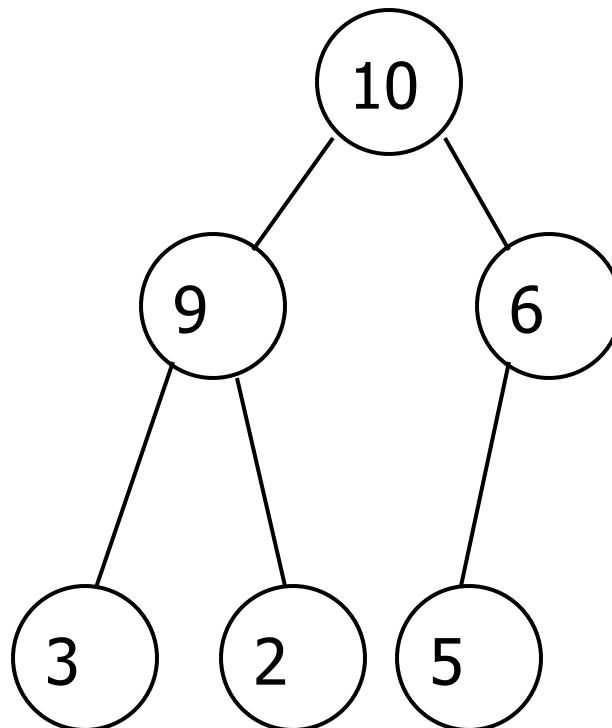


# (Binary) Heaps

- Similar to a Binary Search Tree (BST)
- Heap is sorted in a weaker sense than BST
- Tree is a complete tree
  - Bottom level may not be filled but fills from left to right

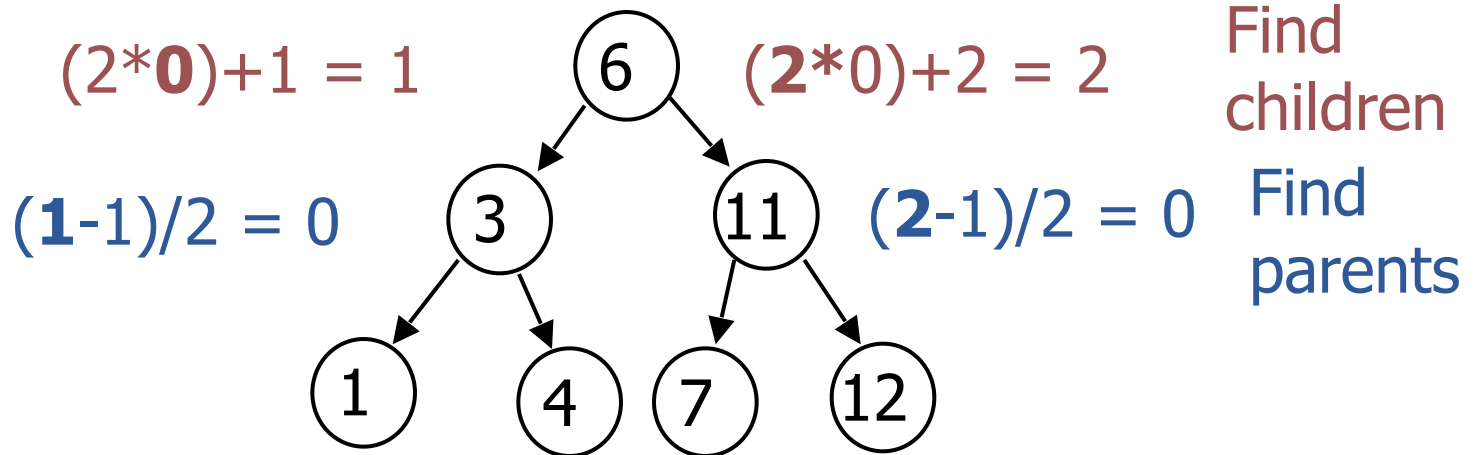


# (Binary) Heaps

- Binary trees are easily stored in an array
  - Array with index started at 0
    - Parent =  $(i-1)/2$
    - Left Child =  $2i + 1$
    - Right Child =  $2i + 2$
  - Array with index started at 1
    - Parent =  $i/2$
    - Left Child =  $2i$
    - Right Child =  $2i + 1$

# Binary Tree – Array

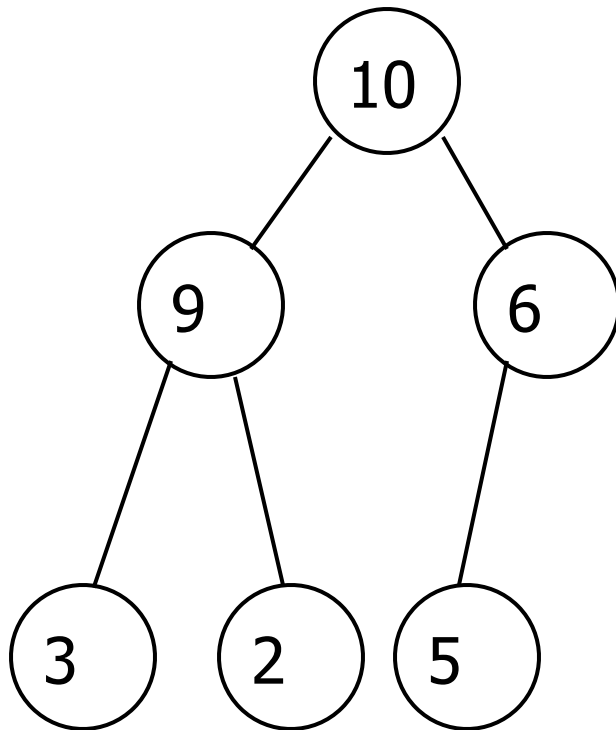
0	1	2	3	4	5	6	7
6	3	11	1	4	7	12	



# Heap Order Properties

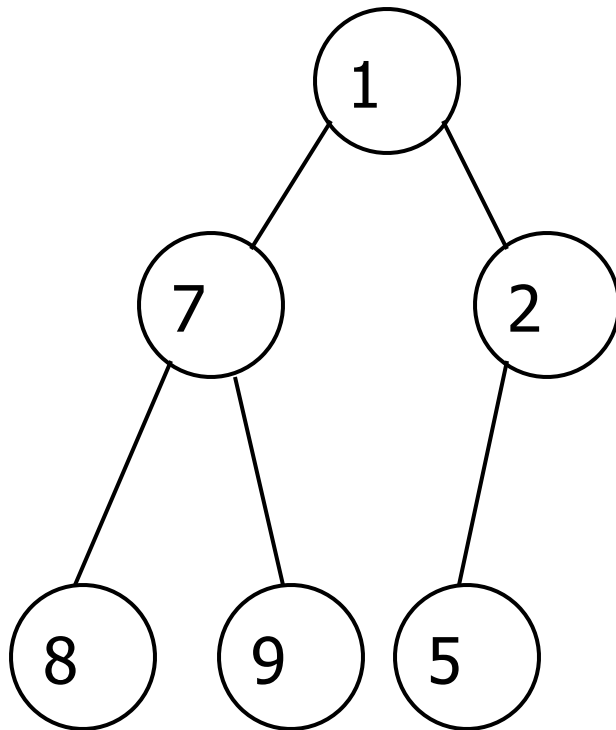
- **Min Heap**
  - For each node  $X$ ,  $X.\text{key} > (X.\text{parent}).\text{key}$ 
    - Smallest number is at the root
- **Max Heap**
  - For each node  $X$ ,  $X.\text{key} < (X.\text{parent}).\text{key}$ 
    - Largest number is at the root

