# ICPC Team Reference

# University of Brasilia

| Contents |      |                              |    | 4 | Strings    |  | 15       |
|----------|------|------------------------------|----|---|------------|--|----------|
|          |      |                              |    |   | 4.1        |  | 15       |
| 1        | Date | Structures                   | 2  |   | 4.2        | Suffix Array                                 |          |
| 1        |      |                              |    |   | 4.3        | Z Algorithm                                  |          |
|          | 1.1  | Merge Sort Tree              | 2  |   | 4.4        | •  | 16       |
|          | 1.2  | Wavelet Tree                 | 2  |   | 4.5        | Min rotation                                 | 16       |
|          | 1.3  | Order Set                    | 2  |   | 4.6        | All palindrome                               | 16       |
|          | 1.4  | Hash table                   | 3  |   | 4.7<br>4.8 | Suffix Automaton                             | 16<br>16 |
|          | 1.5  | Convex Hull Trick Simple     | 3  |   | 4.0        | Sum Tree                                     | 10       |
|          | 1.6  | Convex Hull Trick            | 3  | 5 | Geo        | metry  | 17       |
|          | 1.7  | Min queue                    | 3  |   | 5.1        | 2D basics                                    | 17       |
|          | 1.8  | Sparse Table                 | 4  |   | 5.2        | Circle line intersection                     | 19       |
|          | 1.9  | Treap                        | 4  |   | 5.3        | Half plane intersection                      | 19       |
|          | 1.10 | , <del>-</del>               | 4  |   | 5.4        | Detect empty Half plane intersection         | 19       |
|          | 1.11 | Heavy Light Decomposition    | 5  |   | 5.5        | Circle Circle intersection                   | 20       |
|          |      | Iterative Segtree            | 5  |   | 5.6        | Tangents of two circles                      | 20       |
|          |      | LiChao's Segtree             | 5  |   | 5.7        | Convex Hull                                  | 20       |
|          |      | Palindromic tree             | 5  |   | 5.8        | Check point inside polygon                   | 20       |
|          | 1.17 | Tamaronne acc                | 9  |   | 5.9        | Check point inside polygon without lower/up- |          |
| 2        | Mat  | h                            | 6  |   |            | per hull                                     | 20       |
| _        | 2.1  | Extended Euclidean Algorithm | 6  |   |            | Minkowski sum                                | 20       |
|          | 2.2  | Chinese Remainder Theorem    | 6  |   | 5.11       | Geo Notes                                    | 21       |
|          |      |                              |    |   |            | 5.11.1 Center of mass                        | 21       |
|          | 2.3  | Preffix inverse              | 6  |   |            | 5.11.2 Pick's Theorem                        | 21       |
|          | 2.4  | Pollard Rho                  | 6  | 6 | Mic        | cellaneous                                   | 21       |
|          | 2.5  | Miller Rabin                 | 7  | ١ | 6.1        | LIS  |          |
|          | 2.6  | Totiente                     | 7  |   | 6.2        | DSU rollback                                 |          |
|          | 2.7  | Primitive root               | 7  |   | 6.3        | Buildings                                    |          |
|          | 2.8  | Mobius Function              | 7  |   | 6.4        | Rand   | 22       |
|          | 2.9  | Mulmod TOP                   | 7  |   | 6.5        | Klondike                                     | 22       |
|          |      | Matrix Determinant           | 7  |   | 6.6        | Hilbert Order                                | 22       |
|          | 2.11 | Simplex Method               | 8  |   | 6.7        | Modular Factorial                            | 22       |
|          | 2.12 | FFT                          | 8  |   | 6.8        | Enumeration all submasks of a bitmask        | 23       |
|          | 2.13 | FFT                          | 9  |   | 6.9        | Slope Trick                                  | 23       |
|          | 2.14 | NTT                          | 10 |   |            | Fast IO                                      |          |
|          | 2.15 | Gauss                        | 10 |   |            | Knapsack Bounded with Cost                   |          |
|          | 2.16 | Gauss Xor                    | 11 |   |            | LCA <o(nlgn), o(1)=""></o(nlgn),>            |          |
|          |      | Simpson                      | 11 |   |            | Buffered reader                              |          |
|          |      |                              |    |   |            | Modular summation                            |          |
| 3        | Gra  | ohs                          | 11 |   |            |  | 24       |
|          | 3.1  | Dinic                        | 11 |   |            |  | 24       |
|          | 3.2  | Push relabel                 | 11 |   |            | Fibonacci                                    | 24       |
|          | 3.3  | Min Cost Max Flow            | 12 |   | 0.10       | 6.18.1 Multigraphs                           | 24<br>24 |
|          | 3.4  | Small to Large               | 13 |   |            | 6.18.2 Directed multigraphs                  | 24       |
|          | 3.5  | Centroid Decomposition       | 13 |   | 6 19       | Matroid                                      | 24       |
|          |      | -                            | 13 |   | 0.17       | 6.19.1 Matroid intersection                  | 25       |
|          | 3.6  | Kosaraju                     | 13 |   |            |  | 25       |
|          | 3.7  | Tarjan                       |    |   | 6.20       |  | 25       |
|          | 3.8  | Max Clique                   | 14 |   |            | Notes  |          |
|          | 3.9  | Dominator Tree               | 14 |   |            | Misra Gries edge coloring algorithm notes    |          |
|          | 3.10 | Min Cost Matching            | 14 |   |            |  |          |

University of Brasilia Data Structures, 2

```
set ts=4 sw=4 sta nu rnu sc stl+=%F cindent
set bg=dark ruler clipboard=unnamed,unnamedplus
  timeoutlen=100
imap {<CR> {<CR>}<Esc>0}
nmap <F2> 0V$%d
nmap <C-down> :m+1<CR>
nmap <C-up> :m-2<CR>
vmap < C-c > "+y
nmap <C-a> ggVG
syntax on
alias cmp='g++ -Wall -Wformat=2 -Wshadow -Wconversion -
  fsanitize=address -fsanitize=undefined -fno-sanitize-
  recover -std=c++14'
```

# **Data Structures**

# **Merge Sort Tree**

```
struct MergeTree{
   int n:
   vector<vector<int>> st;
   void build(int p, int L, int R, const int v[]){
          st[p].push_back(v[L]);
          return;
       }
      int mid = (L+R)/2;
      build(2*p, L, mid, v);
      build(2*p+1, mid+1, R, v);
       st[p].resize(R-L+1);
      merge(st[2*p].begin(), st[2*p].end(),
              st[2*p+1].begin(), st[2*p+1].end(),
              st[p].begin());
   }
   int query(int p, int L, int R, int i, int j, int x)
     const{
       if(L > j || R < i) return 0;</pre>
       if(L >= i \&\& R <= j){
          int id = lower_bound(st[p].begin(), st[p].end
            (), x) - st[p].begin();
          return int(st[p].size()) - id;
       int mid = (L+R)/2;
       return query(2*p, L, mid, i, j, x) +
          query(2*p+1, mid+1, R, i, j, x);
   }
public:
   MergeTree(int sz, const int v[]): n(sz), st(4*sz){
      build(1, 1, n, v);
   //number of elements >= x on segment [i, j]
   int query(int i, int j, int x) const{
       if(i > j) swap(i, j);
       return query(1, 1, n, i, j, x);
   }
};
Wavelet Tree
template<typename T>
```

```
class wavelet{
   T L, R;
   vector<int> 1;
    vector<T> sum; // <<</pre>
```

```
wavelet *lef, *rig;
   int r(int i) const{ return i - l[i]; }
public:
   template<typename ITER>
   wavelet(ITER bg, ITER en){
       lef = rig = nullptr;
       L = *bg, R = *bg;
       for(auto it = bg; it != en; it++)
          L = min(L, *it), R = max(R, *it);
       if(L == R) return;
       T mid = L + (R - L)/2;
       1.reserve(std::distance(bg, en) + 1);
       sum.reserve(std::distance(bg, en) + 1);
       1.push_back(0), sum.push_back(0);
       for(auto it = bg; it != en; it++)
          1.push_back(1.back() + (*it <= mid)),</pre>
          sum.push_back(sum.back() + *it);
       auto tmp = stable_partition(bg, en, [mid](T x){
          return x <= mid;</pre>
       });
       if(bg != tmp) lef = new wavelet(bg, tmp);
       if(tmp != en) rig = new wavelet(tmp, en);
   }
    ~wavelet(){
       delete lef;
       delete rig;
   }
   // 1 index, first is 1st
   T kth(int i, int j, int k) const{
       if(L >= R) return L;
       int c = l[j] - l[i-1];
       if(c \ge k) return lef \ge kth(l[i-1]+1, l[j], k);
       else return rig->kth(r(i-1)+1, r(j), k - c);
   }
   // # elements > x on [i, j]
   int cnt(int i, int j, T x) const{
       if(L > x) return j - i + 1;
       if(R <= x || L == R) return 0;</pre>
       int ans = 0;
       if(lef) ans += lef->cnt(l[i-1]+1, l[j], x);
       if(rig) ans += rig->cnt(r(i-1)+1, r(j), x);
       return ans;
   }
   // sum of elements <= k on [i, j]</pre>
   T sumk(int i, int j, T k){
       if(L == R) return R <= k ? L * (j - i + 1) : 0;
       if(R <= k) return sum[j] - sum[i-1];</pre>
       int ans = 0;
       if(lef) ans += lef->sumk(l[i-1]+1, l[j], k);
       if(rig) ans += rig->sumk(r(i-1)+1, r(j), k);
       return ans;
   // swap (i, i+1) just need to update "array" l[i]
};
```

University of Brasilia Data Structures, 3

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
                                                            11 binary_search(const vector<Line> &cht, 11 x){
                                                               int L = 0, R = (int)cht.size()-2;
#include <ext/pb_ds/detail/standard_policies.hpp>
                                                               int bans = (int)cht.size()-1;
                                                               while(L <= R){</pre>
using namespace __gnu_pbds; // or pb_ds;
                                                                   int mid = (L+R)/2;
                                                                   if(cht[mid].eval(x) >= cht[mid+1].eval(x)) // <<<</pre>
template<typename T, typename B = null_type>
                                                                      L = mid + 1;
using oset = tree<T, B, less<T>, rb_tree_tag,
                                                                   else bans = mid, R = mid - 1;
  tree_order_statistics_node_update>;
// find_by_order / order_of_key
                                                               return cht[bans].eval(x);
Hash table
                                                            Convex Hull Trick
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
                                                            * Author: Simon Lindholm
struct custom_hash {
                                                            * source: https://github.com/kth-competitive-
   static uint64_t splitmix64(uint64_t x) {
                                                               programming/kactl/blob/master/content/data-structures
       // http://xorshift.di.unimi.it/splitmix64.c
                                                               /LineContainer.h
                                                             * License: CC0
      x += 0x9e3779b97f4a7c15;
      x = (x \hat{ } (x >> 30)) * 0xbf58476d1ce4e5b9;
      x = (x \hat{ } (x >> 27)) * 0x94d049bb133111eb;
      return x \hat{ } (x >> 31);
                                                            struct Line {
                                                               mutable 11 m, b, p;
   }
                                                               bool operator<(const Line& o) const { return m < o.m</pre>
   size_t operator()(uint64_t x) const {
                                                                 ; }
       static const uint64_t FIXED_RANDOM = chrono::
                                                               bool operator<(ll x) const { return p < x; }</pre>
         steady_clock::now().time_since_epoch().count();
                                                            };
       return splitmix64(x + FIXED_RANDOM);
   }
                                                            struct LineContainer : multiset<Line, less<>> { // CPP14
};
                                                               // (for doubles, use inf = 1/.0, div(a,b) = a/b)
gp_hash_table<long long, int, custom_hash> table;
                                                               const 11 inf = LLONG_MAX;
unordered_map<long long, int, custom_hash> uhash;
                                                               ll div(ll a, ll b) { // floored division
                                                                   return a / b - ((a ^ b) < 0 && a % b); }
uhash.reserve(1 << 15);</pre>
uhash.max_load_factor(0.25);
                                                               bool isect(iterator x, iterator y) {
                                                                   if (y == end()) { x->p = inf; return false; }
Convex Hull Trick Simple
                                                                   if (x->m == y->m) x->p = x->b > y->b? inf: -inf
struct Line{
                                                                   else x->p = div(y->b - x->b, x->m - y->m);
   11 m, b;
                                                                   return x->p >= y->p;
   inline ll eval(ll x) const{
                                                               }
      return x * m + b;
                                                               void add(ll m, ll b) {
   }
                                                                   auto z = insert(\{m, b, 0\}), y = z++, x = y;
};
                                                                   while (isect(y, z)) z = erase(z);
                                                                   if (x != begin() \&\& isect(--x, y)) isect(x, y =
// min => cht.back().m >= L.m
                                                                     erase(y));
// max => cht.back().m <= L.m
                                                                   while ((y = x) != begin() && (--x)->p >= y->p)
void push_line(vector<Line> &cht, Line L){
                                                                      isect(x, erase(y));
 while((int)cht.size() >= 2){
   int sz = (int)cht.size();
                                                               11 query(ll x) {
   if((long double)(L.b-cht[sz-1].b)*(cht[sz-2].m-L.m)
                                                                   assert(!empty());
  <= (long double)(L.b-cht[sz-2].b)*(cht[sz-1].m-L.m)){</pre>
                                                                   auto 1 = *lower_bound(x);
     cht.pop_back();
                                                                   return 1.m * x + 1.b;
   }
                                                               }
   else break;
                                                            };
 cht.push_back(L);
                                                            Min queue
                                                            template<typename T>
// x increasing; pos = 0 in first call
                                                            class minQ{
11 linear_search(const vector<Line> &cht,ll x,int &pos){
                                                               deque<tuple<T, int, int> > p;
   while(pos+1 < (int)cht.size()){</pre>
                                                               T delta:
 *>>*/ if(cht[pos].eval(x) >= cht[pos+1].eval(x)) pos++;
                                                               int sz:
                                                            public:
      else break;
   }
                                                               minQ() : delta(0), sz(0) {}
   return cht[pos].eval(x);
                                                               inline int size() const{ return sz; }
                                                               inline void add(T x){ delta += x; }
}
```

University of Brasilia Data Structures, 4

