ICPC Team Reference

University of Brasilia

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University of Brasilia Data Structures, 2

```
set ts=4 sw=4 sta nu rnu sc stl+=%F cindent
set bg=dark ruler clipboard=unnamed,unnamedplus
  timeoutlen=100
imap {<CR> {<CR>}}<Esc>0
nmap <F2> 0V$%d
nmap <C-down> :m+1<CR>
nmap <C-up> :m-2<CR>
vmap < C-c > "+y
nmap <C-a> ggVG
syntax on
alias cmp='g++ -Wall -Wformat=2 -Wshadow -Wconversion -
  fsanitize=address -fsanitize=undefined -fno-sanitize-
  recover -std=c++14'
```

Data Structures

Merge Sort Tree

```
struct MergeTree{
   int n:
   vector<vector<int>> st;
   void build(int p, int L, int R, const int v[]){
          st[p].push_back(v[L]);
          return;
       }
      int mid = (L+R)/2;
      build(2*p, L, mid, v);
      build(2*p+1, mid+1, R, v);
       st[p].resize(R-L+1);
      merge(st[2*p].begin(), st[2*p].end(),
              st[2*p+1].begin(), st[2*p+1].end(),
              st[p].begin());
   }
   int query(int p, int L, int R, int i, int j, int x)
     const{
       if(L > j || R < i) return 0;</pre>
       if(L >= i \&\& R <= j){
          int id = lower_bound(st[p].begin(), st[p].end
            (), x) - st[p].begin();
          return int(st[p].size()) - id;
       int mid = (L+R)/2;
       return query(2*p, L, mid, i, j, x) +
          query(2*p+1, mid+1, R, i, j, x);
   }
public:
   MergeTree(int sz, const int v[]): n(sz), st(4*sz){
      build(1, 1, n, v);
   //number of elements >= x on segment [i, j]
   int query(int i, int j, int x) const{
       if(i > j) swap(i, j);
       return query(1, 1, n, i, j, x);
   }
};
Wavelet Tree
template<typename T>
```

```
class wavelet{
   T L, R;
   vector<int> 1;
    vector<T> sum; // <<</pre>
```

```
wavelet *lef, *rig;
   int r(int i) const{ return i - l[i]; }
public:
   template<typename ITER>
   wavelet(ITER bg, ITER en){
       lef = rig = nullptr;
       L = *bg, R = *bg;
       for(auto it = bg; it != en; it++)
          L = min(L, *it), R = max(R, *it);
       if(L == R) return;
       T mid = L + (R - L)/2;
       1.reserve(std::distance(bg, en) + 1);
       sum.reserve(std::distance(bg, en) + 1);
       1.push_back(0), sum.push_back(0);
       for(auto it = bg; it != en; it++)
          1.push_back(1.back() + (*it <= mid)),</pre>
          sum.push_back(sum.back() + *it);
       auto tmp = stable_partition(bg, en, [mid](T x){
          return x <= mid;</pre>
       });
       if(bg != tmp) lef = new wavelet(bg, tmp);
       if(tmp != en) rig = new wavelet(tmp, en);
   }
    ~wavelet(){
       delete lef;
       delete rig;
   }
   // 1 index, first is 1st
   T kth(int i, int j, int k) const{
       if(L >= R) return L;
       int c = l[j] - l[i-1];
       if(c \ge k) return lef \ge kth(l[i-1]+1, l[j], k);
       else return rig->kth(r(i-1)+1, r(j), k - c);
   }
   // # elements > x on [i, j]
   int cnt(int i, int j, T x) const{
       if(L > x) return j - i + 1;
       if(R <= x || L == R) return 0;</pre>
       int ans = 0;
       if(lef) ans += lef->cnt(l[i-1]+1, l[j], x);
       if(rig) ans += rig->cnt(r(i-1)+1, r(j), x);
       return ans;
   }
   // sum of elements <= k on [i, j]</pre>
   T sumk(int i, int j, T k){
       if(L == R) return R <= k ? L * (j - i + 1) : 0;
       if(R <= k) return sum[j] - sum[i-1];</pre>
       int ans = 0;
       if(lef) ans += lef->sumk(l[i-1]+1, l[j], k);
       if(rig) ans += rig->sumk(r(i-1)+1, r(j), k);
       return ans;
   // swap (i, i+1) just need to update "array" l[i]
};
```

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```
#include <ext/pb_ds/assoc_container.hpp>
                                                            11 binary_search(const vector<Line> &cht, 11 x){
#include <ext/pb_ds/tree_policy.hpp>
                                                                int L = 0, R = (int)cht.size()-2;
#include <ext/pb_ds/detail/standard_policies.hpp>
                                                                int bans = (int)cht.size()-1;
                                                                while(L <= R){</pre>
using namespace __gnu_pbds; // or pb_ds;
                                                                   int mid = (L+R)/2:
                                                                   if(cht[mid].eval(x) >= cht[mid+1].eval(x)) // <<<</pre>
template<typename T, typename B = null_type>
                                                                       L = mid + 1;
using oset = tree<T, B, less<T>, rb_tree_tag,
                                                                   else bans = mid, R = mid - 1;
  tree_order_statistics_node_update>;
// find_by_order / order_of_key
                                                                return cht[bans].eval(x);
Hash table
                                                            Convex Hull Trick
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
                                                            const ll is_query = -(1LL<<62);</pre>
                                                            struct Line{
struct custom_hash {
                                                               11 m. b:
   static uint64_t splitmix64(uint64_t x) {
                                                                mutable function<const Line*()> succ;
       // http://xorshift.di.unimi.it/splitmix64.c
                                                                bool operator<(const Line& rhs) const{</pre>
      x += 0x9e3779b97f4a7c15;
                                                                   if(rhs.b != is_query) return m < rhs.m;</pre>
      x = (x \hat{ } (x >> 30)) * 0xbf58476d1ce4e5b9;
                                                                   const Line* s = succ();
      x = (x \hat{ } (x >> 27)) * 0x94d049bb133111eb;
                                                                   if(!s) return 0;
                                                                   11 x = rhs.m;
      return x \hat{ } (x >> 31);
                                                                   return b - s->b < (s->m - m) * x;
   }
   size_t operator()(uint64_t x) const {
                                                            };
       static const uint64_t FIXED_RANDOM = chrono::
                                                            struct Cht : public multiset<Line>{ // maintain max
         steady_clock::now().time_since_epoch().count();
                                                                bool bad(iterator y){
       return splitmix64(x + FIXED_RANDOM);
                                                                   auto z = next(y);
   }
                                                                   if(y == begin()){
};
                                                                       if(z == end()) return 0;
                                                                       return y->m == z->m \&\& y->b <= z->b;
gp_hash_table<long long, int, custom_hash> table;
                                                                   }
unordered_map<long long, int, custom_hash> uhash;
                                                                   auto x = prev(y);
uhash.reserve(1 << 15);</pre>
                                                                   if(z == end()) return y->m == x->m && y->b <= x->
uhash.max_load_factor(0.25);
                                                                   return (long double) (x->b - y->b)*(z->m - y->m)
Convex Hull Trick Simple
                                                                     >= (long double)(y->b - z->b)*(y->m - x->m);
struct Line{
                                                                void insert_line(ll m, ll b){
   11 m, b;
                                                                   auto y = insert({ m, b });
   inline 11 eval(11 x) const{
                                                                   y->succ = [=]{ return next(y) == end() ? 0 : &*
      return x * m + b;
                                                                     next(y); };
   }
                                                                   if(bad(y)){ erase(y); return; }
};
                                                                   while(next(y) != end() && bad(next(y))) erase(
                                                                     next(v)):
// min => cht.back().m >= L.m
                                                                   while(y != begin() && bad(prev(y))) erase(prev(y)
// max => cht.back().m <= L.m
                                                                     );
void push_line(vector<Line> &cht, Line L){
 while((int)cht.size() >= 2){
                                                                11 eval(11 x){
   int sz = (int)cht.size();
                                                                   auto 1 = *lower_bound((Line) { x, is_query });
   if((long double)(L.b-cht[sz-1].b)*(cht[sz-2].m-L.m)
                                                                   return 1.m * x + 1.b;
  <= (long double)(L.b-cht[sz-2].b)*(cht[sz-1].m-L.m)){</pre>
     cht.pop_back();
                                                            };
   }
                                                            Convex Hull Trick
   else break;
 cht.push_back(L);
                                                             * Author: Simon Lindholm
}
                                                             * source: https://github.com/kth-competitive-
// x increasing; pos = 0 in first call
                                                               programming/kactl/blob/master/content/data-structures
11 linear_search(const vector<Line> &cht,ll x,int &pos){
                                                               /LineContainer.h
                                                             * License: CC0
   while(pos+1 < (int)cht.size()){</pre>
 *>>*/ if(cht[pos].eval(x) >= cht[pos+1].eval(x)) pos++;
      else break;
   }
                                                            struct Line {
                                                               mutable 11 m, b, p;
   return cht[pos].eval(x);
}
                                                                bool operator<(const Line& o) const { return m < o.m</pre>
```

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```
}
     : }
   bool operator<(11 x) const { return p < x; }</pre>
                                                           int lg(int x){ return 31 - __builtin_clz(x); }
};
struct LineContainer : multiset<Line, less<>>> { // CPP14
                                                           int getmn(int 1, int r)\{ // [1, r]
  only
                                                               int 1z = 1g(r - 1 + 1);
                                                               return min(fn(1, lz), fn(r - (1 << lz) + 1, lz));
   // (for doubles, use inf = 1/.0, div(a,b) = a/b)
   const 11 inf = LLONG_MAX;
                                                           }
   11 div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b); }</pre>
                                                           Treap
   bool isect(iterator x, iterator y) {
                                                           // source: https://github.com/victorsenam/caderno/blob/
       if (y == end()) { x->p = inf; return false; }
                                                             master/code/treap.cpp
       if (x->m == y->m) x->p = x->b > y->b ? inf : -inf
                                                           //const int N = ; typedef int num;
                                                           num X[N]; int en = 1, Y[N], sz[N], L[N], R[N];
       else x->p = div(y->b - x->b, x->m - y->m);
                                                           void calc (int u) { // update node given children info
      return x->p >= y->p;
                                                               if(!u) return:
                                                               sz[u] = sz[L[u]] + 1 + sz[R[u]];
   void add(ll m, ll b) {
      auto z = insert(\{m, b, 0\}), y = z++, x = y;
                                                               // code here, no recursion
                                                           }
       while (isect(y, z)) z = erase(z);
                                                           void unlaze (int u) {
       if (x != begin() \&\& isect(--x, y)) isect(x, y =
                                                               if(!u) return;
         erase(v)):
                                                               // code here, no recursion
       while ((y = x) != begin() \&\& (--x)->p >= y->p)
          isect(x, erase(y));
                                                           void split_val(int u, num x, int &l, int &r) { // l gets
                                                               <= x, r gets > x
   11 query(11 x) {
                                                               unlaze(u); if(!u) return (void) (1 = r = 0);
      assert(!empty());
       auto 1 = *lower_bound(x);
                                                               if(X[u] <= x) { split_val(R[u], x, 1, r); R[u] = 1;</pre>
       return 1.m * x + 1.b;
                                                                 1 = u; }
                                                               else { split_val(L[u], x, 1, r); L[u] = r; r = u; }
   }
};
                                                               calc(u);
Min queue
                                                           void split_sz(int u, int s, int &l, int &r) { // 1 gets
                                                             first s, r gets remaining
template<typename T>
                                                               unlaze(u); if(!u) return (void) (1 = r = 0);
class minQ{
                                                               if(sz[L[u]] < s)  { split_sz(R[u], s - sz[L[u]] - 1,
   deque<tuple<T, int, int> > p;
                                                                 1, r); R[u] = 1; 1 = u; }
   T delta;
                                                               else { split_sz(L[u], s, l, r); L[u] = r; r = u; }
   int sz;
                                                               calc(u);
public:
   minQ() : delta(0), sz(0) {}
                                                           int merge(int 1, int r) { // els on 1 <= els on r</pre>
   inline int size() const{ return sz; }
                                                               unlaze(1); unlaze(r); if(!1 || !r) return 1 + r; int
   inline void add(T x){ delta += x; }
                                                                  11:
   inline void push(T x, int id){
                                                               if(Y[1] > Y[r]) { R[1] = merge(R[1], r); u = 1; }
      x -= delta, sz++;
                                                               else { L[r] = merge(1, L[r]); u = r; }
       int t = 1;
                                                               calc(u); return u;
       while(p.size() > 0 \& get<0>(p.back()) >= x)
          t += get<1>(p.back()), p.pop_back();
                                                           void init(int n=N-1) { // XXX call before using other
      p.emplace_back(x, t, id);
                                                               for(int i = en = 1; i \le n; i++) { Y[i] = i; sz[i] =
   inline void pop(){
                                                                  1; L[i] = R[i] = 0; }
      get<1>(p.front())--, sz--;
                                                               random_shuffle(Y + 1, Y + n + 1);
       if(!get<1>(p.front())) p.pop_front();
                                                           }
                                                           void insert(int &u, int it){
   T getmin() const{ return get<0>(p.front())+delta; }
                                                               unlaze(u);
   int getid() const{ return get<2>(p.front()); }
                                                               if(!u) u = it;
}:
                                                               else if(Y[it] > Y[u]) split_val(u, X[it], L[it], R[
Sparse Table
                                                                 it]), u = it;
                                                               else insert(X[it] < X[u] ? L[u] : R[u], it);</pre>
const int N = 100005;
                                                               calc(u);
int v[N], n;
                                                           void erase(int &u, num key){
int dn[N][20];
                                                               unlaze(u);
int fn(int i, int j){
                                                               if(!u) return;
   if(j == 0) return v[i];
                                                               if(X[u] == key) u = merge(L[u], R[u]);
   if(~dn[i][j]) return dn[i][j];
                                                               else erase(key < X[u] ? L[u] : R[u], key);</pre>
   calc(u);
     ), j-1));
```

