Cryptography

Lecture 6

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September 16, 2024

Outline

Lecture 5 Review

- Security of PRG+OTP
- Quiz on reductions

Outline

Problem 1

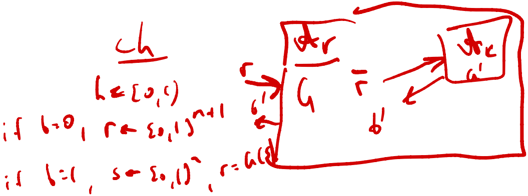
Let $G : \{0,1\}^n \to \{0,1\}^{n+1}$ be a PRG. Prove that

$$G'(s) = \overline{G(s)}$$

is a secure PRG

1. Assume PPT $\mathcal{A}_c$ that breaks G'

2. Build $\mathcal{A}_r$ using $\mathcal{A}_c$ s.t. $\mathcal{A}_r$ breaks G



$b=0$, $\bar{r}$ is random

$b=1$ $\bar{r} = G'(s)$

$\Pr[\mathcal{A}_r \text{ wins vs } G] = \Pr[\mathcal{A}_c \text{ wins vs } G'] \geq \frac{1}{2} + \frac{1}{poly(n)}$

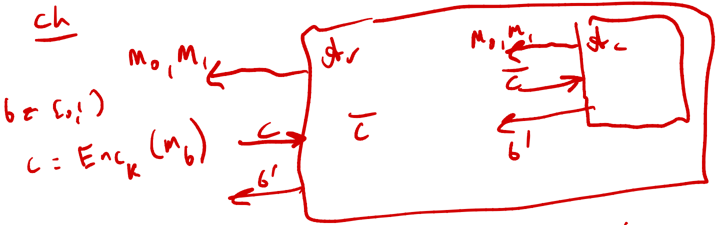
Problem 2

Let $\Pi = (\mathsf{Gen}, \mathsf{Enc}, \mathsf{Dec})$ be an encryption scheme secure vs. eavesdropper. Prove that

$$\mathsf{Enc}'_k(m) = \overline{\mathsf{Enc}_k(m)}$$

is also secure

1. Assume $\mathcal{A}_1$ that breaks Enc'
2. Construct $\mathcal{A}_2$ that breaks Enc



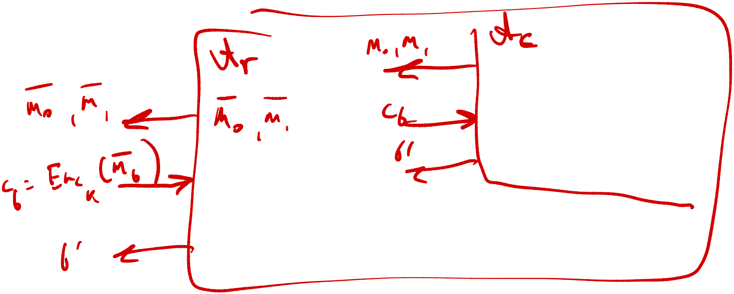
if $b=0$, $c = \mathsf{Enc}(m_0)$

$\bar{c} = \mathsf{Enc}'(m_0)$

$\Pr[\mathcal{A}_2 \text{ wins}] = \Pr[\mathcal{A}_1 \text{ wins}] \geq \frac{1}{2} + \frac{1}{poly(n)}$

Problem 3

What would change if we defined $\mathsf{Enc}'_k(m) = \mathsf{Enc}_k(\overline{m})$



Some Notes on Pseudorandomness

What is a PRG (Informal)

PRG says the following two distributions are indistinguishable

- $s \leftarrow \{0,1\}^n$, output $G(s)$
- $r \leftarrow \{0,1\}^{l(n)}$

Some Notes on Pseudorandomness

if $x = G(s)$ for some s $\quad Pr[x] = \frac{1}{2^n}$ $\quad$ o.w $Pr[x] = 0$

not true

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- $r \leftarrow \{0, 1\}^{l(n)}$ $\quad \forall\ x \in \{0,1\}^{\ell(n)}$ $\quad Pr[x] = \frac{1}{2^{\ell(n)}}$

Observations:

- This does not mean that $G(s) = r$. Equality of distributions, not strings

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Observations:

- This does not mean that $G(s) = r$. Equality of distributions, not strings
- In particular, incorrect to say string w is pseudorandom
- Indistinguishability only holds for PPT adversaries
- Most easily captured in a game

Outline

Security of PRG+OTP: Intuition

PRG+OTP Encryption

- $\text{Gen}(1^n)$: $k \leftarrow \{0,1\}^n$
- $\text{Enc}(k, m)$: $c = G(k) \oplus m$
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Security of PRG+OTP: Intuition

Assumption: $G : \{0,1\}^n \to \{0,1\}^{l(n)}$ is PRG

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Proof:

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- Assume there exists PPT $\mathcal{A}_c$ that breaks Π $(\Pr[PrivK^{eav}_{\mathcal{A}_c,\Pi}(1^n) = 1] > 1/2 + 1/\text{poly}(n))$
- Construct $\mathcal{A}_r$ that breaks G:

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Intuition

- $\mathcal{A}_r$ receives either $r \leftarrow \{0,1\}^{l(n)}$ or $r = G(s)$, uses r as mask

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- $\mathcal{A}_r$ receives either $r \leftarrow \{0,1\}^{l(n)}$ or $r = G(s)$, uses r as mask
- If $r \leftarrow \{0,1\}^{l(n)}$, Π is just OTP ($\Pr[\mathcal{A}_c \text{ WINS}] = 1/2$)

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Intuition

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- If $r = G(s)$, Π is PRG+OTP (by assumption, $\Pr[\mathcal{A}_c$ WINS$] > 1/2 + 1/\text{poly}(n)$)
- $\mathcal{A}_r$ runs $\mathcal{A}_c$ generating challenge c using r, observes if $\mathcal{A}_c$ wins, and if so outputs "PRG".

Security of PRG+OTP: Building $\mathcal{A}_r$

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$PRG_{\mathcal{D},G}(n)$

- The challenger chooses $b \leftarrow \{0,1\}$.
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- On input r, the distinguisher $\mathcal{D}$ outputs a guess b'
- $PRG_{\mathcal{D},G}(n) = 1$ (i.e., $\mathcal{D}$ wins) if $b' = b$

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- $\mathcal{A}$ outputs two messages $m_0, m_1 \in \mathcal{M}$
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  - $\mathcal{A}_r$ outputs 1 ("PRG") if $b = b'$ and 0 otherwise

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  - $\Pr[\mathcal{A}_r(r) = 1] = \Pr[\mathcal{A}_c$ outputs $b' = b] = \Pr[PrivK^{eav}_{\mathcal{A}_c,\Pi}(1^n) = 1] \geq 1/2 + 1/\text{poly}(n)$
- Summing these together, we get

r is random

$$\begin{aligned}\Pr[PRG_{\mathcal{A}_r,G}(1^n) = 1] &\geq 1/2 \cdot 1/2 + 1/2 \cdot (1/2 + 1/\text{poly}(n)) \\ &= 1/2 + 1/(2\text{poly}(n))\end{aligned}$$

Contradiction!

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Where Are We Now

- Features of PRG+OTP encryption
  - Can encrypt messages of arbitrary length, just need PRG with enough stretch.
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- Features of PRG+OTP encryption
  - Can encrypt messages of arbitrary length, just need PRG with enough stretch.
  - Achieve security against an eavesdropper
- Limitations of PRG+OTP encryption
  - Can only see one encryption
  - If see two, can tell whether they are equal

CPA Security Intuition

CPA Security

- $\mathcal{A}$ is allowed to request encryptions (under key k) of any messages of its choice.
- $\mathcal{A}$ still cannot learn any information about encrypted message when seeing challenge ciphertext c.

Why We Need CPA Security - A Historical Motivation

British Mines:

- British would bury a mine at specific latitude, longitude
- When Germans would find the mine, they would encrypt location and send back to HQ
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Battle of Midway:

- US forces intercepted and partially decrypted Japanese message
- Learned that Japan was going to attack location "AF" - wanted to confirm that this was Midway Island
- US sent out a message that "Midway is low on water" making sure Japanese intercepted it
- Japanese forces send message "AF is low on water" to HQ