```
inline void push(T x, int id){
                                                               if(Y[1] > Y[r]) { R[1] = merge(R[1], r); u = 1; }
      x \rightarrow delta, sz++;
                                                                else { L[r] = merge(1, L[r]); u = r; }
       int t = 1:
                                                                calc(u); return u;
      while(p.size() > 0 \& get<0>(p.back()) >= x)
          t += get<1>(p.back()), p.pop_back();
                                                            void init(int n=N-1) { // XXX call before using other
      p.emplace_back(x, t, id);
                                                                for(int i = en = 1; i <= n; i++) { Y[i] = i; sz[i] =</pre>
   }
   inline void pop(){
                                                                   1; L[i] = R[i] = 0; }
                                                                random_shuffle(Y + 1, Y + n + 1);
       get<1>(p.front())--, sz--;
       if(!get<1>(p.front())) p.pop_front();
                                                            void insert(int &u, int it){
   T getmin() const{ return get<0>(p.front())+delta; }
                                                               unlaze(u);
   int getid() const{ return get<2>(p.front()); }
                                                                if(!u) u = it;
                                                                else if(Y[it] > Y[u]) split_val(u, X[it], L[it], R[
                                                                  it]), u = it;
Sparse Table
                                                                else insert(X[it] < X[u] ? L[u] : R[u], it);</pre>
                                                                calc(u);
const int N = 100005;
                                                            void erase(int &u, num key){
int v[N], n;
                                                               unlaze(u);
int dn[N][20];
                                                                if(!u) return;
int fn(int i, int j){
                                                                if(X[u] == key) u = merge(L[u], R[u]);
   if(j == 0) return v[i];
                                                                else erase(key < X[u] ? L[u] : R[u], key);</pre>
   if(~dn[i][j]) return dn[i][j];
                                                                calc(u):
   return dn[i][j] = min(fn(i, j-1), fn(i + (1 << (j-1)))
     ), j-1));
                                                            int create_node(num key){
}
                                                               X[en] = key;
                                                                sz[en] = 1;
int lg(int x){ return 31 - __builtin_clz(x); }
                                                               L[en] = R[en] = 0;
                                                                return en++;
int getmn(int 1, int r) { // [1, r]
   int 1z = 1g(r - 1 + 1);
                                                            int query(int u, int 1, int r){//0 index
   return min(fn(1, lz), fn(r - (1 << lz) + 1, lz));
                                                                unlaze(u);
                                                                if(u! or r < 0 or 1 >= sz[u]) return
Treap
                                                                  identity_element;
                                                                if(1 <= 0 and r >= sz[u] - 1) return subt_data[u];
// source: https://github.com/victorsenam/caderno/blob/
                                                                int ans = query(L[u], 1, r);
 master/code/treap.cpp
                                                                if(1 \le sz[L[u]] and sz[L[u]] \le r)
//const int N = ; typedef int num;
                                                                   ans = max(ans, st[u]);
num X[N]; int en = 1, Y[N], sz[N], L[N], R[N];
                                                                ans = max(ans, query(R[u], l-sz[L[u]]-1, r-sz[L[u]])
void calc (int u) { // update node given children info
                                                                  ]]-1));
   if(!u) return;
                                                                return ans;
   sz[u] = sz[L[u]] + 1 + sz[R[u]];
   // code here, no recursion
                                                            ColorUpdate
void unlaze (int u) {
   if(!u) return;
                                                            // source: https://github.com/tfg50/Competitive-
   // code here, no recursion
                                                              Programming/tree/master/Biblioteca/Data%20Structures
void split_val(int u, num x, int &l, int &r) { // l gets
                                                            #include <set>
                                                            #include <vector>
   <= x, r gets > x
   unlaze(u); if(!u) return (void) (1 = r = 0);
                                                            template <class Info = int>
   if(X[u] <= x) { split_val(R[u], x, 1, r); R[u] = 1;</pre>
     1 = u; }
                                                            class ColorUpdate {
   else { split_val(L[u], x, 1, r); L[u] = r; r = u; }
                                                            public:
   calc(u);
                                                                struct Range {
                                                                   Range(int l = 0) { this->l = 1; }
void split_sz(int u, int s, int &l, int &r) { // l gets
                                                                   Range(int 1, int r, Info v) {
  first s, r gets remaining
                                                                      this \rightarrow l = 1;
   unlaze(u); if(!u) return (void) (1 = r = 0);
                                                                       this->r = r;
   if(sz[L[u]] < s)  { split_sz(R[u], s - sz[L[u]] - 1,
                                                                       this -> v = v;
     1, r); R[u] = 1; 1 = u; }
                                                                   }
   else { split_sz(L[u], s, l, r); L[u] = r; r = u; }
                                                                   int 1, r;
   calc(u);
                                                                   Info v;
int merge(int 1, int r) { // els on l <= els on r</pre>
                                                                   bool operator < (const Range &b) const { return l</pre>
                                                                      < b.1; }
   unlaze(1); unlaze(r); if(!1 || !r) return 1 + r; int
                                                                };
```

University of Brasilia Data Structures, 5

```
std::vector<Range> upd(int 1, int r, Info v) {
                                                           // subtree of u => [ in[u], out[u] )
       std::vector<Range> ans;
                                                           // path from nxt[u] to u => [ in[ nxt[u] ], in[u] ]
       if(1 >= r) return ans;
                                                           Iterative Segtree
       auto it = ranges.lower_bound(1);
       if(it != ranges.begin()) {
                                                           T query(int 1, int r, int &pos){ // [1, r]
          it--;
                                                              T rl, rr;
          if(it->r>1) {
                                                               for (1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1)
              auto cur = *it;
                                                                  if(l & 1) rl = merge(rl, st[l++]);
              ranges.erase(it);
                                                                  if(r & 1) rr = merge(st[--r], rr);
              ranges.insert(Range(cur.1, 1, cur.v));
              ranges.insert(Range(1, cur.r, cur.v));
                                                              return merge(rl, rr);
      it = ranges.lower_bound(r);
                                                           // initially save v[i] in st[n+i] for all i in [0, n)
       if(it != ranges.begin()) {
                                                           void build(){
          it--:
                                                               for(int p = n-1; p > 0; p--)
          if(it->r > r) {
                                                                  st[p] = merge(st[2*p], st[2*p+1]);
              auto cur = *it;
                                                           }
              ranges.erase(it);
              ranges.insert(Range(cur.1, r, cur.v));
                                                           void update(int p, T val){
              ranges.insert(Range(r, cur.r, cur.v));
                                                               st[p += n] = val;
                                                               while(p >>= 1) st[p] = merge(st[2*p], st[2*p+1]);
       for(it = ranges.lower_bound(1); it != ranges.end
                                                           LiChao's Segtree
         () && it->l < r; it++) {
          ans.push_back(*it);
                                                           void add_line(line nw, int v = 1, int l = 0, int r =
                                                             maxn) { // [1, r)}
      ranges.erase(ranges.lower_bound(1), ranges.
                                                               int m = (1 + r) / 2;
         lower_bound(r));
                                                               bool lef = nw.eval(1) < st[v].eval(1);</pre>
       ranges.insert(Range(1, r, v));
                                                               bool mid = nw.eval(m) < st[v].eval(m);</pre>
       return ans;
                                                               if(mid) swap(st[v], nw);
   }
                                                               if(r - 1 == 1) {
private:
                                                                  return;
   std::set<Range> ranges;
                                                               } else if(lef != mid) {
                                                                  add_line(nw, 2 * v, 1, m);
Heavy Light Decomposition
                                                               } else {
                                                                  add_line(nw, 2 * v + 1, m, r);
void dfs_sz(int u){
                                                           }
   sz[u] = 1;
                                                           int get(int x, int v = 1, int l = 0, int r = maxn) {
   for(auto &v : g[u]) if(v == p[u]){
                                                               int m = (1 + r) / 2;
       swap(v, g[u].back());
                                                               if(r - 1 == 1) {
       g[u].pop_back();
                                                                  return st[v].eval(x);
       break;
                                                               else if(x < m) {
                                                                  return min(st[v].eval(x), get(x, 2*v, 1, m));
   for(auto &v : g[u]){
                                                                  return min(st[v].eval(x), get(x, 2*v+1, m, r));
      p[v] = u;
      dfs_sz(v);
       sz[u] += sz[v];
      if(sz[v] > sz[ g[u][0] ])
                                                           Palindromic tree
          swap(v, g[u][0]);
                                                           #include <bits/stdc++.h>
   }
}
                                                           using namespace std;
// nxt[u] = start of path with u
// set nxt[root] = root beforehand
                                                           const int maxn = 3e5 + 1, sigma = 26;
void dfs_hld(int u){
                                                           int len[maxn], link[maxn], to[maxn][sigma];
                                                           int slink[maxn], diff[maxn], series_ans[maxn];
   in[u] = t++;
   rin[in[u]] = u;
                                                           int sz, last, n;
                                                           char s[maxn];
   for(auto v : g[u]){
      nxt[v] = (v == g[u][0] ? nxt[u] : v);
      dfs_hld(v);
                                                           void init()
                                                           {
   out[u] = t;
                                                               s[n++] = -1;
                                                               link[0] = 1;
}
```

```
len[1] = -1;
   sz = 2;
}
int get_link(int v)
   while(s[n - len[v] - 2] != s[n - 1]) v = link[v];
   return v;
void add_letter(char c)
   s[n++] = c -= 'a';
   last = get_link(last);
   if(!to[last][c])
       len[sz] = len[last] + 2;
       link[sz] = to[get_link(link[last])][c];
       diff[sz] = len[sz] - len[link[sz]];
       if(diff[sz] == diff[link[sz]])
          slink[sz] = slink[link[sz]];
          slink[sz] = link[sz];
       to[last][c] = sz++;
   }
   last = to[last][c];
}
int main()
{
   ios::sync_with_stdio(0);
   cin.tie(0);
   init();
   string s;
   cin >> s;
   int n = s.size();
   int ans[n + 1];
   memset(ans, 63, sizeof(ans));
   ans [0] = 0;
   for(int i = 1; i <= n; i++)</pre>
   {
       add_letter(s[i - 1]);
       for(int v = last; len[v] > 0; v = slink[v])
          series_ans[v] = ans[i - (len[slink[v]] + diff
            [v])];
          if(diff[v] == diff[link[v]])
              series_ans[v] = min(series_ans[v],
                series_ans[link[v]]);
          ans[i] = min(ans[i], series_ans[v] + 1);
       cout << ans[i] << "\n";</pre>
   return 0;
```

# Math

# **Extended Euclidean Algorithm**

```
// a*x + b*y = gcd(a, b), < gcd, x, y>
tuple<int, int, int> gcd(int a, int b) {
   if(b == 0) return make_tuple(a, 1, 0);
   int q, w, e;
   tie(q, w, e) = gcd(b, a % b);
   return make_tuple(q, e, w - e * (a / b));
}
```

#### **Chinese Remainder Theorem**

```
// x = vet[i].first (mod vet[i].second)
11 crt(vector<pair<11, 11>> vet){
   11 ans = vet[0].first, lcm = vet[0].second;
   11 a, b, g, x, y;
   for(int i = 1; i < (int)vet.size(); i++){</pre>
       tie(a, b) = vet[i];
       tie(g, x, y) = gcd(lcm, b);
       ans = ans + x * (a - ans) / g % (b / g) * lcm;
       lcm = lcm * b / g;
       ans = (ans \% lcm + lcm) \% lcm;
   return ans;
Preffix inverse
inv[1] = 1;
for(int i = 2; i < p; i++)</pre>
   inv[i] = (p - (p/i) * inv[p%i] % p) % p;
Pollard Rho
11 rho(11 n){
   if(n \% 2 == 0) return 2;
   ll d, c, x, y;
   do{
       c = 11rand() \% n, x = 11rand() \% n, y = x;
          x = add(mul(x, x, n), c, n);
          y = add(mul(y, y, n), c, n);
          y = add(mul(y, y, n), c, n);
           d = \_gcd(abs(x - y), n);
       }while(d == 1);
   }while(d == n);
   return d;
}
11 pollard_rho(11 n){
   ll x, c, y, d, k;
   int i;
   do{
       i = 1;
       x = 11rand() % n, c = 11rand() % n;
       y = x, k = 4;
       do{
          if(++i == k) y = x, k *= 2;
          x = add(mul(x, x, n), c, n);
          d = \_gcd(abs(x - y), n);
       }while(d == 1);
   }while(d == n);
   return d;
void factorize(ll val, map<ll, int> &fac){
   if(rabin(val)) fac[ val ]++;
   else{
       11 d = pollard_rho(val);
       factorize(d, fac);
       factorize(val / d, fac);
   }
}
map<ll, int> factor(ll val){
```

