0);
   }
   cod dist(const line &o) const{
       if(!parallel(o)) return 0;
       if(!eq(o.a * b, o.b * a)) return 0;
       if(!eq(a, 0))
          return abs(c - o.c * a / o.a) / n.len();
       if(!eq(b, 0))
          return abs(c - o.c * b / o.b) / n.len();
       return abs(c - o.c);
   bool parallel(const line &o) const{
       return eq(n ^ o.n, 0);
   bool operator==(const line &o) const{
       if(!eq(a*o.b, b*o.a)) return false;
       if(!eq(a*o.c, c*o.a)) return false;
       if(!eq(c*o.b, b*o.c)) return false;
       return true;
   bool intersect(const line &o) const{
       return !parallel(o) || *this == o;
   vec inter(const line &o) const{
       if(parallel(o)){
          if(*this == o){ }
          else{ /* dont intersect */ }
       auto tmp = n \hat{o.n};
       return {(o.c*b -c*o.b)/tmp, (o.a*c -a*o.c)/tmp};
   }
```

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```
vec at_x(cod x) const{
      return \{x, (-c-a*x)/b\};
   }
   vec at_y(cod y) const{
       return \{(-c-b*y)/a, y\};
   }
   cod eval(const vec &o) const{
      return a * o.x + b * o.y + c;
};
struct segment{
   vec p, q;
   segment(vec a = vec(), vec b = vec()): p(a), q(b) {}
   bool onstrip(const vec &o) const{ // onstrip strip
       return p.dot(o, q) >= -eps && q.dot(o, p) >= -eps
   }
   cod len() const{
      return (p-q).len();
   cod dist(const vec &o) const{
       if(onstrip(o)) return line(p, q).dist(o);
      return min((o-q).len(), (o-p).len());
   bool contains(const vec &o) const{
      return eq(p.cross(q, o), 0) && onstrip(o);
   }
   bool intersect(const segment &o) const{
       if(contains(o.p)) return true;
       if(contains(o.q)) return true;
       if(o.contains(q)) return true;
       if(o.contains(p)) return true;
       return p.ccw(q, o.p) * p.ccw(q, o.q) == -1
      && o.p.ccw(o.q, q) * o.p.ccw(o.q, p) == -1;
   bool intersect(const line &o) const{
      return o.eval(p) * o.eval(q) <= 0;</pre>
   cod dist(const segment &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q)
          || o.onstrip(p) || o.onstrip(q))
             return line(p, q).dist(line(o.p, o.q));
       else if(intersect(o)) return 0;
       return min(min(dist(o.p), dist(o.q)),
                min(o.dist(p), o.dist(q)));
   }
   cod dist(const line &o) const{
       if(line(p, q).parallel(o))
          return line(p, q).dist(o);
       else if(intersect(o)) return 0;
       return min(o.dist(p), o.dist(q));
};
struct hray{
   vec p, q;
   hray(vec a = vec(), vec b = vec()): p(a), q(b){}
   bool onstrip(const vec &o) const{ // onstrip strip
      return p.dot(q, o) >= -eps;
   cod dist(const vec &o) const{
       if(onstrip(o)) return line(p, q).dist(o);
      return (o-p).len();
   }
```

```
bool intersect(const segment &o) const{
       if(!o.intersect(line(p,q))) return false;
       if(line(o.p, o.q).parallel(line(p,q)))
          return contains(o.p) || contains(o.q);
       return contains(line(p,q).inter(line(o.p,o.q)));
   bool contains(const vec &o) const{
       return eq(line(p, q).eval(o), 0) && onstrip(o);
   cod dist(const segment &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q))
              return line(p, q).dist(line(o.p, o.q));
          return o.dist(p);
       else if(intersect(o)) return 0;
       return min(min(dist(o.p), dist(o.q)),
                o.dist(p));
   bool intersect(const hray &o) const{
       if(!line(p, q).parallel(line(o.p, o.q)))
          return false;
       auto pt = line(p, q).inter(line(o.p, o.q));
       return contains(pt) && o.contains(pt); // <<</pre>
   bool intersect(const line &o) const{
       if(line(p, q).parallel(o)) return line(p, q)== o;
       if(o.contains(p) || o.contains(q)) return true;
       return (o.eval(p) >= -eps)^(o.eval(p)<o.eval(q));</pre>
       return contains(o.inter(line(p, q)));
   cod dist(const line &o) const{
       if(line(p,q).parallel(o))
          return line(p,q).dist(o);
       else if(intersect(o)) return 0;
       return o.dist(p);
   }
   cod dist(const hray &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || o.onstrip(p))
              return line(p,q).dist(line(o.p, o.q));
          return (p-o.p).len();
       else if(intersect(o)) return 0;
       return min(dist(o.p), o.dist(p));
   }
};
double heron(cod a, cod b, cod c){
   cod s = (a + b + c) / 2;
   return sqrt(s * (s - a) * (s - b) * (s - c));
Circle line intersection
// intersection of line a * x + b * y + c = 0
// and circle centered at the origin with radius r
double r, a, b, c; // given as input
double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
if(c*c > r*r*(a*a+b*b)+EPS)
   puts("no points");
else if(abs(c*c - r*r*(a*a+b*b)) < EPS){
   puts("1 point");
   cout << x0 << ' ' << y0 << ' 'n';
}
else {
   double d = r*r - c*c/(a*a+b*b);
   double mult = sqrt (d / (a*a+b*b));
```