# Max Heap



10
9
6
3
2
5

# Min Heap



1
7
2
8
9
5

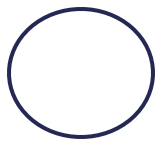
# Heap - Insert

- Create a hole (empty node) in the next available complete tree location
  - Remember to fill bottom layer from left to right
- If the item can be inserted into the hole without violation of the heap property, insert item
- Otherwise, copy the hole's parent item into the hole then trickle up

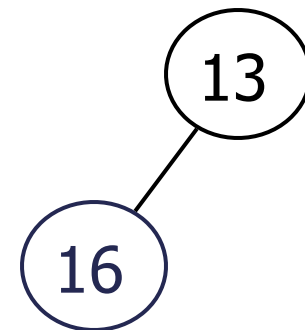
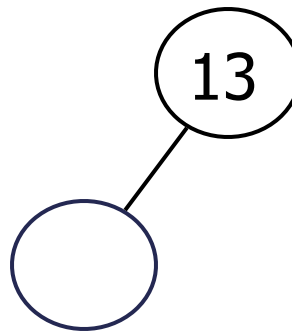
# Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

Insert 13



Insert 16

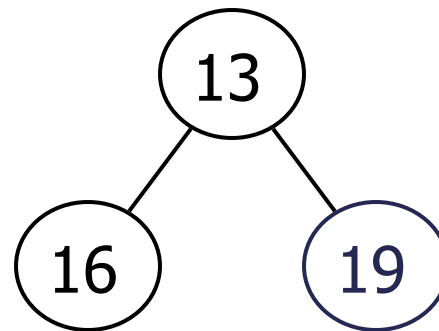
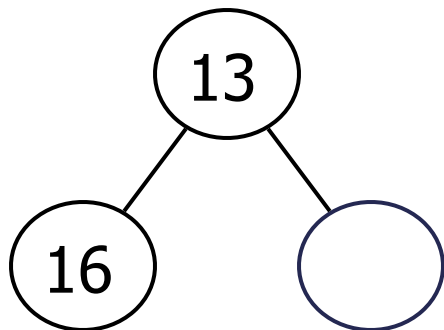




# Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

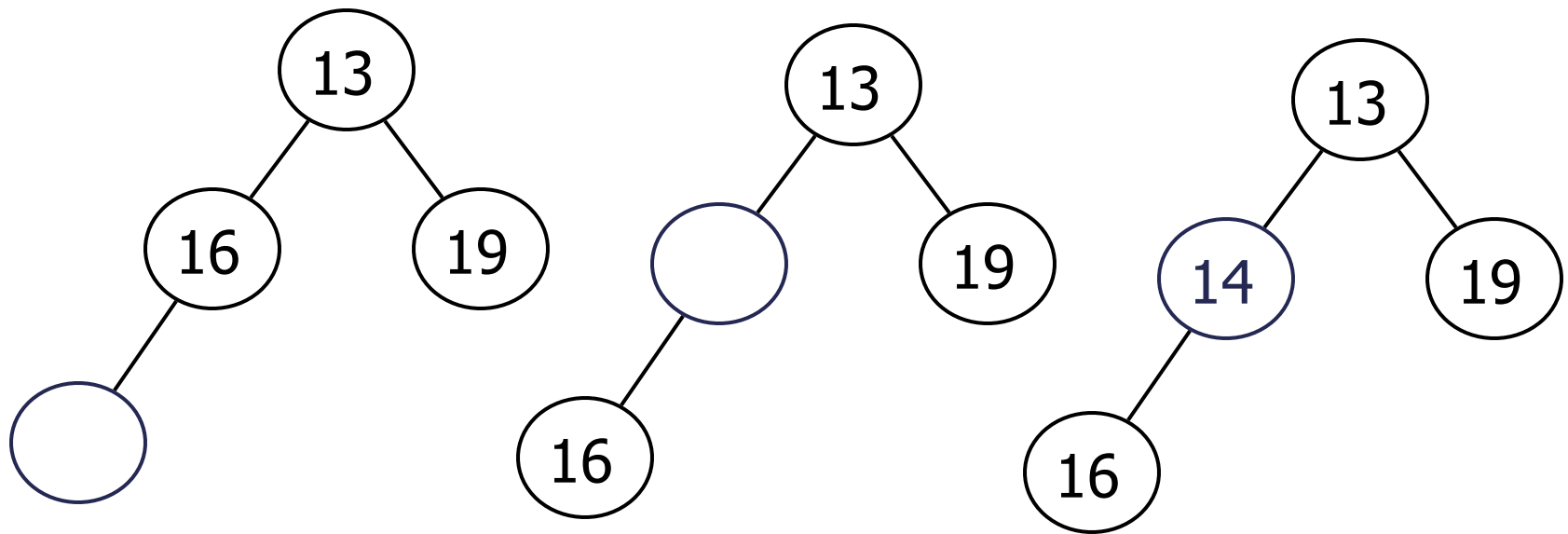
Insert 19



# Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

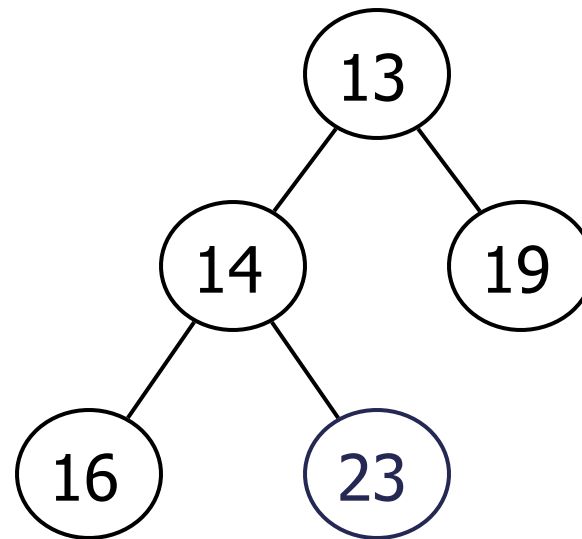
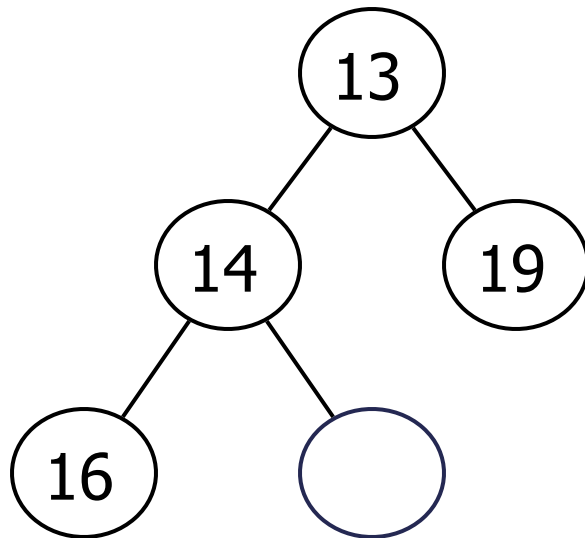
Insert 14



# Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

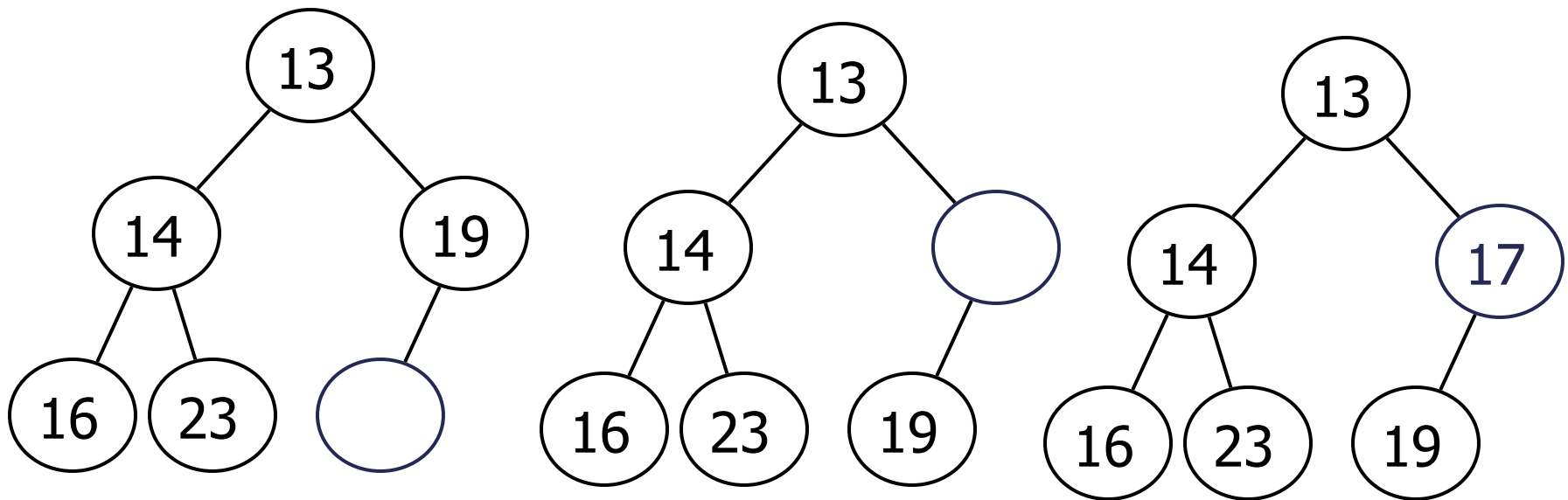
Insert 23



# Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

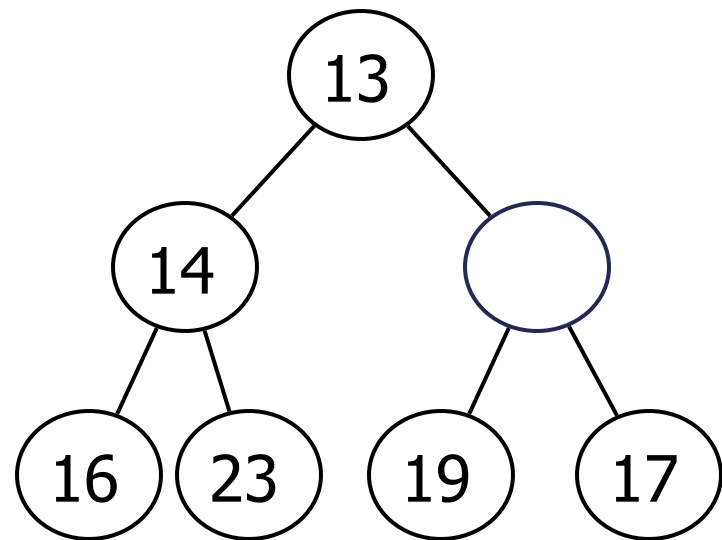
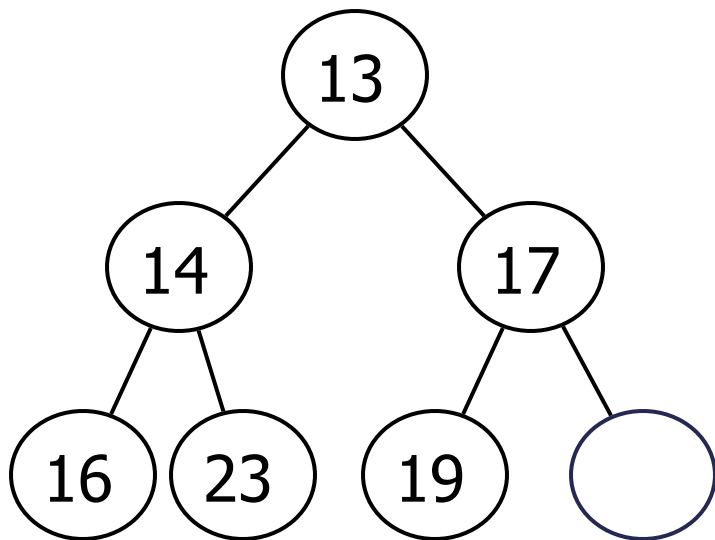
Insert 17



# Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

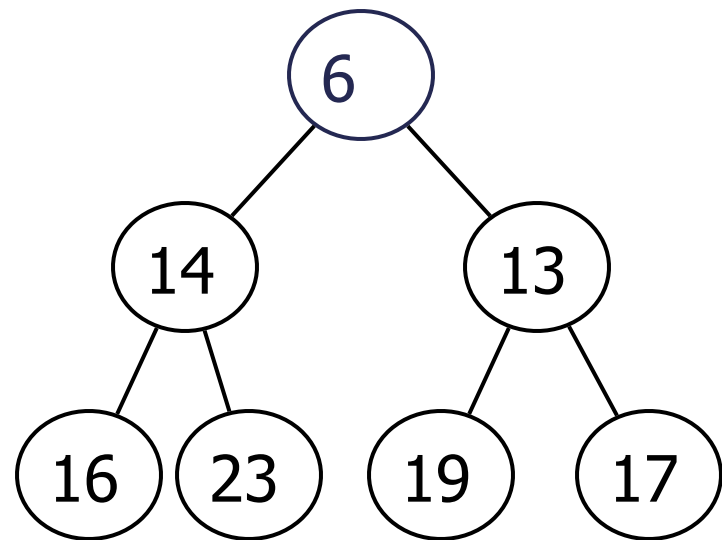
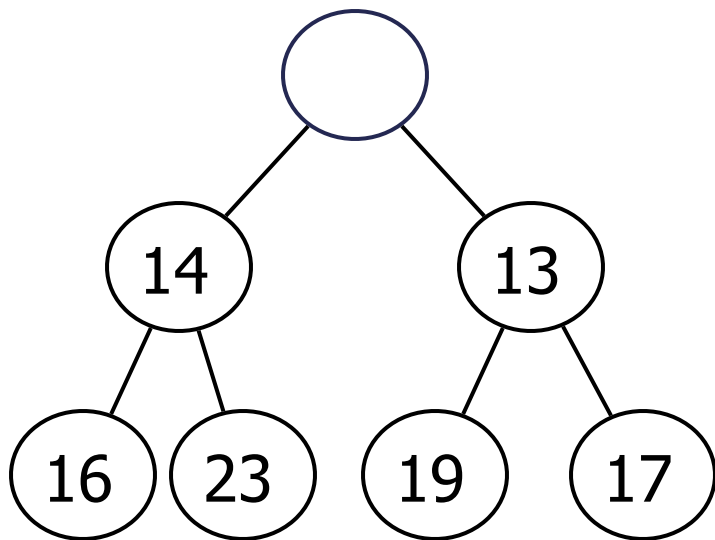
Insert 6



# Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

Insert 6



# Min Heap – Extract Min

6, 14, 13, 16, 23, 19, 17

6,14,13,16 23,19,17

13,14,17,16,23,19

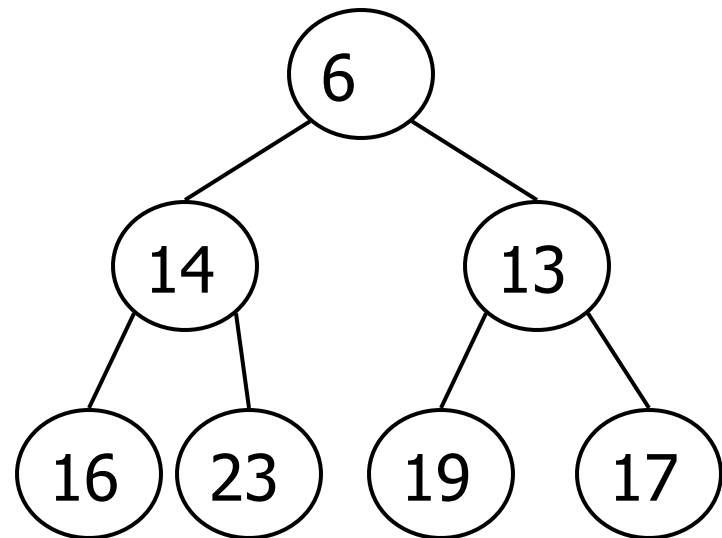
14,16,17,19,23

16,19,17,23

17,19,23

19,23

23



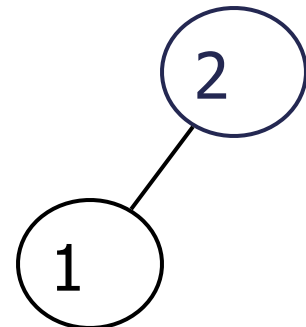
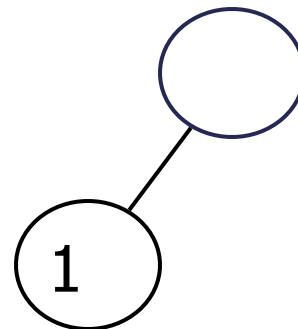
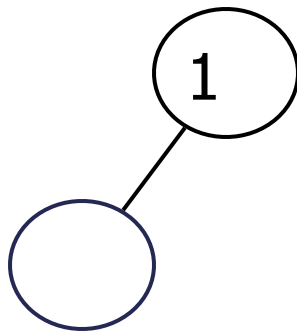
# Max Heap - Insert

1, 2, 3, 4, 5, 6

Insert 1



Insert 2

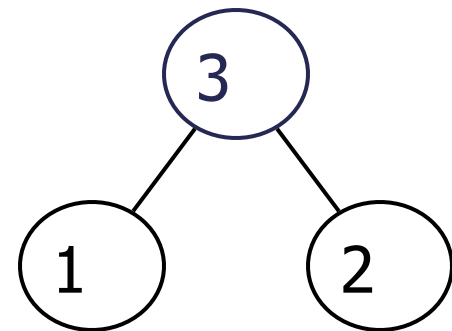
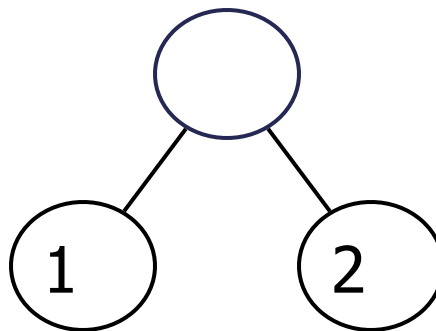
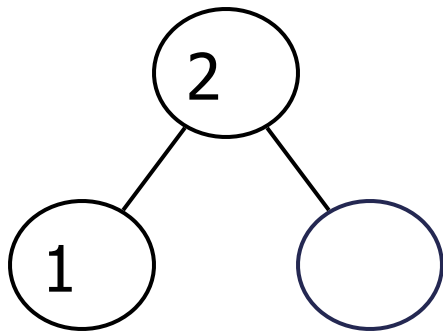




