

haskell系列教程 II

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- 函数式编程基本套路
- 函子抽象
- 透镜组

函数组合

```
(.) :: (b -> c) -> (a -> b) -> a -> c
f . g x = f (g x)
infixr 9 .
```

```
($) :: (a -> b) -> a -> b
f $ x = f x
infixr 0 $
```

-- 方便利用优先级省略括号，下面写法都可以

```
unlines (map reverse (lines ( "hello\nworld" )))
(unlines . map reverse . lines) "hello\nworld"
unlines $ map reverse $ lines $ "hello\nworld"
unlines . map reverse . lines $ "hello\nworld"
-- "olleh\ndlrow\n"
```

```
reverseByLine :: String -> String
reverseByLine = unlines . map reverse . lines
```

常见函数

```
id :: a -> a      -- 直接返回参数
id x = x
```

```
const :: a -> b -> a      -- 忽略第二个参数
const x y = x
```

```
flip :: (a -> b -> c) -> b -> a -> c  -- 交换参数的顺序
flip f x y = f y x
```


```
flip map [1..10] $ \ x -> ...
-- (`map` [1..10]) $ \ x -> ...
```

```
(&) :: a -> (a -> b) -> b      -- flip ($)
x & f = f x
infixl 1 &
```

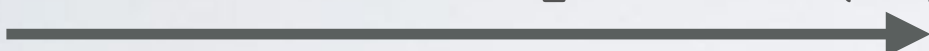
```
"hello\nworld" & lines & map reverse & unlines
-- "olleh\ndlrow\n"
```

递归

```
map :: (a -> b) -> [a] -> [b]
map f [] = []
map f (x:xs) = f x : map f xs
```

map f  map f xs

```
xs :: [Int]
xs = 1 : 1 : zipWith (+) xs (tail xs)
```

xs  xs tail xs

```
data BinTree a = Nil | Node a (BinTree a) (BinTree a)
```

```
countNode :: BinTree a -> Int
countNode Nil = 0
countNode (Node _ left right) =
    countNode left + 1 + countNode right
```

```
countNode
    (Node 1 (Node 2 Nil Nil) (Node 3 Nil (Node 4 Nil Nil)))
-- 4
```

foldr

```
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr _ acc [] = acc
foldr f acc (x:xs) = x `f` foldr f acc xs
```

```
foldr (+) 0 [1,2,3]
-- 1 + (foldr (+) 0 [2,3])
-- 1 + (2 + foldr (+) 0 [3])
-- 1 + (2 + (3 + foldr (+) 0 []))
-- 1 + ( 2 + (3 + 0))
```

```
-----
max :: (Ord a) => a -> a -> a
max x y | x >= y      = x
        | otherwise = y
```

```
maximum :: (Ord a) => [a] -> a
maximum [] = error "empty list"
maximum (x:xs) = foldr max x xs
```

foldr

```
length :: [a] -> Int
```

```
length [] = 0
```

```
length (_:xs) = 1 + count xs
```

```
length = foldr (const (+1)) 0
```

```
const :: a -> b -> a
```

```
const :: (Int -> Int) -> b -> Int -> Int
```

```
const (+1) == \ x y -> y + 1
```

```
-----  
all :: (a -> Bool) -> [a] -> Bool
```

```
all f = foldr ((&&) . f) True
```

```
all f xs = foldr (\ x acc -> f x && acc) True xs
```

```
all even [1..]
```

```
-- even 1 && (even 2 && (even 3 && ..
```

```
-- False
```

foldr

```
map f = foldr ((:) . f) []
```

```
map id [1,2,3]
```

```
-- foldr ((:) . id) []
```