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```
double ax, ay, bx, by;
ax = x0 + b * mult;
bx = x0 - b * mult;
ay = y0 - a * mult;
by = y0 + a * mult;
puts ("2 points");
cout<<ax<<' '<<ay<<'\n'<<bx<<' '<<by<<'\n';</pre>
```

Circle Circle intersection

Assume that the first circle is centered at the origin and second at (x2, y2). Find circle line intersection of first circle and line Ax + By + C = 0, where $A = -2x_2$, $B = -2y_2$, $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$.

Be aware of corner case with two circles centered at the same point.

Tangents of two circles

```
// solve first for same circle(and infinitely many
  tangents)
// Find up to four tangents of two circles
void tangents(pt c, double r1, double r2, vector<line> &
   ans){
   double r = r2 - r1;
   double z = c.x * c.x + c.y * c.y;
   double d = z - r * r;
   if(d < -EPS) return;</pre>
   d = sqrt(abs(d));
   line 1;
   1.a = (c.x * r + c.y * d) / z;
   1.b = (c.y * r - c.x * d) / z;
   1.c = r1;
   ans.push_back (1);
vector<line> tangents(circle a, circle b){
   vector<line> ans;
   pt aux = a.center - b.center;
   for(int i = -1; i \le 1; i += 2)
       for(int j = -1; j \le 1; j += 2)
          tangents(aux, a.r * i, b.r * j, ans);
   for(size_t i = 0; i < ans.size(); ++i)</pre>
       ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
   return ans;
Convex Hull
vector<vec> monotone_chain_ch(vector<vec> P){
   sort(P.begin(), P.end());
   vector<vec> L, U;
   for(auto p : P){
      while(L.size() >= 2 && L[L.size() - 2].cross(L.
         back(), p) < 0)
          L.pop_back();
      L.push_back(p);
   }
   reverse(P.begin(), P.end());
   for(auto p : P){
      while(U.size() >= 2 && U[U.size() - 2].cross(U.
         back(), p) < 0
          U.pop_back();
```

```
U.push_back(p);
}
L.pop_back(), U.pop_back();
L.reserve(L.size() + U.size());
L.insert(L.end(), U.begin(), U.end());
return L;
}
```

Check point inside polygon

```
bool below(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) <= 0;</pre>
}
bool above(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) >= 0;
}
// lowerhull, upperhull and point, borders included
bool inside_poly(const vector<vec> &lo, const vector<vec</pre>
  > &hi, vec p){
   return below(hi, p) && above(lo, p);
```

Check point inside polygon without lower/upper hull

```
// borders included
// must not have 3 colinear consecutive points
bool inside_poly(const vector<vec> &v, vec p){
   if(v[0].ccw(v[1], p) < 0) return false;</pre>
   if(v[0].ccw(v.back(), p) > 0) return 0;
   if(v[0].ccw(v.back(), p) == 0)
       return v[0].dot(p, v.back()) >= 0
          && v.back().dot(p, v[0]) >= 0;
   int L = 1, R = (int)v.size() - 1, ans = 1;
   while(L <= R){</pre>
       int mid = (L+R)/2;
       if(v[0].ccw(v[mid], p) >= 0) ans = mid, L = mid
         +1;
       else R = mid-1;
   }
   return v[ans].ccw(v[(ans+1)%v.size()], p) >= 0;
```

Minkowski sum

