

# Tree / Heap

Hsuan-Tien Lin

Dept. of CSIE, NTU

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# What We Have Done

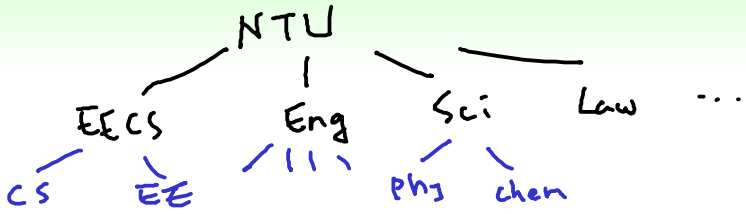
- stack (LIFO): infix to postfix, postfix evaluation
- queue (FIFO): maze solving (how different data structures affect algorithm behavior), implementation by circular array
- deque: stack + queue + push\_front

# Nature of Data Structures

data structure	nature
array	indexed access
linked list	sequential access
stack/queue/deque	restricted (boundary) access
tree	hierarchical access

next: tree

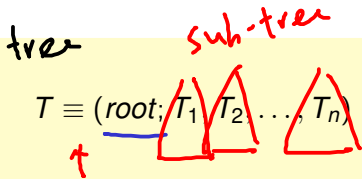
# Trees



- parent-child relationship: file system, hierarchical structure

dir  
file

# Formal Definition of Trees



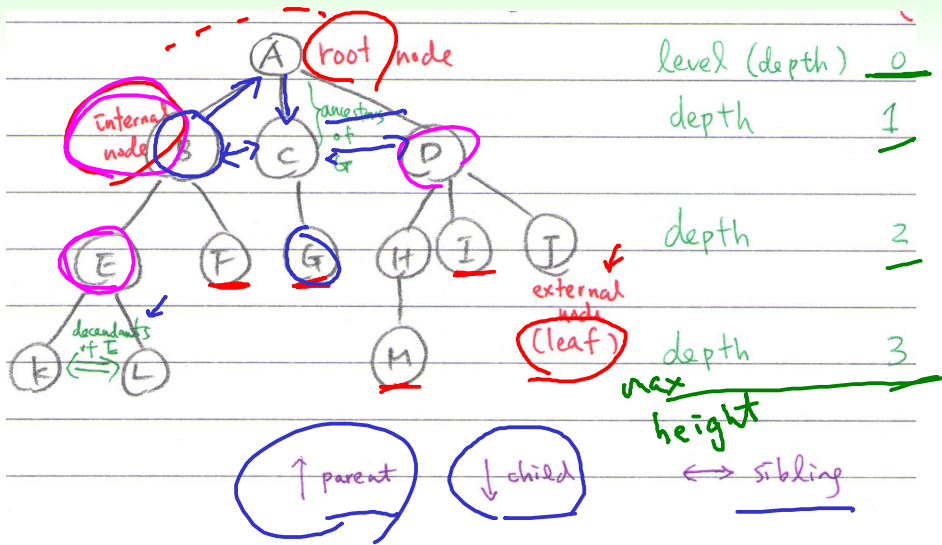
- recursive definition
- sub-trees  $(T_1, \dots, T_n)$  disjoint
- recursion termination?



$T \equiv (\text{root}; \text{nil})$

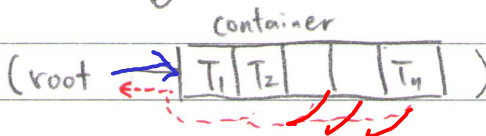
leaf

# Names in Trees

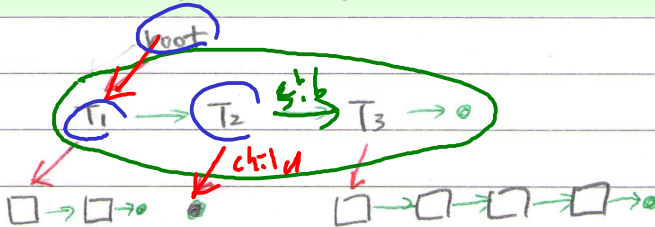


# Representing Trees

$$T \equiv (\text{root} ; T_1, T_2, \dots, T_n)$$



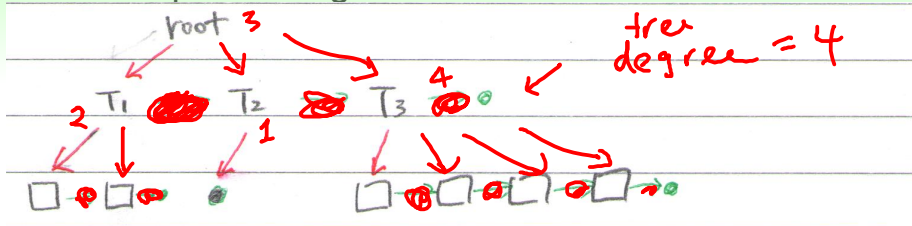
# Representing Trees with Linked Lists



called left-child right-sibling  
# link per node?



