

# Heap / Binary Search Tree

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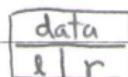
Dept. of CSIE, NTU

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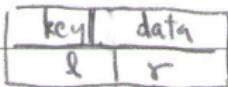
# What We Have Done

- tree: hierarchical access
- binary tree: “simplest” tree, where (complete binary tree  $\sim$  array)
- binary max-tree: binary tree + max-at-root, for priority queue

# Need: Priority Queue (with Binary Tree)



so far



next

- key: priority
  - data: item in todo list
- goal: get the node with **largest** key

## Design Thoughts of Priority Queue with Tree (2/4)

removeLargest needs 2nd largest to replace

- to allow fastest removal, put **2nd largest** close to next entry points
- next entry points of tree?

max tree:

- root key  $\geq$  other node's key
- every sub-tree also max tree

)

removeLargest (version 0): recursive **duel** of sub-tree roots

-  $O(h^2)$   $O(h)$   $O(n^2)$   
 $T O(\log n)$  for  $n$  elements

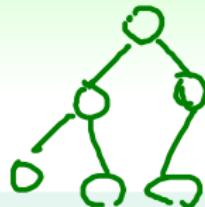
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  - what is shortest binary tree with  $n$  nodes?

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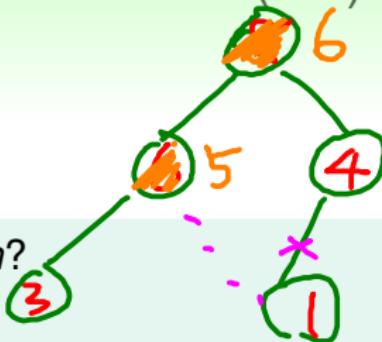
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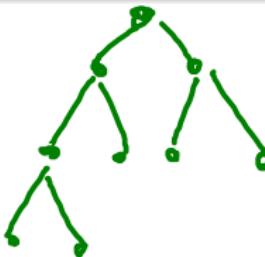


8



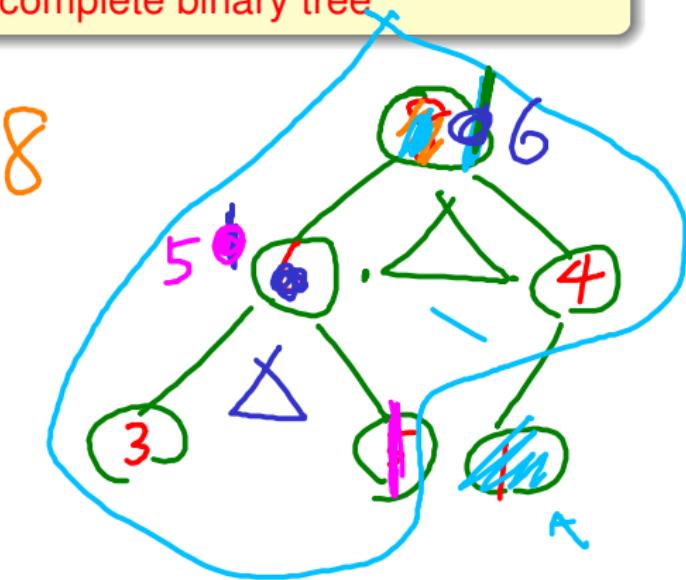
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- (binary) max-heap: max-tree + **complete binary tree**

does `removeLargest` (version 0) work with binary heap?



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- removeLargest (version 1)
  - move last node to root
  - three-way-duel (root & two sub-roots) in each sub-tree recursively:  
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# Design Thoughts of Priority Queue with Tree (4/4)