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```
int create_node(num key){
                                                                     () && it->l < r; it++) {
   X[en] = key;
                                                                      ans.push_back(*it);
                                                                  }
   sz[en] = 1;
   L[en] = R[en] = 0;
                                                                  ranges.erase(ranges.lower_bound(1), ranges.
   return en++;
                                                                     lower_bound(r));
                                                                  ranges.insert(Range(1, r, v));
int query(int u, int 1, int r){//0 index
                                                                  return ans;
   unlaze(u);
                                                               }
   if(u! or r < 0 or 1 >= sz[u]) return
                                                           private:
     identity_element;
                                                               std::set<Range> ranges;
   if(1 <= 0 and r >= sz[u] - 1) return subt_data[u];
   int ans = query(L[u], 1, r);
                                                           Heavy Light Decomposition
   if(1 \le sz[L[u]] and sz[L[u]] \le r)
       ans = max(ans, st[u]);
                                                           void dfs_sz(int u){
   ans = max(ans, query(R[u], l-sz[L[u]]-1, r-sz[L[u]])
                                                               sz[u] = 1;
     11-1)):
   return ans;
                                                               for(auto &v : g[u]) if(v == p[u]){
}
                                                                  swap(v, g[u].back());
ColorUpdate
                                                                  g[u].pop_back();
                                                                  break:
// source: https://github.com/tfg50/Competitive-
                                                               }
  Programming/tree/master/Biblioteca/Data%20Structures
                                                               for(auto &v : g[u]){
                                                                  p[v] = u;
#include <set>
                                                                  dfs_sz(v);
#include <vector>
                                                                  sz[u] += sz[v];
                                                                  if(sz[v] > sz[ g[u][0] ])
template <class Info = int>
                                                                      swap(v, g[u][0]);
class ColorUpdate {
                                                               }
public:
                                                           }
   struct Range {
      Range(int l = 0) { this->l = 1; }
                                                           // nxt[u] = start of path with u
       Range(int 1, int r, Info v) {
                                                           // set nxt[root] = root beforehand
          this -> 1 = 1;
                                                           void dfs_hld(int u){
          this->r = r;
          this->v = v;
                                                               in[u] = t++;
      }
                                                               rin[in[u]] = u;
                                                               for(auto v : g[u]){
      int 1, r;
                                                                  nxt[v] = (v == g[u][0] ? nxt[u] : v);
      Info v;
                                                                  dfs_hld(v);
      bool operator < (const Range &b) const { return 1</pre>
                                                               out[u] = t;
          < b.1; }
                                                           }
   };
                                                           // subtree of u => [ in[u], out[u] )
   std::vector<Range> upd(int 1, int r, Info v) {
                                                           // path from nxt[u] to u \Rightarrow [in[nxt[u]], in[u]]
       std::vector<Range> ans;
       if(1 >= r) return ans;
                                                           Iterative Segtree
       auto it = ranges.lower_bound(1);
       if(it != ranges.begin()) {
                                                           T query(int 1, int r, int &pos){ // [1, r]
          it--:
                                                               T rl, rr;
          if(it->r>1) {
                                                               for(1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1){
              auto cur = *it;
                                                                  if(1 & 1) rl = merge(rl, st[l++]);
             ranges.erase(it);
                                                                  if(r & 1) rr = merge(st[--r], rr);
             ranges.insert(Range(cur.1, 1, cur.v));
                                                               }
             ranges.insert(Range(1, cur.r, cur.v));
                                                               return merge(rl, rr);
                                                           }
       it = ranges.lower_bound(r);
                                                           // initially save v[i] in st[n+i] for all i in [0, n)
       if(it != ranges.begin()) {
                                                           void build(){
          it--;
                                                               for(int p = n-1; p > 0; p--)
          if(it->r > r) {
                                                                  st[p] = merge(st[2*p], st[2*p+1]);
              auto cur = *it;
              ranges.erase(it);
             ranges.insert(Range(cur.1, r, cur.v));
                                                           void update(int p, T val){
              ranges.insert(Range(r, cur.r, cur.v));
                                                               st[p += n] = val;
          }
                                                               while(p \gg 1) st[p] = merge(st[2*p], st[2*p+1]);
                                                           }
       for(it = ranges.lower_bound(1); it != ranges.end
```

LiChao's Segtree

```
void add_line(line nw, int v = 1, int l = 0, int r =
  maxn) \{ // [1, r) \}
   int m = (1 + r) / 2;
   bool lef = nw.eval(1) < st[v].eval(1);</pre>
   bool mid = nw.eval(m) < st[v].eval(m);</pre>
   if(mid) swap(st[v], nw);
   if(r - 1 == 1) {
       return;
   } else if(lef != mid) {
       add_line(nw, 2 * v, 1, m);
   } else {
       add_line(nw, 2 * v + 1, m, r);
}
int get(int x, int v = 1, int l = 0, int r = maxn) {
   int m = (1 + r) / 2;
   if(r - 1 == 1) {
      return st[v].eval(x);
   else if(x < m) 
      return min(st[v].eval(x), get(x, 2*v, 1, m));
      return min(st[v].eval(x), get(x, 2*v+1, m, r));
```

Palindromic tree

```
#include <bits/stdc++.h>
using namespace std;
const int maxn = 3e5 + 1, sigma = 26;
int len[maxn], link[maxn], to[maxn][sigma];
int slink[maxn], diff[maxn], series_ans[maxn];
int sz, last, n;
char s[maxn];
void init()
   s[n++] = -1;
   link[0] = 1;
   len[1] = -1;
   sz = 2;
int get_link(int v)
   while(s[n - len[v] - 2] != s[n - 1]) v = link[v];
   return v;
}
void add_letter(char c)
   s[n++] = c -= 'a';
   last = get_link(last);
   if(!to[last][c])
      len[sz] = len[last] + 2;
      link[sz] = to[get_link(link[last])][c];
      diff[sz] = len[sz] - len[link[sz]];
       if(diff[sz] == diff[link[sz]])
          slink[sz] = slink[link[sz]];
          slink[sz] = link[sz];
       to[last][c] = sz++;
```

```
last = to[last][c];
}
int main()
{
   ios::sync_with_stdio(0);
   cin.tie(0);
   init();
   string s;
   cin >> s;
   int n = s.size();
   int ans[n + 1];
   memset(ans, 63, sizeof(ans));
   ans[0] = 0;
   for(int i = 1; i <= n; i++)</pre>
       add_letter(s[i - 1]);
       for(int v = last; len[v] > 0; v = slink[v])
           series_ans[v] = ans[i - (len[slink[v]] + diff
             [v])]:
           if(diff[v] == diff[link[v]])
              series_ans[v] = min(series_ans[v],
                series_ans[link[v]]);
           ans[i] = min(ans[i], series_ans[v] + 1);
       }
       cout << ans[i] << "\n";</pre>
   }
   return 0;
Math
```

Extended Euclidean Algorithm

```
// a*x + b*y = gcd(a, b), <gcd, x, y>
tuple<int, int, int> gcd(int a, int b) {
   if(b == 0) return make_tuple(a, 1, 0);
   int q, w, e;
   tie(q, w, e) = gcd(b, a % b);
   return make_tuple(q, e, w - e * (a / b));
}
```

Chinese Remainder Theorem

Preffix inverse

```
inv[1] = 1;
for(int i = 2; i < p; i++)
  inv[i] = (p - (p/i) * inv[p%i] % p) % p;</pre>
```

```
Pollard Rho
11 rho(11 n){
   if(n % 2 == 0) return 2;
   11 d, c, x, y;
   do{
       c = 11rand() % n, x = 11rand() % n, y = x;
      do{
          x = add(mul(x, x, n), c, n);
          y = add(mul(y, y, n), c, n);
          y = add(mul(y, y, n), c, n);
          d = \_gcd(abs(x - y), n);
       }while(d == 1);
   }while(d == n);
   return d;
ll pollard_rho(ll n){
   ll x, c, y, d, k;
   int i;
   do{
      i = 1;
      x = 11rand() % n, c = 11rand() % n;
      y = x, k = 4;
       do{
          if(++i == k) y = x, k *= 2;
          x = add(mul(x, x, n), c, n);
          d = \_gcd(abs(x - y), n);
       }while(d == 1);
   }while(d == n);
   return d;
}
void factorize(ll val, map<ll, int> &fac){
   if(rabin(val)) fac[ val ]++;
   else{
       11 d = pollard_rho(val);
       factorize(d, fac);
       factorize(val / d, fac);
   }
}
map<ll, int> factor(ll val){
   map<ll, int> fac;
   if(val > 1) factorize(val, fac);
   return fac;
Miller Rabin
bool rabin(ll n){
   if(n <= 1) return 0:
   if(n <= 3) return 1;
   11 s = 0, d = n - 1;
   while(d % 2 == 0) d /= 2, s++;
   for(int k = 0; k < 64; k++){
      11 a = (11rand() \% (n - 3)) + 2;
      11 x = fexp(a, d, n);
       if(x != 1 \&\& x != n-1){
          for(int r = 1; r < s; r++){
             x = mul(x, x, n);
             if(x == 1) return 0;
             if(x == n-1) break;
          if(x != n-1) return 0;
       }
```

}