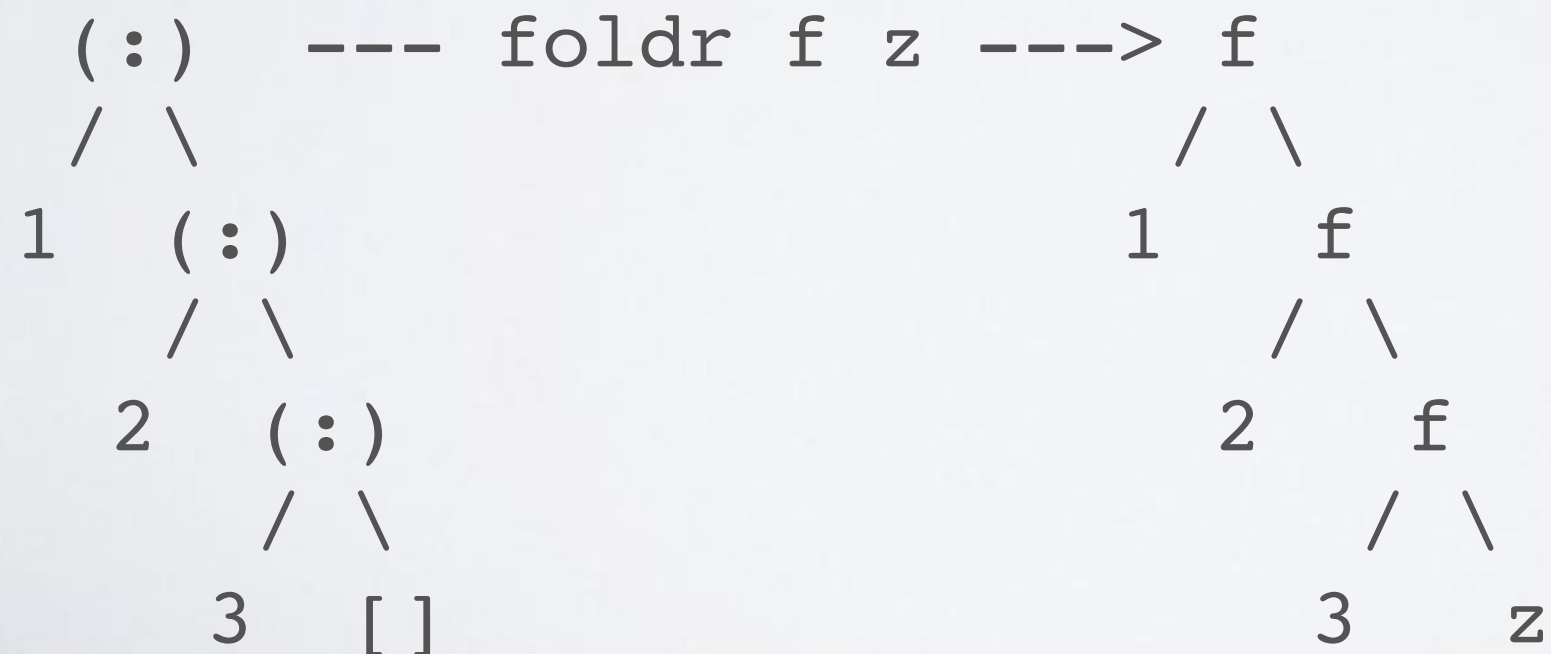
```
-- foldr (:) [] [1,2,3]
```

```
-- 1 : (foldr (:) [] [2,3])
```

```
-- 1 : (2 : foldr (:) [] [3])
```

```
-- 1 : (2 : (3 : foldr (:) [] []))
```

