CS 014: Intro to Data Structures and Algorithms

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Lists

Lists

- Lists
 - List of students
 - List of games
 - List of assignments to complete
 - Etc.
- Is a <u>collection</u> of elements.
- One of most fundamental/simple data structures.
- Implementation:
 - Array-based
 - Node-based

Lists Implemented with Arrays

- A list can be implemented with an array.
 - Static memory allocation (space allocated at compile time)

```
int MAX_ARRAY_SIZE = 50;
int list[MAX_ARRAY_SIZE];
```

Dynamic memory allocation (space allocated at runtime)

```
int* list = NULL;
int MAX_ARRAY_SIZE;
cin >> MAX_ARRAY_SIZE;
list = new int[MAX_ARRAY_SIZE];
  // Initialize all elements to zero.
  delete[] list;
  list = NULL;
```

Lists Implemented with Arrays

Strengths:

- + Direct access
 - + Using subscript operator []
- + Good choice where list is built by inserts at end, and no deletes occur, only array accesses, i.e. find(pos)
- + Can be very fast and efficient.

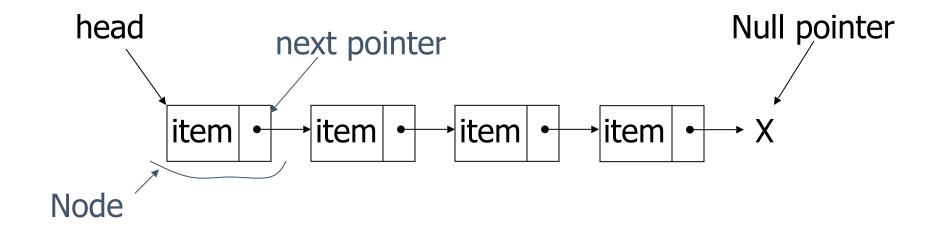
• Weaknesses:

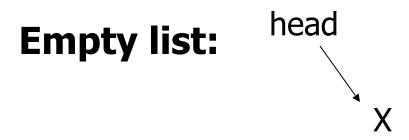
- Wasted space or run out of space.
 - Mostly empty (waste)
 - Full (need to resize, typically double—making list half full)
- Costly to insert in the middle, or beginning of array-based list.
 - Need to shift items up.

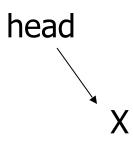
Linked List

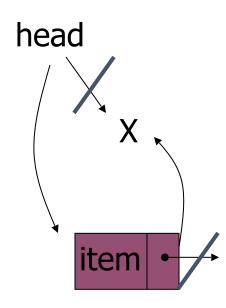
- Composed of nodes.
- Each node holds data (item, value, element).
- Each node has a pointer (next) that connects to the next node in the list.
- Special pointer (head) to beginning of list. (Access through list ID->head)
- Size of the list is flexible (grows on inserts, shrinks on deletes)
- List operations:
 - Insert item
 - Remove item
 - Read from item (access)
 - Write to item (mutate)
 - Traversal (visit each item: search, print)