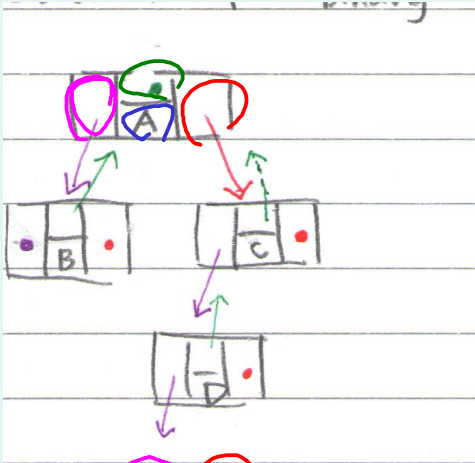
# Representing Trees with Sub-Tree Links



sibling links  $\rightarrow$  child links

'easiest' visual representation of trees  
# link per node?

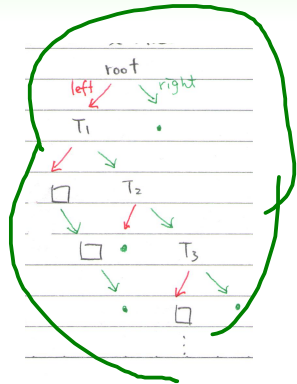
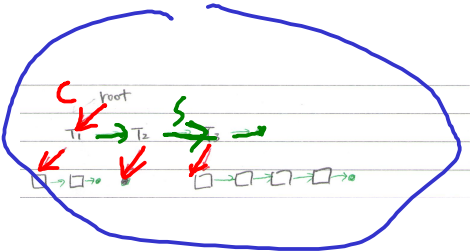
# Binary Trees



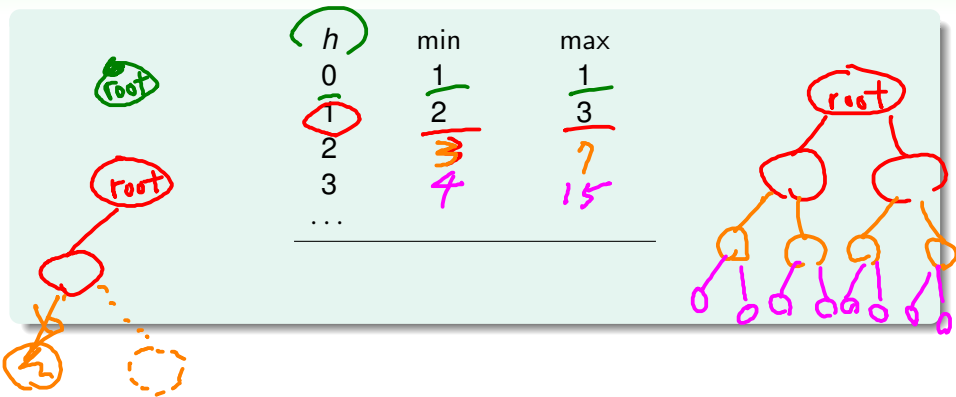
left, right, (parent)

$\leq 2$  ordered children per node

# General Tree + Left-Child Right-Sibling $\equiv$ Binary Tree



# How Many Nodes in Binary Trees



# How Many Nodes in Binary Trees

$$n=10 \quad \log_2(n+1) - 1 \leq h \leq n-1$$

$h$	min	max
0	1	1
1	2	3
2		
3		
...		

