

A Type-Theoretical system for the FraCaS test suite: Grammatical Framework meets Coq

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September 20, 2017

Brief summary of the talk

- Present a type-theoretical framework for formal semantics leveraging two well-studied tools
 - ▶ Grammatical Framework (GF, Ranta 2004, 2011)
 - ▶ Coq
- Providing a compositional resource semantics for GF
- Evaluation on the FraCaS test suite
- State-of-the-art results

Structure the talk

- Brief intro to the systems used, GF and Coq
- Presenting the FraCoq system
 - ▶ We concentrate on the most linguistically relevant features and also the features relevant for the FraCaS
- Evaluation against the FraCaS test suite
 - ▶ Some brief remarks about the FraCaS and NLI platforms
 - ★ Results and comparison with previous logical approaches
 - ★ The issue of automation
- Conclusions and Future work

Background: Grammatical Framework (GF)

- Functional programming language for writing multi-lingual grammars
- Involves an abstract syntax, comprised of:
 - ▶ A number of syntactic categories
 - ▶ A number of syntactic construction functions, which provide the means to compose basic syntactic categories into more complex ones
 - ★ $AdjCN:AP \rightarrow CN \rightarrow CN$ (appending an adjectival phrase to a common noun and obtaining a new common noun)
- GF comes with a library of mappings from abstract syntax to concrete
 - ▶ These mappings can be inverted by GF, thus offering parsers from natural text into abstract syntax
 - ▶ We use the parse trees constructed by Ljunglöf (2012) thereby avoiding any syntactic ambiguity (GF FraCaS treebank).

Background: Type-Theoretical Semantics

- We use the type of logics that have been traditionally dubbed as constructive
 - ▶ Initiated by the work of Martin-Löf (1971, 1984)
 - ▶ In linguistics this types of logics go back to Ranta's seminal work (Ranta 1994) or even earlier to Sundholm (1988)
 - ★ More recent approaches can be found as well. Please see Chatzikyriakidis and Luo (2017) for a collection of papers on constructive type theories for NL semantics

Background: Type Theoretical Semantics

- Main features of MTTs

- ▶ Type many-sortedness.
- ▶ Dependent sum and product types
 - ★ Σ -types, often written $\sum_{x:A} B[x]$ and which have product types $A \times B$ as a special case when B does not depend on x .
 - ★ Dependent product, Π -types, often written $(\prod_{x:A} B[x])$, and which have arrow-types $A \rightarrow B$ as a special case
 - ★ They generalize universal quantification and function types and they offer type polymorphism
- ▶ Proof-theoretical specification and support for effective reasoning.
 - ★ Most powerful proof assistants implement MTTs (e.g. Coq, Agda)

Background: Coq

- Proof assistant based on the calculus of inductive constructions (CiC, Coquand and Paulin-Mohring 1988)
 - ▶ Arguably one of the leading proof assistants
 - ★ a proof of the four-color theorem (Gonthier 2008)
 - ★ a proof of the odd order theorem (Gonthier et al. 2013)
 - ★ developing CompCert, a formally verified compiler for C (Leroy 2013)
 - ★ One of the assistants used in the Univalent Foundations project (Homotopy Type Theory, The Univalent Foundations Program 2013)

Background: Coq

- Important features used

- ▶ Π types

- ★ in Coq: $\prod_{x:A} B[x]$ is written `forall (x:A), B` or (simply $A \rightarrow B$ when B does not depend on x)

- ▶ Record types

- ★ Generalization of Σ -types and are encoded as (trivial) inductive types with a single constructor.
 - ★ $\Sigma x:A. B(x)$ can be expressed as a dependent record type in Coq:

`Record AB : Type := mkAB {x :> A ; P : B x}`

The FraCoq system

- We use Ljunglöf's FraCaS treebank and take these trees to their semantic counterparts
- The structure of the semantic representation
 - 1 Every GF syntactic category C is mapped to a Coq *Set*, noted $\llbracket C \rrbracket$.
 - 2 GF Functional types are mapped compositionally : $\llbracket A \rightarrow B \rrbracket = \llbracket A \rrbracket \rightarrow \llbracket B \rrbracket$
 - 3 Every GF syntactic construction function $f:X$ is mapped to a function $\llbracket f \rrbracket$ such that $\llbracket f \rrbracket : \llbracket X \rrbracket$.
 - 4 GF function applications are mapped compositionally:
 $\llbracket t(u) \rrbracket = \llbracket t \rrbracket(\llbracket u \rrbracket)$.