# Max Heap - Insert

1, 2, 3, 4, 5, 6

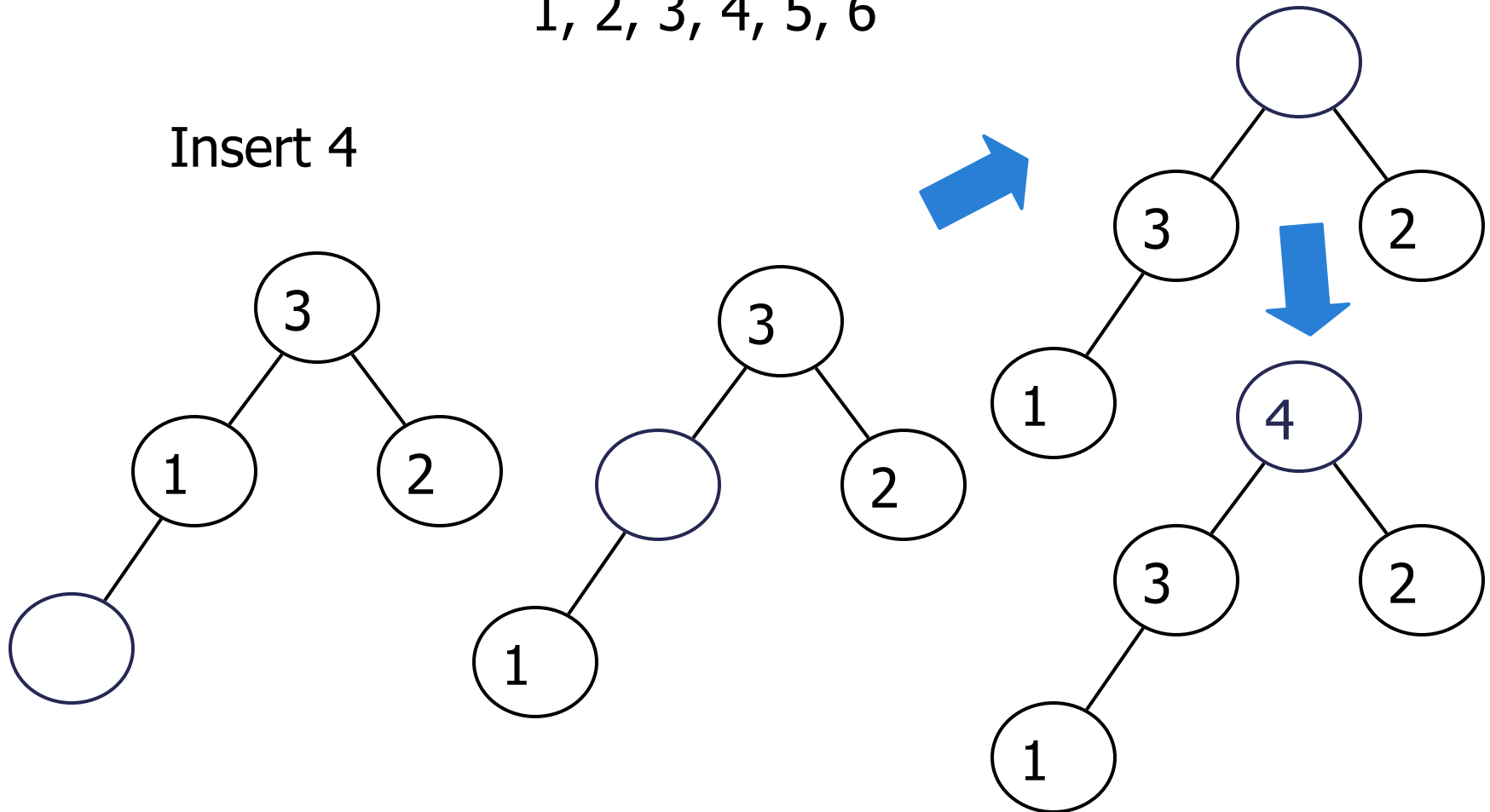
Insert 3



# Max Heap - Insert

1, 2, 3, 4, 5, 6

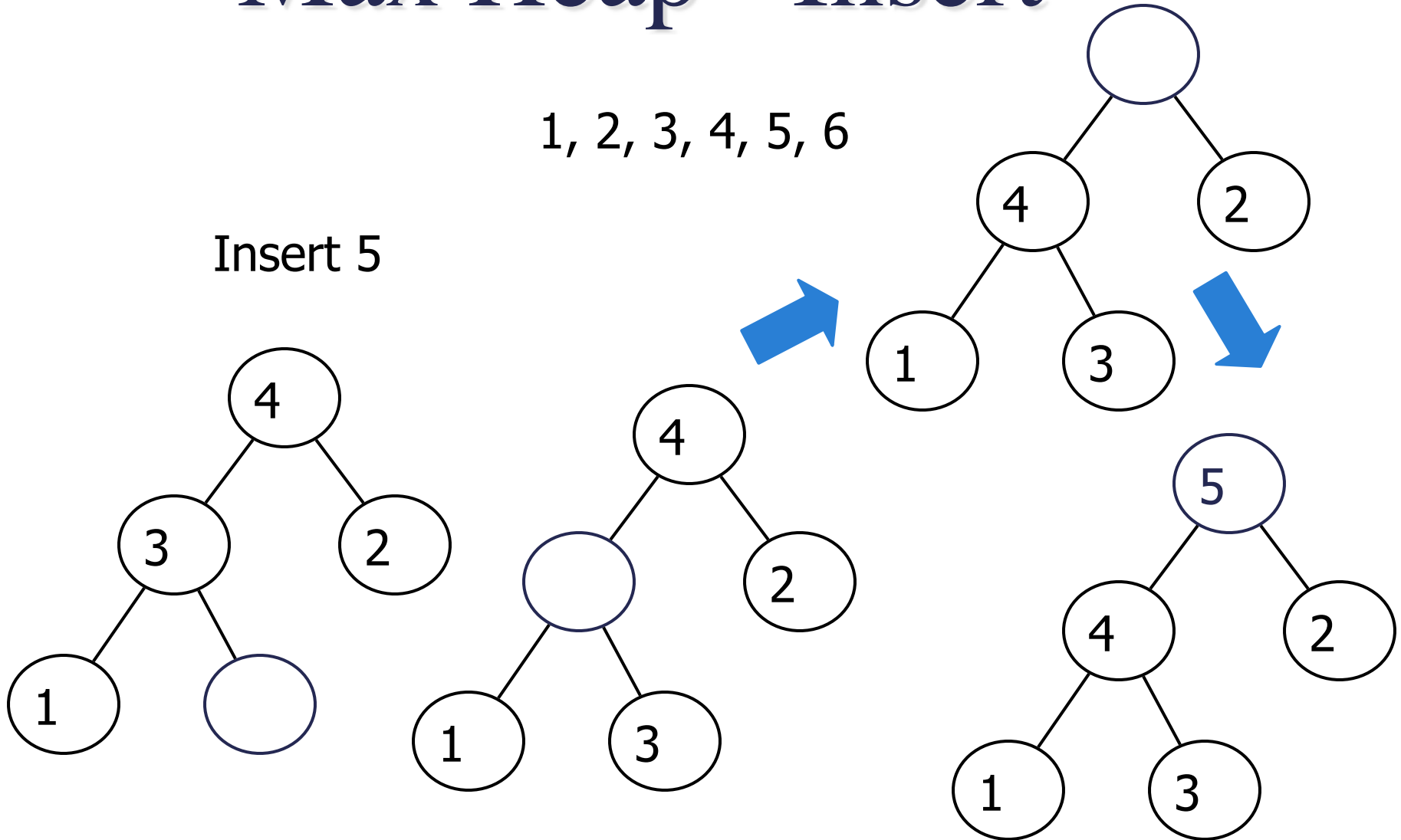
Insert 4



# Max Heap - Insert

1, 2, 3, 4, 5, 6

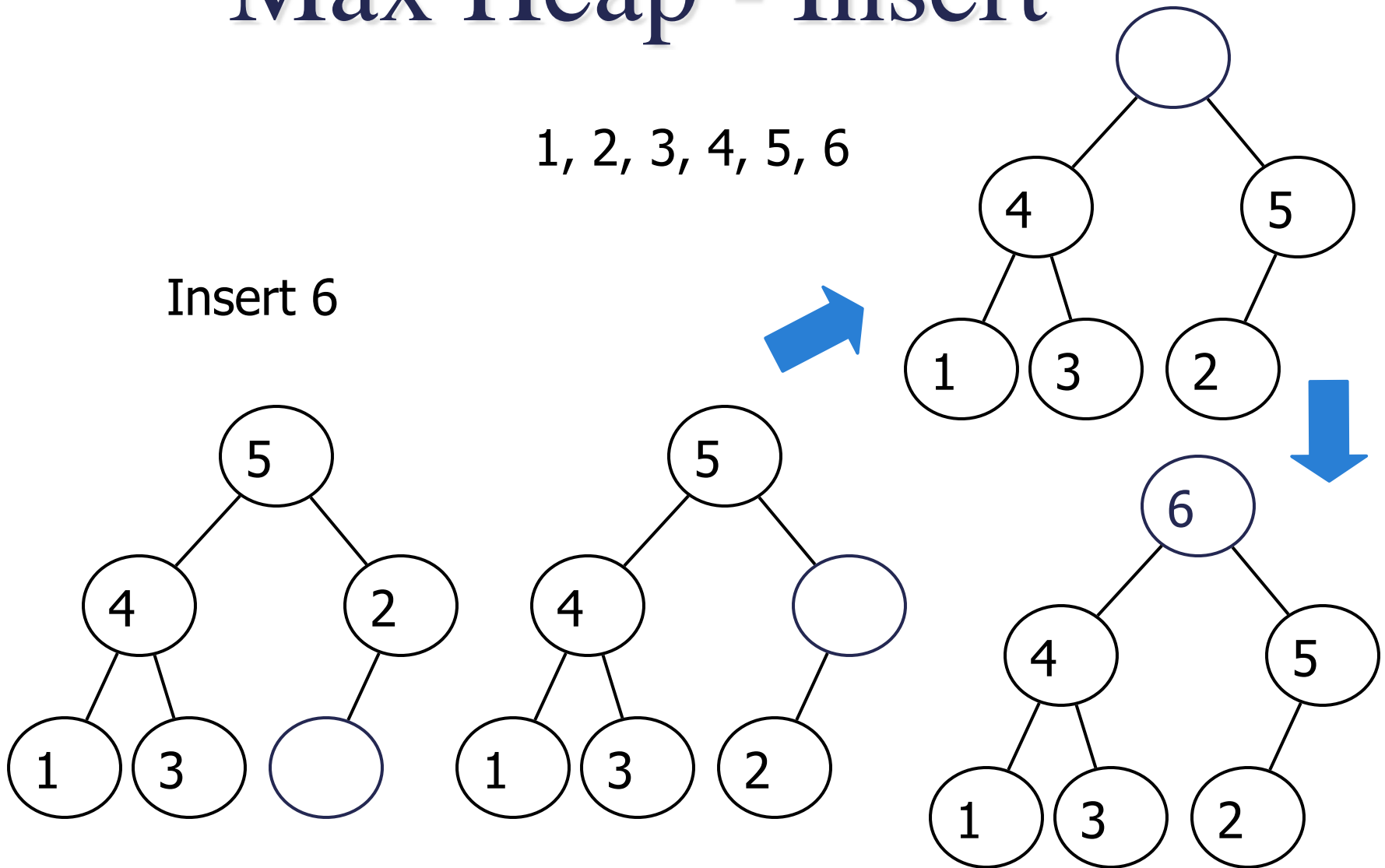
Insert 5



# Max Heap - Insert

1, 2, 3, 4, 5, 6

Insert 6



# Min Heap Insert

Running time = $O(\log n)$
----------------------------

```
void Min-Heap-Insert ( itemtype item )  
    if heap full  
        throw heap full exception  
    // Trickle up  
    x = numNodes  
    while (x > 0 && arr[(x-1)/2] > item)  
        arr[x] = arr[(x-1)/2]  