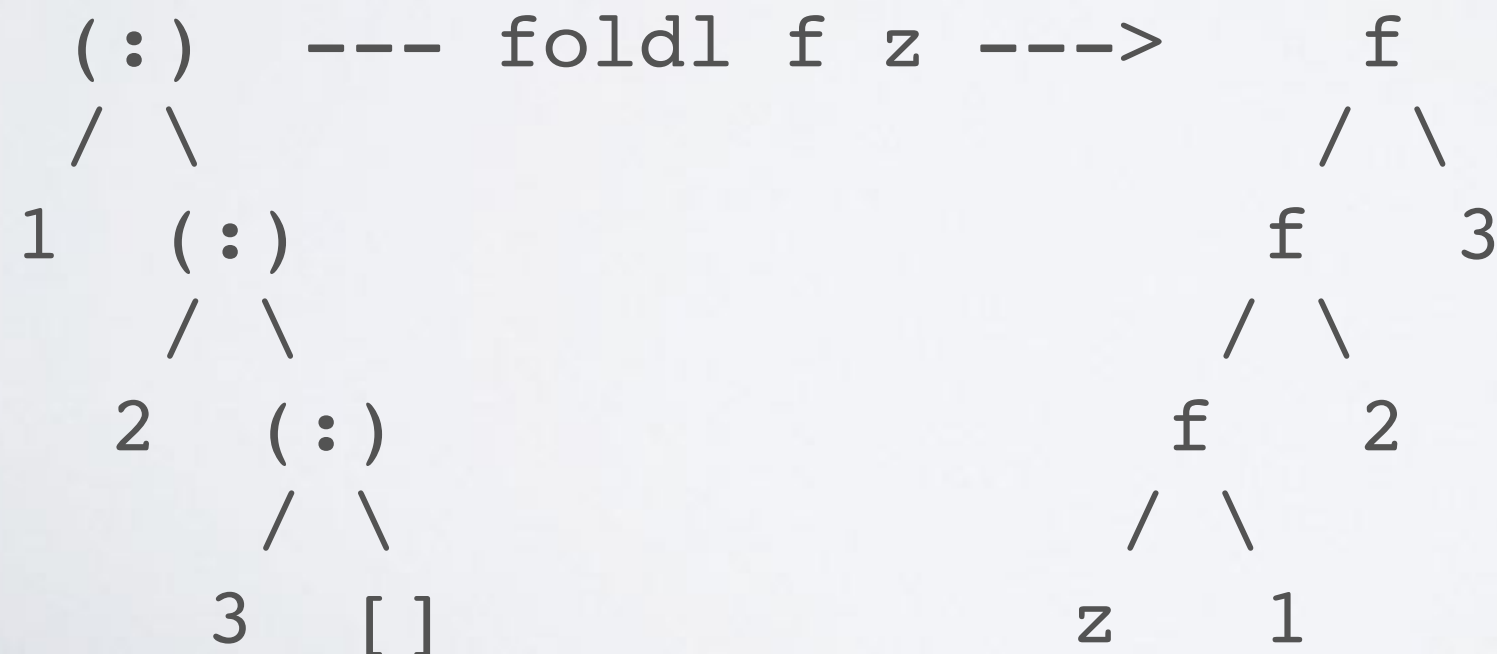
```
-- 1 : ( 2 : (3 : []))
```



foldl

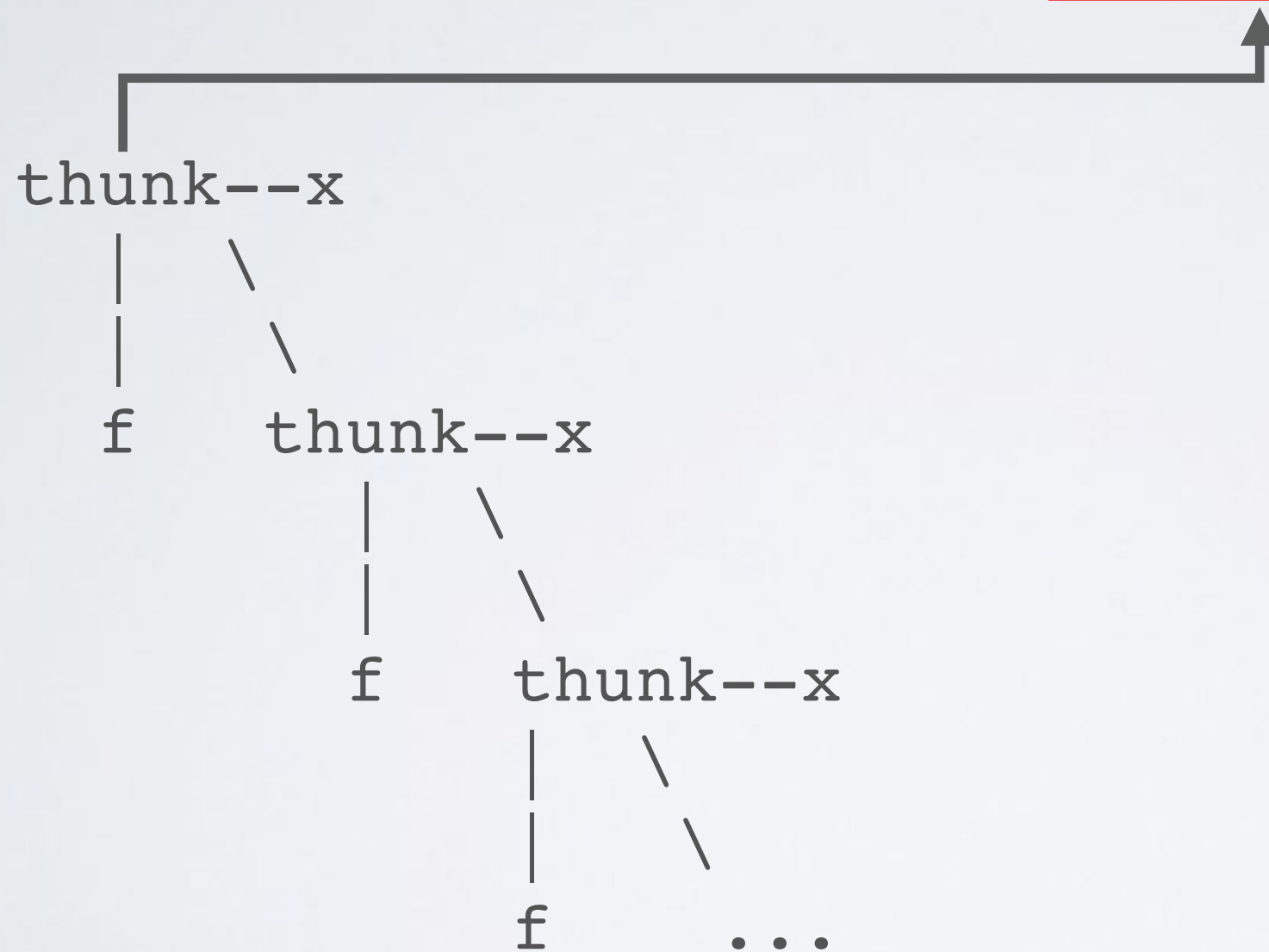
```
foldl :: (b -> a -> b) -> b -> [a] -> b
foldl f acc [] = acc
foldl f acc (x:xs) = foldl f (acc `f` x) xs
```

```
foldl (+) 0 [1,2,3]
-- foldl (+) (0 + 1) [2,3]
-- foldl (+) (0 + 1 + 2) [3]
-- foldl (+) (0 + 1 + 2 + 3) []
-- 0 + 1 + 2 + 3
```



thunk leak

```
foldl :: (b -> a -> b) -> b -> [a] -> b
foldl f acc [] = acc
foldl f acc (x:xs) = foldl f (acc `f` x) xs
```



`seq :: a -> b -> b` -- 把对b的求值过程转化为对a和b的求值过程

foldl'