The FraCoq system

- Sentences

- ▶ We interpret sentences as propositions: $\llbracket S \rrbracket = Prop$.
- ▶ To verify that P entails H , we prove the proposition $\llbracket P \rrbracket \rightarrow \llbracket H \rrbracket$.

Definition S := Prop.

- Common Nouns

- ▶ Predicates over an abstract object type

Parameter object : Set.

Definition CN := object->Prop.

The FraCoq system

- Verb phrases

- ▶ Parameterize over the *noun* of the subject (using Π types)

Definition VP := forall (subjectClass : CN)
object -> Prop.

The FraCoq system

- Adjectives

- ▶ Functions from cn to cn (predicates to predicates)

Definition A := CN -> CN.

- ▶ Different classes of adjectives are captured using coercions (subtyping).
All special classes of adjectives are subtypes of A.

Definition IntersectiveA := object -> Prop.

Definition wkIntersectiveA : IntersectiveA -> A
:= fun a cn (x:object) => a x /\ cn x.

Coercion wkIntersectiveA : IntersectiveA >-> A.

- ▶ Provision is made for intersective, subjective, privative and non-committal adjectives

The FraCoq system

- Adverbs

- ▶ Similar method to adjectives but instead the modification is on verbal predicates
- ▶ The adverb cases in the FraCaS are all veridical and covariant.
- ▶ We define such a subclass *VeridicalAdv* and declare it as a coercion *Adv*
 - ★ Adverbs of type *VeridicalAdv* are also of type *Adv*

Parameter on_time_Adv : VeridicalAdv .

The FraCoq system

- Noun Phrases and Predeterminers

- ▶ A clean definition of NPs as functions from predicates to truth values will not work
 - ★ Problem with GF's abstract syntax: existence of pre-determiners which include cases like *most*, *all* among others and are defined as functions from NPs to NPs
 - ★ In general, the category includes elements that are naturally interpreted as GQs
 - ★ Solution: Remember the components of NPs (number, quantifier and common noun)
 - ★ Pre-determiners then are able to substitute the quantifier part with the appropriate quantifier
 - ★ This has to be done, otherwise pre-determiners introduce a dummy indefinite in these cases

The FraCoq system

- Generalized Quantifiers

- ▶ They turn a number and a common noun into a noun-phrase (which we call *NPO*).

Definition $\text{Quant} := \text{Num} \rightarrow \text{CN} \rightarrow \text{NPO}$.

- ▶ Some quantifiers ignore the number and are given usual definitions (e.g. *some* or *all*), whereas others make essential use of number (e.g. *at most*)
 - ★ In the latter case, the function *CARD*, a context-dependent abstract function which turns a predicate into a natural number is used to get the correct semantics

The FraCoq system

- The definite article
 - ▶ Checks for plural noun phrases
 - ★ If found, then universal quantification
 - ★ If not, it looks up the object of discourse in an abstract *environment*, which is a function which turns a common noun into an object

```
Definition DefArt:Quant:= fun (num : Num) (P:CN)=> fun Q:VP
=> match num with plural => (forall x, P x -> Q P x)
  /\ Q P (environment P) /\ P (environment P) |
_ => Q P (environment P) /\ P (environment P) end.
```


Evaluation: The FraCaS test suite

- A test suite for NLI

- ▶ 346 NLI examples in the form of one or more premises followed by a question along with an answer to that question
- ▶ Three potential answers
 - ★ YES: The declarative sentence formed out of the question follows from the premises
 - ★ NO: The declarative sentence does not follow from the premises
 - ★ UNK: The declarative sentence neither follows nor does not follow from the premises

Evaluation: The FraCaS test suite

- (1) A Swede won the Nobel Prize.
Every Swede is Scandinavian.
Did a Scandinavian win the Nobel prize? [Yes, FraCas 049]
- (2) No delegate finished the report on time..
Did any Scandinavian delegate finish the report on time? [No, FraCaS 070]
- (3) A Scandinavian won the Nobel Prize.
Every Swede is Scandinavian.
Did a Swede win the Nobel prize? [UNK, FraCaS 065]

Evaluation: The FraCaS test suite

- The FraCaS has considerable weaknesses
 - ▶ Small size
 - ▶ Artificial nature of the examples
- However, it covers a lot of phenomena associated with NLI
- It is still a very good suite to test logical approaches