```
return 1;
}
Totiente
11 totiente(ll n){
   11 \text{ ans} = n;
   for(ll i = 2; i*i <= n; i++){
       if(n \% i == 0){
          ans = ans / i * (i - 1);
          while(n % i == 0) n /= i;
       }
   }
   if(n > 1) ans = ans / n * (n - 1);
   return ans;
Primitive root
// a primitive root modulo n is any number g such that
 any c coprime to n is congruent to a power of g modulo
bool exists_root(ll n){
   if(n == 1 || n == 2 || n == 4) return true;
   if(n \% 2 == 0) n /= 2;
   if(n % 2 == 0) return false;
   // test if n is a power of only one prime
   for(ll i = 3; i * i <= n; i += 2) if(n % i == 0){
       while(n % i == 0) n /= i;
       return n == 1;
   }
   return true;
}
11 primitive_root(ll n){
   if(n == 1 || n == 2 || n == 4) return n - 1;
   if(not exists_root(n)) return -1;
   11 x = phi(n);
   auto pr = factorize(x);
   auto check = [x, n, pr](11 m){
       for(11 p : pr) if(fexp(m, x / p, n) == 1)
          return false;
      return true;
   };
   for(ll m = 2; ; m++) if(__gcd(m, n) == 1)
       if(check(m)) return m;
}
// Let's denote R(n) as the set of primitive roots
  modulo n, p is prime
// g \in R(p) => (pow(g, p-1, p * p) == 1 ? g+p : g) \in
   R(pow(p, k)), for all k > 1
// g \text{ in } R(pow(p, k)) \Rightarrow (g \% 2 == 1 ? g : g + pow(p, k))
   Mobius Function
memset(mu, 0, sizeof mu);
mu[1] = 1;
for(int i = 1; i < N; i++)
   for(int j = i + i; j < N; j += i)
      mu[j] -= mu[i];
// g(n) = sum{f(d)} => f(n) = sum{mu(d)*g(n/d)}
Mulmod TOP
constexpr uint64_t mod = (1ull<<61) - 1;</pre>
uint64_t modmul(uint64_t a, uint64_t b){
```

```
uint64_t 11 = (uint32_t)a, h1 = a>>32, 12 = (
                                                                          b[i] -= A[i][y] * b[x];
     uint32_t)b, h2 = b >> 32;
                                                                           for(int j = 0; j < n; j++) if(j != y)
   uint64_t l = 11*12, m = 11*h2 + 12*h1, h = h1*h2;
                                                                              A[i][j] -= A[i][y] * A[x][j];
   uint64_t ret = (1\&mod) + (1>>61) + (h << 3) + (m >>
                                                                           A[i][y] = -A[i][y] * A[x][y];
     29) + (m << 35 >> 3) + 1;
                                                                       }
                                                                   ans += c[y] * b[x];
   ret = (ret & mod) + (ret>>61);
   ret = (ret & mod) + (ret>>61);
                                                                   for(int i = 0; i < n; i++)</pre>
                                                                       if(i != y)
   return ret-1;
                                                                          c[i] -= c[y] * A[x][i];
                                                                   c[y] = -c[y] * A[x][y];
Matrix Determinant
                                                                }
int n:
                                                                // maximiza sum(x[i] * c[i])
long double a[n][n];
                                                                // sujeito a
                                                                // sum(a[i][j] * x[j]) <= b[i] para 0 <= i < m (Ax)
long double gauss(){
   long double det = 1;
                                                                // x[i] >= 0 para 0 <= i < n (x >= 0)
   for(int i = 0; i < n; i++){
                                                                // (n variaveis, m restricoes)
       int q = i;
                                                                // guarda a resposta em ans e retorna o valor otimo
       for(int j = i+1; j < n; j++){
                                                                dbl solve(int _n, int _m) {
          if(abs(a[j][i]) > abs(a[q][i]))
                                                                   this->n = _n; this->m = _m;
              q = j;
       }
                                                                    for(int i = 1; i < m; i++){
       if(abs(a[q][i]) < EPS){
                                                                       int id = uniform_int_distribution<int>(0, i)(
          det = 0;
                                                                         rng);
          break;
                                                                       swap(b[i], b[id]);
                                                                       for(int j = 0; j < n; j++)
       if(i != q){
                                                                           swap(A[i][j], A[id][j]);
          for(int w = 0; w < n; w++)
                                                                   }
              swap(a[i][w], a[q][w]);
          det = -det;
                                                                   ans = 0.;
                                                                   for(int i = 0; i < n; i++) X[i] = i;</pre>
       det *= a[i][i];
                                                                    for(int i = 0; i < m; i++) Y[i] = i + n;
       for(int j = i+1; j < n; j++) a[i][j] /= a[i][i];</pre>
                                                                   while(true) {
                                                                       int x = min_element(b, b + m) - b;
       for(int j = 0; j < n; j++) if(j != i){
                                                                       if(b[x] >= -eps)
          if(abs(a[j][i]) > EPS)
                                                                          break:
              for(int k = i+1; k < n; k++)
                                                                       int y = find_if(A[x], A[x] + n, [](dbl d) {
                  a[j][k] -= a[i][k] * a[j][i];
                                                                         return d < -eps; }) - A[x];</pre>
       }
                                                                       if(y == n) throw 1; // no solution
   }
                                                                       pivot(x, y);
   return det;
                                                                   while(true) {
                                                                       int y = max_{element}(c, c + n) - c;
Simplex Method
                                                                       if(c[y] <= eps) break;</pre>
                                                                       int x = -1;
typedef long double dbl;
                                                                       dbl mn = 1. / 0.;
const dbl eps = 1e-6;
                                                                       for(int i = 0; i < m; i++)</pre>
const int N = , M = ;
                                                                           if(A[i][y] > eps \&\& b[i] / A[i][y] < mn)
                                                                              mn = b[i] / A[i][y], x = i;
mt19937 rng(chrono::steady_clock::now().time_since_epoch
                                                                       if(x == -1) throw 2; // unbounded
  ().count()):
                                                                       pivot(x, y);
struct simplex {
   int X[N], Y[M];
                                                                   memset(sol, 0, sizeof(dbl) * n);
   dbl A[M][N], b[M], c[N];
                                                                   for(int i = 0; i < m; i++)
   dbl ans:
                                                                       if(Y[i] < n)
   int n, m;
                                                                           sol[Y[i]] = b[i];
   dbl sol[N];
                                                                   return ans;
   void pivot(int x, int y){
                                                            };
       swap(X[y], Y[x]);
                                                            FFT
       b[x] /= A[x][y];
       for(int i = 0; i < n; i++)
          if(i != y)
                                                            struct base{
              A[x][i] /= A[x][y];
                                                                double r, i;
       A[x][y] = 1. / A[x][y];
                                                                base(double _r = 0, double _i = 0) : r(_r), i(_i) {}
       for(int i = 0; i < m; i++)</pre>
                                                                base operator*(base &o) const{
          if(i != x && abs(A[i][y]) > eps) {
                                                                   return {r*o.r - i*o.i, r*o.i + o.r*i};
```

```
double real() const{ return r; }
   void operator*=(const base &o){
       (*this) = \{r*o.r-i*o.i, r*o.i+o.r*i\};
   void operator+=(const base &o)\{r \neq o.r, i \neq o.i; \}
   void operator/=(const double &o){ r /= o, i /= o; }
   void operator-=(const base &o){r -= o.r, i -= o.i; }
   base operator+(const base &o){return {r+o.r,i+o.i};}
   base operator-(const base &o){return {r-o.r,i-o.i};}
double PI = acos(-1);
void fft(vector<base> &a, bool inv){
   int n = (int)a.size();
   for(int i = 1, j = 0; i < n; i++){
       int bit = n \gg 1;
       for(; j >= bit; bit >>= 1) j -= bit;
       j += bit;
       if(i < j) swap(a[i], a[j]);
   }
   for(int sz = 2; sz <= n; sz <<= 1) {</pre>
       double ang = 2*PI/sz * (inv ? -1 : 1);
       base wlen(cos(ang), sin(ang));
       for(int i = 0; i < n; i += sz){
          base w(1);
          for(int j = 0; j < sz/2; j++){
              base u = a[i+j], v = a[i+j+sz/2] * w;
              a[i+j] = u + v;
              a[i+j+sz/2] = u - v;
              w *= wlen;
          }
       }
   }
   if(inv) for(int i = 0; i < n; i++) a[i] /= 1.0 * n;
}
void multiply(const vector<int> &a, const vector<int> &b
  , vector<int> &res){
   vector<base> fa(a.begin(), a.end());
   vector<base> fb(b.begin(), b.end());
   size_t n = 1;
   while(n < a.size()) n <<= 1;
   while(n < b.size()) n <<= 1;</pre>
   n \ll 1:
   fa.resize(n), fb.resize(n);
   fft(fa, false), fft(fb, false);
   for(size_t i = 0; i < n; i++)</pre>
       fa[i] *= fb[i];
   fft(fa, true);
   res.resize (n);
   for(size_t i = 0; i < n; ++i)</pre>
       res[i] = int(fa[i].real() + 0.5);
}
NTT
const int mod = 7340033;
const int root = 5;
const int root_1 = 4404020;
const int root_pw = 1<<20;</pre>
void fft (vector<int> & a, bool invert) {
```

```
int n = (int) a.size();
   for (int i=1, j=0; i<n; ++i) {</pre>
       int bit = n \gg 1;
       for (; j>=bit; bit>>=1)
          j -= bit;
       j += bit;
       if (i < j)
           swap (a[i], a[j]);
   for (int len=2; len<=n; len<<=1) {</pre>
       int wlen = invert ? root_1 : root;
       for (int i=len; i<root_pw; i<<=1)</pre>
           wlen = int (wlen * 111 * wlen % mod);
       for (int i=0; i<n; i+=len) {</pre>
          int w = 1;
           for (int j=0; j<len/2; ++j) {</pre>
              int u = a[i+j], v = int (a[i+j+len/2] * 1
                11 * w % mod);
              a[i+j] = u+v < mod ? u+v : u+v-mod;
              a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
              w = int (w * 111 * wlen % mod);
          }
       }
   }
   if (invert) {
       int nrev = reverse (n, mod);
       for (int i=0; i<n; ++i)</pre>
           a[i] = int (a[i] * 111 * nrev % mod);
   }
}
Gauss
// Solves systems of linear equations.
// To use, build a matrix of coefficients and call run(
  mat, R, C). If the i-th variable is free, row[i] will
  be -1, otherwise it's value will be ans[i].
namespace Gauss {
 const int MAXC = 1001;
 int row[MAXC];
 double ans[MAXC];
 void run(double mat[][MAXC], int R, int C) {
   REP(i, C) row[i] = -1;
   int r = 0;
   REP(c, C) {
     int k = r;
     FOR(i, r, R) if(fabs(mat[i][c]) > fabs(mat[k][c]))
       k = i:
     if(fabs(mat[k][c]) < eps) continue;</pre>
     REP(j, C+1) swap(mat[r][j], mat[k][j]);
     REP(i, R) if (i != r) {
       double w = mat[i][c] / mat[r][c];
       REP(j, C+1) mat[i][j] -= mat[r][j] * w;
     row[c] = r++;
   REP(i, C) {
     int r = row[i];
     ans[i] = r == -1 ? 0 : mat[r][C] / mat[r][i];
```

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```
Gauss Xor
const 11 MAX = 1e9;
const ll LOG_MAX = 64 - __builtin_clzll((11)MAX);
struct gauss{
   vector<11> vet;
   gauss(){}
   gauss(ll val){
       if(val) vet.push_back(val);
   void add(ll val){
       int sig = LOG_MAX;
       for(int i = 0; i < (int)vet.size(); i++){</pre>
          while(!(vet[i] & (1LL << sig))) sig--;</pre>
          if(val & (1LL << sig)) val ^= vet[i];</pre>
       if(!val) return;
       sig = LOG_MAX;
       while(!(val & (1 << sig))) sig--;</pre>
       for(auto &x : vet) if(x & (1LL << sig)) x ^= val;</pre>
       vet.push_back(val);
       for(int i = (int)vet.size() - 2; i >= 0 && vet[i]
          < vet[i + 1]; i--){
          swap(vet[i], vet[i + 1]);
   }
};
Graphs
Dinic
const int N = 100005;
const int E = 2000006;
vector<int> g[N];
int ne;
struct Edge{
   int from, to;
   11 flow, cap;
} edge[E];
int lvl[N], vis[N], pass, start = N-2, target = N-1;
int qu[N], qt, px[N];
11 run(int s, int sink, ll minE){
   if(s == sink) return minE;
   11 \text{ ans} = 0;
   for(; px[s] < (int)g[s].size(); px[s]++){</pre>
       int e = g[s][ px[s] ];
       auto &v = edge[e], &rev = edge[e^1];
       if(lvl[v.to] != lvl[s]+1 || v.flow >= v.cap)
          continue; // v.cap - v.flow < lim</pre>
       11 tmp = run(v.to, sink,min(minE, v.cap-v.flow));
       v.flow += tmp, rev.flow -= tmp;
       ans += tmp, minE -= tmp;
       if(minE == 0) break;
   3
   return ans:
}
bool bfs(int source, int sink){
```