-- 每次迭代的同时对累计值求值

```
foldl' :: (b -> a -> b) -> b -> [a] -> b
```

```
foldl' f acc [] = acc
```

```
foldl' f acc (x:xs) =
```

```
    let acc' = acc `f` x
```

```
    in acc' `seq` foldl f acc' xs
```

```
foldl' (+) 0 [1,2,3]
```

```
-- foldl' (+) 1 [2,3]
```

```
-- foldl' (+) 3 [3]
```

```
-- foldl (+) 6 []
```

```
-- 6
```

```
-----  
all' f = foldl' (\ acc x -> acc && f x) True
```

```
all' even [1..]
```

```
-- (...((True && even 1) && even 2) && even 3 ...
```

```
-- Never stop...
```

类型类语法

```
class Eq a where
```

```
    (==), (/=) :: a -> a -> Bool
```

```
    x /= y      = not (x == y)    -- 默认实现
```

```
    x == y      = not (x /= y)
```

```
instance Eq Double where
```

```
    -- (==) :: Double -> Double -> Bool
```

```
    a == b = eqDouble# a b
```

```
instance Eq a => Eq (Maybe a) where
```

```
    (Just a) == (Just b) = a == b
```

```
    Nothing == Nothing   = True
```

```
    _       == _         = False
```

类型类语法

```
data Ordering = LT | EQ | GT -- 小于|等于|大于
```

```
class Eq a => Ord a where
```

```
    compare                :: a -> a -> Ordering
```

```
    (<), (<=), (>), (>=) :: a -> a -> Bool
```

```
    max, min                :: a -> a -> a
```

```
    compare x y = if x == y then EQ
```

```
                  else if x <= y then LT
```

```
                  else GT
```

```
x < y = case compare x y of { LT -> True; _ -> False }
```

```
x <= y = case compare x y of { GT -> False; _ -> True }
```

```
x > y = case compare x y of { GT -> True; _ -> False }
```

```
x >= y = case compare x y of { LT -> False; _ -> True }
```

```
max x y = if x <= y then y else x
```

```
min x y = if x <= y then x else y
```

```
class Num a where
  (+), (-), (*) :: a -> a -> a
  negate, abs, signum :: a -> a
  fromInteger :: Integer -> a

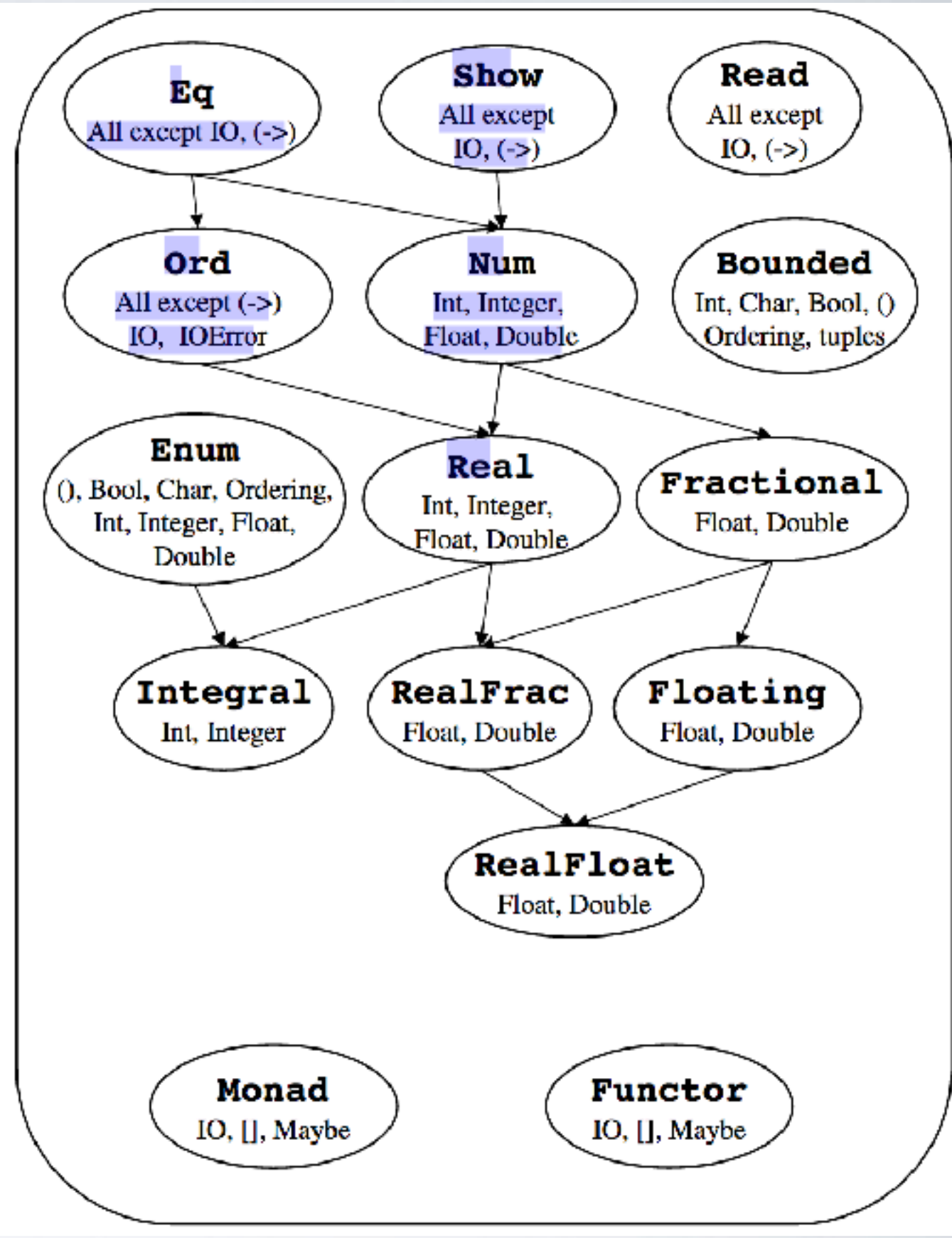
class (Num a, Ord a) => Real a where
  toRational :: a -> Rational

class (Real a, Enum a) => Integral a
where
  quot, rem, div, mod :: a -> a -> a
  divMod, quotRem :: a -> a -> (a, a)
  toInteger :: a -> Integer

class Num a => Fractional a where
  (/) :: a -> a -> a
  recip :: a -> a
  fromRational :: Rational -> a

class Fractional a => Floating a where
  pi :: a
  exp, log, sqrt :: a -> a
  (**), logBase :: a -> a -> a
  sin, cos, tan :: a -> a
  asin, acos, atan :: a -> a
  sinh, cosh, tanh :: a -> a
  asinh, acosh, atanh :: a -> a
```