Evaluation

- Evaluation against 5 sections of the FraCaS
 - ▶ Total of 174 examples
 - ▶ Excluded sections where a lot of context-dependency has to be taken into consideration (e.g. the section on ellipsis)
- YES: a proof can be constructed from the premises to the hypothesis
- NO: a proof of the negated hypothesis can be constructed
- UNK: otherwise

Evaluation

- The following table presents the results (Ours) as well as a comparison with the approach in Mineshima et al. (MINE, 2015), Bos (Nut, 2008) and Abzianidze (Langpro, 2015)

	Section	# examples	Ours	MINE	Nut	Langpro
1	Quantifiers	75	.96	.77	.53	.93 (44)
2	Plurals	33	.76	.67	.52	.73 (24)
3	Adjectives	22	.95	.68	.32	.73 (12)
4	Comparatives	31	.56	.48	.45	-
5	Attitudes	13	.85	.77	.46	.92 (9)
6	Total	174 (181)	0.83	0.69	0.50	0.85

- Our approach outperforms Mineshema et al. by 13 percentage points.
- The approach by Abzianidze has an accuracy of 0.85 without involving the comparative section. If this section is taken out, our system's accuracy rises to 0.88

Error Analysis

- Improvement over earlier approaches. Still, there were a couple of difficulties
 - ▶ Comparatives: cases that needed one to provide adequate semantics for *more* but also to take care of ellipsis
 - (4) ITEL won more orders than APCOM.
ITEL won some orders.
Did ITEL win some orders? [Yes, FraCaS 233]
 - ▶ Definite Plurals: Universal reading was captured. Cases of existential readings were not
 - (5) The inhabitants of Cambridge voted for a Labour MP.
very inhabitant of Cambridge voted for a Labour MP.
Did every inhabitant of Cambridge vote for a Labour MP?
[UNK, FraCaS 094]

Automation

- So far, our proofs are not automated
 - ▶ A couple of steps (usually very few) to reach a proof
 - ▶ Earlier approaches using Coq (e.g. Chatzikyriakidis 2014 and Mineshima et al. 2015) use Coq's tactical language LTac to define automated macros of actions
 - ★ This is not difficult to do in our case as well
 - ★ Just go through all the proof tactics or observe the tactics that are used in the proofs to create a macro that will automate the proofs
 - ★ The question remains: can that macro of tactics generalize outside the suite?
 - ★ Answer: only to a limited extent, i.e. when exactly the same set of tactics yields a proof
 - ★ For this reason, we have not automated proof search to obtain the results presented in this paper, even though this can be done easily

Automation

- Automating would also make an unprincipled use of higher-order logic (HOL)
 - ▶ No algorithm which can decide if a proposition has a proof or not
 - ★ We must use heuristics both to search for proofs and to decide when to give up searching
- Most problems have either obvious proofs or obviously lack a proof (fortunately)
 - ▶ Due to its heuristic nature the proof search necessarily contains a human component
 - ★ Problematic to make a statement about the suitability of FraCoq outside FraCas
 - ★ Small dataset and lack of separation between a development and a test set does not help the situation either
 - ★ Related shortcoming: specialized semantics for specific lexical entries

Future Work

- Address the issue of automation
 - ▶ Define a decidable fragment of the logic and only work within such fragment
 - ★ Possible to concisely characterize how the approach generalises
 - ▶ Train a neural network on a body of freely available proofs on the net and see whether it can generalize to automatically provide the proof tactics for the cases interested
- Improvement at the GF level: make the abstract syntax more compatible with compositional semantics
 - ▶ For example, do something about problematic syntactic categories like pre-determiners or cases where the syntax makes it impossible to recover elliptical fragments
- Extend into the whole suite (first attempt to do anaphora using monads on the way!)

Conclusions

- We have connected two well-defined systems based on type-theory
 - ▶ GF and Coq
 - ▶ Providing resource semantics for GF
- The issue of generalization remains a shortcoming
 - ▶ It is possible to achieve very precise semantics for specific domains
 - ★ Our system outperforms previous logical systems w.r.t accuracy
- Useful in performing inference tasks on controlled natural language domains
- Hybrid NLI systems

Conclusions

- The system can be found here:
<https://github.com/GU-CLASP/FraCoq>