```
vector<vec> mk(const vector<vec>&a,const vector<vec>&b){
   int i = 0, j = 0;
   for(int k = 0; k < (int)a.size(); k++)if(a[k] <a[i])
        i = k;
   for(int k = 0; k < (int)b.size(); k++)if(b[k] <b[j])
        j = k;

   vector<vec> c;
   c.reserve(a.size() + b.size());
   for(int k = 0; k < int(a.size()+b.size()); k++){</pre>
```

```
vec pt{a[i] + b[j]};
if((int)c.size() >= 2
    && c[c.size()-2].ccw(c.back(), pt) == 0)
        c.pop_back();
c.push_back(pt);
int q = i+1, w = j+1;
if(q == int(a.size())) q = 0;
if(w == int(b.size())) w = 0;
if(c.back().ccw(a[i]+b[w], a[q]+b[j]) < 0) i = q;
else j = w;
}
c.shrink_to_fit();</pre>
```

Geo Notes

Center of mass

System of points(2D/3D): Mass weighted average of points. **Frame(2D/3D):** Get middle point of each segment solve as previously.

Triangle: Average of vertices.

2D Polygon: Compute **signed** area and center of mass of triangle $((0,0), p_i, p_{i+1})$. Then solve as system of points.

Polyhedron surface: Solve each face as a 2D polygon(be aware of (0, 0)) then replace each face with its center of mass and solve as system of points.

Tetrahedron(Triangular pyramid): As triangles, its the average of points.

Polyhedron: Can be done as 2D polygon, but with tetrahedralization intead of triangulation.

Pick's Theorem

Given a polygon without self-intersections and all its vertices on integer coordinates in some 2D grid. Let A be its area, I the number of points with interger coordinates stricly inside the polygon and B the number of points with interger coordinates in the border of the polygon. The following formula holds: $A = I + \frac{B}{2} - 1$.

Miscellaneous

multiset<int> S;

LIS

```
for(int i = 0; i < n; i++){
    auto it = S.upper_bound(a[i]); // low for inc
    if(it != S.end()) S.erase(it);
    S.insert(a[i]);
}
ans = S.size();

DSU rollback

struct DSU{
    vector<int> sz, p, change;
    vector<tuple<int, int, int>> modifications;
    vector<size_t> saves;
    bool bipartite;

DSU(int n): sz(n+1, 1), p(n+1), change(n+1),
    bipartite(true){
        iota(p.begin(), p.end(), 0);
    }
```

```
void add_edge(int u, int v){
       if(!bipartite) return;
       int must_change = get_colour(u) == get_colour(v);
       int a = rep(u), b = rep(v);
       if(sz[a] < sz[b]) swap(a, b);
       if(a != b){
          p[b] = a;
          modifications.emplace_back(b, change[b],
            bipartite);
          change[b] ^= must_change;
          sz[a] += sz[b];
       }
       else if(must_change){
          modifications.emplace_back(0, change[0],
            bipartite);
          bipartite = false;
       }
   }
   int rep(int u){
       return p[u] == u ? u : rep(p[u]);
   }
   int get_colour(int u){
       if(p[u] == u) return change[u];
       return change[u] ^ get_colour(p[u]);
   }
   void reset(){
       modifications.clear();
       saves.clear();
       iota(p.begin(), p.end(), 0);
       fill(sz.begin(), sz.end(), 1);
       fill(change.begin(), change.end(), 0);
       bipartite = true;
   }
   void rollback(){
       int u = get<0>(modifications.back());
       tie(ignore, change[u], bipartite) = modifications
         .back();
       sz[ p[u] ] -= sz[u];
       p[u] = u;
       modifications.pop_back();
   }
   void reload(){
       while(modifications.size() > saves.back())
          rollback();
       saves.pop_back();
   }
   void save(){
       saves.push_back(modifications.size());
};
Buildings
// count the number of circular arrays
// of size m, with elements on range
// [1, c**(x*x)]
#include<bits/stdc++.h>
using namespace std;
#define debug(x) cerr << fixed << #x << " = " << x <<
  endl:
```