```
map<ll, int> fac;
   if(val > 1) factorize(val, fac);
                                                               for(11 m = 2; ; m++) if(\_gcd(m, n) == 1)
   return fac;
                                                                   if(check(m)) return m;
                                                           }
Miller Rabin
                                                            // Let's denote R(n) as the set of primitive roots
                                                              modulo n, p is prime
bool rabin(ll n){
                                                            // g \ln R(p) => (pow(g, p-1, p * p) == 1 ? g+p : g) \ln r
   if(n <= 1) return 0;</pre>
                                                               R(pow(p, k)), for all k > 1
   if(n <= 3) return 1;
                                                            // g \text{ in } R(pow(p, k)) \Rightarrow (g \% 2 == 1 ? g : g + pow(p, k))
   11 s = 0, d = n - 1;
                                                               while(d % 2 == 0) d /= 2, s++;
   for(int k = 0; k < 64; k++){
                                                            Mobius Function
      11 a = (11rand() \% (n - 3)) + 2;
      11 x = fexp(a, d, n);
                                                            memset(mu, 0, sizeof mu);
       if(x != 1 \&\& x != n-1){
                                                            mu[1] = 1;
          for(int r = 1; r < s; r++){
                                                            for(int i = 1; i < N; i++)</pre>
              x = mul(x, x, n);
                                                               for(int j = i + i; j < N; j += i)
              if(x == 1) return 0;
                                                                  mu[j] -= mu[i];
              if(x == n-1) break;
                                                            // g(n) = sum{f(d)} => f(n) = sum{mu(d)*g(n/d)}
                                                            Mulmod TOP
          if(x != n-1) return 0;
                                                            constexpr uint64_t mod = (1ull<<61) - 1;</pre>
   }
                                                            uint64_t modmul(uint64_t a, uint64_t b){
   return 1;
                                                               uint64_t 11 = (uint32_t)a, h1 = a>>32, 12 = (
                                                                 uint32_t)b, h2 = b>>32;
Totiente
                                                               uint64_t 1 = 11*12, m = 11*h2 + 12*h1, h = h1*h2;
                                                               uint64_t ret = (1&mod) + (1>>61) + (h << 3) + (m >>
11 totiente(11 n){
                                                                 29) + (m \ll 35 \gg 3) + 1;
   11 \text{ ans} = n;
                                                               ret = (ret & mod) + (ret>>61);
   for(ll i = 2; i*i <= n; i++){</pre>
                                                               ret = (ret & mod) + (ret>>61);
       if(n \% i == 0){
                                                               return ret-1;
          ans = ans / i * (i - 1);
          while(n \% i == 0) n /= i;
                                                            Matrix Determinant
   }
                                                            int n:
                                                            long double a[n][n];
   if(n > 1) ans = ans / n * (n - 1);
   return ans:
                                                            long double gauss(){
                                                               long double det = 1;
Primitive root
                                                               for(int i = 0; i < n; i++){
                                                                   int q = i;
// a primitive root modulo n is any number g such that
                                                                   for(int j = i+1; j < n; j++){
  any c coprime to n is congruent to a power of g modulo
                                                                      if(abs(a[j][i]) > abs(a[q][i]))
  n.
                                                                          q = j;
bool exists_root(ll n){
                                                                   if(abs(a[q][i]) < EPS){
   if(n == 1 || n == 2 || n == 4) return true;
                                                                      det = 0;
   if(n \% 2 == 0) n /= 2;
                                                                      break;
   if(n % 2 == 0) return false;
                                                                   if(i != q){
   // test if n is a power of only one prime
                                                                      for(int w = 0; w < n; w++)
   for(11 i = 3; i * i <= n; i += 2) if(n % i == 0){
                                                                          swap(a[i][w], a[q][w]);
      while(n \% i == 0) n /= i;
                                                                      det = -det;
      return n == 1;
                                                                   det *= a[i][i];
   return true;
                                                                   for(int j = i+1; j < n; j++) a[i][j] /= a[i][i];</pre>
                                                                   for(int j = 0; j < n; j++) if(j != i){
11 primitive_root(ll n){
                                                                      if(abs(a[j][i]) > EPS)
   if(n == 1 || n == 2 || n == 4) return n - 1;
                                                                          for(int k = i+1; k < n; k++)
   if(not exists_root(n)) return -1;
                                                                             a[j][k] = a[i][k] * a[j][i];
   11 x = phi(n);
                                                                   }
   auto pr = factorize(x);
                                                               }
   auto check = [x, n, pr](ll m){
       for(ll p : pr) if(fexp(m, x / p, n) == 1)
                                                               return det;
          return false;
                                                           }
       return true;
```

# **Simplex Method**

```
typedef long double dbl;
const dbl eps = 1e-6;
const int N = , M = ;
mt19937 rng(chrono::steady_clock::now().time_since_epoch
  ().count()):
struct simplex {
   int X[N], Y[M];
   dbl A[M][N], b[M], c[N];
   dbl ans;
   int n, m;
   dbl sol[N];
   void pivot(int x, int y){
       swap(X[y], Y[x]);
      b[x] /= A[x][y];
       for(int i = 0; i < n; i++)</pre>
          if(i != y)
              A[x][i] /= A[x][y];
      A[x][y] = 1. / A[x][y];
       for(int i = 0; i < m; i++)</pre>
          if(i != x && abs(A[i][y]) > eps) {
              b[i] -= A[i][y] * b[x];
              for(int j = 0; j < n; j++) if(j != y)
                 A[i][j] -= A[i][y] * A[x][j];
              A[i][y] = -A[i][y] * A[x][y];
          }
       ans += c[y] * b[x];
       for(int i = 0; i < n; i++)</pre>
          if(i != y)
              c[i] -= c[y] * A[x][i];
      c[y] = -c[y] * A[x][y];
   }
   // maximiza sum(x[i] * c[i])
   // sujeito a
   // sum(a[i][j] * x[j]) <= b[i] para 0 <= i < m (Ax)
   // x[i] >= 0 para 0 <= i < n (x >= 0)
   // (n variaveis, m restricoes)
   // guarda a resposta em ans e retorna o valor otimo
   dbl solve(int _n, int _m) {
       this->n = _n; this->m = _m;
       for(int i = 1; i < m; i++){
          int id = uniform_int_distribution<int>(0, i)(
            rng);
          swap(b[i], b[id]);
          for(int j = 0; j < n; j++)
              swap(A[i][j], A[id][j]);
      }
       ans = 0.;
       for(int i = 0; i < n; i++) X[i] = i;
       for(int i = 0; i < m; i++) Y[i] = i + n;
          int x = min_element(b, b + m) - b;
          if(b[x] >= -eps)
             break:
          int y = find_if(A[x], A[x] + n, [](dbl d) {
            return d < -eps; }) - A[x];</pre>
          if(y == n) throw 1; // no solution
          pivot(x, y);
      while(true) {
          int y = max_element(c, c + n) - c;
```

```
if(c[y] <= eps) break;</pre>
          int x = -1;
          dbl mn = 1. / 0.;
          for(int i = 0; i < m; i++)
              if(A[i][y] > eps && b[i] / A[i][y] < mn)</pre>
                 mn = b[i] / A[i][y], x = i;
          if(x == -1) throw 2; // unbounded
          pivot(x, y);
       memset(sol, 0, sizeof(dbl) * n);
       for(int i = 0; i < m; i++)
          if(Y[i] < n)
              sol[Y[i]] = b[i];
       return ans;
};
FFT
struct base{
   double r, i;
   base(double _r = 0, double _i = 0) : r(_r), i(_i) {}
   base operator*(base &o) const{
       return {r*o.r - i*o.i, r*o.i + o.r*i};
   double real() const{ return r; }
   void operator*=(const base &o){
       (*this) = \{r*o.r-i*o.i, r*o.i+o.r*i\};
   }
   void operator+=(const base &o){r += o.r, i += o.i; }
   void operator/=(const double &o){ r /= o, i /= o; }
   void operator==(const base &o){r == o.r, i == o.i; }
   base operator+(const base &o){return {r+o.r,i+o.i};}
   base operator-(const base &o){return {r-o.r,i-o.i};}
};
double PI = acos(-1);
void fft(vector<base> &a, bool inv){
   int n = (int)a.size();
   for(int i = 1, j = 0; i < n; i++){
       int bit = n \gg 1;
       for(; j >= bit; bit >>= 1) j -= bit;
       j += bit;
       if(i < j) swap(a[i], a[j]);
   for(int sz = 2; sz <= n; sz <<= 1) {
       double ang = 2*PI/sz * (inv ? -1 : 1);
       base wlen(cos(ang), sin(ang));
       for(int i = 0; i < n; i += sz){
          base w(1);
          for(int j = 0; j < sz/2; j++){
              base u = a[i+j], v = a[i+j+sz/2] * w;
              a[i+j] = u + v;
              a[i+j+sz/2] = u - v;
              w *= wlen;
          }
       }
   if(inv) for(int i = 0; i < n; i++) a[i] /= 1.0 * n;
}
void multiply(const vector<int> &a, const vector<int> &b
  , vector<int> &res){
   vector<base> fa(a.begin(), a.end());
   vector<base> fb(b.begin(), b.end());
```

```
size_t n = 1;
   while(n < a.size()) n <<= 1;</pre>
   while(n < b.size()) n <<= 1;</pre>
   n <<= 1:
   fa.resize(n), fb.resize(n);
   fft(fa, false), fft(fb, false);
   for(size_t i = 0; i < n; i++)</pre>
       fa[i] *= fb[i];
   fft(fa, true);
   res.resize (n);
   for(size_t i = 0; i < n; ++i)</pre>
       res[i] = int(fa[i].real() + 0.5);
FFT
namespace fft {
 typedef double dbl;
 struct num {
   dbl x, y;
   num() \{ x = y = 0; \}
   num(dbl x, dbl y) : x(x), y(y) {}
 inline num operator+(num a, num b) { return num(a.x +
   b.x, a.y + b.y); }
 inline num operator-(num a, num b) { return num(a.x -
   b.x, a.y - b.y); }
 inline num operator*(num a, num b) { return num(a.x *
   b.x - a.y * b.y, a.x * b.y + a.y * b.x); }
 inline num conj(num a) { return num(a.x, -a.y); }
 int base = 1;
 vector<num> roots = \{\{0, 0\}, \{1, 0\}\};
 vector < int > rev = \{0, 1\};
 const dbl PI = acosl(-1.0);
 void ensure_base(int nbase) {
   if(nbase <= base) return;</pre>
   rev.resize(1 << nbase);</pre>
   for(int i = 0; i < (1 << nbase); i++) {</pre>
     rev[i] = (rev[i >> 1] >> 1) + ((i \& 1) << (nbase -
   roots.resize(1 << nbase);</pre>
   while(base < nbase) {</pre>
     dbl angle = 2*PI / (1 << (base + 1));
     for(int i = 1 << (base - 1); i < (1 << base); i++)</pre>
       {
       roots[i << 1] = roots[i];</pre>
       dbl \ angle_i = angle * (2 * i + 1 - (1 << base));
       roots[(i << 1) + 1] = num(cos(angle_i), sin(
         angle_i));
     base++;
   }
 }
 void fft(vector<num> &a, int n = -1) {
   if(n == -1) {
     n = a.size();
```

```
assert((n & (n-1)) == 0);
  int zeros = __builtin_ctz(n);
  ensure_base(zeros);
 int shift = base - zeros;
  for(int i = 0; i < n; i++) {</pre>
   if(i < (rev[i] >> shift)) {
     swap(a[i], a[rev[i] >> shift]);
  for(int k = 1; k < n; k <<= 1) {
   for(int i = 0; i < n; i += 2 * k) {</pre>
     for(int j = 0; j < k; j++) {
       num z = a[i+j+k] * roots[j+k];
       a[i+j+k] = a[i+j] - z;
       a[i+j] = a[i+j] + z;
   }
 }
}
vector<num> fa, fb;
vector<int> multiply(vector<int> &a, vector<int> &b) {
  int need = a.size() + b.size() - 1;
  int nbase = 0;
 while((1 << nbase) < need) nbase++;</pre>
  ensure_base(nbase);
  int sz = 1 << nbase;</pre>
  if(sz > (int) fa.size()) {
   fa.resize(sz);
  for(int i = 0; i < sz; i++) {
   int x = (i < (int) a.size() ? a[i] : 0);</pre>
   int y = (i < (int) b.size() ? b[i] : 0);</pre>
   fa[i] = num(x, y);
 fft(fa, sz);
 num r(0, -0.25 / sz);
  for(int i = 0; i <= (sz >> 1); i++) {
   int j = (sz - i) & (sz - 1);
   num z = (fa[j] * fa[j] - conj(fa[i] * fa[i])) * r;
   if(i != j) {
     fa[j] = (fa[i] * fa[i] - conj(fa[j] * fa[j])) * r
   fa[i] = z;
  fft(fa, sz);
 vector<int> res(need);
  for(int i = 0; i < need; i++) {</pre>
   res[i] = fa[i].x + 0.5;
 return res;
vector<int> multiply_mod(vector<int> &a, vector<int> &
  b, int m, int eq = 0) {
  int need = a.size() + b.size() - 1;
 int nbase = 0:
 while ((1 << nbase) < need) nbase++;</pre>
  ensure base(nbase):
  int sz = 1 << nbase;</pre>
  if (sz > (int) fa.size()) {
   fa.resize(sz);
  for (int i = 0; i < (int) a.size(); i++) {</pre>
   int x = (a[i] \% m + m) \% m;
   fa[i] = num(x & ((1 << 15) - 1), x >> 15);
```