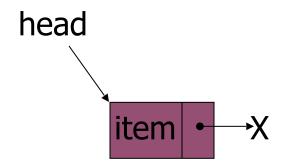
Linked List

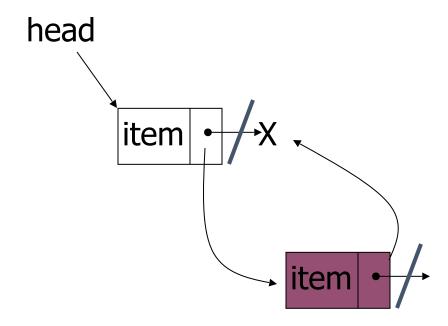


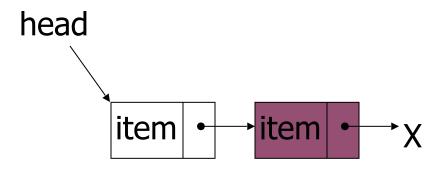


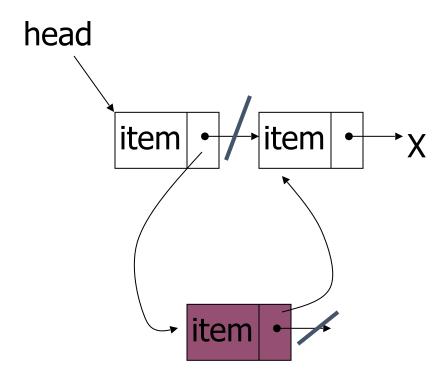


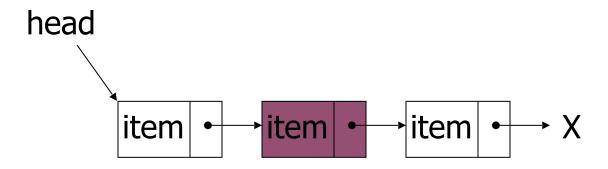




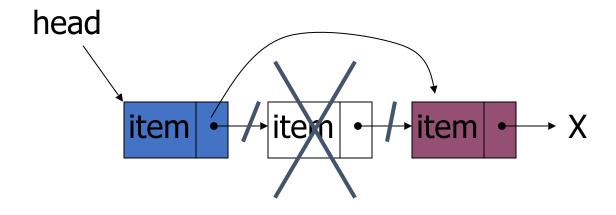








Linked list - Deleting Nodes



Linked List Implementation

```
class List {
private:
       Node* head;
                                    class Node {
public:
                                       friend class List;
      List ( ) {
                                    private:
         head = NULL;
                                       itemtype Item;
                                       Node* next;
      // Member functions
      // List Interface
};
```

```
Node* a = new Node;

Node* b = new Node;

Node* c = new Node;

a->next = b;

b->next = c;

b->item = X;

c->item = Y;

a->item = Z
```

```
Node* a = new Node;
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Node* c = new Node;
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Node* c = new Node;

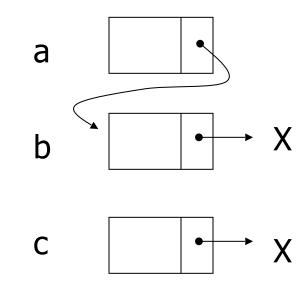
a->next = b;

b->next = c;

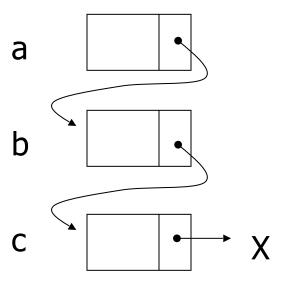
b->item = X;

c->item = Y;

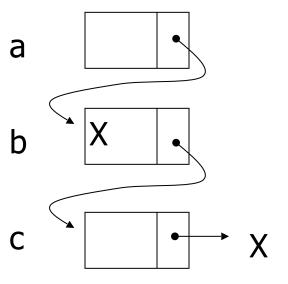
a->item = Z
```