```
Klondike
const int MOD = 1e9 + 7, MAX = 1e5 + 5;
int dp[MAX];
                                                            // minimum number of moves to make
                                                            // all elements equal
inline int add(int a, int b) {
 return a + b >= MOD ? a + b - MOD : a;
                                                            // move: change a segment of equal value
                                                            // elements to any value
inline int sub(int a, int b) {
 return a - b < 0? a - b + MOD: a;
                                                            int v[305];
                                                            int dp[305][305];
                                                            int rec[305][305];
inline int mult(int a, int b) {
 return (1LL * a * b) % MOD;
                                                            int f(int 1, int r){
                                                              if(r == 1) return 1;
int f_exp(int x, int exp) {
                                                              if(r < 1) return 0;</pre>
 if(exp == 0) return 1;
                                                              if(dp[l][r] != -1) return dp[l][r];
 else if(exp & 1) return mult(x, f_exp(x, exp - 1));
 return f_exp(mult(x, x), exp / 2);
                                                              int ans = f(1+1, r) + 1;
                                                              for(int i = l+1; i <= r; i++)</pre>
inline int inv(int x) {
                                                               if(v[i] == v[1])
                                                                 ans = min(ans, f(1, i - 1) + f(i+1, r));
 return f_exp(x, MOD - 2);
int main() {
                                                              return dp[l][r] = ans;
 ios::sync_with_stdio(false);
 cin.tie(NULL); cout.tie(NULL);
                                                            int main() {
                                                              int n. m:
 int n, m, c;
                                                              memset(dp, -1, sizeof dp);
 cin >> n >> m >> c;
                                                              scanf("%d %d",&n , &m);
 int x = f_{exp}(c, n * n);
                                                              for(int i = 0; i < n; i++){
 int ans = f_{exp}(x, m);
                                                               scanf("%d",v+i);
 for(int i = 1; i <= m; i++) {</pre>
                                                               if(i \&\& v[i] == v[i-1]){
   if(m % i == 0) {
                                                                 i--;
     int y = f_{exp}(x, i);
                                                                 n--;
     for(int j = 1; j < i; j++) {
                                                               }
                                                              }
      if(i % j == 0)
                                                             printf("%d\n",f(0, n-1) - 1);
        y = sub(y, mult(j, dp[j]));
                                                              // printf("%d\n",rec[0][n-1] );
                                                              // printf("%d\n",rec[1][n-1] );
     dp[i] = mult(y, inv(i));
     ans = sub(ans, mult(i - 1, dp[i]));
                                                              // printf("%d\n",rec[2][n-3] );
 }
                                                            Hilbert Order
 cout << ans << '\n';</pre>
                                                            // maybe use B = n / sqrt(q)
                                                            inline int64_t hilbertOrder(int x, int y, int pow = 21,
 return 0;
                                                              int rotate = 0) {
                                                               if(pow == 0) return 0;
                                                                int hpow = 1 \ll (pow-1);
Rand
                                                                int seg = (x < hpow) ? (
                                                                   (y < hpow) ? 0 : 3
#include <random>
                                                               ):(
#include <chrono>
                                                                   (y < hpow) ? 1 : 2
                                                               );
cout << RAND_MAX << endl;</pre>
                                                               seg = (seg + rotate) & 3;
                                                               const int rotateDelta[4] = {3, 0, 0, 1};
mt19937 rng(chrono::steady_clock::now().time_since_epoch
                                                               int nx = x & (x \hat{pow}), ny = y & (y \hat{pow});
  ().count());
                                                               int nrot = (rotate + rotateDelta[seg]) & 3;
vector<int> permutation(N);
                                                               int64_t subSquareSize = int64_t(1) << (2*pow - 2);</pre>
                                                                int64_t ans = seg * subSquareSize;
iota(permutation.begin(), permutation.end(), 0);
                                                                int64_t add = hilbertOrder(nx, ny, pow-1, nrot);
shuffle(permutation.begin(), permutation.end(), rng);
                                                                ans += (seg == 1 || seg == 2) ? add : (subSquareSize
                                                                   - add - 1);
iota(permutation.begin(), permutation.end(), 0);
                                                               return ans;
for(int i = 1; i < N; i++){
                                                            Modular Factorial
   swap(permutation[i], permutation[
     uniform_int_distribution<int>(0, i)(rng)]);
                                                            // Compute (1*2*...*(p-1)*1*(p+1)*(p+2)*..*n) % p
}
                                                            // in O(p*lg(n))
```