```
for (int i=1, j=0; i<n; ++i) {</pre>
   fill(fa.begin() + a.size(), fa.begin() + sz, num {0,
                                                                    int bit = n \gg 1;
      0}):
                                                                    for (; j>=bit; bit>>=1)
   fft(fa, sz);
                                                                       j -= bit;
   if (sz > (int) fb.size()) {
                                                                    i += bit;
     fb.resize(sz);
                                                                    if (i < j)
                                                                        swap (a[i], a[j]);
   if (eq) {
                                                                }
     copy(fa.begin(), fa.begin() + sz, fb.begin());
                                                                for (int len=2; len<=n; len<<=1) {</pre>
   } else {
     for (int i = 0; i < (int) b.size(); i++) {</pre>
                                                                    int wlen = invert ? root_1 : root;
       int x = (b[i] \% m + m) \% m;
                                                                    for (int i=len; i<root_pw; i<<=1)</pre>
                                                                        wlen = int (wlen * 111 * wlen % mod);
       fb[i] = num(x & ((1 << 15) - 1), x >> 15);
                                                                    for (int i=0; i<n; i+=len) {</pre>
     fill(fb.begin() + b.size(), fb.begin() + sz, num
                                                                        int w = 1;
                                                                        for (int j=0; j<len/2; ++j) {</pre>
       \{0.0\}\}:
     fft(fb, sz);
                                                                           int u = a[i+j], v = int (a[i+j+len/2] * 1
                                                                             11 * w % mod);
                                                                           a[i+j] = u+v < mod ? u+v : u+v-mod;
   dbl ratio = 0.25 / sz;
                                                                           a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
   num r2(0, -1);
   num r3(ratio, 0);
                                                                           w = int (w * 111 * wlen % mod);
   num r4(0, -ratio);
                                                                       }
                                                                    }
   num r5(0, 1);
   for (int i = 0; i \le (sz >> 1); i++) {
                                                                if (invert) {
     int j = (sz - i) & (sz - 1);
     num a1 = (fa[i] + conj(fa[j]));
                                                                    int nrev = reverse (n, mod);
     num a2 = (fa[i] - conj(fa[j])) * r2;
                                                                    for (int i=0; i<n; ++i)</pre>
     num b1 = (fb[i] + conj(fb[j])) * r3;
                                                                        a[i] = int (a[i] * 111 * nrev % mod);
     num b2 = (fb[i] - conj(fb[j])) * r4;
                                                                }
     if (i != j) {
                                                            }
       num c1 = (fa[j] + conj(fa[i]));
                                                             Gauss
       num c2 = (fa[j] - conj(fa[i])) * r2;
       num d1 = (fb[j] + conj(fb[i])) * r3;
       num d2 = (fb[j] - conj(fb[i])) * r4;
                                                             // Solves systems of linear equations.
       fa[i] = c1 * d1 + c2 * d2 * r5;
                                                             // To use, build a matrix of coefficients and call run(
       fb[i] = c1 * d2 + c2 * d1;
                                                               mat, R, C). If the i-th variable is free, row[i] will
     }
                                                               be -1, otherwise it's value will be ans[i].
     fa[j] = a1 * b1 + a2 * b2 * r5;
     fb[j] = a1 * b2 + a2 * b1;
                                                            namespace Gauss {
                                                              const int MAXC = 1001;
   fft(fa, sz);
                                                              int row[MAXC];
   fft(fb, sz);
                                                              double ans[MAXC];
   vector<int> res(need);
   for (int i = 0; i < need; i++) {</pre>
                                                              void run(double mat[][MAXC], int R, int C) {
     long long aa = fa[i].x + 0.5;
                                                                REP(i, C) row[i] = -1;
     long long bb = fb[i].x + 0.5;
     long long cc = fa[i].y + 0.5;
                                                                int r = 0;
     res[i] = (aa + ((bb \% m) << 15) + ((cc \% m) << 30))
                                                                REP(c, C) {
                                                                  int k = r;
   }
                                                                  FOR(i, r, R) if(fabs(mat[i][c]) > fabs(mat[k][c]))
   return res;
                                                                  if(fabs(mat[k][c]) < eps) continue;</pre>
 vector<int> square_mod(vector<int> &a, int m) {
                                                                  REP(j, C+1) swap(mat[r][j], mat[k][j]);
   return multiply_mod(a, a, m, 1);
                                                                  REP(i, R) if (i != r) {
                                                                    double w = mat[i][c] / mat[r][c];
                                                                    REP(j, C+1) mat[i][j] -= mat[r][j] * w;
NTT
                                                                  row[c] = r++;
const int mod = 7340033;
const int root = 5;
                                                                REP(i, C) {
const int root_1 = 4404020;
                                                                  int r = row[i];
const int root_pw = 1<<20;</pre>
                                                                  ans[i] = r == -1 ? 0 : mat[r][C] / mat[r][i];
void fft (vector<int> & a, bool invert) {
                                                              }
   int n = (int) a.size();
                                                            }
```

```
Gauss Xor
const 11 MAX = 1e9;
const 11 LOG_MAX = 64 - __builtin_clzll((11)MAX);
struct gauss{
   vector<11> vet;
   gauss(){}
   gauss(ll val){
       if(val) vet.push_back(val);
   void add(ll val){
       int sig = LOG_MAX;
       for(int i = 0; i < (int)vet.size(); i++){</pre>
          while(!(vet[i] & (1LL << sig))) sig--;</pre>
          if(val & (1LL << sig)) val ^= vet[i];</pre>
       if(!val) return;
       sig = LOG_MAX;
       while(!(val & (1 << sig))) sig--;</pre>
       for(auto &x : vet) if(x & (1LL << sig)) x ^= val;</pre>
       vet.push_back(val);
       for(int i = (int)vet.size() - 2; i >= 0 && vet[i]
          < vet[i + 1]; i--){
          swap(vet[i], vet[i + 1]);
       }
   }
};
Simpson
inline double simpson(double fl,double fr,double fmid,
  double 1,double r) { return (fl+fr+4.0*fmid)*(r-1)
double rsimpson(double slr, double fl, double fr, double
  fmid,double 1,double r)
{
   double mid = (1+r)*0.5;
   double fml = f((1+mid)*0.5);
   double fmr = f((mid+r)*0.5);
   double slm = simpson(fl,fmid,fml,l,mid);
   double smr = simpson(fmid, fr, fmr, mid, r);
   if(fabs(slr-slm-smr) < eps) return slm+smr;</pre>
   return rsimpson(slm,fl,fmid,fml,l,mid)+rsimpson(smr,
     fmid, fr, fmr, mid, r);
double integrate(double 1,double r) {
   double mid = (1+r)*0.5;
   double fl = f(1);
   double fr = f(r);
   double fmid = f(mid);
   return rsimpson(simpson(fl,fr,fmid,l,r),fl,fr,fmid,l
      ,r);
Graphs
Dinic
const int N = 100005;
const int E = 2000006;
vector<int> g[N];
int ne:
struct Edge{
   int from, to;
   11 flow, cap;
```

} edge[E];

```
int lvl[N], vis[N], pass, start = N-2, target = N-1;
int qu[N], qt, px[N];
11 run(int s, int sink, ll minE){
   if(s == sink) return minE;
   11 \text{ ans} = 0;
   for(; px[s] < (int)g[s].size(); px[s]++){</pre>
       int e = g[s][ px[s] ];
       auto &v = edge[e], &rev = edge[e^1];
       if(lvl[v.to] != lvl[s]+1 || v.flow >= v.cap)
           continue; // v.cap - v.flow < lim</pre>
       11 tmp = run(v.to, sink,min(minE, v.cap-v.flow));
       v.flow += tmp, rev.flow -= tmp;
       ans += tmp, minE -= tmp;
       if(minE == 0) break;
   }
   return ans;
bool bfs(int source, int sink){
   qt = 0;
   qu[qt++] = source;
   lvl[source] = 1;
   vis[source] = ++pass;
   for(int i = 0; i < qt; i++){</pre>
       int u = qu[i];
       px[u] = 0;
       if(u == sink) return true;
       for(int e : g[u]){
           auto v = edge[e];
           if(v.flow >= v.cap || vis[v.to] == pass)
              continue; // v.cap - v.flow < lim</pre>
          vis[v.to] = pass;
          lvl[v.to] = lvl[u]+1;
           qu[qt++] = v.to;
       }
   }
   return false;
11 flow(int source = start, int sink = target){
   11 \text{ ans} = 0;
   //for(lim = (1LL << 62); lim >= 1; lim /= 2)
   while(bfs(source, sink))
       ans += run(source, sink, oo);
   return ans;
void addEdge(int u, int v, ll c = 1, ll rc = 0){
   edge[ne] = \{u, v, 0, c\};
   g[u].push_back(ne++);
   edge[ne] = \{v, u, 0, rc\};
   g[v].push_back(ne++);
}
void reset_flow(){
   for(int i = 0; i < ne; i++)
       edge[i].flow = 0;
Push relabel
// Push relabel in O(V^2 E^0.5) with gap heuristic
// It's quite fast
```

```
template<typename flow_t = long long>
                                                            struct Edge{
struct PushRelabel {
                                                                int from, to;
   struct Edge { int to, rev; flow_t f, c; };
                                                                11 cap, cost;
   vector<vector<Edge> > g;
                                                            } edge[E];
   vector<flow_t> ec;
   vector<Edge*> cur;
                                                            int start = N-1, target = N-2, p[N];
   vector<vector<int> > hs;
                                                            int inqueue[N];
   vector<int> H;
                                                            11 d[N];
   PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H
                                                            bool spfa(int source, int sink){
   void add_edge(int s, int t, flow_t cap, flow_t rcap
                                                                for(int i = 0; i < N; i++) d[i] = oo, inqueue[i] =</pre>
     =0) {
                                                                  0;
       if (s == t) return;
      Edge a = \{t, (int)g[t].size(), 0, cap\};
                                                                d[source] = 0;
      Edge b = \{s, (int)g[s].size(), 0, rcap\};
                                                                queue<int> q;
                                                                q.push(source);
      g[s].push_back(a);
      g[t].push_back(b);
                                                                inqueue[source] = 1;
   void add_flow(Edge& e, flow_t f) {
                                                                while(!q.empty()){
                                                                   int u = q.front(); q.pop();
      Edge &back = g[e.to][e.rev];
       if (!ec[e.to] && f)
                                                                   inqueue[u] = 0;
          hs[H[e.to]].push_back(e.to);
       e.f += f, ec[e.to] += f;
                                                                   for(int e : g[u]){
      back.f -= f, ec[back.to] -= f;
                                                                       auto v = edge[e];
                                                                       if(v.cap > 0 \text{ and } d[u] + v.cost < d[v.to]){
                                                                           d[v.to] = d[u] + v.cost;
   flow_t max_flow(int s, int t) {
                                                                           p[v.to] = e;
      int v = g.size();
      H[s] = v; ec[t] = 1;
                                                                           if(!inqueue[v.to]){
       vector<int> co(2 * v);
                                                                              q.push(v.to);
       co[0] = v-1;
                                                                              inqueue[v.to] = 1;
       for(int i = 0; i < v; ++i) cur[i] = g[i].data();</pre>
       for(auto &e : g[s]) add_flow(e, e.c);
                                                                       }
                                                                   }
       if(hs[0].size())
                                                                }
       for (int hi = 0; hi >= 0;) {
                                                                return d[sink] != oo;
          int u = hs[hi].back();
                                                            }
          hs[hi].pop_back();
          while (ec[u] > 0) // discharge u
                                                            // <max flow, min cost>
                                                            pair<11, 11> mincost(int source = start, int sink =
              if (cur[u] == g[u].data() + g[u].size()) {
                 H[u] = 1e9;
                                                              target){
                 for(auto &e:g[u])
                                                                11 ans = 0, mf = 0;
                     if (e.c - e.f \&\& H[u] > H[e.to]+1)
                                                                while(spfa(source, sink)){
                        H[u] = H[e.to]+1, cur[u] = &e;
                                                                   11 f = oo;
                 if (++co[H[u]], !--co[hi] && hi < v)</pre>
                     for(int i = 0; i < v; ++i)
                                                                   for(int u = sink; u != source; u = edge[ p[u] ].
                         if (hi < H[i] && H[i] < v){</pre>
                            --co[H[i]];
                                                                       f = min(f, edge[p[u]].cap);
                            H[i] = v + 1;
                                                                   for(int u = sink; u != source; u = edge[ p[u] ].
                 hi = H[u];
                                                                      from) {
              } else if (cur[u]->c - cur[u]->f \&\& H[u]
                                                                       edge[p[u]].cap -= f;
                == H[cur[u]->to]+1)
                                                                       edge[p[u] ^1].cap += f;
                 add_flow(*cur[u], min(ec[u], cur[u]->c
                                                                   }
                    - cur[u]->f));
              else ++cur[u];
                                                                   mf += f;
                                                                   ans += f * d[sink];
          while (hi \ge 0 \& hs[hi].empty()) --hi;
                                                                }
                                                                return {mf, ans};
      return -ec[s];
   }
                                                            }
}:
                                                            void addEdge(int u, int v, ll c, ll cost){
Min Cost Max Flow
                                                                edge[ne] = \{u, v, c, cost\};
                                                                g[u].push_back(ne++);
const 11 oo = 1e18;
                                                                edge[ne] = \{v, u, 0, -cost\};
const int N = 222, E = 2 * 1000006;
                                                                g[v].push_back(ne++);
                                                            }
vector<int> g[N];
int ne;
```

# **Small to Large**

```
void cnt_sz(int u, int p = -1){
   sz[u] = 1;
   for(int v : g[u]) if(v != p)
       cnt_sz(v, u), sz[u] += sz[v];
}
void add(int u, int p, int big = -1){
   // Update info about this vx in global answer
   for(int v : g[u]) if(v != p && v != big)
       add(v, u);
void dfs(int u, int p, int keep){
   int big = -1, mmx = -1;
   for(int v : g[u]) if(v != p \&\& sz[v] > mmx)
       mmx = sz[v], big = v;
   for(int v : g[u]) if(v != p && v != big)
       dfs(v, u, 0);
   if(big != -1) dfs(big, u, 1);
   add(u, p, big);
   for(auto x : q[u]){
       // answer all queries for this vx
   if(!keep){
       // Remove data from this subtree
```

# **Centroid Decomposition**