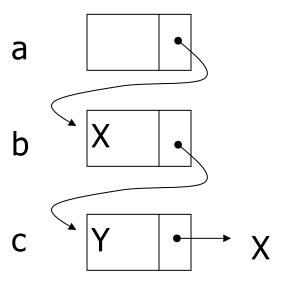
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b->next = c;
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c->item = Y;
a->item = Z
```



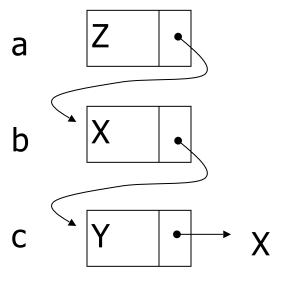
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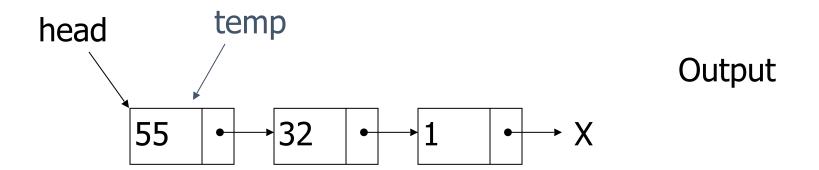
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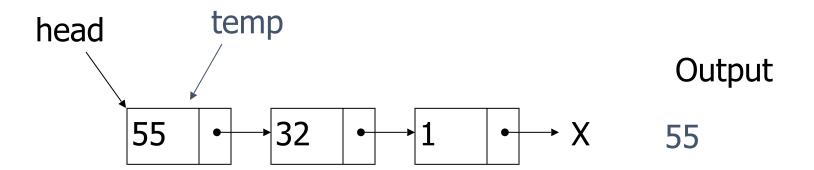


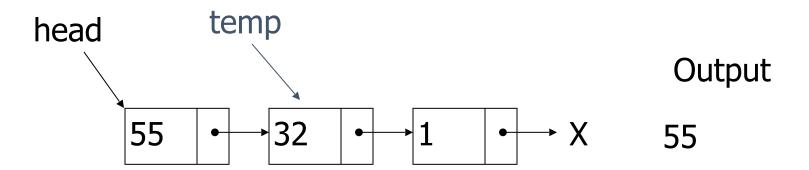
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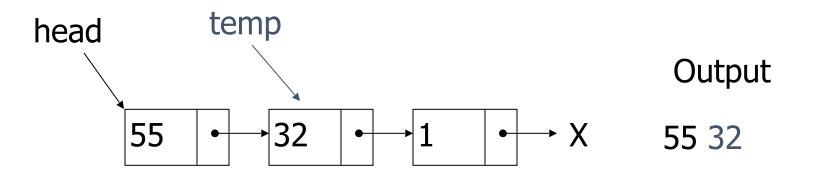


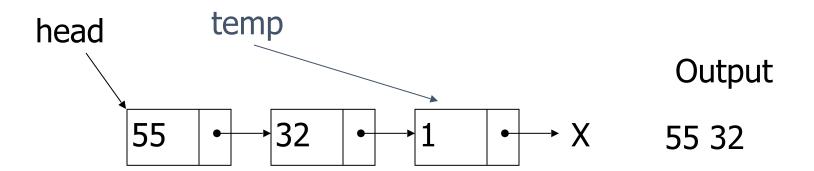
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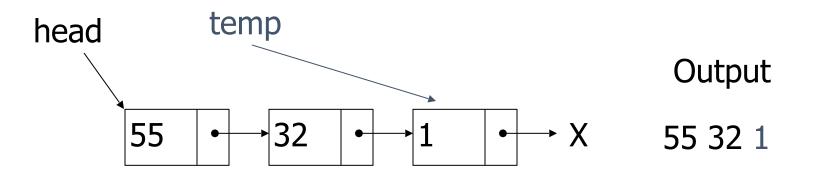




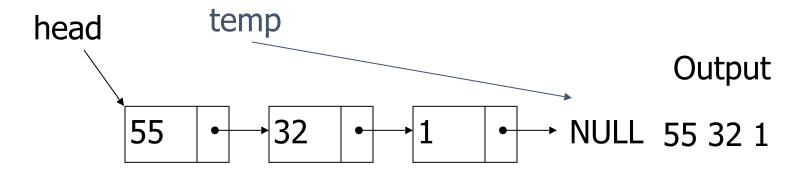




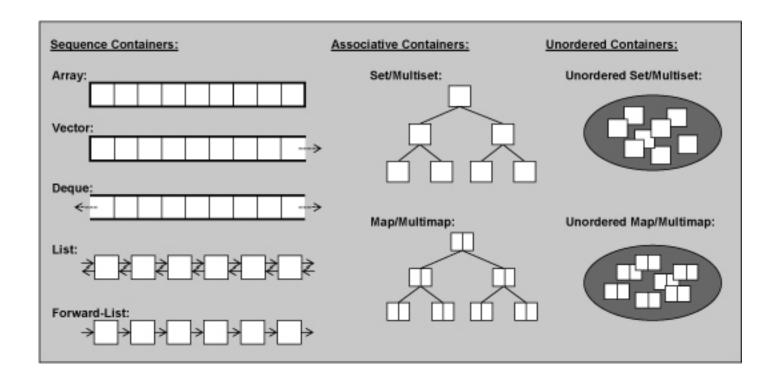




