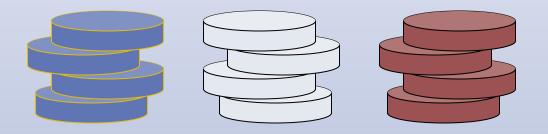
Stacks



The Stack

- The concept of **stack** is derived from the metaphor of a stack of plates in a spring-loaded cafeteria dispenser.
- If you want to remove a plate, you pop the a plate off the top of stack.
- If you want to replace a plate or insert more plates, you
 push onto the top of the stack.
- To check if ALL plates were clean, you would need to check the top plate, remove that plate, and repeat the process until the entire stack was inspected.

Stack ADT

- Definition: A stack is a collection of objects that are inserted and removed according to the last-in-first-out (LIFO) principle.
 - Objects are inserted (as long as stack not full) onto the top of the stack.
 - Objects can ONLY be removed from the top of the stack.
- All stack operations are O(1).

Stack ADT

- Main stack operations:
 - push(o): inserts object o on top of stack
 - STL operation push(o)
 - pop(): <u>removes</u> element from the top of the stack
 - STL operation *pop()*
 - An error occurs if the queue is empty. (exception)
 - o *top()*: examines the top object on the stack **without** removing it.
 - Used in combination with pop()
 - top() to inspect element, pop() to remove top element

- Auxiliary **stack** operations:
 - o **size**(): returns the number of objects in a stack. Either store as a variable counter or calculate it.
 - isEmpty(): returns true if the stack is empty, else false

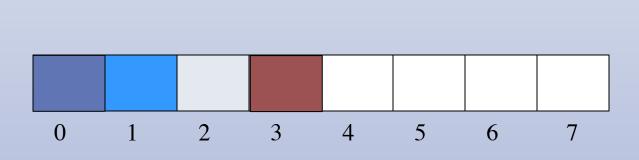
Array-based Stack

- Store the elements in an N-element array S
- Have an integer variable t that gives the index of the top element in the array S
- The top element in the array S is stored in the cell S[t]
- See an example...

Array-based Stack

- A simple way of implementing the Stack ADT uses an array
- We push (add) elements from left to right
- A variable keeps track of the index of the last item pushed

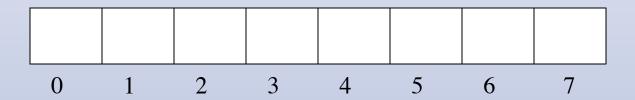
Top = 3



Array-based Stack

• We pop (remove) elements from right to left

$$Top = -1$$



Stack ADT Pseudocode

```
Algorithm size():
  return t+1
Algorithm isEmpty():
  return (t<0)
Algorithm top():
  if isEmpty() then
      throw a StackEmptyException
  return S/t/
```

Stack ADT Pseudocode

```
Algorithm push(o):

if size() == N then

throw a StackFullException

t \leftarrow t+1

S[t] \leftarrow o
```

```
Algorithm pop():

if isEmpty() then

throw a StackEmptyException
t \leftarrow t-1
```

Evaluation

Performance

- Let *n* be the number of elements in the stack
- \circ The space used is O(n)
- Space complexity = memory required new!
 Each operation runs in time *O*(1)
- - Time complexity = running time

Limitations

- o The maximum size of the stack must be defined a priori, and cannot be changed
- Trying to push a new element into a full stack causes an implementation-specific exception

Stack Data Structure

```
class Stack
{
private:
    objectType stack[MAX_STACK_SIZE];
    int top;
public:
    functions for stack manipulation
    constructor sets top to -1
};
```

Stack Implementation - Push

- The array storing the stack elements may become full
 - O Limitation of the array-based implementation

```
void push ( const objectType & o )
{
    if ( top + 1 == MAX_STACK_SIZE )
        throw FullStackException
    else
        S[++top] = o;
}
```

Stack Implementation-Pop

- (Stop here) Quick exercise- write pop and getTop functions
 - Array may be empty when pop
 - o getTop will return top item/object

```
void pop(){
    if ( isEmpty ( ) )
        throw EmptyStackException
    else
    --top;
```

Stack Implementation-Top

- (Stop here) Quick exercise- write pop and getTop functions
 - Array may be empty when pop
 - o getTop will return top item/object

```
objectType getTop(){
    if ( isEmpty ( ) )
        throw EmptyStackException
    else
        return S[top];
```

Stack Applications

Checking for balanced symbols in a program.

}

- Evaluating postfix (Reverse Polish Notation) expressions.
- Infix to Postfix expression conversion.
- Managing function calls in a program.

Reverse Polish Notation (RPN): Postfix

- Operators * , / , + , follow their operands:
 - 0 3 + 8 (in infix)
 0 3 8 + (in postfix)
- For expressions with multiple operands, operator occurs immediately after its second operand.

```
0 40 4 5 * - (in postfix)
0 40 (4*5) - , 40 20 - , 40 - 20 , 20
```

- Eliminates need for parentheses to force operator precedence.
- Used widely for computation in early desktop calculators.

Algorithm PostfixEvaluation

```
Process infix expression one item (p) at a time, left-to-right if (p == operand) // examples: 5, 7, 77, 2

push(p)

if (p == operator) // examples: *, /, +, -

top/pop and write to b

top/pop and write to a

push(a operator b)
```

Algorithm PostfixEvaluation

```
3 * (5 + ((2 + 3) * 8) + 5) => 3 5 2 3 + 8 * + 5 + *
```

Current Symbol	Stack	
3	3	
5	3 5	
2	3 5 2	
3	3 5 2 3	
· ·	2 F F	

Algorithm PostfixEvaluation

```
3 * (5 + ((2 + 3) * 8) + 5) => 3 5 2 3 + 8 * + 5 + *
```

Current Symbol	Stack
8	3 5 5 8
*	3 5 40
+	3 45
5	3 45 5
+	3 50
*	150

Algorithm Infix2Postfix

```
Process infix expression one item (p) at a time, left-to-right
     if (p == operand) // examples: 5, 7, 77, 2
        write to output
     if (p == operator) // examples: *, /, +, -
        top/pop and write to output
        until top of stack is ( or item of lower precedence than p
        push(p)
     if (p == '(')
        push(p)
     if (p == ') ')
        top/pop and write to output
        until (, pop but do NOT write (
     if (p == NULL) // p is empty
       top/pop and write to output
       until stack is empty
```

if (top)

else

return top->obj;

throw stack_empty exception;

```
bool isEmpty(){
                                         objectType getTop(){
   if (top == NULL)
     return true;
   else
     return false;
void push ( const objectType & obj ) {
  Node * newNode = new Node;
   newNode \rightarrow obj = obj;
   newNode \rightarrow next = top;
   top = newNode;
```

Top $\longrightarrow X$