```
vector<pair<int,int>> G[500005];
int subtree[500005], k;
bool erased[500005];
ll dist[500005], ans;
int calc_sz(int v, int p){
   subtree[v] = 1;
   for(auto x : G[v]) if(!erased[x.ff] && x.ff != p){
       subtree[v] += calc_sz(x.ff,v);
   return subtree[v];
}
int centroid(int v, int p, int treesize){
   for(auto x : G[v]) if(!erased[x.ff] && x.ff != p){
      if(subtree[x.ff] * 2 > treesize)
          return centroid(x.ff, v, treesize);
   }
   return v;
}
void procurar_ans(int v, int p, int d_atual, ll custo){
   ans = min(ans, dist[k - d_atual] + custo);
   if(d_atual == k) return;
   for(auto x : G[v]) if(!erased[x.ff] && x.ff != p)
      procurar_ans(x.ff, v, d_atual+1, custo+x.ss);
}
```

```
void atualiza_dist(int v, int p, int d_atual, ll custo){
   dist[d_atual] = min(dist[d_atual], custo);
   if(d_atual == k) return;
   for(auto x : G[v]) if(!erased[x.ff] && x.ff != p)
       atualiza_dist(x.ff,v,d_atual+1,custo+x.ss);
}
void decomp(int v, int p){
   int treesize = calc_sz(v, v);
   if(treesize < k) return;</pre>
   int cent = centroid(v, v, treesize);
   erased[cent] = 1;
   for(int i = 1; i <= treesize; i++) dist[i] = 1e18;</pre>
   for(pair<int,int> x : G[cent]) if(!erased[x.ff]){
      procurar_ans(x.ff, cent, 1, x.ss);
       atualiza_dist(x.ff, cent, 1, x.ss);
   for(pair<int,int> x : G[cent]) if(!erased[x.ff])
       decomp(x.ff, cent);
Kosaraju
vector<int> g[N], gt[N], S;
int vis[N], cor[N];
void dfs(int u){
   vis[u] = 1;
   for(int v : g[u]) if(!vis[v]) dfs(v);
   S.push_back(u);
}
void dfst(int u, int e){
   cor[u] = e:
   for(int v : gt[u]) if(!cor[v]) dfst(v, e);
void kosaraju(){
   for(int i = 1; i <= n; i++) if(!vis[i]) dfs(i);</pre>
   for(int i = 1; i \le n; i++) for(int j : g[i])
       gt[j].push_back(i);
   int e = 0;
   reverse(S.begin(), S.end());
   for(int u : S) if(!cor[u])
       dfst(u, ++e);
Tarjan
int cnt = 0, root;
void dfs(int u, int p = -1){
   low[u] = num[u] = ++t;
   for(int v : g[u]){
       if(!num[v]){
          dfs(v, u);
              if(u == root) cnt++;
          if(low[v] >= num[u]) u PONTO DE ARTICULAÇÃO;
          if(low[v] > num[u]) ARESTA u->v PONTE;
          low[u] = min(low[u], low[v]);
       else if(v != p) low[u] = min(low[u], num[v]);
   }
}
root PONTO DE ARTICULAÇÃO <=> cnt > 1
void tarjanSCC(int u){
```

```
low[u] = num[u] = ++cnt;
                                                                S.push_back(u);
   vis[u] = 1;
                                                                id[u] = ++dfs_time;
                                                                label[u] = sdom[u] = dsu[u] = u;
   S.push_back(u);
   for(int v : g[u]){
       if(!num[v]) tarjanSCC(v);
                                                                for(int v : g[u]){
       if(vis[v]) low[u] = min(low[u], low[v]);
                                                                   if(!id[v])
                                                                       prep(v), down[u].push_back(v);
   if(low[u] == num[u]){
                                                                   gt[v].push_back(u);
       ssc[u] = ++ssc_cnt; int v;
                                                                }
                                                            }
          v = S.back(); S.pop_back(); vis[v] = 0;
          ssc[v] = ssc_cnt;
                                                            int fnd(int u, int flag = 0){
       }while(u != v);
                                                                if(u == dsu[u]) return u;
                                                                int v = fnd(dsu[u], 1), b = label[ dsu[u] ];
   }
                                                                if(id[ sdom[b] ] < id[ sdom[ label[u] ] ])</pre>
                                                                   label[u] = b;
Max Clique
                                                                dsu[u] = v;
                                                                return flag ? v : label[u];
long long adj[N], dp[N];
for(int i = 0; i < n; i++){
                                                            void build_dominator_tree(int root, int sz){
   for(int j = 0; j < n; j++){
                                                                // memset(id, 0, sizeof(int) * (sz + 1));
      int x;
                                                                // for(int i = 0; i <= sz; i++) T[i].clear();
       scanf("%d",&x);
                                                                prep(root);
       if(x \mid | i == j)
                                                                reverse(S.begin(), S.end());
          adj[i] |= 1LL << j;
   }
                                                                int w:
}
                                                                for(int u : S){
                                                                    for(int v : gt[u]){
int resto = n - n/2;
                                                                       w = fnd(v);
int C = n/2;
                                                                       if(id[ sdom[w] ] < id[ sdom[u] ])</pre>
for(int i = 1; i < (1 << resto); i++){</pre>
                                                                           sdom[u] = sdom[w];
   int x = i;
   for(int j = 0; j < resto; j++)
                                                                   gt[u].clear();
       if(i & (1 << j))</pre>
          x \&= adj[j + C] >> C;
                                                                   if(u != root) bucket[ sdom[u] ].push_back(u);
   if(x == i){
      dp[i] = __builtin_popcount(i);
                                                                    for(int v : bucket[u]){
   }
                                                                       w = fnd(v);
}
                                                                       if(sdom[w] == sdom[v]) idom[v] = sdom[v];
                                                                       else idom[v] = w;
for(int i = 1; i < (1 << resto); i++)</pre>
   for(int j = 0; j < resto; j++)
                                                                   bucket[u].clear();
      if(i & (1 << j))</pre>
          dp[i] = max(dp[i], dp[i ^ (1 << j)]);
                                                                    for(int v : down[u]) dsu[v] = u;
                                                                    down[u].clear();
int maxCliq = 0;
                                                                }
for(int i = 0; i < (1 << C); i++){
   int x = i, y = (1 << resto) - 1;
                                                                reverse(S.begin(), S.end());
   for(int j = 0; j < C; j++)
                                                                for(int u : S) if(u != root){
       if(i & (1 << j))</pre>
                                                                    if(idom[u] != sdom[u]) idom[u] = idom[ idom[u] ];
          x \&= adj[j] \& ((1 << C) - 1), y \&= adj[j] >>
                                                                   T[ idom[u] ].push_back(u);
            С;
   if(x != i) continue;
                                                                S.clear();
   maxCliq = max(maxCliq, __builtin_popcount(i) + dp[y
     1);
}
                                                            Min Cost Matching
Dominator Tree
                                                            // Min cost matching
vector<int> g[N], gt[N], T[N];
                                                            // O(n^2 * m)
vector<int> S;
                                                            // n == nro de linhas
int dsu[N], label[N];
                                                            // m == nro de colunas
int sdom[N], idom[N], dfs_time, id[N];
                                                            // n <= m | flow == n
                                                            // a[i][j] = custo pra conectar i a j
vector<int> bucket[N];
                                                            vector < int > u(n + 1), v(m + 1), p(m + 1), way(m + 1);
vector<int> down[N];
                                                            for(int i = 1; i <= n; ++i){</pre>
                                                                p[0] = i;
void prep(int u){
                                                                int j0 = 0;
```

University of Brasilia Strings, 15

```
vector<int> minv(m + 1 , oo);
   vector<char> used(m + 1 , false);
   do{
      used[j0] = true;
       int i0 = p[j0] , delta = oo, j1;
       for(int j = 1; j \le m; ++j)
          if(! used[j]){
              int cur = a[i0][j] - u[i0] - v[j];
              if(cur < minv[j])</pre>
                 minv[j] = cur, way[j] = j0;
              if(minv[j] < delta)</pre>
                 delta = minv[j], j1 = j;
       for(int j = 0; j \le m; ++j)
          if(used[i])
              u[p[j]] += delta, v[j] -= delta;
              minv[j] -= delta;
       j0 = j1;
   }while(p[j0] != 0);
   do{
       int j1 = way[j0];
      p[j0] = p[j1];
       j0 = j1;
   }while(j0);
}
// match[i] = coluna escolhida para linha i
vector<int> match(n + 1);
for(int j = 1; j \ll m; ++j)
   match[p[j]] = j;
int cost = -v[0];
Strings
Aho Corasick
map<char, int> to[N];
int ne = 1, term[N], fail[N];
```

```
void add_string(char *str){
   int p = 0;
   for(int i = 0; str[i]; i++){
       if(!to[p][ str[i] ]) to[p][ str[i] ] = ne++;
      p = to[p][ str[i] ];
   term[p] = 1;
}
int go(int s, char c){
   while(s && !to[s].count(c)) s = fail[s];
   if(to[s].count(c)) return to[s][c];
   return s;
}
void init(){
   queue<int> q;
   q.push(0);
   int u, v; char c;
   while(!q.empty()){
      u = q.front(); q.pop();
       for(auto w : to[u]){
```

```
tie(c, v) = w;
          q.push(v);
          if(u){
              fail[v] = go(fail[u], c);
              term[v] |= term[ fail[v] ];
          }
      }
   }
}
Suffix Array
int lcp[N], c[N];
// Caractere final da string '\0' esta sendo considerado
   parte da string s
void build_sa(char s[], int n, int a[]){
   const int A = 300; // Tamanho do alfabeto
   int c1[n], a1[n], h[n + A];
   memset(h, 0, sizeof h);
   for(int i = 0; i < n; i++) {
      c[i] = s[i];
      h[c[i] + 1]++;
   partial_sum(h, h + A, h);
   for(int i = 0; i < n; i++)
       a[h[c[i]]++] = i;
   for(int i = 0; i < n; i++)</pre>
      h[c[i]]--;
   for(int L = 1; L < n; L <<= 1) {
       for(int i = 0; i < n; i++) {
          int j = (a[i] - L + n) \% n;
          a1[h[c[j]]++] = j;
       }
       int cc = -1;
       for(int i = 0; i < n; i++) {</pre>
          if(i == 0 || c[a1[i]] != c[a1[i-1]] || c[(a1[
             i] + L) % n] != c[(a1[i-1] + L) % n])
              h[++cc] = i;
          c1[a1[i]] = cc;
       }
       memcpy(a, a1, sizeof a1);
       memcpy(c, c1, sizeof c1);
       if(cc == n-1) break;
   }
}
void build_lcp(char s[], int n, int a[]){ // lcp[i] =
  lcp(s[:i], s[:i+1])
   int k = 0;
   //memset(lcp, 0, sizeof lcp);
   for(int i = 0; i < n; i++){
       if(c[i] == n-1) continue;
       int j = a[c[i]+1];
       while(i+k < n && j+k < n && s[i+k] == s[j+k]) k
         ++;
       lcp[c[i]] = k;
       if(k) k--;
   }
}
```

University of Brasilia Strings, 16

```
int comp_lcp(int i, int j){
                                                             int x = 0;
   if(i == j) return n - i;
                                                             FOR(i, 1, m) {
   if(c[i] > c[j]) swap(i, j);
                                                              int &r = rad[i] = 0;
   return min(lcp[k] for k in [c[i], c[j]-1]);
                                                              if (i \le x+rad[x]) r = min(rad[x+x-i], x+rad[x]-i);
                                                              while (i-r-1 >= 0 \& i+r+1 < m \& t[i-r-1] == t[i+r]
                                                                +1]) ++r;
Z Algorithm
                                                              if (i+r >= x+rad[x]) x = i;
vector<int> z_algo(const string &s) {
   int n = s.size(), L = 0, R = 0;
                                                             REP(i, m) if (i-rad[i] == 0 || i+rad[i] == m-1) ++rad[
   vector<int> z(n, 0);
   for(int i = 1; i < n; i++){
                                                             REP(i, m) rad[i] /= 2;
       if(i \le R) z[i] = min(z[i-L], R - i + 1);
      while(z[i]+i < n \& s[z[i]+i] == s[z[i]])
          z[i]++:
                                                           Suffix Automaton
      if(i+z[i]-1 > R) L = i, R = i + z[i] - 1;
   }
                                                           map<char, int> to[2*N];
   return z:
                                                           int link[2*N], len[2*N], last = 0, sz = 1;
Prefix function/KMP
                                                           void add_letter(char c){
                                                              int p = last;
vector<int> preffix_function(const string &s){
                                                              last = sz++;
   int n = s.size();
                                                              len[last] = len[p] + 1;
   vector<int> b(n+1);
                                                               for(; !to[p][c]; p = link[p]) to[p][c] = last;
   b[0] = -1;
                                                              if(to[p][c] == last){
   int i = 0, j = -1;
                                                                  link[last] = 0;
   while(i < n){
                                                                  return;
      while(j >= 0 \&\& s[i] != s[j]) j = b[j];
      b[++i] = ++j;
                                                              int u = to[p][c];
                                                              if(len[u] == len[p]+1){
   return b;
                                                                  link[last] = u;
                                                                  return;
                                                              }
void kmp(const string &t, const string &p){
                                                              int c1 = sz++;
   vector<int> b = preffix_function(p);
                                                              to[c1] = to[u];
   int n = t.size(), m = p.size();
                                                              link[c1] = link[u];
   int j = 0;
                                                              len[c1] = len[p]+1;
   for(int i = 0; i < n; i++){
                                                              link[last] = link[u] = c1;
      while(j >= 0 && t[i] != p[j]) j = b[j];
                                                               for(; to[p][c] == u; p = link[p]) to[p][c] = c1;
      if(j == m){
                                                           Suffix Tree
          //patern of p found on t
          j = b[j];
                                                           namespace sf {
      }
                                                           // const int NS = ; const int N = * 2;
   }
                                                           int cn, cd, ns, en = 1, lst;
}
                                                           string S[NS]; int si = -1;
Min rotation
                                                           vector<int> sufn[N]; // sufn[si][i] no do sufixo S[si][i
int min_rotation(int *s, int N) {
                                                           struct node {
 REP(i, N) s[N+i] = s[i];
                                                              int 1, r, si, p, suf;
                                                              map<char, int> adj;
 int a = 0;
                                                              node() : l(0), r(-1), suf(0), p(0) {}
 REP(b, N) REP(i, N) {
                                                              node(int L, int R, int S, int P) : 1(L), r(R), si(S)
   if (a+i == b \mid \mid s[a+i] < s[b+i]) \{ b += max(0, i-1);
                                                               inline int len() { return r - l + 1; }
   if (s[a+i] > s[b+i]) { a = b; break; }
                                                              inline int operator[](int i) { return S[si][l + i];
 return a;
                                                              inline int& operator()(char c) { return adj[c]; }
All palindrome
                                                           inline int new_node(int L, int R, int S, int P) { t[en]
                                                             = node(L, R, S, P); return en++; }
void manacher(char *s, int N, int *rad) {
                                                           void add_string(string s) {
 static char t[2*MAX];
                                                              s += '; S[++si] = s; sufn[si].resize(s.size() + 1)
 int m = 2*N - 1;
                                                                ; cn = cd = 0;
                                                              int i = 0; const int n = s.size();
                                                               for(int j = 0; j < n; j++)
 REP(i, m) t[i] = -1;
 REP(i, N) t[2*i] = s[i];
                                                                  for(; i <= j; i++) {</pre>
                                                                      if(cd == t[cn].len() \&\& t[cn](s[j])) { cn = t}
```