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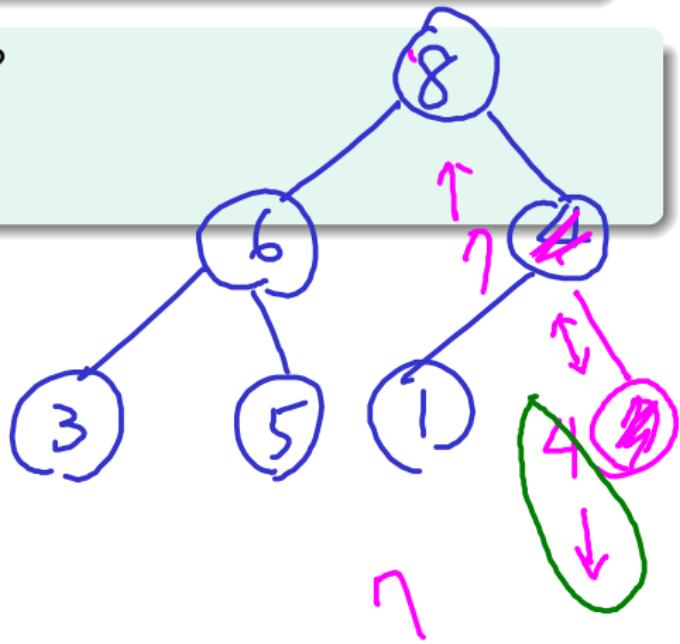
lesson learned: preserve **harder constraints** first

## Priority Queue Insertion

8, 6, 4, 3, 5, 1

binary max-heap: max-tree + **complete binary tree**

- insertion: preserve which first?



# Priority Queue Insertion

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binary max-heap: max-tree + **complete binary tree**

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- how? insert where first?
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insert

- insert at last position 
- duel with parents until satisfied:  $O(h)$

# Binary Max-Heap in Array

complete binary tree  
max-heap



consecutive array

partially-ordered consecutive array



# Binary Max-Heap in Array

complete binary tree  
max-heap

consecutive array  
partially-ordered consecutive array

unordered array  
selection sort  $O(n) \cdot O(n)$

partially-ordered array (max-heap)  
heap sort  $O(n) \cdot O(\log n)$

先有 heap  
run sel sort w/ heap

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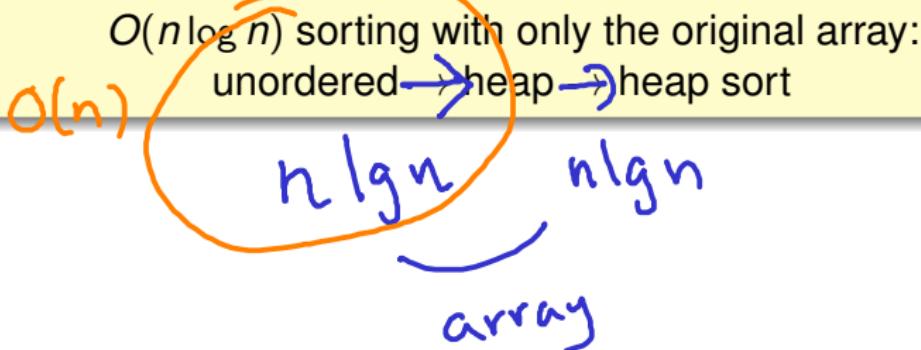
consecutive array  
partially-ordered consecutive array

unordered array  
selection sort  $O(n) \cdot O(n)$

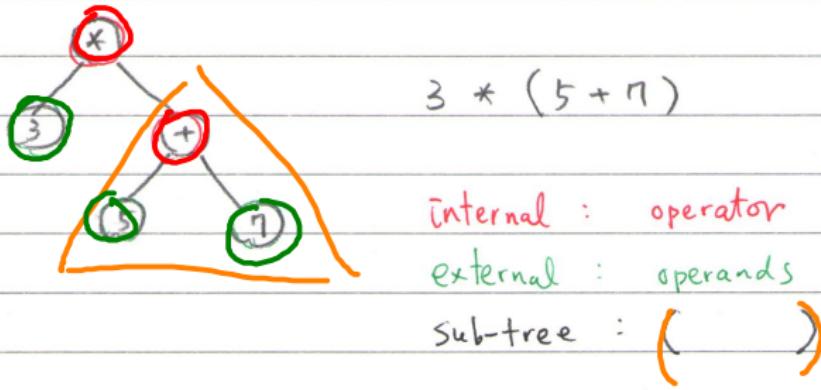
partially-ordered array (max-heap)  
heap sort  $O(n) \cdot O(\log n)$