qt = 0;

qu[qt++] = source;

```
lvl[source] = 1;
   vis[source] = ++pass;
   for(int i = 0; i < qt; i++){
       int u = qu[i];
       px[u] = 0;
       if(u == sink) return true;
       for(int e : g[u]){
           auto v = edge[e];
           if(v.flow >= v.cap || vis[v.to] == pass)
              continue; // v.cap - v.flow < lim</pre>
           vis[v.to] = pass;
          lvl[v.to] = lvl[u]+1;
           qu[qt++] = v.to;
       }
   }
   return false;
}
11 flow(int source = start, int sink = target){
   11 \text{ ans} = 0;
   //for(lim = (1LL << 62); lim >= 1; lim /= 2)
   while(bfs(source, sink))
       ans += run(source, sink, oo);
   return ans;
}
void addEdge(int u, int v, ll c = 1, ll rc = 0){
   edge[ne] = \{u, v, 0, c\};
   g[u].push_back(ne++);
   edge[ne] = \{v, u, 0, rc\};
   g[v].push_back(ne++);
}
void reset_flow(){
   for(int i = 0; i < ne; i++)
       edge[i].flow = 0;
Min Cost Max Flow
const 11 oo = 1e18;
const int N = 222, E = 2 * 1000006;
vector<int> g[N];
int ne;
struct Edge{
   int from, to;
   ll cap, cost;
} edge[E];
int start = N-1, target = N-2, p[N];
int inqueue[N];
11 d[N];
bool spfa(int source, int sink){
   for(int i = 0; i < N; i++) d[i] = oo, inqueue[i] =</pre>
     0;
   d[source] = 0;
   queue<int> q;
   q.push(source);
   inqueue[source] = 1;
   while(!q.empty()){
       int u = q.front(); q.pop();
       inqueue[u] = 0;
```

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```
for(int e : g[u]){
                                                               for(int v : g[u]) if(v != p && v != big)
          auto v = edge[e];
                                                                   dfs(v, u, 0);
          if(v.cap > 0 \text{ and } d[u] + v.cost < d[v.to]){
              d[v.to] = d[u] + v.cost;
                                                               if(big != -1) dfs(big, u, 1);
              p[v.to] = e;
              if(!inqueue[v.to]){
                                                               add(u, p, big);
                 q.push(v.to);
                 inqueue[v.to] = 1;
                                                               for(auto x : q[u]){
                                                                   // answer all queries for this vx
          }
       }
   }
                                                               if(!keep){
                                                                   // Remove data from this subtree
   return d[sink] != oo;
}
                                                            }
// <max flow, min cost>
                                                            Junior e Falta de Ideias
pair<11, 11> mincost(int source = start, int sink =
  target){
                                                            #include <bits/stdc++.h>
   11 ans = 0, mf = 0;
   while(spfa(source, sink)){
                                                            #define ff first
       11 f = oo;
                                                            #define ss second
                                                            #define mp make_pair
       for(int u = sink; u != source; u = edge[ p[u] ].
                                                            using namespace std;
          f = min(f, edge[p[u]].cap);
       for(int u = sink; u != source; u = edge[ p[u] ].
                                                            typedef long long 11;
         from) {
                                                            vector<pair<int,int>> G[500005];
          edge[p[u]].cap -= f;
                                                            int subtree[500005], treesize, k;
          edge[ p[u] ^ 1 ].cap += f;
                                                            bool vis[500005]:
                                                            ll dist[500005], ans;
       mf += f;
                                                            int dfs(int v, int p){
       ans += f * d[sink];
                                                               subtree[v] = 1;
                                                               for(pair<int,int> x : G[v])
   return {mf, ans};
                                                                   if(x.ff != p \&\& !vis[x.ff]) subtree[v] += dfs(x.
}
                                                                     ff,v);
                                                               return subtree[v];
void addEdge(int u, int v, ll c, ll cost){
                                                            }
   edge[ne] = \{u, v, c, cost\};
   g[u].push_back(ne++);
                                                            int centroid(int v, int p){
   edge[ne] = \{v, u, 0, -cost\};
                                                               for(pair<int,int> x : G[v]){
   g[v].push_back(ne++);
                                                                   if(x.ff == p || vis[x.ff]) continue;
                                                                   if(subtree[x.ff]*2 > treesize) return centroid(x.
Small to Large
                                                                     ff.v):
void cnt_sz(int u, int p = -1){
                                                               return v;
   sz[u] = 1;
   for(int v : g[u]) if(v != p)
                                                            void procurar_ans(int v, int p, int d_atual, ll custo){
       cnt_sz(v, u), sz[u] += sz[v];
                                                               ans = min(ans, dist[k-d_atual] + custo);
}
                                                               if(d_atual == k) return;
                                                               for(pair<int,int> x : G[v]){
void add(int u, int p, int big = -1){
                                                                   if(!vis[x.ff] && x.ff != p)
                                                                      procurar_ans(x.ff,v,d_atual+1,custo+x.ss);
   // Update info about this vx in global answer
   for(int v : g[u]) if(v != p && v != big)
                                                            }
       add(v, u);
}
                                                            void atualiza_distancia(int v, int p, int d_atual, 11
                                                              custo){
void dfs(int u, int p, int keep){
                                                               dist[d_atual] = min(dist[d_atual], custo);
                                                               if(d_atual == k) return;
   int big = -1, mmx = -1;
                                                               for(pair<int,int> x : G[v]){
                                                                   if(!vis[x.ff] && x.ff != p)
   for(int v : g[u]) if(v != p \&\& sz[v] > mmx)
                                                                      atualiza_distancia(x.ff,v,d_atual+1,custo+x.
       mmx = sz[v], big = v;
```

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```
}
                                                                return 0;
}
                                                             }
                                                             Tarjan
void decomp(int v, int p){
   treesize = dfs(v,v);
                                                             int cnt = 0, root;
   // if(treesize < k) return;</pre>
                                                             void dfs(int u, int p = -1){
   int cent = centroid(v,v);
                                                                low[u] = num[u] = ++t;
   vis[cent] = 1;
                                                                 for(int v : g[u]){
                                                                    if(!num[v]){
   for(int i = 1; i <= treesize; i++)</pre>
                                                                        dfs(v, u);
       dist[i] = 1e18;
                                                                           if(u == root) cnt++;
                                                                        if(low[v] >= num[u]) u PONTO DE ARTICULAÇÃO;
   for(pair<int,int> x : G[cent]){
                                                                        if(low[v] > num[u]) ARESTA u->v PONTE;
       if(!vis[x.ff]){
                                                                       low[u] = min(low[u], low[v]);
          procurar_ans(x.ff,cent,1,x.ss);
          atualiza_distancia(x.ff,cent,1,x.ss);
                                                                    else if(v != p) low[u] = min(low[u], num[v]);
                                                                }
   }
                                                             }
   for(pair<int,int> x : G[cent]){
                                                             root PONTO DE ARTICULAÇÃO <=> cnt > 1
       if(!vis[x.ff])
           decomp(x.ff, cent);
                                                             void tarjanSCC(int u){
   }
                                                                low[u] = num[u] = ++cnt;
}
                                                                vis[u] = 1;
                                                                S.push_back(u);
int main(){
                                                                 for(int v : g[u]){
   int n,i,a,b;
                                                                    if(!num[v]) tarjanSCC(v);
                                                                    if(vis[v]) low[u] = min(low[u], low[v]);
   scanf("%d%d", &n,&k);
   for(i = 2; i \le n; i++){
                                                                if(low[u] == num[u]){
       scanf("%d%d", &a,&b);
                                                                    ssc[u] = ++ssc_cnt; int v;
       G[i].push_back(mp(a,b));
                                                                    do{
       G[a].push_back(mp(i,b));
                                                                        v = S.back(); S.pop_back(); vis[v] = 0;
   }
                                                                        ssc[v] = ssc_cnt;
   ans = 1e18;
                                                                    }while(u != v);
   decomp(1,-1);
                                                                }
   printf("%lld\n", ans == 1e18 ? -1 : ans);
                                                             Max Clique
   return 0;
                                                             long long adj[N], dp[N];
Kosaraju
                                                             for(int i = 0; i < n; i++){
                                                                 for(int j = 0; j < n; j++){
vector<int> g[N], gt[N], S;
                                                                    int x;
                                                                    scanf("%d",&x);
int vis[N], cor[N], tempo = 1;
                                                                    if(x \mid | i == j)
                                                                        adj[i] |= 1LL << j;
void dfs(int u){
   vis[u] = 1;
                                                             }
   for(int v : g[u]) if(!vis[v]) dfs(v);
   S.push_back(u);
                                                             int resto = n - n/2;
}
                                                             int C = n/2;
                                                             for(int i = 1; i < (1 << resto); i++){</pre>
int e:
void dfst(int u){
                                                                 int x = i;
                                                                 for(int j = 0; j < resto; j++)
   cor[u] = e;
   for(int v : gt[u]) if(!cor[v]) dfst(v);
                                                                    if(i & (1 << j))
                                                                       x \&= adj[j + C] >> C;
                                                                if(x == i){
int main(){
                                                                    dp[i] = __builtin_popcount(i);
                                                                }
   for(int i = 1; i <= n; i++) if(!vis[i]) dfs(i);</pre>
                                                             }
   e = 0:
                                                             for(int i = 1; i < (1 << resto); i++)</pre>
   reverse(S.begin(), S.end());
                                                                for(int j = 0; j < resto; j++)
   for(int u : S) if(!cor[u])
                                                                    if(i & (1 << j))</pre>
       e++, dfst(u);
                                                                        dp[i] = max(dp[i], dp[i ^ (1 << j)]);
```