```
Node* temp;
for ( temp = head; temp != NULL; temp = temp->next )
      cout << temp->item << " ";</pre>
```



STL Containers



[&]quot;Container" Classes manage a collection of elements. - N.M Josuttis

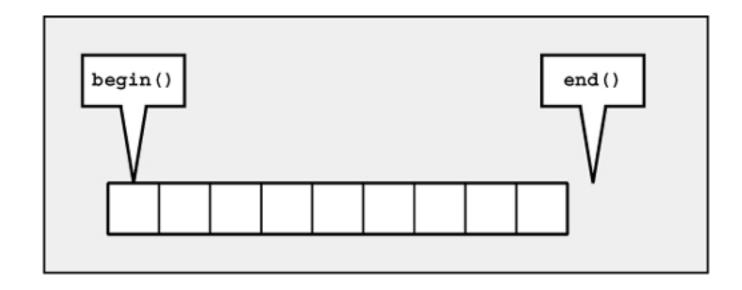
Iterators

- Some operations on lists (container), such as insert and remove require a notion of position.
- http://www.cplusplus.com/reference/stl/
- STL Standard Template Library provides an **iterator** class for representing position in containers.
- Examples...
 - list<string>::iterator it;
 - map<char, int>::iterator it;

Iterators

- iterator begin() returns an appropriate iterator that represents the beginning of elements in the container.
- iterator end() returns the appropriate iterator that represents the end of the elements in container, i.e., position past the last item(valid position).
- itr++ advances the iterator itr to the next position in container.
 (itr--)
- *itr returns the element of the current position of itr. If element has members, use -> to access members.
- itr1 == itr2 returns TRUE if itr1 and itr2 refer to same position.
- itrl != itr2 returns TRUE if itrl and itr2 refer to different positions.