```
[cn](s[j]); cd = 0; 
   if(cd < t[cn].len() && t[cn][cd] == s[j]) {
       cd++;
       if(j < s.size() - 1) break;</pre>
       else {
          if(i) t[lst].suf = cn;
          for(; i <= j; i++) { sufn[si][i] = cn;</pre>
            cn = t[cn].suf; }
   } else if(cd == t[cn].len()) {
       sufn[si][i] = en;
       if(i) t[lst].suf = en; lst = en;
       t[cn](s[j]) = new_node(j, n - 1, si, cn);
       cn = t[cn].suf; cd = t[cn].len();
   } else {
       int mid = new_node(t[cn].1, t[cn].1 + cd -
          1, t[cn].si, t[cn].p);
       t[t[cn].p](t[cn][0]) = mid;
       if(ns) t[ns].suf = mid;
       if(i) t[lst].suf = en; lst = en;
       sufn[si][i] = en;
       t[mid](s[j]) = new_node(j, n - 1, si, mid)
       t[mid](t[cn][cd]) = cn;
       t[cn].p = mid; t[cn].l += cd; cn = t[mid].
       int g = cn? j - cd : i + 1; cn = t[cn].suf
       while(g < j \&\& g + t[t[cn](S[si][g])].len
         () <= j) {
          cn = t[cn](S[si][g]); g += t[cn].len();
       if(g == j) \{ ns = 0; t[mid].suf = cn; cd =
          t[cn].len(); }
       else { ns = mid; cn = t[cn](S[si][g]); cd
         = j - g; }
   }
}
```

# Geometry

# 2D basics

}

```
typedef double cod;
double eps = 1e-7;
bool eq(cod a, cod b){ return abs(a - b) <= eps; }</pre>
struct vec{
   cod x, y; int id;
   vec(cod \ a = 0, \ cod \ b = 0) : x(a), y(b) {}
   vec operator+(const vec &o) const{
       return \{x + o.x, y + o.y\};
   vec operator-(const vec &o) const{
       return \{x - o.x, y - o.y\};
   vec operator*(cod t) const{
       return {x * t, y * t};
   }
   vec operator/(cod t) const{
       return {x / t, y / t};
   }
   cod operator*(const vec &o) const{ // cos
       return x * o.x + y * o.y;
   cod operator^(const vec &o) const{ // sin
```

```
return x * o.y - y * o.x;
   }
   bool operator==(const vec &o) const{
       return eq(x, o.x) && eq(y, o.y);
   bool operator<(const vec &o) const{</pre>
       if(!eq(x, o.x)) return x < o.x;</pre>
       return y < o.y;</pre>
    cod cross(const vec &a, const vec &b) const{
  return (a-(*this)) ^ (b-(*this));
    int ccw(const vec &a, const vec &b) const{
       cod tmp = cross(a, b);
       return (tmp > eps) - (tmp < -eps);</pre>
   }
   cod dot(const vec &a, const vec &b) const{
       return (a-(*this)) * (b-(*this));
    cod len() const{
       return sqrt(x * x + y * y); // <</pre>
    double angle(const vec &a, const vec &b) const{
       return atan2(cross(a, b), dot(a, b));
   }
    double tan(const vec &a, const vec &b) const{
       return cross(a, b) / dot(a, b);
   }
   vec unit() const{
       return operator/(len());
    }
    int quad() const{
       if(x > 0 \& y >= 0) return 0;
       if(x \le 0 \&\& y > 0) return 1;
       if(x < 0 \& y <= 0) return 2;
       return 3;
   }
   bool comp(const vec &a, const vec &b) const{
       return (a - *this).comp(b - *this);
   bool comp(vec b){
       if(quad() != b.quad()) return quad() < b.quad();</pre>
       if(!eq(operator^(b), 0)) return operator^(b) > 0;
       return (*this) * (*this) < b * b;</pre>
   template<class T>
   void sort_by_angle(T first, T last) const{
       std::sort(first, last, [=](const vec &a, const
         vec &b){
           return comp(a, b);
       });
   vec rot90() const{ return {-y, x}; }
   vec rot(double a) const{
       return \{\cos(a)*x - \sin(a)*y, \sin(a)*x + \cos(a)*y\};
   vec proj(const vec &b) const{ // proj of *this onto
       cod k = operator*(b) / (b * b);
       return b * k;
   // proj of (*this) onto the plane orthogonal to b
   vec rejection(vec b) const{
       return (*this) - proj(b);
   }
};
```

```
struct line{
   cod a, b, c; vec n;
   line(vec q, vec w){ // q.cross(w, (x, y)) = 0
       a = -(w.y-q.y);
      b = w.x-q.x;
      c = -(a * q.x + b * q.y);
      n = \{a, b\};
   cod dist(const vec &o) const{
       return abs(eval(o)) / n.len();
   bool contains(const vec &o) const{
       return eq(a * o.x + b * o.y + c, \emptyset);
   }
   cod dist(const line &o) const{
      if(!parallel(o)) return 0;
       if(!eq(o.a * b, o.b * a)) return 0;
       if(!eq(a, 0))
          return abs(c - o.c * a / o.a) / n.len();
       if(!eq(b, 0))
          return abs(c - o.c * b / o.b) / n.len();
       return abs(c - o.c);
   }
   bool parallel(const line &o) const{
      return eq(n ^ o.n, 0);
   }
   bool operator==(const line &o) const{
       if(!eq(a*o.b, b*o.a)) return false;
       if(!eq(a*o.c, c*o.a)) return false;
       if(!eq(c*o.b, b*o.c)) return false;
       return true;
   bool intersect(const line &o) const{
       return !parallel(o) || *this == o;
   vec inter(const line &o) const{
      if(parallel(o)){
          if(*this == o){ }
          else{ /* dont intersect */ }
       auto tmp = n \cdot o.n;
       return {(o.c*b -c*o.b)/tmp, (o.a*c -a*o.c)/tmp};
   }
   vec at_x(cod x) const{
       return \{x, (-c-a*x)/b\};
   }
   vec at_y(cod y) const{
      return \{(-c-b*y)/a, y\};
   cod eval(const vec &o) const{
      return a * o.x + b * o.y + c;
struct segment{
   vec p, q;
   segment(vec a = vec(), vec b = vec()): p(a), q(b) {}
   bool onstrip(const vec &o) const{ // onstrip strip
       return p.dot(o, q) >= -eps && q.dot(o, p) >= -eps
   cod len() const{
      return (p-q).len();
   cod dist(const vec &o) const{
       if(onstrip(o)) return line(p, q).dist(o);
```

```
return min((o-q).len(), (o-p).len());
   }
   bool contains(const vec &o) const{
      return eq(p.cross(q, o), 0) && onstrip(o);
   bool intersect(const segment &o) const{
       if(contains(o.p)) return true;
       if(contains(o.q)) return true;
       if(o.contains(q)) return true;
       if(o.contains(p)) return true;
       return p.ccw(q, o.p) * p.ccw(q, o.q) == -1
      && o.p.ccw(o.q, q) * o.p.ccw(o.q, p) == -1;
   bool intersect(const line &o) const{
      return o.eval(p) * o.eval(q) <= 0;</pre>
   cod dist(const segment &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q)
          || o.onstrip(p) || o.onstrip(q))
              return line(p, q).dist(line(o.p, o.q));
       else if(intersect(o)) return 0;
      return min(min(dist(o.p), dist(o.q)),
                min(o.dist(p), o.dist(q)));
   cod dist(const line &o) const{
      if(line(p, q).parallel(o))
          return line(p, q).dist(o);
       else if(intersect(o)) return 0;
       return min(o.dist(p), o.dist(q));
   }
};
struct hray{
   vec p, q;
   hray(vec a = vec(), vec b = vec()): p(a), q(b){}
   bool onstrip(const vec &o) const{ // onstrip strip
      return p.dot(q, o) >= -eps;
   }
   cod dist(const vec &o) const{
       if(onstrip(o)) return line(p, q).dist(o);
       return (o-p).len();
   bool intersect(const segment &o) const{
      if(!o.intersect(line(p,q))) return false;
       if(line(o.p, o.q).parallel(line(p,q)))
          return contains(o.p) || contains(o.q);
      return contains(line(p,q).inter(line(o.p,o.q)));
   bool contains(const vec &o) const{
      return eq(line(p, q).eval(o), 0) && onstrip(o);
   cod dist(const segment &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q))
              return line(p, q).dist(line(o.p, o.q));
          return o.dist(p);
       else if(intersect(o)) return 0;
       return min(min(dist(o.p), dist(o.q)),
                o.dist(p));
   bool intersect(const hray &o) const{
      if(!line(p, q).parallel(line(o.p, o.q)))
          return false;
       auto pt = line(p, q).inter(line(o.p, o.q));
```

```
return contains(pt) && o.contains(pt); // <<</pre>
   }
                                                               pi slope(){ return pi(a, b); }
                                                           };
   bool intersect(const line &o) const{
       if(line(p, q).parallel(o)) return line(p, q)== o;
                                                           pi cross(line a, line b){
                                                               long double det = a.a * b.b - b.a * a.b;
       if(o.contains(p) || o.contains(q)) return true;
                                                               return pi((a.c * b.b - a.b * b.c) / det, (a.a * b.c
       return (o.eval(p) >= -eps)^(o.eval(p)<o.eval(q));</pre>
       return contains(o.inter(line(p, q)));
                                                                 - a.c * b.a) / det);
   cod dist(const line &o) const{
                                                           bool bad(line a, line b, line c){
                                                               if(ccw(pi(0, 0), a.slope(), b.slope()) <= 0) return</pre>
       if(line(p,q).parallel(o))
          return line(p,q).dist(o);
       else if(intersect(o)) return 0;
                                                               pi crs = cross(a, b);
       return o.dist(p);
                                                               return crs.first * c.a + crs.second * c.b >= c.c;
   }
   cod dist(const hray &o) const{
                                                           bool solve(vector<line> v, vector<pi> &solution){ // ax
       if(line(p, q).parallel(line(o.p, o.q))){
                                                              + bv <= c:
          if(onstrip(o.p) || o.onstrip(p))
                                                               sort(v.begin(), v.end());
              return line(p,q).dist(line(o.p, o.q));
                                                               deque<line> dq;
          return (p-o.p).len();
                                                               for(auto &i : v){
                                                                   if(!dq.empty() && z(ccw(pi(0, 0), dq.back().slope
                                                                     (), i.slope()))) continue;
       else if(intersect(o)) return 0;
       return min(dist(o.p), o.dist(p));
                                                                   while(dq.size() >= 2 && bad(dq[dq.size()-2], dq.
   }
                                                                     back(), i)) dq.pop_back();
};
                                                                   while(dq.size() \geq 2 && bad(i, dq[0], dq[1])) dq.
                                                                     pop_front();
double heron(cod a, cod b, cod c){
                                                                   dq.push_back(i);
   cod s = (a + b + c) / 2;
   return sqrt(s * (s - a) * (s - b) * (s - c));
                                                               while(dq.size() > 2 && bad(dq[dq.size()-2], dq.back
                                                                  (), dq[0])) dq.pop_back();
                                                               while(dq.size() > 2 && bad(dq.back(), dq[0], dq[1]))
Circle line intersection
                                                                  dq.pop_front();
                                                               vector<pi> tmp;
// intersection of line a * x + b * y + c = 0
                                                               for(int i=0; i<dq.size(); i++){</pre>
// and circle centered at the origin with radius r
                                                                   line cur = dq[i], nxt = dq[(i+1)%dq.size()];
double r, a, b, c; // given as input
                                                                   if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps</pre>
double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
                                                                     ) return false;
if(c*c > r*r*(a*a+b*b)+EPS)
                                                                   tmp.push_back(cross(cur, nxt));
   puts("no points");
                                                               }
else if(abs(c*c - r*r*(a*a+b*b)) < EPS){
                                                               solution = tmp:
   puts("1 point");
                                                               return true;
   cout << x0 << ' ' << y0 << ' 'n';
else {
                                                            Detect empty Half plane intersection
   double d = r*r - c*c/(a*a+b*b);
   double mult = sqrt (d / (a*a+b*b));
                                                            // abs(point a) = absolute value of a
   double ax, ay, bx, by;
                                                            // ccw(a, b, c) = a.ccw(b, c)
   ax = x0 + b * mult;
                                                           pair<bool, point> half_inter(vector<pair<point,point> >
   bx = x0 - b * mult;
                                                              &vet){
   ay = y0 - a * mult;
                                                               random_shuffle(all(vet));
   by = y0 + a * mult;
                                                               point p:
   puts ("2 points");
                                                               rep(i,0,sz(vet)) if(ccw(vet[i].x,vet[i].y,p) != 1){
   cout<<ax<<' '<<ay<<'\n'<<bx<<' '<<by<<'\n';
                                                                  point dir = (vet[i].y - vet[i].x) / abs(vet[i].y
                                                                     - vet[i].x);
Half plane intersection
                                                                  point l = vet[i].x - dir*1e15;
                                                                  point r = vet[i].x + dir*1e15;
const double eps = 1e-8;
                                                                   if(r < 1) swap(1, r);
typedef pair<long double, long double> pi;
                                                                   rep(j, 0, i){
bool z(long double x){ return fabs(x) < eps; }</pre>
                                                                      if(ccw(point(), vet[i].x-vet[i].y, vet[j].x-
struct line{
                                                                        vet[j].y) == 0){
   long double a, b, c;
                                                                          if(ccw(vet[j].x, vet[j].y, p) == 1)
   bool operator<(const line &l)const{</pre>
                                                                             continue;
       bool flag1 = pi(a, b) > pi(0, 0);
                                                                          return mp(false, point());
       bool flag2 = pi(1.a, 1.b) > pi(0, 0);
       if(flag1 != flag2) return flag1 > flag2;
                                                                      if(ccw(vet[j].x, vet[j].y, 1) != 1)
       long double t = ccw(pi(0, 0), pi(a, b), pi(l.a, l)
                                                                          1 = max(1, line_intersect(vet[i].x,vet[i].
         .b));
                                                                            y,vet[j].x,vet[j].y));
       return z(t) ? c * hypot(l.a, l.b) < l.c * hypot(a
                                                                      if(ccw(vet[j].x, vet[j].y, r) != 1)
                                                                          r = min(r, line_intersect(vet[i].x,vet[i].
         , b) : t > 0;
```