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```
int maxCliq = 0;
                                                                   down[u].clear();
                                                               }
for(int i = 0; i < (1 << C); i++){
   int x = i, y = (1 << resto) - 1;
   for(int j = 0; j < C; j++)
                                                               reverse(S.begin(), S.end());
       if(i & (1 << j))
                                                                for(int u : S) if(u != root){
          x \&= adj[j] \& ((1 << C) - 1), y \&= adj[j] >>
                                                                   if(idom[u] != sdom[u]) idom[u] = idom[ idom[u] ];
            С;
                                                                   T[ idom[u] ].push_back(u);
   if(x != i) continue;
   maxCliq = max(maxCliq, __builtin_popcount(i) + dp[y
                                                               S.clear();
     ]);
                                                            Min Cost Matching
Dominator Tree
                                                            // Min cost matching
                                                            // O(n^2 * m)
vector<int> g[N], gt[N], T[N];
                                                            // n == nro de linhas
vector<int> S;
                                                            // m == nro de colunas
int dsu[N], label[N];
                                                            // n <= m | flow == n
int sdom[N], idom[N], dfs_time, id[N];
                                                            // a[i][j] = custo pra conectar i a j
                                                            vector < int > u(n + 1), v(m + 1), p(m + 1), way(m + 1);
vector<int> bucket[N];
                                                            for(int i = 1; i <= n; ++i){</pre>
vector<int> down[N];
                                                               p[0] = i;
                                                               int j0 = 0;
void prep(int u){
                                                               vector<int> minv(m + 1 , oo);
   S.push_back(u);
                                                               vector<char> used(m + 1 , false);
   id[u] = ++dfs_time;
   label[u] = sdom[u] = dsu[u] = u;
                                                                   used[j0] = true;
                                                                   int i0 = p[j0], delta = oo, j1;
   for(int v : g[u]){
                                                                   for(int j = 1; j \le m; ++j)
      if(!id[v])
                                                                       if(! used[j]){
          prep(v), down[u].push_back(v);
                                                                          int cur = a[i0][j] - u[i0] - v[j];
       gt[v].push_back(u);
                                                                          if(cur < minv[j])</pre>
   }
                                                                              minv[j] = cur, way[j] = j0;
}
                                                                          if(minv[j] < delta)</pre>
                                                                              delta = minv[j] , j1 = j;
int fnd(int u, int flag = 0){
                                                                      }
   if(u == dsu[u]) return u;
                                                                   for(int j = 0; j \le m; ++j)
   int v = fnd(dsu[u], 1), b = label[ dsu[u] ];
                                                                      if(used[j])
   if(id[ sdom[b] ] < id[ sdom[ label[u] ] ])</pre>
                                                                          u[p[j]] += delta, v[j] -= delta;
       label[u] = b;
                                                                       else
   dsu[u] = v;
                                                                          minv[j] -= delta;
   return flag ? v : label[u];
                                                                   j0 = j1;
}
                                                               }while(p[j0] != 0);
void build_dominator_tree(int root, int sz){
   // memset(id, 0, sizeof(int) * (sz + 1));
                                                                   int j1 = way[j0];
   // for(int i = 0; i <= sz; i++) T[i].clear();
                                                                   p[j0] = p[j1];
   prep(root);
                                                                   j0 = j1;
   reverse(S.begin(), S.end());
                                                               }while(j0);
                                                            }
   int w:
   for(int u : S){
                                                            // match[i] = coluna escolhida para linha i
       for(int v : gt[u]){
                                                            vector<int> match(n + 1);
          w = fnd(v);
                                                            for(int j = 1; j \le m; ++j)
          if(id[ sdom[w] ] < id[ sdom[u] ])</pre>
                                                               match[p[j]] = j;
              sdom[u] = sdom[w];
                                                            int cost = -v[0];
      gt[u].clear();
                                                            Strings
      if(u != root) bucket[ sdom[u] ].push_back(u);
                                                            Aho Corasick
       for(int v : bucket[u]){
                                                            map<char, int> to[N];
          w = fnd(v);
          if(sdom[w] == sdom[v]) idom[v] = sdom[v];
                                                            int ne = 1, term[N], fail[N];
          else idom[v] = w;
                                                            void add_string(char *str){
      bucket[u].clear();
                                                               int p = 0;
```

for(int i = 0; str[i]; i++){

for(int v : down[u]) dsu[v] = u;

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```
if(!to[p][ str[i] ]) to[p][ str[i] ] = ne++;
      p = to[p][ str[i] ];
                                                           void build_lcp(){ // lcp[i] = lcp(s[:i], s[:i+1])
   }
                                                               int k = 0;
   term[p] = 1;
}
                                                               for(int i = 0; i < n; i++) r[ sa[i] ] = i;
int go(int s, char c){
                                                               for(int i = 0; i < n; i++){
   while(s && !to[s].count(c)) s = fail[s];
                                                                  if(r[i] == n-1) k = 0;
   if(to[s].count(c)) return to[s][c];
   return s;
                                                                      int j = sa[r[i]+1];
                                                                      while(i+k < n && j+k < n && s[i+k] == s[j+k])
void init(){
                                                                  lcp[r[i]] = k;
   queue<int> q;
                                                                  if(k) k--;
   q.push(0);
                                                              }
                                                           }
   int u, v; char c;
   while(!q.empty()){
                                                           int comp_lcp(int i, int j){
      u = q.front(); q.pop();
                                                               if(i == j) return n - i;
                                                               if(r[i] > r[j]) swap(i, j);
       for(auto w : to[u]){
                                                               return min(lcp[k] for k in [r[i], r[j]-1]);
          tie(c, v) = w;
          q.push(v);
                                                           Z Algorithm
          if(u){
              fail[v] = go(fail[u], c);
                                                           vector<int> z_algo(const string &s) {
              term[v] |= term[ fail[v] ];
                                                               int n = s.size(), L = 0, R = 0;
          }
                                                               vector<int> z(n, 0);
      }
                                                               for(int i = 1; i < n; i++){</pre>
   }
                                                                  if(i \le R) z[i] = min(z[i-L], R - i + 1);
}
                                                                  while(z[i]+i < n \& s[z[i]+i] == s[z[i]])
                                                                      z[i]++;
Suffix Array
                                                                  if(i+z[i]-1 > R) L = i, R = i + z[i] - 1;
                                                               }
                                                               return z;
int n, sa[N], tsa[N], lcp[N], r[N], nr[N], c[N];
                                                           Prefix function/KMP
void sort(int k, int mx){
                                                           vector<int> preffix_function(const string &s){
   mx += 2:
                                                               int n = s.size();
   memset(c, 0, sizeof(int) * mx);
                                                               vector<int> b(n+1);
   for(int i = 0; i < n; i++) c[i + k < n ? r[i+k]+2 :
                                                              b[0] = -1;
     11++:
                                                               int i = 0, j = -1;
   partial_sum(c, c+mx, c);
                                                               while(i < n){
   int t;
                                                                  while(j >= 0 \&\& s[i] != s[j]) j = b[j];
   for(int i = 0; i < n; i++)</pre>
                                                                  b[++i] = ++j;
       t = sa[i]+k < n ? r[ sa[i]+k ]+1 : 0,
                                                               }
       tsa[c[t]++] = sa[i];
                                                               return b;
   memcpy(sa, tsa, sizeof(int) * n);
                                                           }
                                                           void kmp(const string &t, const string &p){
void build_sa(){
                                                               vector<int> b = preffix_function(p);
                                                               int n = t.size(), m = p.size();
   for(int i = 0; i < n; i++) sa[i] = i, r[i] = s[i];
                                                               int j = 0;
                                                               for(int i = 0; i < n; i++){
   int t = 300, a, b;
                                                                  while(j >= 0 && t[i] != p[j]) j = b[j];
   for(int sz = 1; sz < n; sz *= 2){
                                                                  j++;
       sort(sz, t), sort(0, t);
                                                                  if(j == m){
      t = nr[ sa[0] ] = 0;
                                                                      //patern of p found on t
       for(int i = 1; i < n; i++){
                                                                      j = b[j];
          a = sa[i]+sz < n ? r[ sa[i]+sz ] : -1;
                                                                  }
          b = sa[i-1]+sz < n ? r[ sa[i-1]+sz ] : -1;
                                                              }
          nr[ sa[i] ] = r[ sa[i] ] == r[ sa[i-1] ] && a
             == b ? t : ++t;
                                                           Min rotation
       if(t == n-1) break;
                                                           int min_rotation(int *s, int N) {
      memcpy(r, nr, sizeof(int) * n);
                                                             REP(i, N) s[N+i] = s[i];
   }
```

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```
int sufflink;
 int a = 0;
                                                               int num;
                                                            };
 REP(b, N) REP(i, N) {
   if (a+i == b \mid \mid s[a+i] < s[b+i]) \{ b += max(0, i-1);
                                                            int len;
      break; }
   if (s[a+i] > s[b+i]) \{ a = b; break; \}
                                                            char s[MAXN];
                                                            node tree[MAXN];
                                                            int num; // node 1 - root with len -1, node 2 - root
 return a:
                                                              with len 0
                                                            int suff; // max suffix palindrome
All palindrome
                                                            long long ans;
void manacher(char *s, int N, int *rad) {
                                                            bool addLetter(int pos) {
 static char t[2*MAX];
                                                               int cur = suff, curlen = 0;
 int m = 2*N - 1;
                                                               int let = s[pos] - 'a';
 REP(i, m) t[i] = -1;
                                                               while(true){
 REP(i, N) t[2*i] = s[i];
                                                                   curlen = tree[cur].len;
                                                                   if (pos-1 - curlen) = 0 && s[pos-1 - curlen] == s
 int x = 0;
                                                                     [pos])
 FOR(i, 1, m) {
                                                                      break:
   int &r = rad[i] = 0;
                                                                   cur = tree[cur].sufflink;
   if (i <= x+rad[x]) r = min(rad[x+x-i], x+rad[x]-i);</pre>
   while (i-r-1 >= 0 \& i+r+1 < m \& t[i-r-1] == t[i+r]
                                                               if (tree[cur].next[let]) {
     +1]) ++r;
                                                                   suff = tree[cur].next[let];
   if (i+r >= x+rad[x]) x = i;
                                                                   return false;
                                                               }
 REP(i, m) if (i-rad[i] == 0 || i+rad[i] == m-1) ++rad[
                                                               num++:
                                                               suff = num;
 REP(i, m) rad[i] /= 2;
                                                               tree[num].len = tree[cur].len + 2;
                                                               tree[cur].next[let] = num;
Manacher
                                                               if (tree[num].len == 1){
// odd[i] = largest palindrome with center in character
                                                                   tree[num].sufflink = 2;
                                                                   tree[num].num = 1;
// even[i] = largest palindrome with center between
                                                                   return true;
  chars i-1 and i(0-based)
                                                               }
// pair<even, odd>
pair<vector<int>, vector<int>> build(const string &s){
                                                               while (true){
  // 0(n)
                                                                   cur = tree[cur].sufflink;
   int n = (int)s.size();
                                                                   curlen = tree[cur].len;
   vector<int> vet[2];
                                                                   if(pos-1 - curlen) == 0 \&\& s[pos-1 - curlen] == s[
   for(int add = 0; add < 2; add++){</pre>
                                                                       tree[num].sufflink = tree[cur].next[let];
      vet[add].resize(n);
                                                                      break;
       for(int i = 0, l = 0, r = -1; i < n; i++){
                                                                   }
          int k = i > r? add : min(vet[add][l + r - i
            + !add], r - i + !add);
          while(i - k - !add >= 0 && i + k < n && s[i -
                                                               tree[num].num = 1 + tree[tree[num].sufflink].num;
             k - !add] == s[i + k]) k++;
          vet[add][i] = k--;
                                                               return true;
          if(i + k > r){
                                                            }
              r = i + k;
              1 = i - k - !add;
                                                            void initTree() {
                                                               num = 2; suff = 2;
      }
                                                               tree[1].len = -1; tree[1].sufflink = 1;
   }
                                                               tree[2].len = 0; tree[2].sufflink = 1;
                                                            }
   return make_pair(vet[0], vet[1]);
                                                            int main() {
Palindromic Tree
                                                               initTree();
const int MAXN = 105000:
                                                               for (int i = 0; i < len; i++) {</pre>
struct node {
                                                                   addLetter(i);
   int next[26];
                                                               }
   int len;
```