...



Monoid

```
class Monoid m where
  mempty :: m
  mappend :: m -> m -> m
  mconcat :: [a] -> a
  mconcat = foldr mappend mempty
```

```
(<>) = mappend
infixr 6 <>
```

-- Monoid的满足的定律

```
mempty <> x == x
x <> mempty == x
(x <> y) <> z == x <> (y <> z)
```

Monoid

```
instance Monoid [a] where
    mempty  = []
    mappend = (++)
```

```
instance Monoid Int where
    mempty  = 0
    mappend = (+)
```

```
"hello" <> "world" -- "helloworld"
0 <> 1 <> 2 -- 3
```

-- 如何选择Int的两种Monoid实例?

```
instance Monoid Int where
    mempty  = 1
    mappend = (*)
0 <> 1 <> 2 -- 0
```

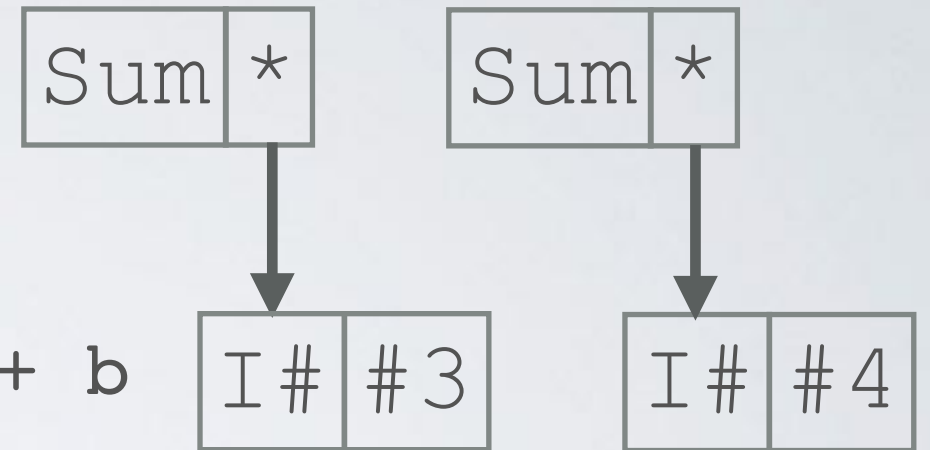
构造新的类型来选择实例

```
data Sum a = Sum {getSum :: a}
```

```
instance Num a => Monoid (Sum a)
```

```
    mempty = Sum 0
```

```
    Sum a `mappend` Sum b = Sum $ a + b
```



```
Sum 3 <> Sum 4 == Sum 7
```

```
getSum . mconcat . map Sum $ [1..100] == 5050
```

```
-----  
data Product a = Product {getProduct :: a}
```

```
instance Num a => Monoid (Product a)
```

```
    mempty = Product 1
```

```
    Product a `mappend` Product b = Product $ a * b
```

```
Product 3 <> Product 4 == Product 12
```