```
for(; ((c < '0' || c > '9') && c != '-'); c = gc());
int factmod(int n, int p){
                                                               if(c == '-'){
   int ans = 1;
                                                                   neg = 1;
   while (n > 1) {
                                                                   c = gc();
       for(int i = 2; i <= n % p; i++)</pre>
          ans = (ans * i) % p;
                                                               for(; c >= '0' && c <= '9'; c = gc())</pre>
      n /= p;
                                                                   x = (x << 1) + (x << 3) + c - '0';
       if(n \% 2) ans = p - ans;
                                                               if(neg) x = -x;
   return ans % p;
                                                            inline void print_int(int n){
                                                               int rev = 0, count = 0, neg;
int fac_pow(int n, int p){
                                                               if(n == 0){
   int ans = 0;
                                                                   pc('0');
   while(n) n \neq p, ans += n;
                                                                   return:
   return ans:
                                                               if(n < 0) n = -n, neg = 1;
                                                               while(n % 10 == 0) count++, n /= 10;
                                                               for(rev = 0; n != 0; n /= 10)
int C(int n, int k, int p){
   if(fac_pow(n, p) > fac_pow(n-k, p) + fac_pow(k, p))
                                                                   rev = (rev << 3) + (rev << 1) + n % 10;
                                                               if(neg) pc('-');
   int tmp = factmod(k, p) * factmod(n-k, p);
                                                               while(rev != 0) pc(rev % 10 + '0'), rev /= 10;
   return (f_exp(tmp, p-2, p) * factmod(n, p)) % p;
                                                               while(count--) pc('0');
                                                               pc('\n');
                                                            }
Enumeration all submasks of a bitmask
                                                            inline void print_string(char *str){
// loop through all submask of a given bitmask
                                                               while(*str) pc(*str++);
// it does not include mask 0
                                                               pc('\n');
for(int sub = mask; sub; sub = (sub-1)&mask){
                                                            Knapsack Bounded with Cost
Slope Trick
                                                            // menor custo para conseguir peso ate {\tt M} usando {\tt N} tipos
                                                              diferentes de elementos, sendo que o i-esimo elemento
///By woqja125, contest: Codeforces Round #371 (Div. 1),
                                                              pode ser usado b[i] vezes, tem peso w[i] e custo c[i]
   problem: (C) Sonya and Problem Wihtout a Legend,
                                                            // O(N * M)
  Accepted, #
#include <stdio.h>
                                                            int b[N], w[N], c[N];
                                                            MinQueue Q[M]
#include <queue>
                                                            int d[M] //d[i] = custo minimo para conseguir peso i
int main() {
                                                            for(int i = 0; i <= M; i++) d[i] = i ? oo : 0;</pre>
   int n, t;
                                                            for(int i = 0; i < N; i++){
   long long ans = 0;
                                                               for(int j = 0; j < w[i]; j++)
   std::priority_queue<int> Q;
                                                                   Q[j].clear();
   scanf("%d%d", &n, &t);
                                                               for(int j = 0; j <= M; j++){</pre>
   Q.push(t);
                                                                   q = Q[j \% w[i]];
   for(int i = 1; i < n; i++) {</pre>
                                                                   if(q.size() >= q) q.pop();
       scanf("%d", &t); t -= i;
                                                                   q.add(c[i]);
      Q.push(t);
                                                                   q.push(d[j]);
       if(Q.top() > t) {
                                                                   d[j] = q.getmin();
          ans += Q.top() - t;
                                                               }
          Q.pop();
          Q.push(t);
                                                            LCA < O(nlgn), O(1) >
   printf("%lld", ans);
                                                            int start[N], dfs_time;
   return 0;
                                                            int tour[2*N], id[2*N];
Fast IO
                                                            void dfs(int u){
                                                               start[u] = dfs_time;
#define pc(x) putchar_unlocked(x)
                                                               id[dfs_time] = u;
#define gc(x) getchar_unlocked(x)
                                                               tour[dfs_time++] = start[u];
                                                               for(int v : g[u]){
                                                                   dfs(v);
inline void scan_int(int &x){
   register int c = gc();
                                                                   id[dfs_time] = u;
   x = 0;
                                                                   tour[dfs_time++] = start[u];
   int neg = 0;
```