```
return 0;
}
Suffix Automaton
map<char, int> to[2*N];
int link[2*N], len[2*N], last = 0, sz = 1;
void add_letter(char c){
   int p = last;
   last = sz++;
   len[last] = len[p] + 1;
   for(; !to[p][c]; p = link[p]) to[p][c] = last;
   if(to[p][c] == last){
       link[last] = 0;
       return;
   }
   int u = to[p][c];
   if(len[u] == len[p]+1){
       link[last] = u;
       return;
   }
   int c1 = sz++;
   to[c1] = to[u];
   link[c1] = link[u];
   len[c1] = len[p]+1;
   link[last] = link[u] = c1;
   for(; to[p][c] == u; p = link[p]) to[p][c] = c1;
Suffix Tree
namespace sf {
// const int NS = ; const int N = * 2;
int cn, cd, ns, en = 1, lst;
string S[NS]; int si = -1;
vector<int> sufn[N]; // sufn[si][i] no do sufixo S[si][i
  . . . ]
struct node {
   int 1, r, si, p, suf;
   map<char, int> adj;
   node() : l(0), r(-1), suf(0), p(0) {}
   node(int L, int R, int S, int P) : 1(L), r(R), si(S)
      p(P) \{ \}
   inline int len() { return r - l + 1; }
   inline int operator[](int i) { return S[si][l + i];
   inline int& operator()(char c) { return adj[c]; }
inline int new_node(int L, int R, int S, int P) { t[en]
  = node(L, R, S, P); return en++; }
void add_string(string s) {
   s += '; S[++si] = s; sufn[si].resize(s.size() + 1)
      : cn = cd = 0:
   int i = 0; const int n = s.size();
   for(int j = 0; j < n; j++)
       for(; i <= j; i++) {</pre>
          if(cd == t[cn].len() \& t[cn](s[j])) { cn = t}
             [cn](s[j]); cd = 0; 
          if(cd < t[cn].len() \&\& t[cn][cd] == s[j]) {
              cd++;
              if(j < s.size() - 1) break;</pre>
              else {
                 if(i) t[lst].suf = cn;
                 for(; i <= j; i++) { sufn[si][i] = cn;</pre>
                   cn = t[cn].suf; }
          } else if(cd == t[cn].len()) {
```

```
sufn[si][i] = en;
              if(i) t[lst].suf = en; lst = en;
              t[cn](s[j]) = new_node(j, n - 1, si, cn);
              cn = t[cn].suf; cd = t[cn].len();
          } else {
              int mid = new_node(t[cn].1, t[cn].1 + cd -
                 1, t[cn].si, t[cn].p);
              t[t[cn].p](t[cn][0]) = mid;
              if(ns) t[ns].suf = mid;
              if(i) t[lst].suf = en; lst = en;
              sufn[si][i] = en;
              t[mid](s[j]) = new_node(j, n - 1, si, mid)
              t[mid](t[cn][cd]) = cn;
              t[cn].p = mid; t[cn].l += cd; cn = t[mid].
                p;
              int g = cn? j - cd : i + 1; cn = t[cn].suf
              while(g < j \&\& g + t[t[cn](S[si][g])].len
                () <= j) {
                 cn = t[cn](S[si][g]); g += t[cn].len();
              if(g == j) \{ ns = 0; t[mid].suf = cn; cd =
                 t[cn].len(); }
              else { ns = mid; cn = t[cn](S[si][g]); cd
                = j - g;  }
          }
      }
   }
};
```

Geometry

2D basics

```
typedef double cod;
double eps = 1e-7;
bool eq(cod a, cod b){ return abs(a - b) <= eps; }</pre>
struct vec{
   cod x, y; int id;
   vec(cod a = 0, cod b = 0) : x(a), y(b) {}
   vec operator+(const vec &o) const{
       return \{x + o.x, y + o.y\};
   vec operator-(const vec &o) const{
       return \{x - o.x, y - o.y\};
   vec operator*(cod t) const{
       return {x * t, y * t};
   vec operator/(cod t) const{
       return {x / t, y / t};
   cod operator*(const vec &o) const{ // cos
       return x * o.x + y * o.y;
   cod operator^(const vec &o) const{ // sin
       return x * o.y - y * o.x;
   bool operator==(const vec &o) const{
       return eq(x, o.x) && eq(y, o.y);
   bool operator<(const vec &o) const{</pre>
       if(!eq(x, o.x)) return x < o.x;
       return y < o.y;</pre>
   cod cross(const vec &a, const vec &b) const{
```

```
return (a-(*this)) ^ (b-(*this));
   }
   int ccw(const vec &a, const vec &b) const{
       cod tmp = cross(a, b);
       return (tmp > eps) - (tmp < -eps);</pre>
   cod dot(const vec &a, const vec &b) const{
       return (a-(*this)) * (b-(*this));
   cod len() const{
       return sqrt(x * x + y * y); // <</pre>
   double angle(const vec &a, const vec &b) const{
       return atan2(cross(a, b), dot(a, b));
   double tan(const vec &a, const vec &b) const{
       return cross(a, b) / dot(a, b);
   }
   vec unit() const{
       return operator/(len());
   int quad() const{
       if(x > 0 \& y >= 0) return 0;
       if(x \le 0 \& y > 0) return 1;
       if(x < 0 \& y <= 0) return 2;
       return 3;
   }
   bool comp(const vec &a, const vec &b) const{
       return (a - *this).comp(b - *this);
   bool comp(vec b){
       if(quad() != b.quad()) return quad() < b.quad();</pre>
       if(!eq(operator^(b), 0)) return operator^(b) > 0;
       return (*this) * (*this) < b * b;</pre>
   template<class T>
   void sort_by_angle(T first, T last) const{
       std::sort(first, last, [=](const vec &a, const
         vec &b){
          return comp(a, b);
       });
   }
   vec rot90() const{ return {-y, x}; }
   vec rot(double a) const{
       return \{\cos(a)*x - \sin(a)*y, \sin(a)*x + \cos(a)*y\};
   }
   vec proj(const vec &b) const{ // proj of *this onto
       cod k = operator*(b) / (b * b);
       return b * k;
   // proj of (*this) onto the plane orthogonal to b
   vec rejection(vec b) const{
       return (*this) - proj(b);
};
struct line{
   cod a, b, c; vec n;
   line(vec q, vec w){ // q.cross(w, (x, y)) = 0
       a = -(w.y-q.y);
       b = w.x-q.x;
       c = -(a * q.x + b * q.y);
       n = \{a, b\};
   }
   cod dist(const vec &o) const{
       return abs(eval(o)) / n.len();
```

```
bool contains(const vec &o) const{
      return eq(a * o.x + b * o.y + c, 0);
   cod dist(const line &o) const{
      if(!parallel(o)) return 0;
       if(!eq(o.a * b, o.b * a)) return 0;
       if(!eq(a, 0))
          return abs(c - o.c * a / o.a) / n.len();
       if(!eq(b, 0))
          return abs(c - o.c * b / o.b) / n.len();
      return abs(c - o.c);
   bool parallel(const line &o) const{
      return eq(n ^ o.n, 0);
   bool operator==(const line &o) const{
       if(!eq(a*o.b, b*o.a)) return false;
       if(!eq(a*o.c, c*o.a)) return false;
       if(!eq(c*o.b, b*o.c)) return false;
       return true;
   bool intersect(const line &o) const{
      return !parallel(o) || *this == o;
   vec inter(const line &o) const{
       if(parallel(o)){
          if(*this == o){ }
          else{ /* dont intersect */ }
      }
       auto tmp = n ^ o.n;
      return {(o.c*b -c*o.b)/tmp, (o.a*c -a*o.c)/tmp};
   vec at_x(cod x) const{
      return \{x, (-c-a*x)/b\};
   }
   vec at_y(cod y) const{
      return {(-c-b*y)/a, y};
   cod eval(const vec &o) const{
      return a * o.x + b * o.y + c;
};
struct segment{
   vec p, q;
   segment(vec a = vec(), vec b = vec()): p(a), q(b) {}
   bool onstrip(const vec &o) const{ // onstrip strip
       return p.dot(o, q) >= -eps && q.dot(o, p) >= -eps
   }
   cod len() const{
      return (p-q).len();
   cod dist(const vec &o) const{
       if(onstrip(o)) return line(p, q).dist(o);
       return min((o-q).len(), (o-p).len());
   bool contains(const vec &o) const{
      return eq(p.cross(q, o), 0) && onstrip(o);
   bool intersect(const segment &o) const{
      if(contains(o.p)) return true;
       if(contains(o.q)) return true;
       if(o.contains(q)) return true;
       if(o.contains(p)) return true;
```

```
return p.ccw(q, o.p) * p.ccw(q, o.q) == -1
      && o.p.ccw(o.q, q) * o.p.ccw(o.q, p) == -1;
   }
   bool intersect(const line &o) const{
      return o.eval(p) * o.eval(q) <= 0;</pre>
   }
   cod dist(const segment &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q)
          || o.onstrip(p) || o.onstrip(q))
              return line(p, q).dist(line(o.p, o.q));
       else if(intersect(o)) return 0;
       return min(min(dist(o.p), dist(o.q)),
                min(o.dist(p), o.dist(q)));
   cod dist(const line &o) const{
       if(line(p, q).parallel(o))
          return line(p, q).dist(o);
       else if(intersect(o)) return 0;
       return min(o.dist(p), o.dist(q));
};
struct hray{
   vec p, q;
   hray(vec a = vec(), vec b = vec()): p(a), q(b){}
   bool onstrip(const vec &o) const{ // onstrip strip
      return p.dot(q, o) >= -eps;
   cod dist(const vec &o) const{
       if(onstrip(o)) return line(p, q).dist(o);
       return (o-p).len();
   bool intersect(const segment &o) const{
       if(!o.intersect(line(p,q))) return false;
       if(line(o.p, o.q).parallel(line(p,q)))
          return contains(o.p) || contains(o.q);
       return contains(line(p,q).inter(line(o.p,o.q)));
   }
   bool contains(const vec &o) const{
      return eq(line(p, q).eval(o), 0) && onstrip(o);
   cod dist(const segment &o) const{
       if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || onstrip(o.q))
              return line(p, q).dist(line(o.p, o.q));
          return o.dist(p);
       else if(intersect(o)) return 0;
       return min(min(dist(o.p), dist(o.q)),
                o.dist(p));
   bool intersect(const hray &o) const{
       if(!line(p, q).parallel(line(o.p, o.q)))
          return false;
       auto pt = line(p, q).inter(line(o.p, o.q));
      return contains(pt) && o.contains(pt); // <<</pre>
   bool intersect(const line &o) const{
       if(line(p, q).parallel(o)) return line(p, q)== o;
       if(o.contains(p) || o.contains(q)) return true;
       return (o.eval(p) >= -eps)^(o.eval(p)<o.eval(q));</pre>
       return contains(o.inter(line(p, q)));
   cod dist(const line &o) const{
       if(line(p,q).parallel(o))
```

```
return line(p,q).dist(o);
       else if(intersect(o)) return 0;
      return o.dist(p);
   cod dist(const hray &o) const{
      if(line(p, q).parallel(line(o.p, o.q))){
          if(onstrip(o.p) || o.onstrip(p))
              return line(p,q).dist(line(o.p, o.q));
          return (p-o.p).len();
       else if(intersect(o)) return 0;
      return min(dist(o.p), o.dist(p));
};
double heron(cod a, cod b, cod c){
   cod s = (a + b + c) / 2;
   return sqrt(s * (s - a) * (s - b) * (s - c));
Circle line intersection
// intersection of line a * x + b * y + c = 0
// and circle centered at the origin with radius r
double r, a, b, c; // given as input
double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
if(c*c > r*r*(a*a+b*b)+EPS)
   puts("no points");
else if(abs(c*c - r*r*(a*a+b*b)) < EPS){
   puts("1 point");
   cout << x0 << '' << y0 << '\n';
}
else {
   double d = r*r - c*c/(a*a+b*b);
   double mult = sqrt (d / (a*a+b*b));
   double ax, ay, bx, by;
   ax = x0 + b * mult;
   bx = x0 - b * mult;
   ay = y0 - a * mult;
   by = y0 + a * mult;
   puts ("2 points");
   cout<<ax<<' '<<ay<<'\n'<<bx<<' '<<by<<'\n';
```

Circle Circle intersection

Assume that the first circle is centered at the origin and second at (x^2, y^2) . Find circle line intersection of first circle and line Ax + By + C = 0, where $A = -2x_2$, $B = -2y_2$, $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$.

Be aware of corner case with two circles centered at the same point.

Tangents of two circles

