

R9 Graphs BFS

A graph $G = (V, E)$ is a mathematical object comprising a set of vertices V (also called nodes) and a set of edges E .

- *directed* (u, v)
- *undirected* $\{u, v\}$

The *in-degree* and *out-degree* of a vertex v denotes the number of incoming and outgoing edges connected to v respectively.

When talk about degree, generally mean out-degree.

- *path*

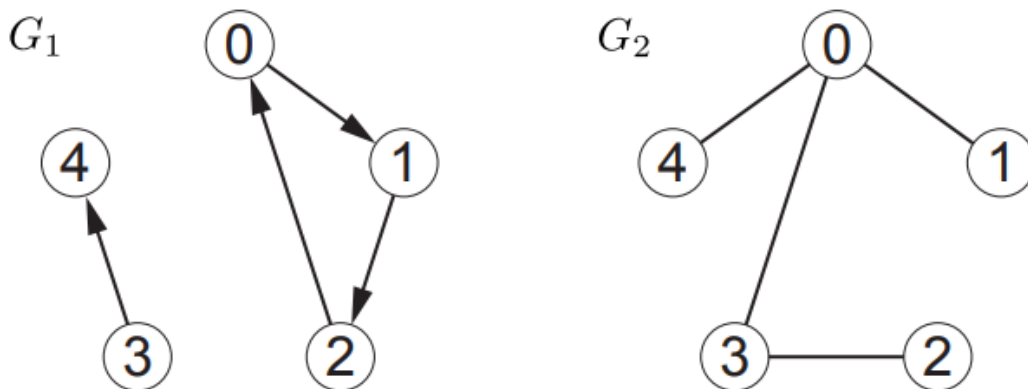
A path¹ in a graph is a sequence of vertices (v_0, \dots, v_k) such that for every ordered pair of vertices (v_i, v_{i+1}) , there exists an outgoing edge in the graph from v_i to v_{i+1} .

The length of a path is the number of edges in the path

Graph Representations

adjacency list

The most common way is to store a Set data structure Adj mapping each vertex u to another data structure Adj(u) storing the adjacencies of u , the set of vertices that are accessible from u via a single outgoing edge.



example of *adjacency list* in python of G_1 & G_2 . Using a direct access array for the top-level Set and an array for each adjacency list.

```
A1 = [[1], [2], [0], [4], []]
A2 = [[1, 3, 4], [0], [3], [0, 2], [0]]
```



example of hash table for outer Adj Set and inner adjacency lists Adj(u), using Python dictionaries

```
S1 = {0:{1}, 1:{2}, 2:{0}, 3:{4}},  
S2 = {0:{1,3,4}, 1:{0}, 2:{3}, 3:{0,2}, 4:{0}}
```

Breadth-First Search

A breadth-first search (BFS) from s discovers the level sets of s : level L_i is the set of vertices reachable from s via a shortest path of length i (not reachable via a path of shorter length).

- level L_i is the set of vertices reachable from s via a shortest path of length i .
- $L_0 = \{s\}$
- So to compute level L_{i+1} , include every vertex with an incoming edge from a vertex in L_i , *that has not already been assigned a level.*, means not repeated vertices.

```
def bfs(Adj, s): # Adj: adjacency list, s: starting vertex  
    parent = [None for v in Adj] # O(V) add a length|Adj| None array.  
    parent[s] = s # root is string vertex  
    level = [[s]] # level 0 is [s]  
    while 0 < len(level[-1]): # while level has vertices  
        level.append([]) # add a new level  
        for u in level[-2]: # loop over last full level  
            for v in Adj[u]: # loop over neighbors  
                if parent[v] is None: # only v is not be visited.  
                    parent[v] = u  
                    level[-1].append(v)  
    return parent, level
```

- runtime
so the runtime is $O(|E| + |V|)$, see in note of LEC9