Foundations of Computing

Lecture 6

Arkady Yerukhimovich

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Outline

1. Lecture 5 Review
2. A Nonregular Language
3. The Pumping Lemma for Regular Languages
4. Using the Pumping Lemma

Lecture 5 Review

- Regular expressions
- Equivalence of regular expressions and NFAs/DFAs

Quiz Solutions

For each of the following languages over $\Sigma = \{a, b\}$, give two strings that are in the language and two strings not in the language.

1. $a^* \cup b^*$

2. $(aa \cup bb)^*$

3. $\Sigma^* a \Sigma^* b \Sigma^* a \Sigma^*$

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2. Languages recognized by a DFA
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Are all languages regular?

Today we will see that there are languages that are not regular

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This means that:

- The number of states is fixed independently of the input size
- An automaton must be able to process strings w s.t. $|w| > |Q|$
- Thus, a finite automaton cannot store its whole input

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$$L = \{0^n 1^n | n \geq 0\}$$

Let's try to build a DFA for L:

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The Problem

We need to count the number of 0s, but this is unbounded so can't have a state for each value

The Need for a Proof

What we just saw

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Consider the following language:

$L = \{w | w \text{ has an equal number of occurrences of } 01 \text{ and } 10 \text{ as substrings}\}$

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2. Pick a string $w \in L$
3. Show that if $M(w) = 1$ then there exists a string $w' \notin L$ s.t. $M(w') = 1$
4. Conclude that L is not regular since any M that accepts all strings in L must also accept strings not in L

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Next steps:

1. Prove the pumping lemma
2. Show how to use the pumping lemma to prove languages nonregular

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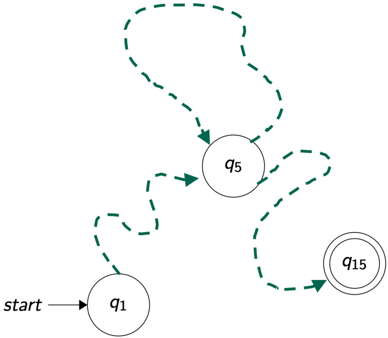

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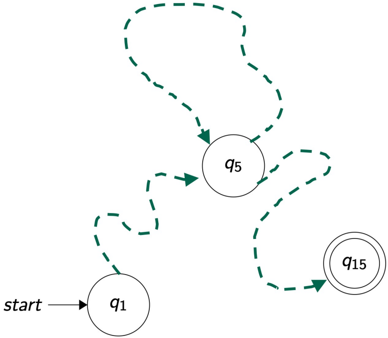

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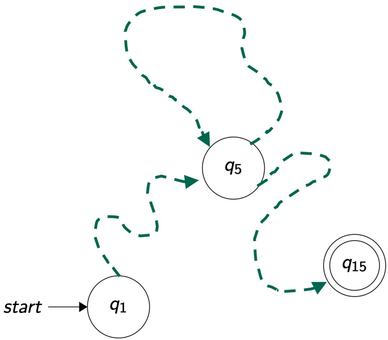

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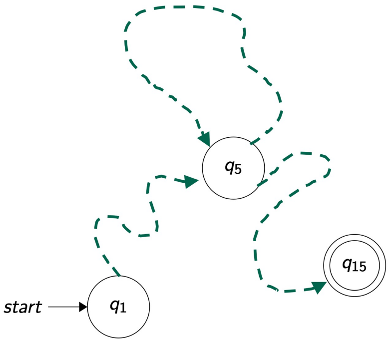

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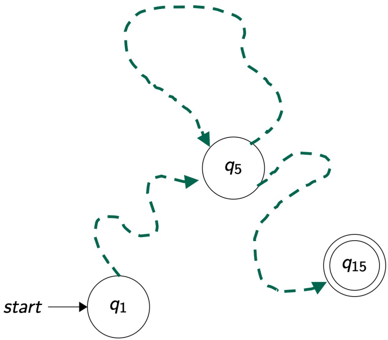

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3. $|xy| \leq p$
   Proof: if q_5 is the first repetition in $M(w)$, then this repetition must occur in the first $p + 1$ states, hence $|xy| \leq p$

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   - But, this means that y must be all 0s

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Consider $L = \{w | w$ has an equal number of 0s and 1s$\}$, prove L is not regular

Proof:

1. Assume L is regular, and let p be the pumping length this implies
2. Choose $w = 0^p 1^p$
3. By pumping lemma, $w = xyz$ s.t. $xy^i z \in L$
4. Problem: If $y = 0^m 1^m$, then w can be pumped – not leading to contradiction
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   - Since $w = 0^p 1^p$ and $|xy| \leq p$, we know that y must be in first p symbols
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7. Contradiction – hence, L is not regular

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3. Since regular languages are closed under $\cap$, if L is regular then L_1 must be regular
4. Since we know L_1 is nonregular, this means that L must be nonregular

Exercise

Prove that the following language is nonregular:

$$L = \{0^i 1^j 2^i 3^j | i, j > 0\}$$

What's Next?

- We will get plenty of practice with proving languages nonregular
- We will add (a small amount of) memory to our machines to recognize a richer class of languages