```
// solve first for same circle(and infinitely many
  tangents)
// Find up to four tangents of two circles
void tangents(pt c, double r1, double r2, vector<line> &
  ans) {
  double r = r2 - r1;
  double z = c.x * c.x + c.y * c.y;
  double d = z - r * r;
  if(d < -EPS) return;
  d = sqrt(abs(d));
  line 1;
  l.a = (c.x * r + c.y * d) / z;
  l.b = (c.y * r - c.x * d) / z;</pre>
```

```
1.c = r1;
   ans.push_back (1);
}
vector<line> tangents(circle a, circle b){
   vector<line> ans;
   pt aux = a.center - b.center;
   for(int i = -1; i \le 1; i += 2)
       for(int j = -1; j <= 1; j += 2)
          tangents(aux, a.r * i, b.r * j, ans);
   for(size_t i = 0; i < ans.size(); ++i)</pre>
       ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
   return ans;
}
Convex Hull
vector<vec> monotone_chain_ch(vector<vec> P){
   sort(P.begin(), P.end());
   vector<vec> L, U;
   for(auto p : P){
       while(L.size() >= 2 && L[L.size() - 2].cross(L.
         back(), p) < 0)
          L.pop_back();
      L.push_back(p);
   }
   reverse(P.begin(), P.end());
   for(auto p : P){
      while(U.size() >= 2 && U[U.size() - 2].cross(U.
         back(), p) < 0)
          U.pop_back();
      U.push_back(p);
   }
   L.pop_back(), U.pop_back();
   L.reserve(L.size() + U.size());
   L.insert(L.end(), U.begin(), U.end());
   return L;
Check point inside polygon
bool below(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) <= 0;</pre>
}
bool above(const vector<vec> &vet, vec p){
   auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
   if(it == vet.begin()) return *it == p;
   return prev(it)->cross(*it, p) >= 0;
// lowerhull, upperhull and point, borders included
bool inside_poly(const vector<vec> &lo, const vector<vec</pre>
  > &hi, vec p){
```

return below(hi, p) && above(lo, p);

Check point inside polygon without lower/upper hull

```
// borders included
// must not have 3 colinear consecutive points
bool inside_poly(const vector<vec> &v, vec p){
   if(v[0].ccw(v[1], p) < 0) return false;</pre>
   if(v[0].ccw(v.back(), p) > 0) return 0;
   if(v[0].ccw(v.back(), p) == 0)
      return v[0].dot(p, v.back()) >= 0
          && v.back().dot(p, v[0]) >= 0;
   int L = 1, R = (int)v.size() - 1, ans = 1;
   while(L <= R){</pre>
       int mid = (L+R)/2;
       if(v[0].ccw(v[mid], p) >= 0) ans = mid, L = mid
       else R = mid-1;
   }
   return v[ans].ccw(v[(ans+1)%v.size()], p) >= 0;
Minkowski sum
vector<vec> mk(const vector<vec>&a,const vector<vec>&b){
   int i = 0, j = 0;
   for(int k = 0; k < (int)a.size(); k++)if(a[k] < a[i])
      i = k;
   for(int k = 0; k < (int)b.size(); k++)if(b[k] < b[j])
      j = k;
   vector<vec> c;
   c.reserve(a.size() + b.size());
   for(int k = 0; k < int(a.size()+b.size()); k++){</pre>
       vec pt{a[i] + b[j]};
       if((int)c.size() >= 2
       && c[c.size()-2].ccw(c.back(), pt) == 0)
          c.pop_back();
      c.push_back(pt);
      int q = i+1, w = j+1;
      if(q == int(a.size())) q = 0;
      if(w == int(b.size())) w = 0;
      if(c.back().ccw(a[i]+b[w], a[q]+b[j]) < 0) i = q;
      else j = w;
   }
   c.shrink_to_fit();
   return c;
```

Geo Notes

Center of mass

System of points(2D/3D): Mass weighted average of points. **Frame(2D/3D):** Get middle point of each segment solve as previously.

Triangle: Average of vertices.

2D Polygon: Compute **signed** area and center of mass of triangle $((0,0), p_i, p_{i+1})$. Then solve as system of points.

Polyhedron surface: Solve each face as a 2D polygon(be aware of (0, 0)) then replace each face with its center of mass and solve as system of points.

Tetrahedron(Triangular pyramid): As triangles, its the average of points.

Polyhedron: Can be done as 2D polygon, but with tetrahe-

dralization intead of triangulation.

Pick's Theorem

Given a polygon without self-intersections and all its vertices on integer coordinates in some 2D grid. Let A be its area, I the number of points with interger coordinates stricly inside the polygon and B the number of points with interger coordinates in the border of the polygon. The following formula holds: $A = I + \frac{B}{2} - 1$.

Miscellaneous

LIS

```
multiset<int> S;
for(int i = 0; i < n; i++){
   auto it = S.upper_bound(a[i]); // low for inc
   if(it != S.end()) S.erase(it);
   S.insert(a[i]);
}
ans = S.size();
DSU rollback
#include <bits/stdc++.h>
using namespace std;
struct DSU{
   vector<int> sz, p, change;
   vector<tuple<int, int, int>> modifications;
   vector<size_t> saves;
   bool bipartite;
   DSU(int n): sz(n+1, 1), p(n+1), change(n+1),
     bipartite(true){
       iota(p.begin(), p.end(), 0);
   }
   void add_edge(int u, int v){
       if(!bipartite) return;
       int must_change = get_colour(u) == get_colour(v);
       int a = rep(u), b = rep(v);
       if(sz[a] < sz[b]) swap(a, b);
       if(a != b){
          p[b] = a;
          modifications.emplace_back(b, change[b],
            bipartite);
          change[b] ^= must_change;
          sz[a] += sz[b];
       else if(must_change){
          modifications.emplace_back(0, change[0],
            bipartite);
          bipartite = false;
   }
   int rep(int u){
      return p[u] == u ? u : rep(p[u]);
   }
   int get_colour(int u){
       if(p[u] == u) return change[u];
       return change[u] ^ get_colour(p[u]);
   }
```

```
void reset(){
       modifications.clear();
       saves.clear();
       iota(p.begin(), p.end(), 0);
       fill(sz.begin(), sz.end(), 1);
       fill(change.begin(), change.end(), 0);
       bipartite = true;
   void rollback(){
       int u = get<0>(modifications.back());
       tie(ignore, change[u], bipartite) = modifications
         .back():
       sz[ p[u] ] -= sz[u];
       p[u] = u;
       modifications.pop_back();
   }
   void reload(){
       while(modifications.size() > saves.back())
          rollback();
       saves.pop_back();
   }
   void save(){
       saves.push_back(modifications.size());
   }
};
const int N = 100005;
const int B = 318;
int n, m, q;
int x[N], y[N], 1[N], r[N], ans[N];
vector<int> qu[N];
int brute(int lef, int rig, DSU &s){
   s.save();
   for(int i = lef; i <= rig; i++)</pre>
       s.add_edge(x[i], y[i]);
   int ret = s.bipartite;
   s.reload();
   return ret;
}
int main(){
   scanf("%d %d %d", &n, &m, &q);
    for(int i = 1; i <= m; i++)</pre>
       scanf("%d %d", x+i, y+i);
   DSU s(n);
   for(int i = 0; i < q; i++){
       scanf("%d %d", l+i, r+i);
       if(r[i] - 1[i] \le B + 10)
           ans[i] = brute(l[i], r[i], s);
       else qu[l[i] / B].push_back(i);
   }
   for(int i = 0; i <= m / B; i++){</pre>
       sort(qu[i].begin(), qu[i].end(),[](int a, int b){
          return r[a] < r[b];</pre>
       });
       s.reset();
```

```
int R = (i+1)*B-1;
                                                             cout << ans << '\n';</pre>
       for(int id : qu[i]){
          while(R < r[id]) ++R, s.add_edge(x[R], y[R]);
                                                             return 0;
          s.save();
          for(int k = 1[id]; k < (i+1)*B; k++)
                                                            Rand
              s.add_edge(x[k], y[k]);
          ans[id] = s.bipartite;
                                                            #include <random>
          s.reload();
                                                            #include <chrono>
   }
                                                            cout << RAND_MAX << endl;</pre>
                                                            mt19937 rng(chrono::steady_clock::now().time_since_epoch
   for(int i = 0; i < q; i++)
                                                              ().count());
      printf("%s\n",ans[i] ? "Possible":"Impossible");
                                                            vector<int> permutation(N);
}
Buildings
                                                            iota(permutation.begin(), permutation.end(), 0);
                                                            shuffle(permutation.begin(), permutation.end(), rng);
// count the number of circular arrays
// of size m, with elements on range
// [1, c**(x*x)]
                                                            iota(permutation.begin(), permutation.end(), 0);
#include<bits/stdc++.h>
using namespace std;
                                                            for(int i = 1; i < N; i++){
                                                               swap(permutation[i], permutation[
                                                                 uniform_int_distribution<int>(0, i)(rng)]);
#define debug(x) cerr << fixed << #x << " = " << x <<
  endl;
                                                            Klondike
const int MOD = 1e9 + 7, MAX = 1e5 + 5;
int dp[MAX];
                                                            // minimum number of moves to make
                                                            // all elements equal
inline int add(int a, int b) {
                                                            // move: change a segment of equal value
 return a + b >= MOD ? a + b - MOD : a;
                                                            // elements to any value
inline int sub(int a, int b) {
                                                            int v[305];
 return a - b < 0? a - b + MOD: a;
                                                            int dp[305][305];
                                                            int rec[305][305];
inline int mult(int a, int b) {
 return (1LL * a * b) % MOD;
                                                            int f(int 1, int r){
                                                             if(r == 1) return 1;
int f_exp(int x, int exp) {
                                                             if(r < 1) return 0;</pre>
 if(exp == 0) return 1;
                                                             if(dp[l][r] != -1) return dp[l][r];
 else if(exp & 1) return mult(x, f_exp(x, exp - 1));
                                                             int ans = f(1+1, r) + 1;
 return f_exp(mult(x, x), exp / 2);
                                                             for(int i = l+1; i <= r; i++)</pre>
                                                               if(v[i] == v[l])
inline int inv(int x) {
                                                                 ans = min(ans, f(1, i - 1) + f(i+1, r));
 return f_exp(x, MOD - 2);
                                                             return dp[l][r] = ans;
int main() {
 ios::sync_with_stdio(false);
 cin.tie(NULL); cout.tie(NULL);
                                                            int main() {
                                                             int n, m;
 int n, m, c;
                                                             memset(dp, -1, sizeof dp);
 cin >> n >> m >> c;
                                                             scanf("%d %d",&n , &m);
                                                             for(int i = 0; i < n; i++){
 int x = f_{exp}(c, n * n);
                                                               scanf("%d",v+i);
 int ans = f_{exp}(x, m);
                                                               if(i && v[i] == v[i-1]){
 for(int i = 1; i <= m; i++) {</pre>
                                                                 i--;
   if(m \% i == 0) {
                                                                 n--;
     int y = f_{exp}(x, i);
                                                               }
     for(int j = 1; j < i; j++) {
      if(i % j == 0)
                                                             printf("%d\n",f(0, n-1) - 1);
        y = sub(y, mult(j, dp[j]));
                                                             // printf("%d\n",rec[0][n-1] );
                                                             // printf("%d\n",rec[1][n-1] );
     dp[i] = mult(y, inv(i));
                                                              // printf("%d\n",rec[2][n-3] );
     ans = sub(ans, mult(i - 1, dp[i]));
```

int main() {