        x = (x-1)/2  
    arr[x] = item  
    numNodes++
```

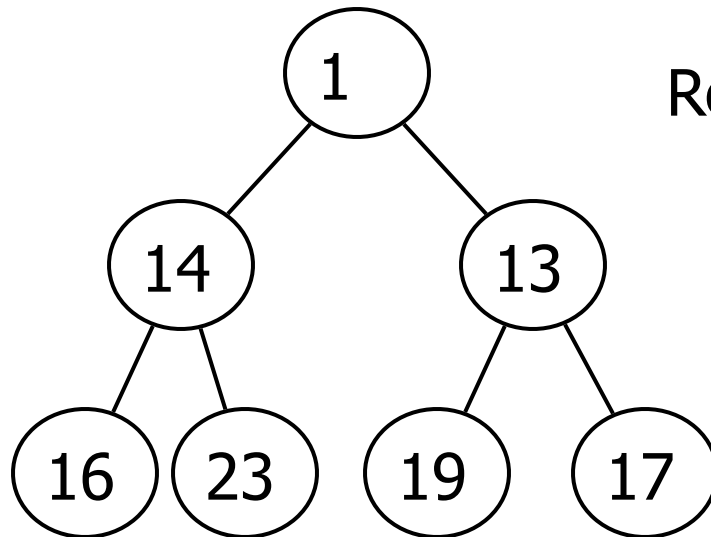
# Heap Find Min/Max

- Finding the node with the highest priority is easy - it is just at the root
  - Running time =  $O(1)$

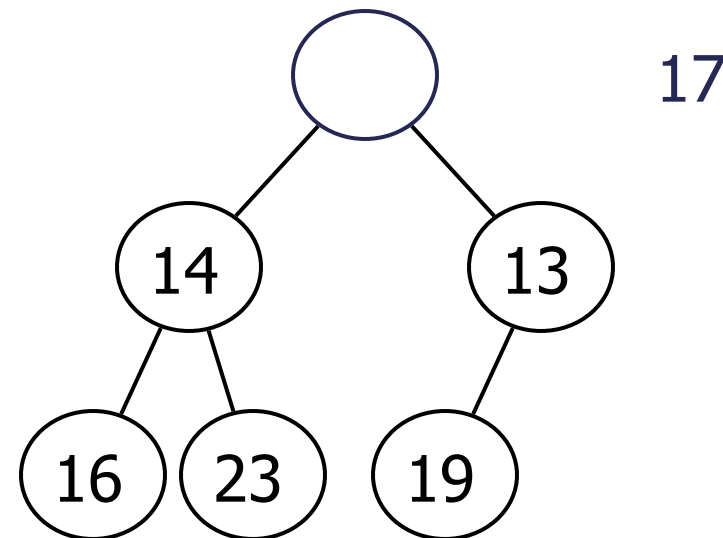
# Heap – Remove (Extract Min or Max)

- Remove highest priority item
  - Makes a hole at the root
  - Want to remain a complete tree, so attempt to place last item in the heap into the hole
    - If item can be placed in hole without violation of the heap property, then done
    - Otherwise, trickle down
    - Pick the child with the highest priority