$h$	$h+1$	$2^{h+1} - 1$
	(skewed)	full

$$h+1 \leq n \leq 2^{h+1} - 1$$

$\Leftrightarrow$

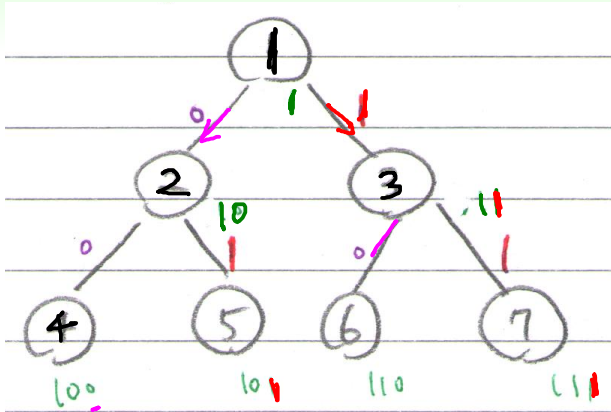
#nodes

# How Many Nodes in Binary Trees

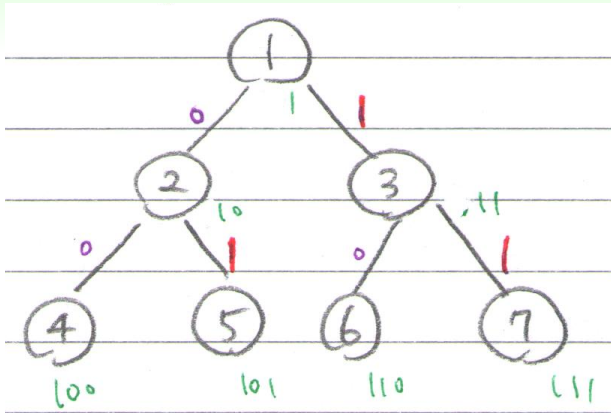
$h$	min	max
0	1	1
1	2	3
2		
3		
...		
$h$	$h + 1$	$2^{h+1} - 1$
	(skewed)	full

$$h + 1 \leq n \leq 2^{h+1} - 1$$
$$\Leftrightarrow \lg(n + 1) - 1 \leq h \leq \frac{n-1}{2} \cdot 2$$

# Node Index in Full Binary Tree



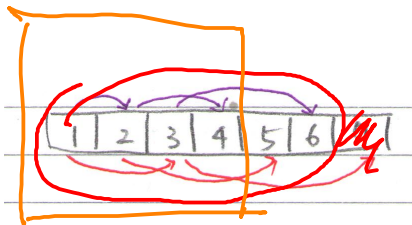
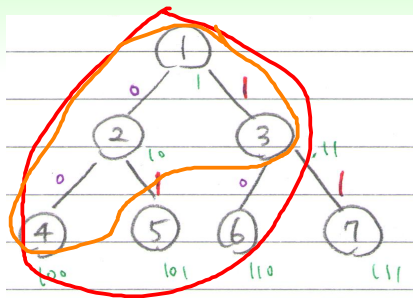
# Node Index in Full Binary Tree



$$\text{node index} = (1\text{path code})_2$$



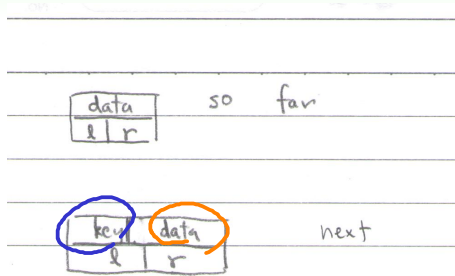
# Representing/Packing Full Binary Tree in Array



- implicit links: no need for explicit pointers
- can similarly pack any binary tree if NIL can represent NO-DATA (with **space wasting** in data field)

complete binary tree: full binary tree with first  $n$  nodes  
(**no waste** with array representation)

# Need: Priority Queue (with Binary Tree)



- key: priority
- data: item in todo list

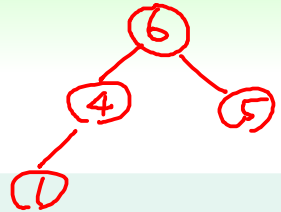
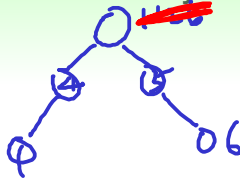
goal: get the node with **largest** key

# Design Thoughts of Priority Queue with Tree (1/4)

getLargest

- entry point of tree? *root*
- to allow fastest access, put **largest** close to entry point

# Design Thoughts of Priority Queue with Tree (1/4)



- entry point of tree?
- to allow fastest access, put **largest** close to entry point

but what if not just `getLargest` but need `removeLargest`?

## Design Thoughts of Priority Queue with Tree (2/4)

removeLargest needs 2nd largest to replace

- to allow fastest removal, put 2nd largest close to next entry points
- next entry points of tree? children of root

# Design Thoughts of Priority Queue with Tree (2/4)

removeLargest needs 2nd largest to replace

- to allow fastest removal, put **2nd largest** close to next entry points
- next entry points of tree?

max tree:

- root key  $\geq$  other node's key
- every sub-tree also max tree



removeLargest (version 0): recursive **duel** of sub-tree roots

$O(h)$