newtype

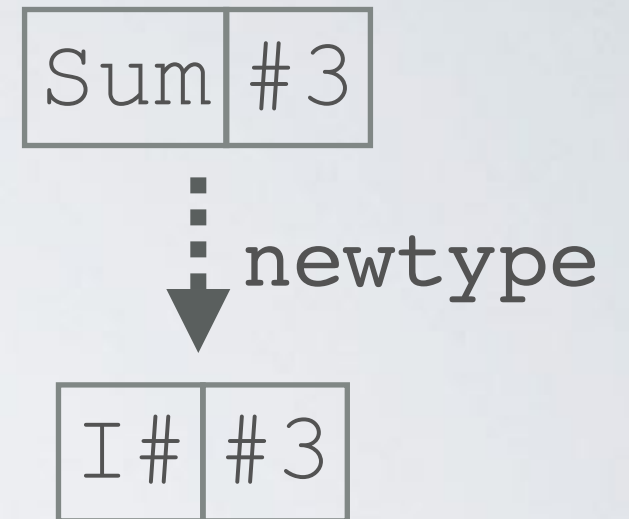
```
newtype Sum a = Sum {getSum :: a}
```

```
instance Num a => Monoid (Sum a)
```

```
    mempty = 0
```

```
    Sum a `mappend` Sum b = Sum $ a + b
```

```
Sum 3 <> Sum 4 == Sum 7
```



`Sum`只存在于编译阶段!

`newtype`的行为和`data`类似，但是省去了一层`boxing`。

因此速度快，而且没有额外的`Laziness` -- `Sum x == x `seq` Sum x`

`newtype`的使用条件：

- 新类型只有一个构造函数
- 构造函数只有一个参数

函子Functor

-- 函子提供了操作容器内部payload的能力

```
class Functor f where
```

```
    fmap :: (a -> b) -> f a -> f b
```

```
instance Functor [] where
```

```
    -- fmap :: (a -> b) -> [a] -> [b]
```

```
    fmap = map
```

```
instance Functor Maybe where
```

```
    -- fmap :: (a -> b) -> Maybe a -> Maybe b
```

```
    fmap f (Just a) = Just (f a)
```

```
    fmap f Nothing = Nothing
```

```
fmap (+1) (Just 0) == Just 1
```

```
fmap (+1) Nothing == Nothing
```

```
fmap even [1,2,3] == [False, True, False]
```

两个极端的函子

-- 什么上下文都没有的Identity

```
newtype Identity a = Identity{runIdentity :: a}
```

```
instance Functor Identity where
```

```
    fmap f (Identity a) = Identity (f a)
```

```
fmap (+1) (Identity 0) == Identity 1
```

-- 只包含上下文，没有payload的Const

```
newtype Const a b = Const{getConst :: a}
```

```
instance Functor Const c where
```

```
    -- fmap :: (a -> b) -> Const c a -> Const c b
```

```
    fmap f (Const c) = Const c
```

```
fmap (+1) (Const 0) == Const 0
```