```
y,vet[j].x,vet[j].y));
    if(!(1 < r)) return mp(false, point());
}
    p = r;
}
return mp(true, p);</pre>
```

### **Circle Circle intersection**

Assume that the first circle is centered at the origin and second at (x2, y2). Find circle line intersection of first circle and line Ax + By + C = 0, where  $A = -2x_2$ ,  $B = -2y_2$ ,  $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$ .

Be aware of corner case with two circles centered at the same point.

# Tangents of two circles

```
// solve first for same circle(and infinitely many
// Find up to four tangents of two circles
void tangents(pt c, double r1, double r2, vector<line> &
   ans){
   double r = r2 - r1;
   double z = c.x * c.x + c.y * c.y;
   double d = z - r * r;
   if(d < -EPS) return;</pre>
   d = sqrt(abs(d));
   line 1;
   1.a = (c.x * r + c.y * d) / z;
   1.b = (c.y * r - c.x * d) / z;
   1.c = r1;
   ans.push_back (1);
vector<line> tangents(circle a, circle b){
   vector<line> ans;
   pt aux = a.center - b.center;
   for(int i = -1; i \le 1; i += 2)
       for(int j = -1; j \le 1; j += 2)
          tangents(aux, a.r * i, b.r * j, ans);
   for(size_t i = 0; i < ans.size(); ++i)</pre>
       ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
   return ans;
}
Convex Hull
vector<vec> monotone_chain_ch(vector<vec> P){
   sort(P.begin(), P.end());
   vector<vec> L, U;
   for(auto p : P){
      while(L.size() >= 2 && L[L.size() - 2].cross(L.
         back(), p) < 0)
          L.pop_back();
      L.push_back(p);
   }
   reverse(P.begin(), P.end());
   for(auto p : P){
      while(U.size() >= 2 && U[U.size() - 2].cross(U.
         back(), p) < 0)
          U.pop_back();
      U.push_back(p);
```

```
}
   L.pop_back(), U.pop_back();
   L.reserve(L.size() + U.size());
   L.insert(L.end(), U.begin(), U.end());
   return L;
Check point inside polygon
bool below(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) <= 0;</pre>
}
bool above(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) >= 0;
// lowerhull, upperhull and point, borders included
bool inside_poly(const vector<vec> &lo, const vector<vec</pre>
  > &hi, vec p){
   return below(hi, p) && above(lo, p);
Check point inside polygon without lower/upper
     hull
// borders included
// must not have 3 colinear consecutive points
bool inside_poly(const vector<vec> &v, vec p){
   if(v[0].ccw(v[1], p) < 0) return false;
   if(v[0].ccw(v.back(), p) > 0) return 0;
   if(v[0].ccw(v.back(), p) == 0)
      return v[0].dot(p, v.back()) >= 0
          && v.back().dot(p, v[0]) >= 0;
   int L = 1, R = (int)v.size() - 1, ans = 1;
   while(L <= R){</pre>
       int mid = (L+R)/2;
       if(v[0].ccw(v[mid], p) >= 0) ans = mid, L = mid
         +1;
       else R = mid-1;
   }
   return v[ans].ccw(v[(ans+1)%v.size()], p) >= 0;
Minkowski sum
vector<vec> mk(const vector<vec>&a,const vector<vec>&b){
   int i = 0, j = 0;
   for(int k = 0; k < (int)a.size(); k++)if(a[k] < a[i])
      i = k:
   for(int k = 0; k < (int)b.size(); k++)if(b[k] < b[j])
      j = k;
```

vector<vec> c:

c.reserve(a.size() + b.size());

vec pt{a[i] + b[j]};

for(int k = 0; k < int(a.size()+b.size()); k++){

#### **Geo Notes**

#### Center of mass

**System of points(2D/3D):** Mass weighted average of points. **Frame(2D/3D):** Get middle point of each segment solve as previously.

**Triangle:** Average of vertices.

**2D Polygon:** Compute **signed** area and center of mass of triangle  $((0,0), p_i, p_{i+1})$ . Then solve as system of points.

**Polyhedron surface:** Solve each face as a 2D polygon(be aware of (0, 0)) then replace each face with its center of mass and solve as system of points.

**Tetrahedron(Triangular pyramid):** As triangles, its the average of points.

**Polyhedron:** Can be done as 2D polygon, but with tetrahedralization intead of triangulation.

#### Pick's Theorem

Given a polygon without self-intersections and all its vertices on integer coordinates in some 2D grid. Let A be its area, I the number of points with interger coordinates stricly inside the polygon and B the number of points with interger coordinates in the border of the polygon. The following formula holds:  $A = I + \frac{B}{2} - 1$ .

## Miscellaneous

multiset<int> S;

#### LIS

```
for(int i = 0; i < n; i++){
   auto it = S.upper_bound(a[i]); // low for inc
   if(it != S.end()) S.erase(it);
   S.insert(a[i]);
}
ans = S.size();
DSU rollback
struct DSU{
   vector<int> sz, p, change;
   vector<tuple<int, int, int>> modifications;
   vector<size_t> saves;
   bool bipartite;
   DSU(int n): sz(n+1, 1), p(n+1), change(n+1),
     bipartite(true){
      iota(p.begin(), p.end(), 0);
   void add_edge(int u, int v){
```

```
if(!bipartite) return;
       int must_change = get_colour(u) == get_colour(v);
      int a = rep(u), b = rep(v);
      if(sz[a] < sz[b]) swap(a, b);
       if(a != b){
          p[b] = a;
          modifications.emplace_back(b, change[b],
            bipartite);
          change[b] ^= must_change;
          sz[a] += sz[b];
       else if(must_change){
          modifications.emplace_back(0, change[0],
            bipartite);
          bipartite = false;
      }
   }
   int rep(int u){
       return p[u] == u ? u : rep(p[u]);
   int get_colour(int u){
       if(p[u] == u) return change[u];
      return change[u] ^ get_colour(p[u]);
   }
   void reset(){
      modifications.clear();
       saves.clear();
       iota(p.begin(), p.end(), 0);
       fill(sz.begin(), sz.end(), 1);
       fill(change.begin(), change.end(), 0);
      bipartite = true;
   }
   void rollback(){
       int u = get<0>(modifications.back());
       tie(ignore, change[u], bipartite) = modifications
         .back();
       sz[p[u]] = sz[u];
      p[u] = u;
      modifications.pop_back();
   }
   void reload(){
      while(modifications.size() > saves.back())
          rollback();
       saves.pop_back();
   }
   void save(){
       saves.push_back(modifications.size());
};
Buildings
// count the number of circular arrays
// of size m, with elements on range
// [1, c**(x*x)]
#include<bits/stdc++.h>
using namespace std;
#define debug(x) cerr << fixed << #x << " = " << x <<
  endl;
const int MOD = 1e9 + 7, MAX = 1e5 + 5;
```

```
int dp[MAX];
                                                            // minimum number of moves to make
                                                            // all elements equal
inline int add(int a, int b) {
                                                            // move: change a segment of equal value
 return a + b >= MOD ? a + b - MOD : a;
                                                            // elements to any value
                                                            int v[305];
inline int sub(int a, int b) {
                                                            int dp[305][305];
 return a - b < 0? a - b + MOD: a;
                                                            int rec[305][305];
inline int mult(int a, int b) {
 return (1LL * a * b) % MOD;
                                                            int f(int 1, int r){
                                                              if(r == 1) return 1;
int f_exp(int x, int exp) {
                                                              if(r < 1) return 0;</pre>
                                                              if(dp[l][r] != -1) return dp[l][r];
 if(exp == 0) return 1;
                                                              int ans = f(1+1, r) + 1;
 else if(exp & 1) return mult(x, f_exp(x, exp - 1));
                                                              for(int i = 1+1; i <= r; i++)</pre>
 return f_exp(mult(x, x), exp / 2);
                                                               if(v[i] == v[1])
inline int inv(int x) {
                                                                 ans = min(ans, f(1, i - 1) + f(i+1, r));
 return f_exp(x, MOD - 2);
                                                              return dp[l][r] = ans;
int main() {
 ios::sync_with_stdio(false);
 cin.tie(NULL); cout.tie(NULL);
                                                            int main() {
                                                              int n, m;
                                                              memset(dp, -1, sizeof dp);
 int n, m, c;
                                                              scanf("%d %d",&n , &m);
 cin >> n >> m >> c;
                                                              for(int i = 0; i < n; i++){
                                                               scanf("%d",v+i);
 int x = f_{exp}(c, n * n);
 int ans = f_{exp}(x, m);
                                                               if(i \& v[i] == v[i-1]){
 for(int i = 1; i <= m; i++) {</pre>
                                                                 i--;
   if(m \% i == 0) {
                                                                 n--;
     int y = f_exp(x, i);
     for(int j = 1; j < i; j++) {
      if(i \% j == 0)
                                                              printf("%d\n",f(0, n-1) - 1);
                                                              // printf("%d\n",rec[0][n-1] );
        y = sub(y, mult(j, dp[j]));
                                                              // printf("%d\n",rec[1][n-1] );
                                                              // printf("%d\n",rec[2][n-3] );
     dp[i] = mult(y, inv(i));
     ans = sub(ans, mult(i - 1, dp[i]));
   }
                                                            Hilbert Order
 }
                                                            // maybe use B = n / sqrt(q)
 cout << ans << '\n';</pre>
                                                            inline int64_t hilbertOrder(int x, int y, int pow = 21,
                                                              int rotate = 0) {
 return 0;
                                                               if(pow == 0) return 0;
}
                                                               int hpow = 1 << (pow-1);</pre>
Rand
                                                               int seg = (x < hpow) ? (
                                                                   (y < hpow) ? 0 : 3
#include <random>
                                                               ):(
#include <chrono>
                                                                   (y < hpow) ? 1 : 2
                                                               );
cout << RAND_MAX << endl;</pre>
                                                               seg = (seg + rotate) & 3;
mt19937 rng(chrono::steady_clock::now().time_since_epoch
                                                               const int rotateDelta[4] = {3, 0, 0, 1};
                                                               int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
  ().count());
vector<int> permutation(N);
                                                               int nrot = (rotate + rotateDelta[seg]) & 3;
                                                               int64_t subSquareSize = int64_t(1) << (2*pow - 2);</pre>
iota(permutation.begin(), permutation.end(), 0);
                                                               int64_t ans = seg * subSquareSize;
                                                               int64_t add = hilbertOrder(nx, ny, pow-1, nrot);
shuffle(permutation.begin(), permutation.end(), rng);
                                                               ans += (seg == 1 || seg == 2) ? add : (subSquareSize
                                                                   - add - 1);
iota(permutation.begin(), permutation.end(), 0);
                                                               return ans;
for(int i = 1; i < N; i++){
                                                            Modular Factorial
   swap(permutation[i], permutation[
     uniform_int_distribution<int>(0, i)(rng)]);
                                                            // Compute (1*2*...*(p-1)*1*(p+1)*(p+2)*..*n) % p
                                                            // in O(p*lg(n))
Klondike
                                                            int factmod(int n, int p){
```

