

Linked List

Hsuan-Tien Lin

Dept. of CSIE, NTU

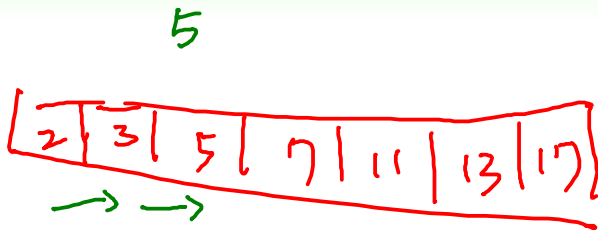
March 9, 2021

What We Have Done

- pseudo code: 'spoken' language of programmers
- data structure: scheme of data organization
- why dsa: proper use of resources;
move from coding to programming
- array: (fast)-index-access data structure
- ordered array: consecutive array with sorted values
 - harder to maintain, e.g. `insert`
 - easier to search by value

More on Ordered Array

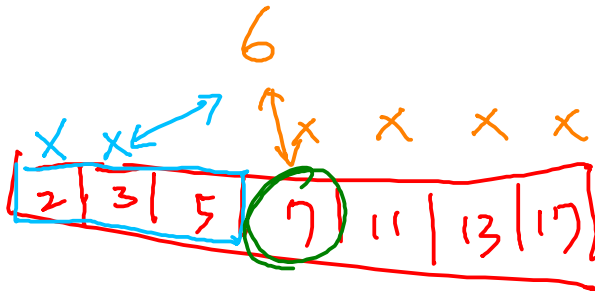
Ordered Array: Sequential Search Algorithm with Cut



→ → → × × ×

ordered: possibly easier to declare **not found**

Ordered Array: Binary Search Algorithm



“cut” multiple times by fast random access to the middle

Singly Linked List

Application: Polynomial Computation

$$f(x) = \underline{5} x^{\underline{9}} - \underline{7} x^{\underline{4}} + \underline{3}$$

$$(5, 9) \quad (-7, 4) \quad (\underline{3}, \underline{0})$$

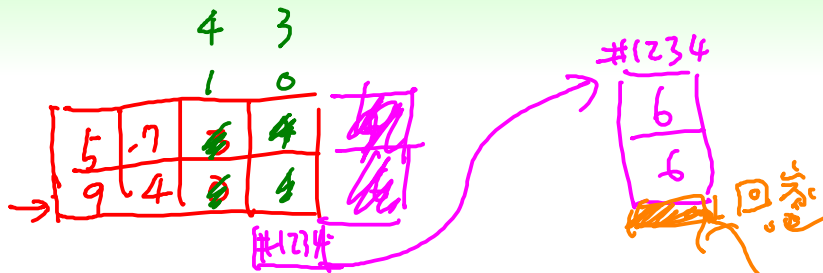
→

5	9
-7	4
3	0

5	9	5	9	5	9	5	9	...
								
								
								
								

solution 0: use ordered array on (exponent, coefficient)

Issues of (Ordered) Array for Polynomial Computation



head

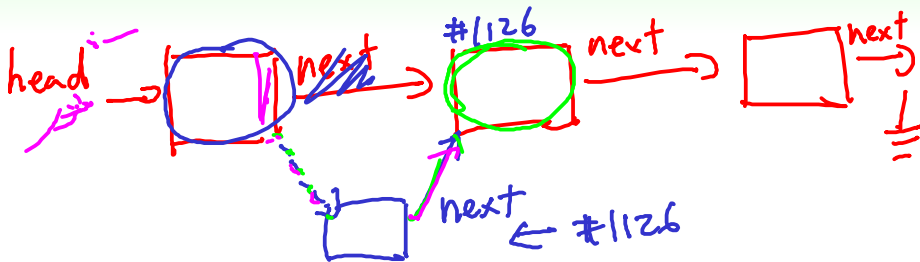
$$f(x) = 5x^9 - 7x^4 + 3 + 4x + 6x^6$$

↓

(5, 9) → (-7, 4) → (4, 1) → (3, 0)

ordered (consecutive) array:
not flexible for resizing/insertion/removal

Solution: Singly Linked List for Flexible Insertion



overhead of next \Leftrightarrow flexible insertAfter

Singly Linked List as Abstract Data Structure: Access access

- data getAt (node) ←
- node getHead() ←
- node getNext (node) ←
- insertAfter (node, data) ←
- insertHead (data) ←

linked list: sequential access; array: random access

Singly Linked List as ADT: Maintenance

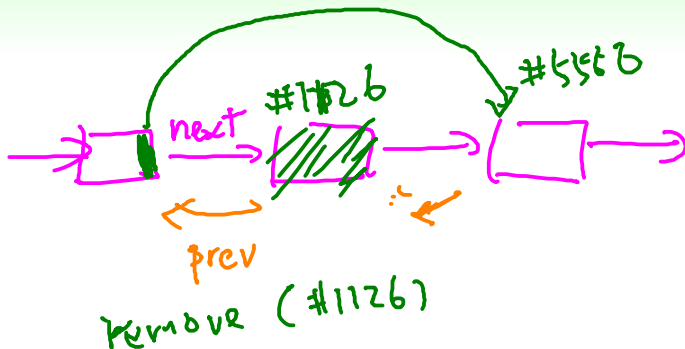
maintenance

- `construct (length) : trivial` ←
- `updateHere (node, data) : trivial` ←
- `removeAfter (node) : simple` ←
- `removeHead : simple` ←

think: dummy head node or not?

Doubly Linked List

removeHere for Singly Linked List



↓ ↓
 removeHere (and insertHere): hard for singly linked list

Doubly Linked List: More Flexible `removeHere`

overhead of `prev` \Leftrightarrow flexible `removeHere` (and flexible traverse backward)

Linked List for Sparse Vectors

Application: Sparse Vector in Scientific Computing



✓

polynomial: can be viewed as special case of sparse vector

Sparse Vector: (Dense) Array versus Linked List

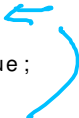
storing only non-zeros can be time/space efficient

Merging Sparse Vectors

“running cursors” algorithm:
similar for other uses, like dot product

Real-World Usage of Sparse Vector: LIBSVM

```
1  double Kernel::dot(const svm_node *px, const svm_node *py){
2      double sum = 0;
3      while(px->index != -1 && py->index != -1){
4          if(px->index == py->index){
5              sum += px->value * py->value;
6              ++px;
7              ++py;
8          }
9          else{
10             if(px->index > py->index)
11                 ++py;
12             else
13                 ++px;
14         }
15     }
16     return sum;
17 }
```



good data structure needed everywhere