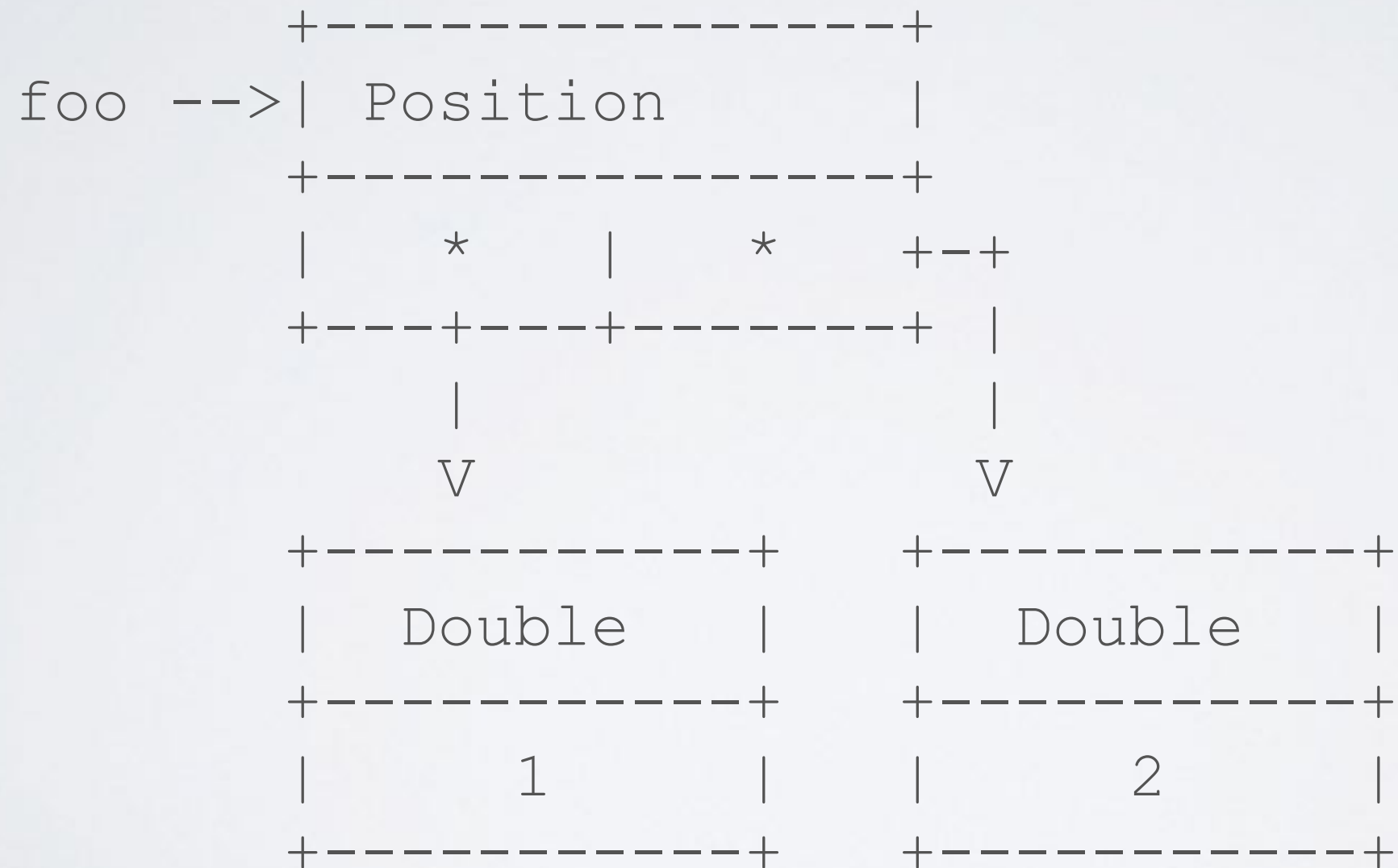


Lens

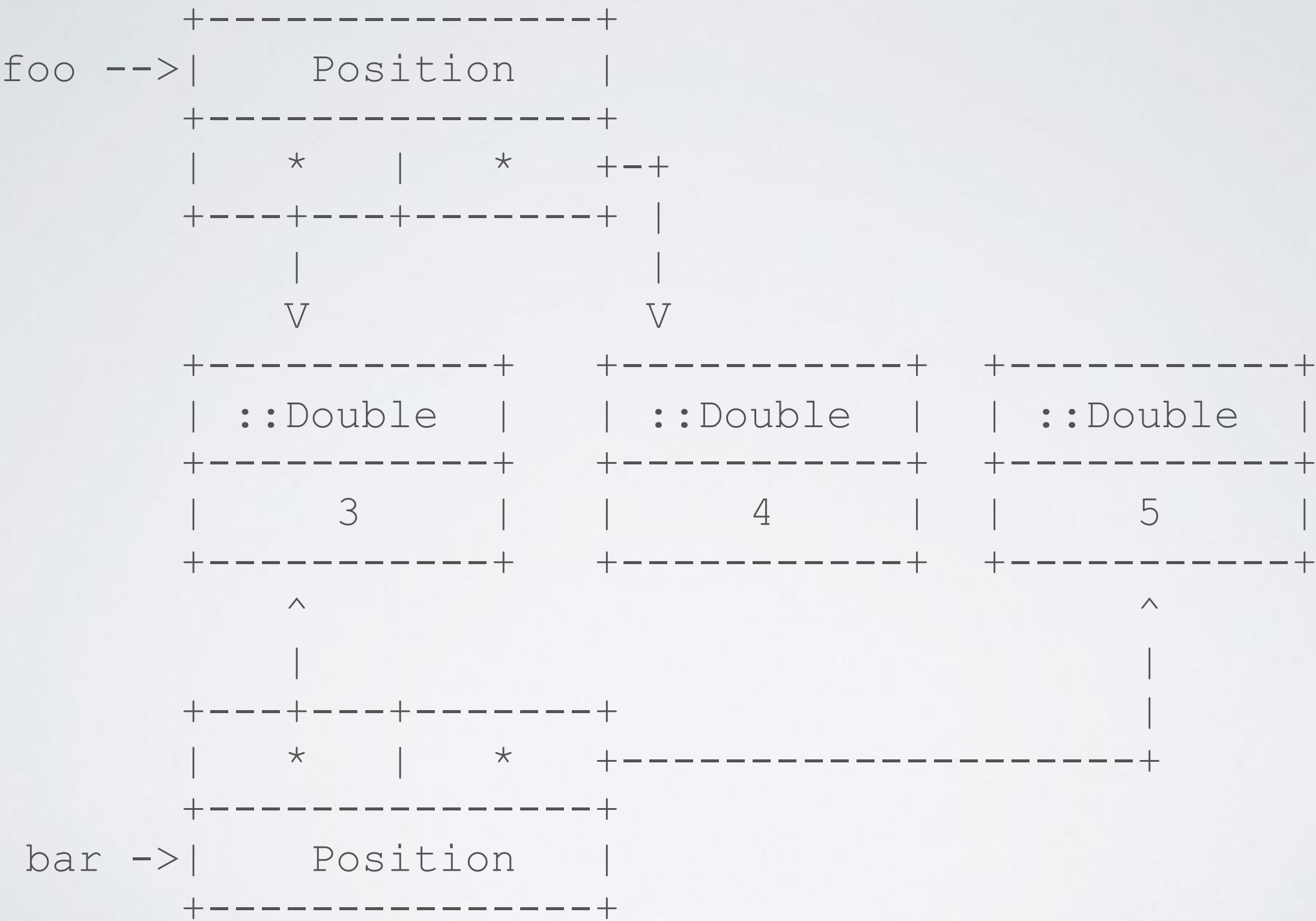
数据操作的艺术

```
data Position =
    Position{posX :: Double, posY :: Double}
```

```
foo :: Position
foo = Position 1 2
```



```
bar = foo{posY = 5}
```