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- Example: Can give oracle access to $\mathsf{Enc}_k(\cdot)$. This allows caller to encrypt m of its choice without learning k.
- Notation:
  - We write $\mathcal{A}^{\mathcal{O}(\cdot)}$ to indicate a party $\mathcal{A}$ given oracle access to some function $\mathcal{O}$.
  - Calls to $\mathcal{O}$ cost 1 computation step
  - Oracles are a useful tool in security definitions and proofs

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Definition: An encryption scheme $\Pi = (\mathsf{Gen}, \mathsf{Enc}, \mathsf{Dec})$ with message space $\mathcal{M}$ is CPA-secure if for all PPT $\mathcal{A}$ it holds that

$$\Pr[\mathsf{PrivK}^{cpa}_{\mathcal{A},\Pi}(n) = 1] \leq 1/2 + \mathsf{negl}(n)$$

Observations

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- She still cannot learn any information about encrypted message.

Benefits of CPA-Security

- Can encrypt many messages
  - $\mathcal{A}$ gets to see encryptions of many messages of its choice, still cannot break security of challenge
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- Can encrypt arbitrarily long messages
  - Break message m into n-bit blocks, $m = m_1 || m_2 || \cdots || m_\ell$
  - To encrypt m separately encrypt each m_i.
  - Secure since this is just encrypting many messages

How to Construct CPA-Secure Encryption

- Recall that PRG+OTP encryption allowed us to encrypt long messages.
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Key Idea

What if encryption (and decryption) could generate a different OTP for each ciphertext?

Note: We need to produce enough OTP's for as many encryptions as $\mathcal{A}$ wants. So, can't just pre-generate them all.

Outline

What Is a Random Function?

Consider a function $f : \{0,1\}^n \to \{0,1\}^n$

x	f(x)
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Question:

How can we get the benefits of a random function without paying the overhead?

PRF Goals

Construct an efficient, keyed function
$F : \{0,1\}^n \times \{0,1\}^n \to \{0,1\}^n$ such that:

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- $F_k(\cdot)$ is efficiently computable
- For a random key $k \leftarrow \{0,1\}^n$, $F_k(\cdot)$ looks like a *random function* from n bits to n bits (to someone who doesn't know k).

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Let $F : \{0,1\}^n \times \{0,1\}^n \to \{0,1\}^n$ be a deterministic, keyed, poly-time function.

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- With access to oracle $\mathcal{O}$, the distinguisher $\mathcal{D}$ outputs a bit b'

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Let $F : \{0,1\}^n \times \{0,1\}^n \to \{0,1\}^n$ be a deterministic, keyed, poly-time function.

$PRF_{\mathcal{D},F}(n)$

- The challenger chooses $b \leftarrow \{0,1\}$.
  If $b = 0$, he chooses $f \leftarrow \mathcal{F}_n$ and gives $\mathcal{D}$ an oracle $\mathcal{O} = f$.
  if $b = 1$, he chooses $k \leftarrow \{0,1\}^n$, and gives $\mathcal{D}$ an oracle $\mathcal{O} = F_k$.
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$\mathcal{D}$ cannot distinguish between oracle access to a random function and oracle access to a PRF (for a key k he doesn't know).

Observations

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- $\mathcal{D}$ can make polynomially many queries to $\mathcal{O}$
- $\mathcal{D}$ can choose its queries adaptively based on results of earlier queries
- The set of polynomially many evaluations of $F_k(\cdot)$ must look random
- Clearly, this is not possible if $\mathcal{D}$ knows k

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  - If $\mathcal{O} = F_k$, then $(y_1 \oplus y_2) = (x_1 \oplus x_2)$ with probability 1
  - If $\mathcal{O} = f$, then $(y_1 \oplus y_2) = (x_1 \oplus x_2)$ with probability $1/2^n$
- So, $\mathcal{D}$ always outputs 1 when $\mathcal{O} = F_k$ and outputs 1 with probability $1/2^n$ when $\mathcal{O} = f$.

$$\Pr[\mathcal{D} \text{ WINS}] = \Pr[b = 1] \cdot 1 + \Pr[b = 0] \cdot (1 - 1/2^n) > 1/2$$

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  - In applied crypto, this is often called a *blockcipher*.

Relationship Between PRG and PRF

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Relationships:

- Not hard to show that a PRF can be used to build a PRG
- In fact, PRG can also be used to build a PRF
- But, important to remember the differences in functionalities and security definitions