# Min Heap – Remove

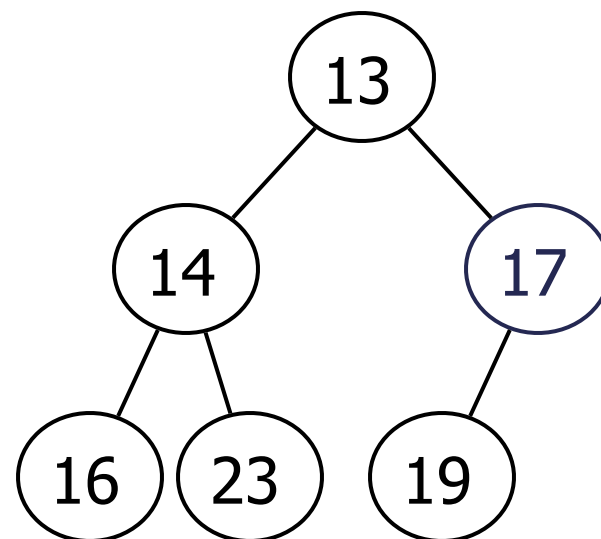
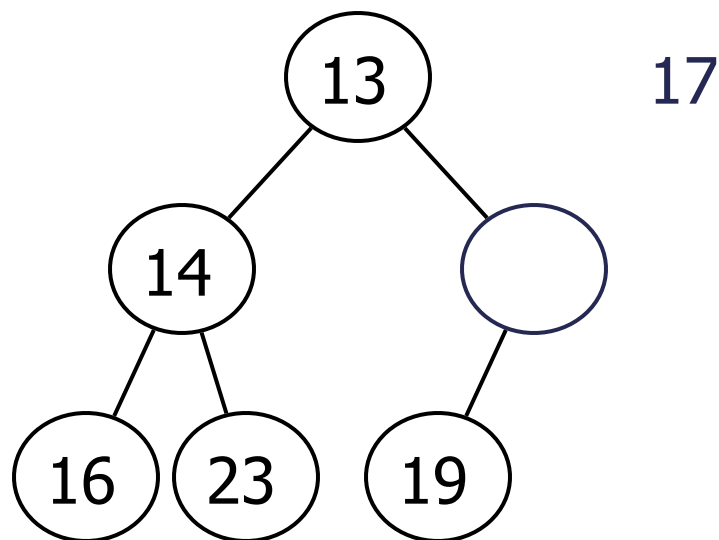


Remove 1

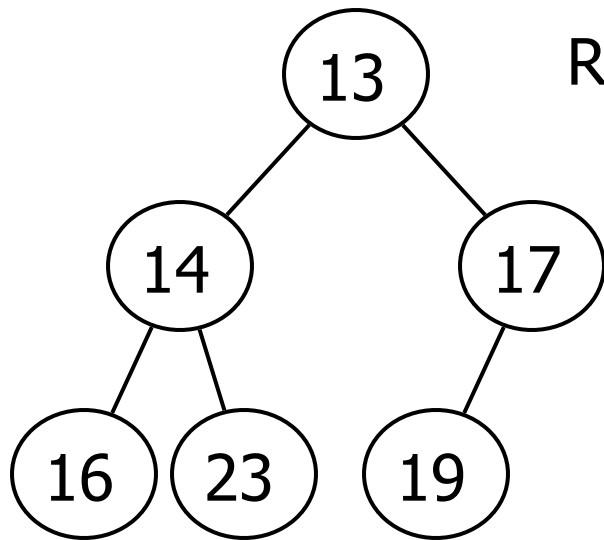




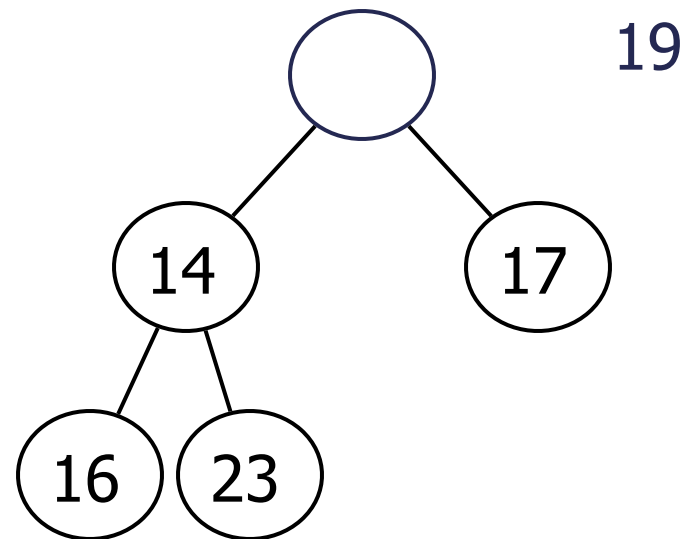
# Min Heap - Remove



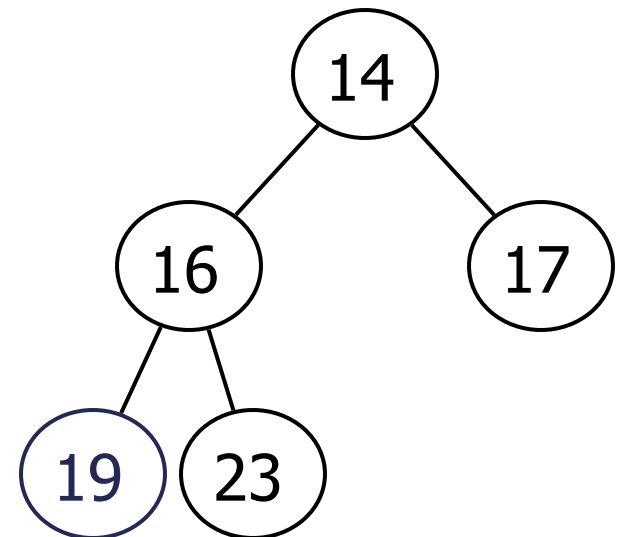
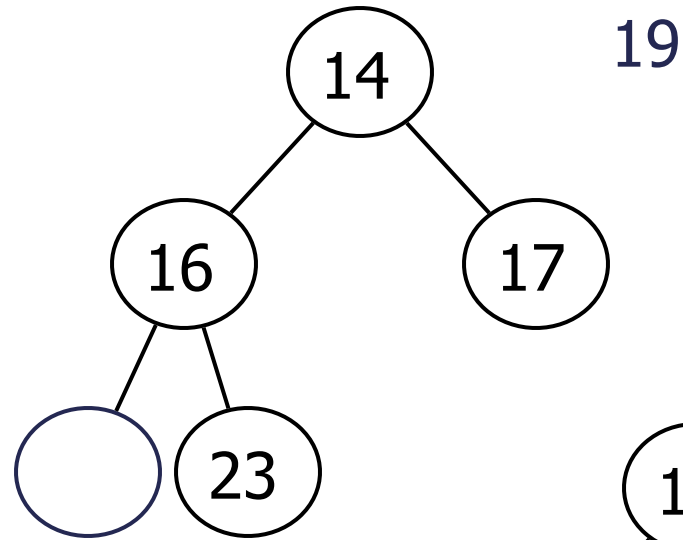
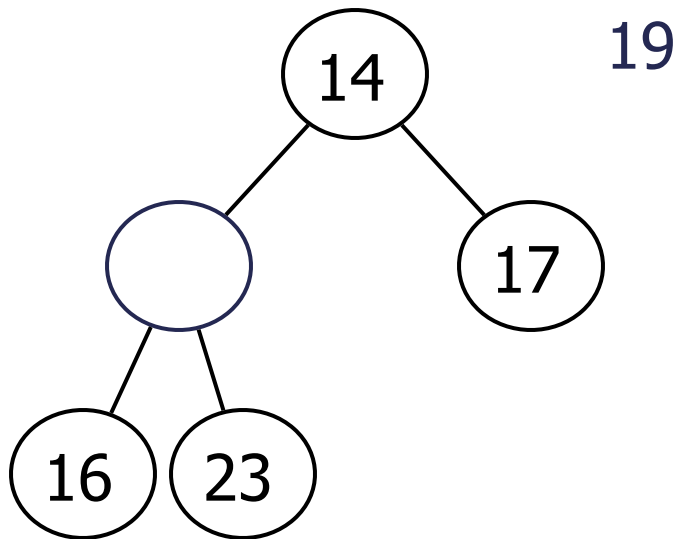
# Min Heap - Remove