问题?

```
data Line = Line{start :: Position, end :: Position}
```

```
lineA = Line (Position 0 0) (Position 3 4)
```

```
-- move end's posY to 5
```

```
lineB =  
    let lineA'sEnd = end lineA  
    in lineA{  
        end = lineA'sEnd{ posY = 5 }  
    }
```

- 深层次嵌套的不可变数据操作如此麻烦?
- 我们希望能够使用类似`lineA.end.posY = 5`的语法
- 同时保持数据不可变和共享的特性

meet lens

```
type Lens b a =
    (Functor f) => (a -> f a) -> b -> f b

posXLens :: Lens Position Double
posXLens :: (Functor f) =>
    (Double -> f Double) -> Position -> f Position

posXLens f p =
    let x = (posX p)    -- x :: Double
        x' = f x        -- x' :: f Double
    in
        -- setter :: Double -> Position
        setter = \x -> p{posX = x}
        fmap setter x'    -- f Position
```


lens与functor

```
posXLens :: Lens Position Double
posXLens :: (Functor f) =>
  (Double -> f Double) -> Position -> f Position
```

```
mirror x :: Double -> [Double]
mirror x = [-x, x]
```

```
foo = Position 3 4
```

```
posXLens mirror foo
== [Position -3 4, Position 3 4]
/= Position [-3,3] 4
```

使用lens获得modifier

```
over :: Lens b a -> (a -> a) -> b -> b
over lens f b =
  let
    idF = Identity . f      -- idF :: a -> Identity a
    idB = lens idF          -- idB :: b -> Identity b
    fb = idB b              -- fb :: Identity b
  in
    runIdentity fb          -- b

over lens f = runIdentity . lens (Identity . f)

foo = Position 3 4
over posXLens (+1) foo == Position 4 4
```

使用lens获得setter

```
set :: Lens b a -> a -> b -> b
```

```
set lens a b =
```

```
  let
```

```
    idF = Identity . (\_ -> a)
```

```
    idB = lens idF
```

```
    fb = idB b
```

```
  in runIdentity fb
```

```
const :: a -> b -> a
```

```
const a _ = a
```

```
set lens a
```

```
  = runIdentity . lens (Identity . const a)
```

```
set posXLens 0 foo == Position 0 4
```

使用lens获得getter

```
view :: Lens b a -> b -> a
```

```
Lens b a :: Functor f => (a -> f a) -> b -> f b
```

从f b类型得到a类型? ? view lens b = ???

```
view lens b =
```

```
  let
```

```
    -- Const :: a -> Const a a
```

```
    cb = lens Const          -- cb :: b -> Const a b
```

```
    ca = cb b               -- ca :: Const a b
```

```
  in getConst ca           -- a
```

```
view lens = getConst . (lens Const)
```

```
view posXLens foo == 3
```

(.)

The power of dot

组合lens

```
posXLens :: Lens Position Double
posXLens :: Functor f =>
  (Double -> f Double) -> Position -> f Position
posXLens f p =
  fmap (\x -> p{posX = x}) (f $posX p)

startLens :: Lens Line Position
startLens :: Functor f =>
  (Position -> f Position) -> Line -> f Line
startLens f l =
  fmap (\s -> l{start = s}) (f $ start l)

(startLens . posXLens) :: Functor f =>
  (Double -> f Double) -> Line -> f Line
(startLens . posXLens) :: Lens Line Double
```

组合lens

```
lineA = Line (Position 0 0) (Position 3 4)

-- move start's posX to 5
set (startLens . posXLens) lineA 5
-- Line (Position 5 0) (Position 3 4)
```

实际上你并不需要书写每一个Lens，真正的使用大约类似这样：

```
data Profile = Profile {
  _name :: String
, _address :: Address
, ...
}
```

```
mkLenses ''Profile
-- now you have name/address lenses
```

中缀函数DSL

```
(%~) = over
-- lens (%~) f = over lens f
infixr 4 %~
lineA & start. posX %~ (+1)

(.~) = set
infixr 4 .~
lineA & start. posX .~ 0

value ^. lens = view lens value
-- (^.) = flip view

infixl 8 ^.
lineA ^. start . posX
```


being first class!

```
triAngleA :: [Line]
triAngleA = [ Line (Position 0 0) (Position 0 4)
              , Line (Position 0 4) (Position 3 0)
              , Line (Position 3 0) (Position 0 0)
              ]

map (^ . start) triAngleA
-- [(Position 0 0),(Position 0 4),(Position 3 0)]

foldl' (\ acc x -> acc + x ^ . start . posX) 0 triAngleA
-- 3

map (start . posX %~ (+5)) triAngleA
-- [ Line (Position 5 0) (Position 0 4)
    , Line (Position 5 4) (Position 3 0)
    , Line (Position 8 0) (Position 0 0)
    ]
```