```
char BUF[BUFSIZE+1], *inp=BUF;
#define GETCHAR(INP) { \
    if(!*inp && !REACHEOF) { \
        memset(BUF,0,sizeof BUF);\
        int inpzzz = fread(BUF,1,BUFSIZE,stdin);\
        if (inpzzz != BUFSIZE) REACHEOF = true;\
        inp=BUF; \
    } \
    INP=*inp++; \
#define DIG(a) (((a)>='0')&&((a)<='9'))
#define GN(j) { \
    AM=0; \
    GETCHAR(INP); while(!DIG(INP) && INP!='-') GETCHAR(
      INP);\
    if (INP=='-') {AM=1;GETCHAR(INP);} \
    j=INP-'0'; GETCHAR(INP); \
    \label{eq:while} \begin{aligned} &\text{while}(\texttt{DIG}(\texttt{INP})) \{ j = 10*j + (\texttt{INP-'0'}); \texttt{GETCHAR}(\texttt{INP}); \} \  \, \\ \end{aligned}
    if (AM) j=-j;\
```

Modular summation

```
//calcula (sum(0 <= i <= n) P(i)) % mod,
//onde P(i) eh uma PA modular (com outro modulo)
namespace sum_pa_mod{
   ll calc(ll a, ll b, ll n, ll mod){
       assert(a&&b);
       if(a >= b){
          ll ret = ((n*(n+1)/2) \mod)*(a/b);
          if(a%b) ret = (ret + calc(a%b,b,n,mod))%mod;
          else ret = (ret+n+1)%mod;
          return ret;
       return ((n+1)*(((n*a)/b+1)%mod) - calc(b,a,(n*a)/b+1)%mod)
         b, mod) + mod + n/b + 1)%mod;
   }
   //P(i) = a*i \mod m
   11 solve(11 a, 11 n, 11 m, 11 mod){
       a = (a\%m + m)\%m;
       if(!a) return 0;
       11 \text{ ret} = (n*(n+1)/2) \text{%mod};
       ret = (ret*a)%mod;
       11 g = \_\_gcd(a,m);
       ret -= m*(calc(a/g,m/g,n,mod)-n-1);
       return (ret%mod + mod)%mod;
   }
   //P(i) = a + r*i \mod m
   11 solve(l1 a, l1 r, l1 n, l1 m, l1 mod){
       a = (a\%m + m)\%m;
       r = (r\%m + m)\%m;
       if(!r) return (a*(n+1))%mod;
       if(!a) return solve(r, n, m, mod);
       11 g, x, y;
```

```
g = gcdExtended(r, m, x, y);
x = (x%m + m)%m;
ll d = a - (a/g)*g;
a -= d;
x = (x*(a/g))%m;
return (solve(r, n+x, m, mod) - solve(r, x-1, m, mod) + mod + d*(n+1))%mod;
};
```

Burnside's Lemma

Let (G, \oplus) be a finite group that acts on a set X. It should hold that $e_g * x = x$ and $g_1 * (g_2 * x) = (g_1 \oplus g_2) * x$, $\forall x \in X, g_1, g_2 \in G$. For each $g \in G$ let $X^g = \{x \in X \mid g * x = x\}$. The number of orbits its given by:

 $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$

Wilson's Theorem

 $(n-1)! = -1 \mod n \iff n \text{ is prime}$

Fibonacci

- $F_{n-1}F_{n+1} F_n^2 = (-1)^n$
- $F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$
- $GCD(F_n, F_m) = F_{GCD(n,m)}$
- $F_n = \frac{(\frac{1+\sqrt{5}}{2})^n (\frac{1-\sqrt{5}}{2})^n}{\sqrt{5}}$

Kirchhoff's Theorem

Laplacian matrix is L = D - A, where D is a diagonal matrix with vertex degrees on the diagonals and A is adjacency matrix.

The number of spanning trees is any cofactor of L. i-th cofactor is determinant of the matrix gotten by removing i-th row and column of L.

Multigraphs

In D[i][i] all loops are excluded. A[i][j] = number of edges from i to j.

Directed multigraphs

D[i][i] = indegree of i minus the number of loops at i. A[i][j] = number of edges from i to j.

The number of oriented spanning trees rooted at a vertex i is the determinant of the matrix gotten by removing the ith row and column of L.

Matroid

Let *X* set of objects, $I \subseteq 2^X$ set of independents sets such that:

- 1. $\emptyset \in I$
- 2. $A \in I, B \subseteq A \implies B \in I$
- 3. Exchange axiom, $A \in I, B \in I, |B| > |A| \implies \exists x \in B \setminus A : A \cup \{x\} \in I$
- 4. $A \subseteq X$ and I and I' are maximal independent subsets of A then |I| = |I'|

Then (X, I) is a matroid. The combinatorial optimization problem associated with it is: Given a weight $w(e) \ge 0 \ \forall e \in X$, find an independet subset that has the largest possible total weight.

Matroid intersection