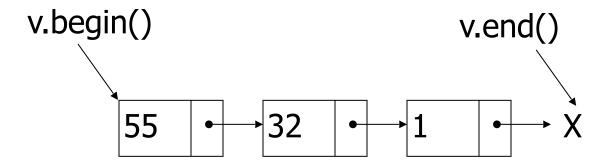
Iterators



The Standard Template Library - A Tutorial and Reference - N.M Josuttis

Printing the list - using an iterator

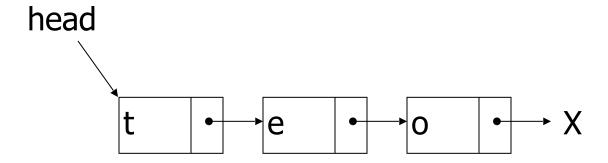
```
list<int>::iterator itr;
for ( itr = v.begin(); itr != v.end(); ++itr) {
     cout << *itr << " ";</pre>
```

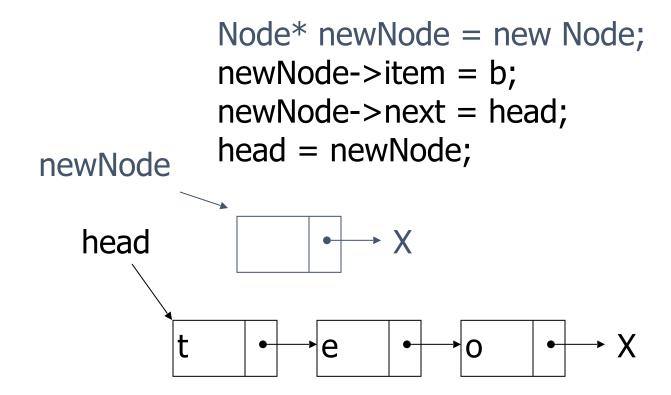


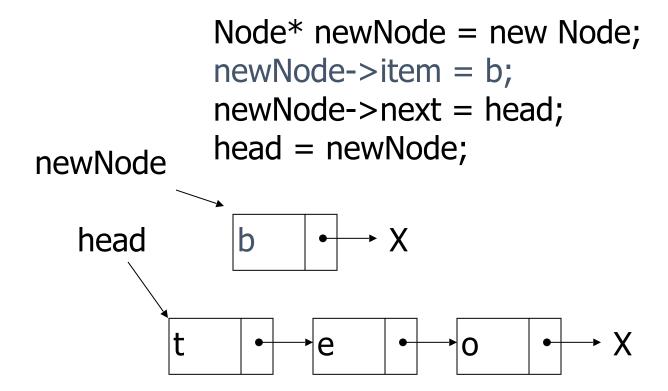
Linked List Code - Get Last

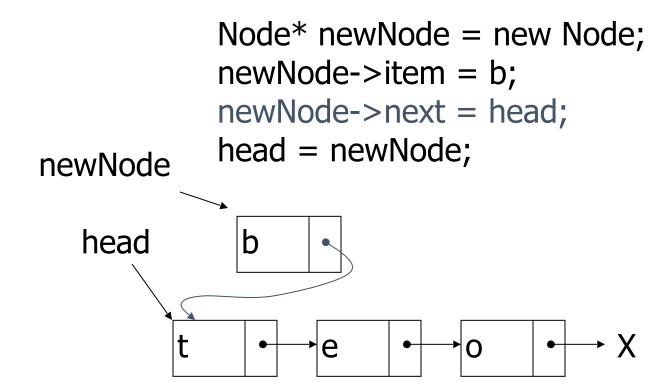
• (Stop here) Quick exercise - write the code to get a pointer to the last node in a list

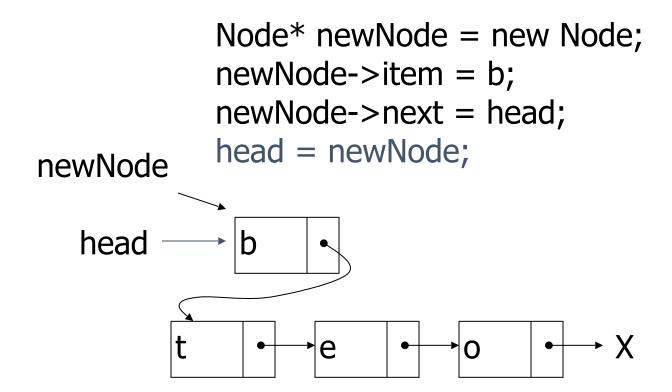
```
Node* newNode = new Node;
newNode->item = b;
newNode->next = head;
head = newNode;
```



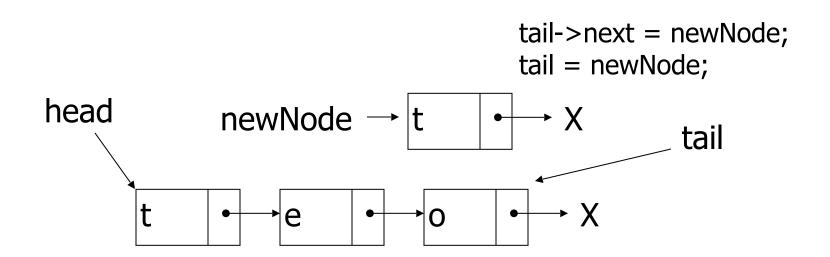




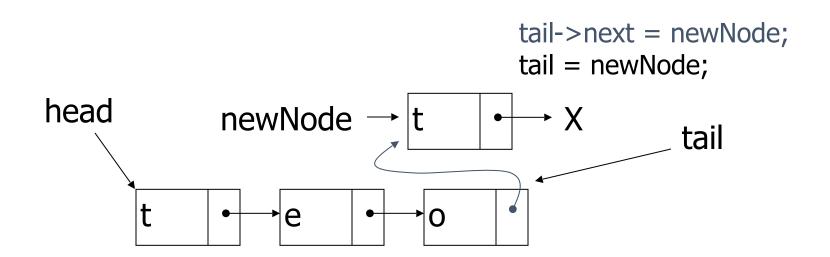




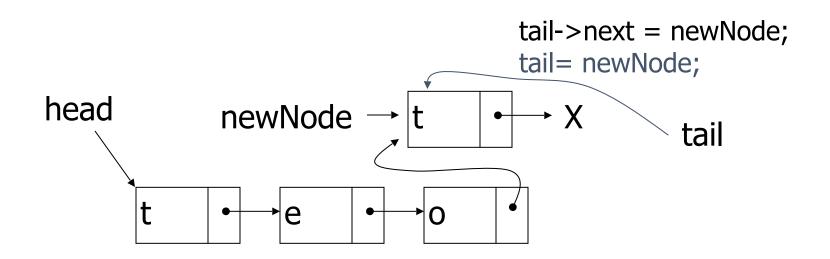
- Inserting at tail/end
 - For now, we assume that we have a pointer to the last node called tail.
 - Assume newNode is created and has its value set.



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Linked List Code - Search

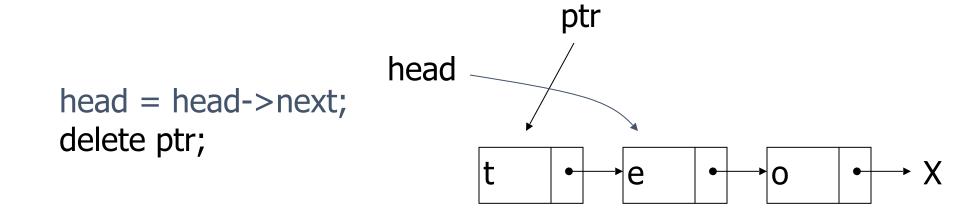
 Search - returns a pointer to the node containing the key or NULL if the key doesn't exit