$O(n)$        $O(n \log n)$  sorting with only the original array:  
unordered  $\rightarrow$  heap  $\rightarrow$  heap sort

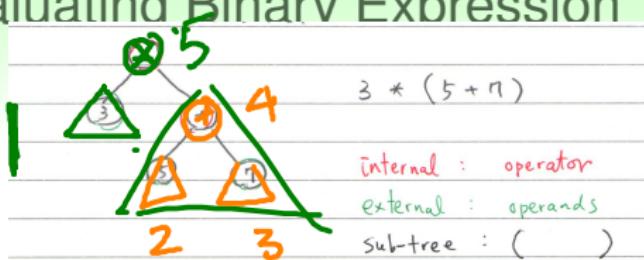
$n \lg n$        $n \lg n$   
array



# More on Binary Trees: (Binary) Expression Tree

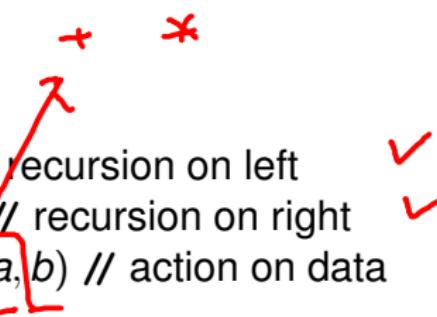


# Evaluating Binary Expression Tree



EVALUATE( $T$ )

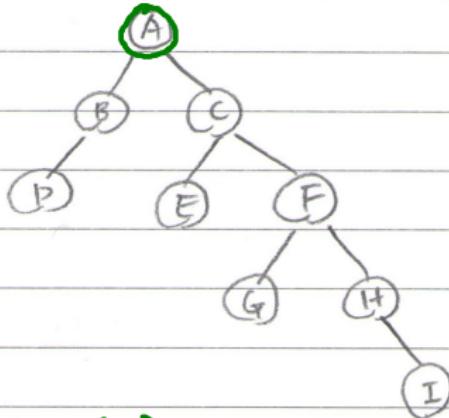
```
1 if  $T.data$  is a number
2   return  $T.data$ 
3 else
4    $a = \text{EVALUATE}(T.left)$  // recursion on left
5    $b = \text{EVALUATE}(T.right)$  // recursion on right
6   return  $\text{compute}(a, T.data, b)$  // action on data
```



traversal: how  $T.data$  is visited

post-order traversal:  $T.data$  visited after sub-tree recursions

# Other Traversals



in : DBAECGFHI  
post : DBE[G][I][H][F]CA ←  
pre : ABDCEFGHI  
.....

- **Post( $T$ )** ↓ ↓ ↓
  - post-order: Post(T.left); Post(T.right); Visit(T.data);  
—expression tree evaluation: Visit by result computation
  - pre-order: Visit(T.data); Pre(T.left); Pre(T.right);  
e.g. comparing if two trees are equal using equality as Visit
  - in-order: In(T.left); Visit(T.data); In(T.right);  
e.g. output ordered data from tree if  $T_L < \text{root} < T_R$  (see next)

# Binary Search Tree



< root <



any subtree

BST-SEARCH( $T, key$ )

```
1
2 while  $T \neq \text{NIL}$ 
3      $mid = T // \text{root}$ 
4     if  $mid.\text{key} = \text{key}$ 
5         return  $mid$ 
6     elseif  $mid.\text{key} < \text{key}$  )T1
7         return  $\text{??}$ 
8     elseif  $mid.\text{key} > \text{key}$  )T2
9         return  $\text{??}$ 
10    return FAIL
complexity??  $O(h)$   $\Leftarrow O(\log n)$ 
```

BINARY-SEARCH( $A, key$ )

```
1  $left = 1, right = A.\text{length}$ 
2 while  $left \leq right$ 
3      $mid = \text{floor}(\frac{left+right}{2})$ 
4     if  $A[mid] = \text{key}$ 
5         return  $mid$ 
6     elseif  $A[mid] < \text{key}$  )T3
7          $left = mid + 1$ 
8     elseif  $A[mid] > \text{key}$  )T4
9          $right = mid - 1$ 
10    return FAIL
complexity:  $O(\log n)$   $\Leftarrow$ 
```