

Foundations of Computing

Lecture 11

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Outline

- 1 Lecture 10 Review
- 2 The CFG Pumping Lemma
- 3 Midterm Review

Lecture 10 Review

- $\text{CFG} == \text{PDA}$
 - Construct PDA from CFG
 - Construct CFG from PDA
- CFG Pumping Lemma

Lecture 10 Review

- CFG \equiv PDA
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Today

- Midterm review

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The CFG Pumping Lemma

Theorem

If L is a CFL, then there exists a pumping length p s.t. for any $s \in L$, with $|s| \geq p$, s can be divided into 5 pieces $s = uvxyz$ satisfying:

- ① For each $i \geq 0$, $uv^i xy^i z \in L$
- ② $|vy| > 0$
- ③ $|vxy| \leq p$

- Last week we used this to prove that

$$L = \{a^n b^n c^n \mid n \geq 0\}$$

is not context-free

Example 2

Consider $L = \{ww \mid w \in \{0,1\}^*\}$, prove L is not CFL

$L = \{ww^R\}$ is CFL

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Proof:

① Assume L is CFL, and let p be the pumping length

② Try 1: Choose $s = 0^p 0^p \in L$

$uvxyz = s$

$v = 00$

$y = 00$

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$\underbrace{\quad}_v \quad \underbrace{\quad}_\delta$

$y = 0^{p-2}$
 $\boxed{v = 0^2}$
 $x = 1$
 $\boxed{z = 0^2}$
 $z = 0^{p-2} \quad 1$

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- ③ Try 2: Choose $s = 0^p 10^p 1 \in L$
- ④ Choose $s = \underbrace{0^p 1^p}_{\text{red bracket}} \underbrace{0^p 1^p}_{\text{red bracket}} \in L$

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- ⑥ Contradiction – Hence L is not CFL

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- Recall closure properties of regular languages (complement, union, intersection, concatenation, $*$ closure)

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- NFA to DFA using the finger method

5 Regular Expressions

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- Understand why it is true (state of NFA must repeat)

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- Understand how to use it.

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- Be able to build an RE for a language
- RE to NFA
- NFA to RE

⑥ Regular Language Pumping Lemma

- Remember statement as sequence of quantifiers
- Understand why it is true (state of NFA must repeat)
- Understand how to use it.
- Also know how to prove languages not regular using closure properties
 - IMPORTANT: Make sure you understand how to make this into a proof by contradiction!

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- 8 Context-free Grammars (CFG)
 - Remember what this means

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- Be able to construct one from language description

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- Remember what a derivation is and what a parse tree is

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- Be able to construct one from language description
- Remember what a derivation is and what a parse tree is
- $\text{PDA} \implies \text{CFG}$ (at a high level)

9 CFL pumping lemma

- There will not be any questions on the CFL pumping lemma on the exam
- But, there will be on the next homework

Exam Format

- 7 questions – most have multiple parts
- Covers most of the material outlined above
- 2 questions requiring proofs, the rest are more constructive
- Some yes/no questions

Don't Forget

- Exam is in class on Thursday 11:10-12:25, don't be late!
- You can bring two 8.5×11 piece of paper
- 5 point bonus for participating in ACM hackathon this weekend

Any Questions?

RE \rightarrow NFA

\downarrow
 $(01)^*$ \cup (1010)

$\{0,1,6\}, 0 \rightarrow \Sigma 7\}$

$\{0,1,6\}, 1 \rightarrow \Sigma 2\}$