```
int ans = 1;
                                                                  neg = 1;
   while (n > 1) {
                                                                  c = gc();
       for(int i = 2; i <= n % p; i++)</pre>
                                                               }
          ans = (ans * i) % p;
                                                               for(; c >= '0' && c <= '9'; c = gc())</pre>
                                                                  x = (x << 1) + (x << 3) + c - '0';
      n /= p;
      if(n \% 2) ans = p - ans;
                                                               if(neg) x = -x;
                                                           }
   return ans % p;
                                                           inline void print_int(int n){
                                                               int rev = 0, count = 0, neg;
int fac_pow(int n, int p){
                                                               if(n == 0){
   int ans = 0;
                                                                  pc('0');
   while(n) n \neq p, ans += n;
                                                                  return;
   return ans;
                                                               if(n < 0) n = -n, neg = 1;
                                                               while(n % 10 == 0) count++, n /= 10;
int C(int n, int k, int p){
                                                               for(rev = 0; n != 0; n /= 10)
   if(fac_pow(n, p) > fac_pow(n-k, p) + fac_pow(k, p))
                                                                  rev = (rev << 3) + (rev << 1) + n % 10;
                                                               if(neg) pc('-');
   int tmp = factmod(k, p) * factmod(n-k, p);
                                                               while(rev != 0) pc(rev % 10 + '0'), rev /= 10;
   return (f_exp(tmp, p-2, p) * factmod(n, p)) % p;
                                                               while(count--) pc('0');
                                                               pc('\n');
                                                           }
Enumeration all submasks of a bitmask
                                                           inline void print_string(char *str){
// loop through all submask of a given bitmask
                                                               while(*str) pc(*str++);
// it does not include mask 0
                                                               pc('\n');
for(int sub = mask; sub; sub = (sub-1)&mask){
                                                           Knapsack Bounded with Cost
Slope Trick
                                                           // menor custo para conseguir peso ate M usando N tipos
                                                             diferentes de elementos, sendo que o i-esimo elemento
///By wogja125, contest: Codeforces Round #371 (Div. 1),
                                                             pode ser usado b[i] vezes, tem peso w[i] e custo c[i]
   problem: (C) Sonya and Problem Wihtout a Legend,
                                                           // O(N * M)
  Accepted, #
                                                           int b[N], w[N], c[N];
#include <stdio.h>
                                                           MinQueue Q[M]
#include <queue>
                                                           int d[M] //d[i] = custo minimo para conseguir peso i
int main() {
                                                           for(int i = 0; i <= M; i++) d[i] = i ? oo : 0;</pre>
   int n, t;
                                                           for(int i = 0; i < N; i++){
   long long ans = 0;
                                                               for(int j = 0; j < w[i]; j++)
   std::priority_queue<int> Q;
                                                                  Q[j].clear();
   scanf("%d%d", &n, &t);
                                                               for(int j = 0; j \le M; j++){
   Q.push(t);
                                                                  q = Q[j \% w[i]];
   for(int i = 1; i < n; i++) {</pre>
                                                                   if(q.size() >= q) q.pop();
       scanf("%d", &t); t -= i;
                                                                  q.add(c[i]);
      Q.push(t);
                                                                  q.push(d[j]);
       if(Q.top() > t) {
                                                                  d[j] = q.getmin();
          ans += Q.top() - t;
          Q.pop();
          Q.push(t);
      }
                                                           LCA < O(nlgn), O(1) >
   printf("%11d", ans);
                                                           int start[N], dfs_time;
   return 0;
                                                           int tour[2*N], id[2*N];
Fast IO
                                                           void dfs(int u){
                                                               start[u] = dfs_time;
#define pc(x) putchar_unlocked(x)
                                                               id[dfs_time] = u;
                                                               tour[dfs_time++] = start[u];
#define gc(x) getchar_unlocked(x)
                                                               for(int v : g[u]){
inline void scan_int(int &x){
                                                                  dfs(v);
   register int c = gc();
                                                                   id[dfs_time] = u;
   x = 0:
                                                                   tour[dfs_time++] = start[u];
   int neg = 0;
   for(; ((c < '0' || c > '9') && c != '-'); c = gc());
   if(c == '-'){
```

```
int LCA(int u, int v){
   if(start[u] > start[v]) swap(u, v);
   return id[min(tour[k]for k in [start[u],start[v]])];
Buffered reader
// source: https://github.com/ngthanhtrung23/
  ACM_Notebook_new/blob/master/buffered_reader.h
int INP,AM,REACHEOF;
#define BUFSIZE (1<<12)</pre>
char BUF[BUFSIZE+1], *inp=BUF;
#define GETCHAR(INP) { \
   if(!*inp && !REACHEOF) { \
      memset(BUF,0,sizeof BUF);\
       int inpzzz = fread(BUF,1,BUFSIZE,stdin);\
       if (inpzzz != BUFSIZE) REACHEOF = true;\
       inp=BUF; \
   } \
   INP=*inp++; \
#define DIG(a) (((a)>='0')&&((a)<='9'))
#define GN(j) { \
   AM=0; \setminus
   GETCHAR(INP); while(!DIG(INP) && INP!='-') GETCHAR(
     INP);\
   if (INP=='-') {AM=1;GETCHAR(INP);} \
   j=INP-'0'; GETCHAR(INP); \
   while(DIG(INP)){j=10*j+(INP-'0');GETCHAR(INP);} \
   if (AM) j=-j;\
Modular summation
//calcula (sum(0 <= i <= n) P(i)) % mod,
//onde P(i) eh uma PA modular (com outro modulo)
namespace sum_pa_mod{
   11 calc(l1 a, l1 b, l1 n, l1 mod){
      assert(a&&b);
       if(a >= b){
          ll ret = ((n*(n+1)/2) \mod)*(a/b);
          if(a%b) ret = (ret + calc(a%b,b,n,mod))%mod;
          else ret = (ret+n+1)%mod;
          return ret;
       return ((n+1)*(((n*a)/b+1))mod) - calc(b,a,(n*a)/a)
         b, mod) + mod + n/b + 1)%mod;
   }
   //P(i) = a*i \mod m
   11 solve(ll a, ll n, ll m, ll mod){
      a = (a\%m + m)\%m;
      if(!a) return 0;
      11 \text{ ret} = (n*(n+1)/2) \% mod;
      ret = (ret*a)%mod;
      11 g = \_gcd(a,m);
      ret -= m*(calc(a/g,m/g,n,mod)-n-1);
      return (ret%mod + mod)%mod;
   }
   //P(i) = a + r*i \mod m
   11 solve(ll a, ll r, ll n, ll m, ll mod){
      a = (a\%m + m)\%m;
      r = (r\%m + m)\%m;
      if(!r) return (a*(n+1))%mod;
       if(!a) return solve(r, n, m, mod);
```

11 g, x, y;

x = (x%m + m)%m;

g = gcdExtended(r, m, x, y);

```
11 d = a - (a/g)*g;
a -= d;
x = (x*(a/g))%m;
return (solve(r, n+x, m, mod) - solve(r, x-1, m, mod) + mod + d*(n+1))%mod;
}
};
```

#### Burnside's Lemma

Let  $(G, \oplus)$  be a finite group that acts on a set X. It should hold that  $e_g * x = x$  and  $g_1 * (g_2 * x) = (g_1 \oplus g_2) * x$ ,  $\forall x \in X, g_1, g_2 \in G$ . For each  $g \in G$  let  $X^g = \{x \in X \mid g * x = x\}$ . The number of orbits its given by:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

#### Wilson's Theorem

 $(n-1)! = -1 \mod n \iff n \text{ is prime}$ 

#### Fibonacci

- $F_{n-1}F_{n+1} F_n^2 = (-1)^n$
- $F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$
- $GCD(F_n, F_m) = F_{GCD(n,m)}$
- $F_n = \frac{(\frac{1+\sqrt{5}}{2})^n (\frac{1-\sqrt{5}}{2})^n}{\sqrt{5}}$

#### Kirchhoff's Theorem

Laplacian matrix is L = D - A, where D is a diagonal matrix with vertex degrees on the diagonals and A is adjacency matrix.

The number of spanning trees is any cofactor of L. i-th cofactor is determinant of the matrix gotten by removing i-th row and column of L.

#### Multigraphs

In D[i][i] all loops are excluded. A[i][j] = number of edges from i to j.

#### Directed multigraphs

D[i][i] = indegree of i minus the number of loops at i. A[i][j] = number of edges from i to j.

The number of oriented spanning trees rooted at a vertex i is the determinant of the matrix gotten by removing the ith row and column of L.

#### Matroid

Let *X* set of objects,  $I \subseteq 2^X$  set of independents sets such that:

- 1.  $\emptyset \in I$
- 2.  $A \in I, B \subseteq A \implies B \in I$
- 3. Exchange axiom,  $A \in I, B \in I, |B| > |A| \implies \exists x \in B \setminus A : A \cup \{x\} \in I$
- 4.  $A \subseteq X$  and I and I' are maximal independent subsets of A then |I| = |I'|

Then (X, I) is a matroid. The combinatorial optimization problem associated with it is: Given a weight  $w(e) \ge 0 \ \forall e \in X$ , find an independet subset that has the largest possible total weight.

#### Matroid intersection

```
// Input two matroids (X, I_a) and (X, I_b)
// output set I of maximum size, I \in I_a and I \in I_b
set<> I;
while(1){
   for(e_i : X \setminus I)
       if(I + e_i \in I_a \text{ and } I + e_i \in I_b)
           I = I + e_i;
   set<> A, T; queue<> Q;
   for(x : X) label[x] = MARK1;
   for(e_i : X \setminus I){
       if(I + e_i \setminus in I_a)
           Q.push(e_i), label[e_i] = MARK2;
       else{
           for (x \text{ such that } I - x + e_i \in I_a)
               A[x].push(e_i);
       if(I + e_i \setminus in I_b)
           T = T + \{e_i\}
       else{
           for (x \text{ such that } I - x + e_i \setminus in I_b)
               A[e_i].push(x);
   }
   if(T.empty()) break;
   bool found = false;
   while(!Q.empty() and !found){
       auto e = Q.front(); Q.pop();
       for(x : A[e]) if(label[x] == MARK1){
           label[x] = e;
           Q.push(x);
           if(x \in T)
               found = true;
               put = 1;
               while(label[x] != MARK2){
                   if(put) I = I + x;
                   else I = I - x;
                   put = 1 - put;
               I = I + x;
               break;
           }
   if(!found) break;
}
return I;
   Where path(e) = [e] if label[e] = MARK2, path(label[e]) +
[e] otherwise.
```

Matroid Union

Given k matroids over the same set of objects  $(X, I_1)$ ,  $(X, I_2)$ , ...,  $(X, I_k)$  find  $A_1 \in I_1$ ,  $A_2 \in I_2$ , ...,  $A_k \in I_k$  such that  $i \neq j, A_i \cap A_j = \emptyset$  and  $|\bigcup_{i=1}^k A_i|$  is maximum. Matroid union

can be reduced to matroid intersection as follows.

Let  $X' = X \times \{1, 2, ..., k\}$ , ie, k copies of each element of X with different colors. M1 = (X', Q) where  $B \in Q \iff \forall 1 \le i \le k$ ,  $\{x \mid (x, i) \in B\} \in I_i$ , ie, for each color, B is independent. M2 = (X', W) where  $B \in W \iff i \ne j \implies \neg((x, i) \in B \land (x, j) \in B)$ , ie, each element is picked by at most one color.

Intersection of *M*1 and *M*2 is the answer for the combinatorial problem of matroid union.

# **Edge coloring**

```
Data: A graph G

Result: A proper coloring c of the edges of G

Let U := E(G);

while U \neq \emptyset do

Let (u,v) be any edge in U;

Let F[1:k] be a maximal fan of u starting at F[1]=v;

Let c be a color that is free on u and d be a color

that is free on F[k];

Invert the cd_u path;

Let w \in V(G) be such that w \in F, F' = [F[1]...w] is a

fan and d is free on w;

Rotate F' and set c(u,w)=d;

U := U - (u,v);
```

### Notes

When we repeat something and each time we have probability p to succeed then the expected number or tries is  $\frac{1}{p}$ , till we succeed.

#### Small to large

**Trick in statement** If k sets are given you should note that the amount of different set sizes is  $O(\sqrt{s})$  where s is total size of those sets. And no more than  $\sqrt{s}$  sets have size greater than  $\sqrt{s}$ . For example, a path to the root in Aho-Corasick through suffix links will have at most  $O(\sqrt{s})$  vertices.

**gcd on subsegment**, we have at most  $log(a_i)$  different values in  $\{gcd(a_j, a_{j+1}, ..., a_i) \text{ for } j < i\}$ .

From static set to expandable. To insert, create a new set with the new element. While there are two sets with same size, merge them. There will be at most  $\log(n)$  disjoints sets.

**Matrix exponentiation optimization**. Save binary power of  $A_{nxn}$  and answer q queries  $b = A^m x$  in  $O((n^3 + qn^2)log(m))$ .

**Ternary search on integers into binary search**, comparing f(mid) and f(mid+1), binary search on derivative

**Dynamic offline set** For each element we will wind segment of time [a, b] such that element is present in the set during this whole segment. Now we can come up with recursive procedure which handles [l, r] time segment considering that all elements such that  $[l, r] \subset [a, b]$  are already included into the set. Now, keeping this invariant we recursively go into [l, m] and [m + 1, r] subsegments. Finally when we come into segment of length 1.

```
a > b \implies a \mod b < \frac{a}{2}
```