```
// Input two matroids (X, I_a) and (X, I_b)
// output set I of maximum size, I \in I_a and I \in I_b
set<> I;
while(1){
   for(e_i : X \setminus I)
       if(I + e_i \in I_a \text{ and } I + e_i \in I_b)
           I = I + e_i;
   set<> A, T; queue<> Q;
   for(x : X) label[x] = MARK1;
   for(e_i : X \setminus I){
       if(I + e_i \setminus in I_a)
           Q.push(e_i), label[e_i] = MARK2;
       else{
           for (x \text{ such that } I - x + e_i \in I_a)
               A[x].push(e_i);
       if(I + e_i \setminus in I_b)
           T = T + \{e_i\}
       else{
           for (x \text{ such that } I - x + e_i \setminus in I_b)
               A[e_i].push(x);
   }
   if(T.empty()) break;
   bool found = false;
   while(!Q.empty() and !found){
       auto e = Q.front(); Q.pop();
       for(x : A[e]) if(label[x] == MARK1){
           label[x] = e;
           Q.push(x);
           if(x \in T)
               found = true;
               put = 1;
               while(label[x] != MARK2){
                   if(put) I = I + x;
                   else I = I - x;
                   put = 1 - put;
               I = I + x;
               break;
           }
   if(!found) break;
}
return I;
   Where path(e) = [e] if label[e] = MARK2, path(label[e]) +
[e] otherwise.
```

Matroid Union

Given k matroids over the same set of objects (X, I_1) , (X, I_2) , ..., (X, I_k) find $A_1 \in I_1$, $A_2 \in I_2$, ..., $A_k \in I_k$ such that $i \neq j, A_i \cap A_j = \emptyset$ and $|\bigcup_{i=1}^k A_i|$ is maximum. Matroid union

can be reduced to matroid intersection as follows.

Let $X' = X \times \{1, 2, ..., k\}$, ie, k copies of each element of X with different colors. M1 = (X', Q) where $B \in Q \iff \forall 1 \le i \le k$, $\{x \mid (x, i) \in B\} \in I_i$, ie, for each color, B is independent. M2 = (X', W) where $B \in W \iff i \ne j \implies \neg((x, i) \in B \land (x, j) \in B)$, ie, each element is picked by at most one color.

Intersection of *M*1 and *M*2 is the answer for the combinatorial problem of matroid union.

Edge coloring

```
Data: A graph G

Result: A proper coloring c of the edges of G

Let U := E(G);

while U \neq \emptyset do

Let (u,v) be any edge in U;

Let F[1:k] be a maximal fan of u starting at F[1]=v;

Let c be a color that is free on u and d be a color

that is free on F[k];

Invert the cd_u path;

Let w \in V(G) be such that w \in F, F' = [F[1]...w] is a

fan and d is free on w;

Rotate F' and set c(u,w)=d;

U := U - (u,v);
```

Notes

When we repeat something and each time we have probability p to succeed then the expected number or tries is $\frac{1}{p}$, till we succeed.

Small to large

Trick in statement If k sets are given you should note that the amount of different set sizes is $O(\sqrt{s})$ where s is total size of those sets. And no more than \sqrt{s} sets have size greater than \sqrt{s} . For example, a path to the root in Aho-Corasick through suffix links will have at most $O(\sqrt{s})$ vertices.

gcd on subsegment, we have at most $log(a_i)$ different values in $\{gcd(a_j, a_{j+1}, ..., a_i) \text{ for } j < i\}$.

From static set to expandable. To insert, create a new set with the new element. While there are two sets with same size, merge them. There will be at most log(n) disjoints sets.

Matrix exponentiation optimization. Save binary power of A_{nxn} and answer q queries $b = A^m x$ in $O((n^3 + qn^2)log(m))$.

Ternary search on integers into binary search, comparing f(mid) and f(mid+1), binary search on derivative

Dynamic offline set For each element we will wind segment of time [a, b] such that element is present in the set during this whole segment. Now we can come up with recursive procedure which handles [l, r] time segment considering that all elements such that $[l, r] \subset [a, b]$ are already included into the set. Now, keeping this invariant we recursively go into [l, m] and [m + 1, r] subsegments. Finally when we come into segment of length 1.

```
a > b \implies a \mod b < \frac{a}{2}
```