```
Hilbert Order
// maybe use B = n / sqrt(q)
inline int64_t hilbertOrder(int x, int y, int pow = 21,
  int rotate = 0) {
   if(pow == 0) return 0;
   int hpow = 1 << (pow-1);</pre>
   int seg = (x < hpow) ? (
       (y < hpow) ? 0 : 3
   ):(
       (y < hpow) ? 1 : 2
   ) •
   seg = (seg + rotate) & 3;
   const int rotateDelta[4] = \{3, 0, 0, 1\};
   int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
   int nrot = (rotate + rotateDelta[seg]) & 3;
   int64_t subSquareSize = int64_t(1) << (2*pow - 2);</pre>
   int64_t ans = seg * subSquareSize;
   int64_t add = hilbertOrder(nx, ny, pow-1, nrot);
   ans += (seg == 1 || seg == 2) ? add : (subSquareSize
      - add - 1):
   return ans;
}
Modular Factorial
// Compute (1*2*...*(p-1)*1*(p+1)*(p+2)*..*n) % p
// in O(p*lg(n))
int factmod(int n, int p){
   int ans = 1;
   while (n > 1) {
       for(int i = 2; i <= n % p; i++)</pre>
          ans = (ans * i) % p;
      n /= p;
      if(n \% 2) ans = p - ans;
   }
   return ans % p;
}
int fac_pow(int n, int p){
   int ans = 0;
   while(n) n \neq p, ans += n;
   return ans;
}
int C(int n, int k, int p){
   if(fac_pow(n, p) > fac_pow(n-k, p) + fac_pow(k, p))
      return 0:
   int tmp = factmod(k, p) * factmod(n-k, p);
   return (f_exp(tmp, p-2, p) * factmod(n, p)) % p;
Enumeration all submasks of a bitmask
// loop through all submask of a given bitmask
// it does not include mask 0
for(int sub = mask; sub; sub = (sub-1)&mask){
Slope Trick
///By woqja125, contest: Codeforces Round #371 (Div. 1),
  problem: (C) Sonya and Problem Wihtout a Legend,
  Accepted. #
```

#include <stdio.h>

#include <queue>

```
int n, t;
   long long ans = 0;
   std::priority_queue<int> Q;
   scanf("%d%d", &n, &t);
   Q.push(t);
   for(int i = 1; i < n; i++) {</pre>
       scanf("%d", &t); t -= i;
       Q.push(t);
       if(Q.top() > t) {
          ans += Q.top() - t;
          Q.pop();
          Q.push(t);
       }
   printf("%lld", ans);
   return 0;
Fast IO
#define pc(x) putchar_unlocked(x)
#define gc(x) getchar_unlocked(x)
inline void scan_int(int &x){
   register int c = gc();
   x = 0;
   int neg = 0;
   for(; ((c < '0' || c > '9') && c != '-'); c = gc());
   if(c == '-'){
      neg = 1;
       c = gc();
   }
   for(; c >= '0' && c <= '9'; c = gc())</pre>
      x = (x << 1) + (x << 3) + c - '0';
   if(neg) x = -x;
}
inline void print_int(int n){
   int rev = 0, count = 0, neg;
   if(n == 0){
      pc('0');
       return;
   if(n < 0) n = -n, neg = 1;
   while(n % 10 == 0) count++, n /= 10;
   for(rev = 0; n != 0; n /= 10)
       rev = (rev << 3) + (rev << 1) + n % 10;
   if(neg) pc('-');
   while(rev != 0) pc(rev % 10 + '0'), rev /= 10;
   while(count--) pc('0');
   pc('\n');
}
inline void print_string(char *str){
   while(*str) pc(*str++);
   pc('\n');
Knapsack Bounded with Cost
// menor custo para conseguir peso ate M usando N tipos
  diferentes de elementos, sendo que o i-esimo elemento
  pode ser usado b[i] vezes, tem peso w[i] e custo c[i]
// O(N * M)
int b[N], w[N], c[N];
MinQueue Q[M]
int d[M] //d[i] = custo minimo para conseguir peso i
```

```
for(int i = 0; i <= M; i++) d[i] = i ? oo : 0;
for(int i = 0; i < N; i++){
    for(int j = 0; j < w[i]; j++)
        Q[j].clear();
    for(int j = 0; j <= M; j++){
        q = Q[j % w[i]];
        if(q.size() >= q) q.pop();
        q.add(c[i]);
        q.push(d[j]);
        d[j] = q.getmin();
    }
}
```

LCA < O(nlgn), O(1)>

```
int start[N], dfs_time;
int tour[2*N], id[2*N];

void dfs(int u){
    start[u] = dfs_time;
    id[dfs_time] = u;
    tour[dfs_time++] = start[u];
    for(int v : g[u]){
        dfs(v);
        id[dfs_time] = u;
        tour[dfs_time++] = start[u];
    }
}

int LCA(int u, int v){
    if(start[u] > start[v]) swap(u, v);
    return id[min(tour[k]for k in [start[u],start[v]])];
}
```

Buffered reader

```
// source: https://github.com/ngthanhtrung23/
 ACM_Notebook_new/blob/master/buffered_reader.h
int INP.AM.REACHEOF;
#define BUFSIZE (1<<12)</pre>
char BUF[BUFSIZE+1], *inp=BUF;
#define GETCHAR(INP) { \
   if(!*inp && !REACHEOF) { \
      memset(BUF,0,sizeof BUF);\
       int inpzzz = fread(BUF,1,BUFSIZE,stdin);\
       if (inpzzz != BUFSIZE) REACHEOF = true;\
      inp=BUF; \
   } \
   INP=*inp++; \
#define DIG(a) (((a)>='0')&&((a)<='9'))
#define GN(j) { \
   AM=0; \
   GETCHAR(INP); while(!DIG(INP) && INP!='-') GETCHAR(
   if (INP=='-') {AM=1;GETCHAR(INP);} \
   j=INP-'0'; GETCHAR(INP); \
   while(DIG(INP)){j=10*j+(INP-'0');GETCHAR(INP);} \
   if (AM) j=-j;\
```

Burnside's Lemma

Let (G, \oplus) be a finite group that acts on a set X. It should hold that $e_g * x = x$ and $g_1 * (g_2 * x) = (g_1 \oplus g_2) * x$, $\forall x \in X, g_1, g_2 \in G$. For each $g \in G$ let $X^g = \{x \in X \mid g * x = x\}$. The number of orbits its given by:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

Wilson's Theorem

 $(n-1)! = -1 \mod n \iff n \text{ is prime}$

Fibonacci

- $F_{n-1}F_{n+1} F_n^2 = (-1)^n$
- $F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$
- $GCD(F_n, F_m) = F_{GCD(n,m)}$
- $F_n = \frac{(\frac{1+\sqrt{5}}{2})^n (\frac{1-\sqrt{5}}{2})^n}{\sqrt{5}}$

Kirchhoff's Theorem

Laplacian matrix is L = D - A, where D is a diagonal matrix with vertex degrees on the diagonals and A is adjacency matrix.

The number of spanning trees is any cofactor of L. i-th cofactor is determinant of the matrix gotten by removing i-th row and column of L.

Multigraphs

In D[i][i] all loops are excluded. A[i][j] = number of edges from i to j.

Directed multigraphs

D[i][i] = indegree of i minus the number of loops at i. A[i][j] = number of edges from i to j.

The number of oriented spanning trees rooted at a vertex i is the determinant of the matrix gotten by removing the ith row and column of L.

Matroid

Let *X* set of objects, $I \subseteq 2^X$ set of independents sets such that:

- 1. $\emptyset \in I$
- $2. \ A \in I, B \subseteq A \implies B \in I$
- 3. Exchange axiom, $A \in I, B \in I, |B| > |A| \implies \exists x \in B \setminus A : A \cup \{x\} \in I$
- 4. $A \subseteq X$ and I and I' are maximal independent subsets of A then |I| = |I'|

Then (X, I) is a matroid. The combinatorial optimization problem associated with it is: Given a weight $w(e) \ge 0 \ \forall e \in X$, find an independet subset that has the largest possible total weight.

Matroid intersection

```
// Input two matroids (X, I_a) and (X, I_b)
// output set I of maximum size, I \in I_a and I \in I_b
set<> I;
while(1){
    for(e_i : X \ I)
        if(I + e_i \in I_a and I + e_i \in I_b)
        I = I + e_i;
    set<> A, T; queue<> Q;
    for(x : X) label[x] = MARK1;
    for(e_i : X \ I){
        if(I + e_i \in I_a)
```

```
Q.push(e_i), label[e_i] = MARK2;
       else{
           for (x \text{ such that } I - x + e_i \setminus in I_a)
              A[x].push(e_i);
       if(I + e_i \setminus in I_b)
           T = T + \{e_i\}
       else{
           for(x such that I - x + e_i \in I_b)
               A[e_i].push(x);
   }
   if(T.empty()) break;
   bool found = false;
   while(!Q.empty() and !found){
       auto e = Q.front(); Q.pop();
       for(x : A[e]) if(label[x] == MARK1){
           label[x] = e;
           Q.push(x);
           if(x \in T)
               found = true;
              put = 1;
               while(label[x] != MARK2){
                  if(put) I = I + x;
                  else I = I - x;
                  put = 1 - put;
              I = I + x;
              break;
           }
   if(!found) break;
return I;
```

Where path(e) = [e] if label[e] = MARK2, path(label[e]) +[e] otherwise.

Matroid Union

}

Given k matroids over the same set of objects (X, I_1) , (X, I_2) , ..., (X, I_k) find $A_1 \in I_1, A_2 \in I_2, ..., A_k \in I_k$ such that $i \neq j, A_i \cap A_j = \emptyset$ and $|\bigcup A_i|$ is maximum. Matroid union can be reduced to matroid intersection as follows.

Let $X' = X \times \{1, 2, ..., k\}$, ie, k copies of each element of Xwith different colors. M1 = (X', Q) where $B \in Q \iff \forall 1 \leq$ $i \le k$, $\{x \mid (x, i) \in B\} \in I_i$, ie, for each color, B is independent. M2 = (X', W) where $B \in W \iff i \neq j \implies \neg((x, i) \in X)$ $B \land (x, j) \in B$), ie, each element is picked by at most one color.

Intersection of M1 and M2 is the answer for the combinatorial problem of matroid union.

Edge coloring

```
Data: A graph G
Result: A proper coloring c of the edges of G
Let U := E(G);
while U \neq \emptyset do
   Let (u,v) be any edge in U;
   Let F[1:k] be a maximal fan of u starting at F[1]=v;
   Let c be a color that is free on u and d be a color
    that is free on F[k];
   Invert the cd_u path;
   Let w \in V(G) be such that w \in F, F' = [F[1]...w] is a
    fan and d is free on w;
   Rotate F' and set c(u,w)=d;
   U := U - (u,v);
end
```

Notes

When we repeat something and each time we have probability *p* to succeed then the expected number or tries is $\frac{1}{n}$, till we succeed.

Small to large

Trick in statement If *k* sets are given you should note that the amount of different set sizes is $O(\sqrt{s})$ where s is total size of those sets. And no more than \sqrt{s} sets have size greater than \sqrt{s} . For example, a path to the root in Aho-Corasick through suffix links will have at most $O(\sqrt{s})$ vertices.

gcd on subsegment, we have at most $log(a_i)$ different values in {gcd(a_i , a_{i+1} , ..., a_i) for j < i}.

From static set to expandable. To insert, create a new set with the new element. While there are two sets with same size, merge them. There will be at most log(n) disjoints sets.

Matrix exponentiation optimization. Save binary power of A_{nxn} and answer q queries $b = A^m x$ in $O((n^3 + qn^2)log(m))$.

Ternary search on integers into binary search, comparing f(mid) and f(mid+1), binary search on derivative

Dynamic offline set For each element we will wind segment of time [a, b] such that element is present in the set during this whole segment. Now we can come up with recursive procedure which handles [l, r] time segment considering that all elements such that $[l,r] \subset [a,b]$ are already included into the set. Now, keeping this invariant we recursively go into [l, m] and [m + 1, r] subsegments. Finally when we come into segment of length 1.

```
a > b \implies a \mod b < \frac{a}{2}
```