```
Node* search ( itemtype key ) {
    Node* temp;
    for ( temp = head; temp != NULL; temp = temp->next ) {
        if ( temp->item == key )
            return temp;
    }
    return NULL;
}
```

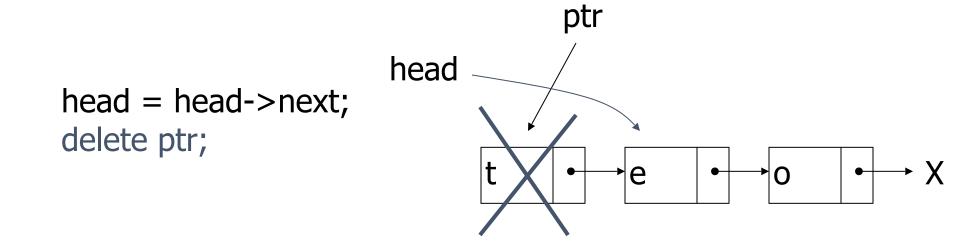
- Removing a node assume ptr points to the node we want to remove
- Remove first node:

head = head->next; delete ptr; head thead = head->next; delete ptr;

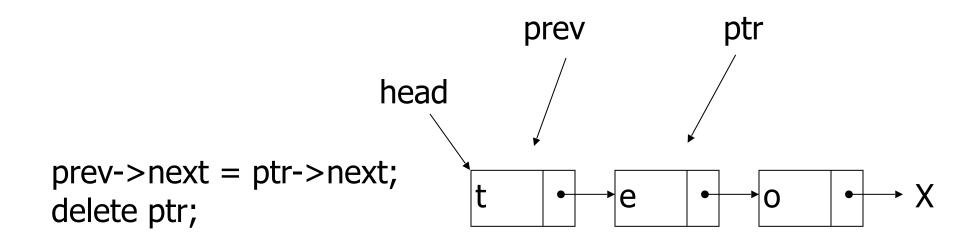
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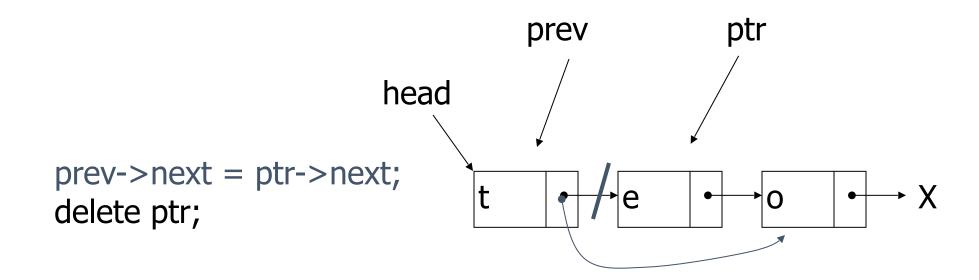
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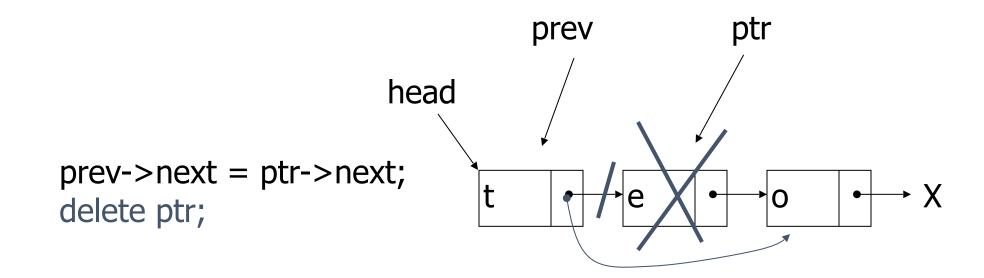
- Removing a node assume ptr points to the node we want to remove
- Remove middle node assume *prev* pointer:



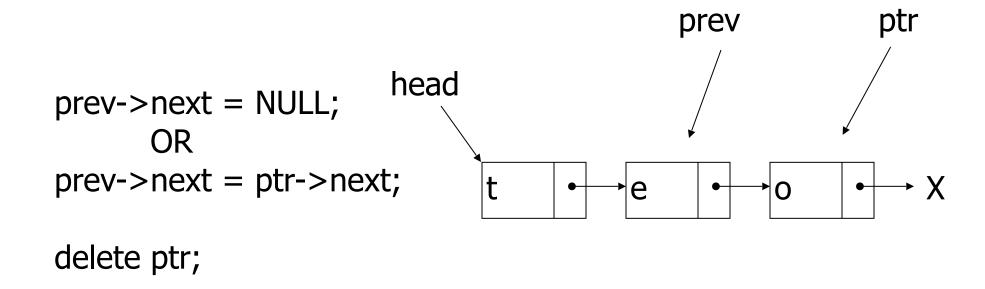
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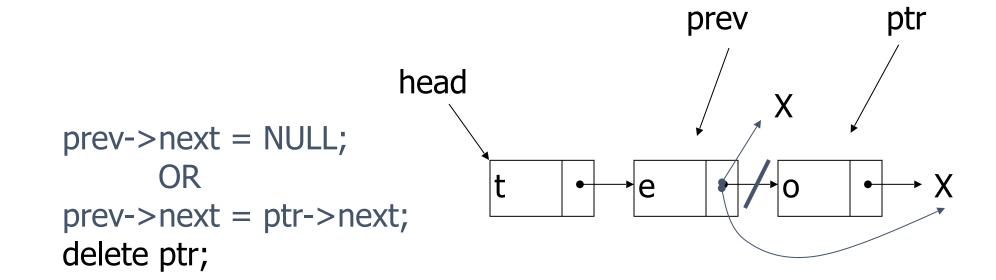
- Removing a node assume ptr points to the node we want to remove
- Remove middle node assume *prev* pointer:



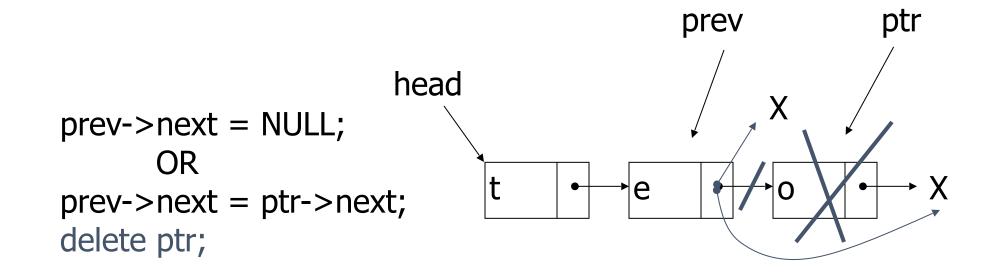
- Removing a node assume ptr points to the node we want to remove
- Remove last node assume *prev* pointer



- Removing a node assume ptr points to the node we want to remove
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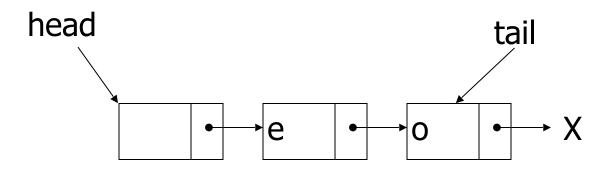
- Removing a node assume ptr points to the node we want to remove
- Remove last node assume *prev* pointer:



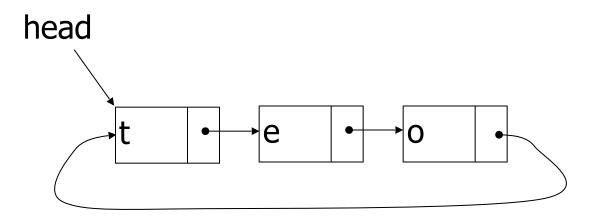
Linked List Code - Destructor

Every new needs to have a corresponding delete.

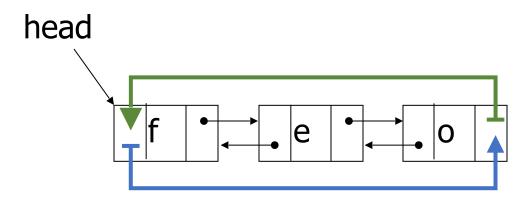
- Keep a head and tail pointer to make tail inserts faster
 - More pointer upkeep on inserts and removes
- Use a sentinel/dummy node
 - Easier removes because there is only one case



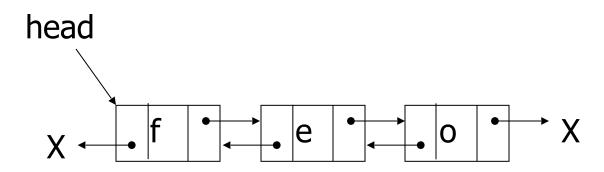
- Circularly linked list
 - No special cases for head or tail



- Doubly-linked list
- Still have to find previous node for inserts and removes
- Solved with a doubly linked list



- Doubly-linked list
- Still have to find previous node for inserts and removes
- Solved with a doubly linked list



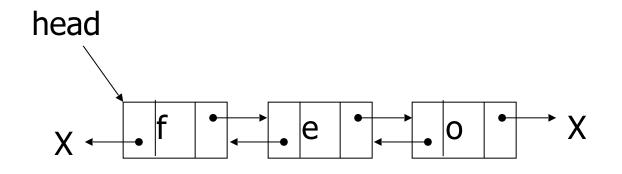
Doubly Linked Lists

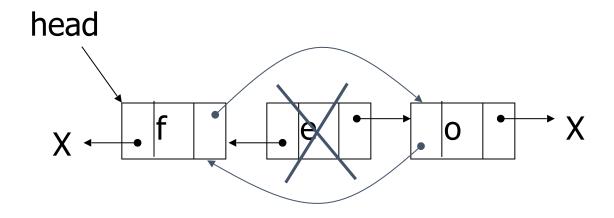
- Required adjustments for doubly linked list:
 - Size increase per node
 - extra *prev* pointer kept in each node
 - Slight increase in time to do inserts and removes
 - more pointer manipulations
 - Increases pointer complexity
 - If you don't like pointers

Doubly Linked List Implementation

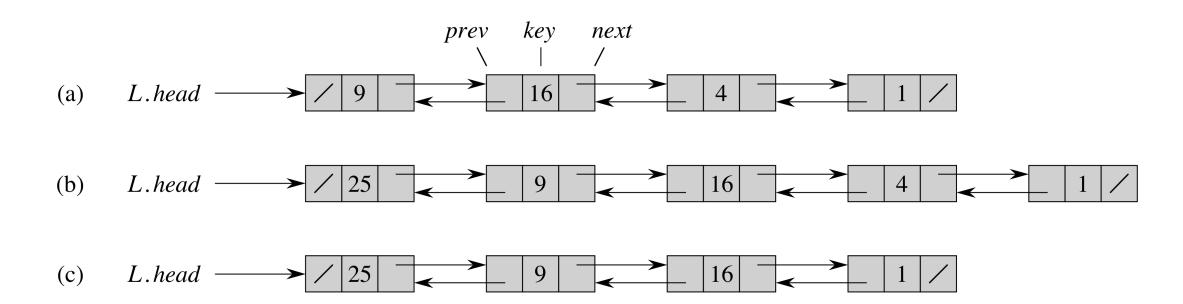
```
class List {
private:
                                    class Node {
       Node* head;
                                    friend class List;
public:
       List ( ) {
                                    private:
              head = NULL;
                                           itemtype Item;
                                           Node* next;
       // Member functions to
                                           Node* prev;
       // interface with list
                                     };
};
```

Doubly linked list - Remove





Doubly Linked List



Doubly Linked List Code - Remove

- In class exercise write remove for a doubly linked list
- Assumptions:
 - Only have key and pointer to head
 - Don't assume item exists
 - Remove first instance

Doubly Linked List Code - Remove

Linked List Programming Notes

- Memory leaks
 - Every insert invokes a new, be sure to have a complimentary delete for each new. Since every node will not necessarily be removed during the program, the destructor should go through and delete all the nodes still left.
- Segmentation faults
 - Be sure to compare all pointers to NULL before accessing any information

Linked List Programming Notes

- Deal with all possible cases
 - Check for empty list
 - Deal with first, middle, and last node cases if they require different code
- Most importantly....be generic.
 - Don't take type of item to store into consideration
 - Plan to reuse code

Linked List Running Times

- Assuming only a head pointer and a singly linked list
 - Insert head
 - O(1)
 - Insert tail
 - O(N)
 - Search
 - O(N)

Remove

Num items

Print