Remove 13

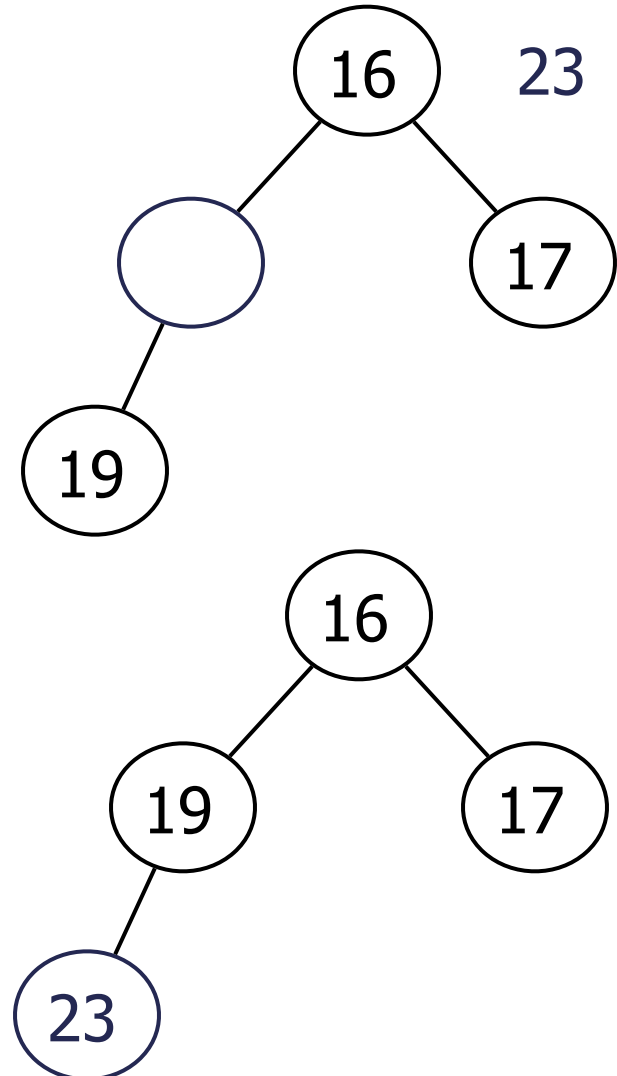
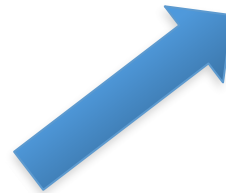
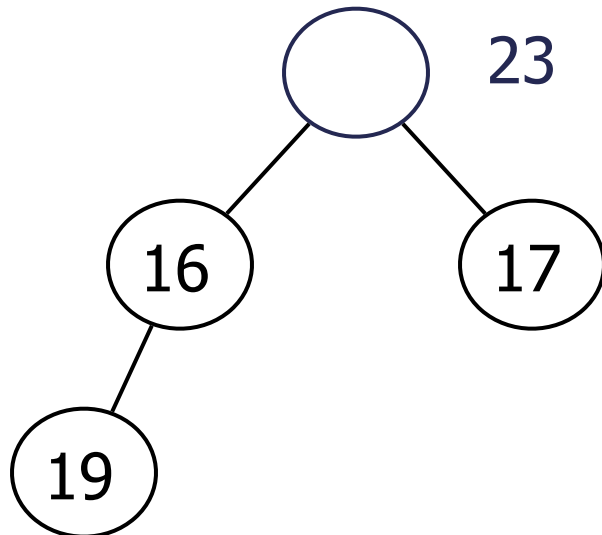
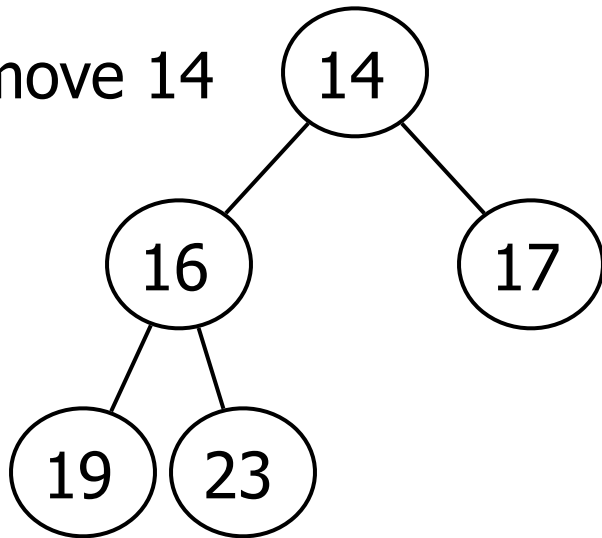


# Min Heap - Remove



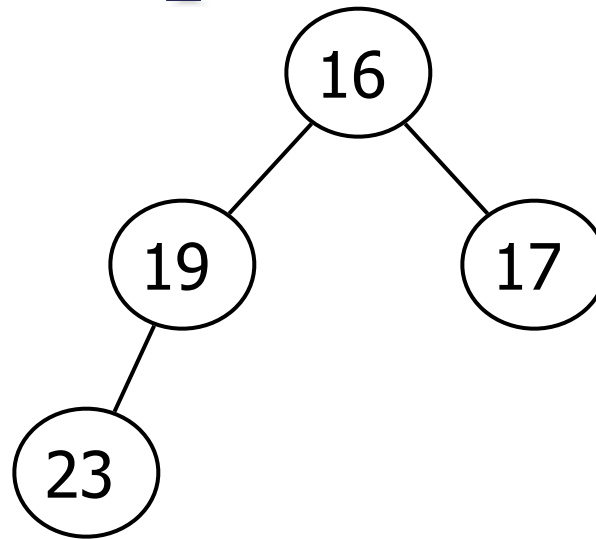
# Min Heap - Remove

Remove 14

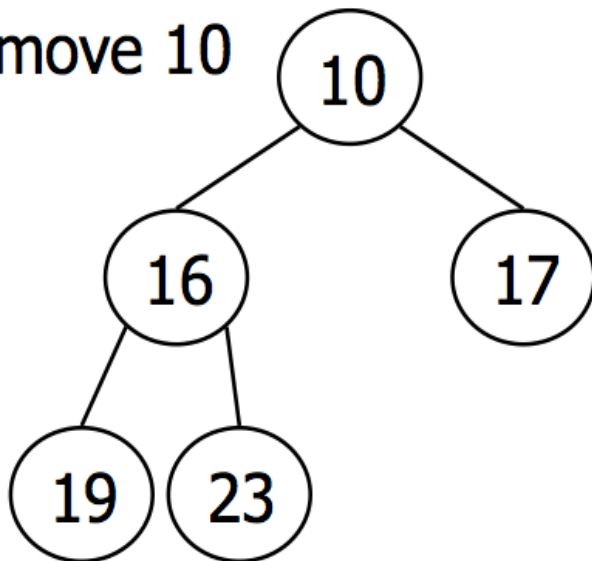


# Exam7 Q5

## Min Heap – Insert



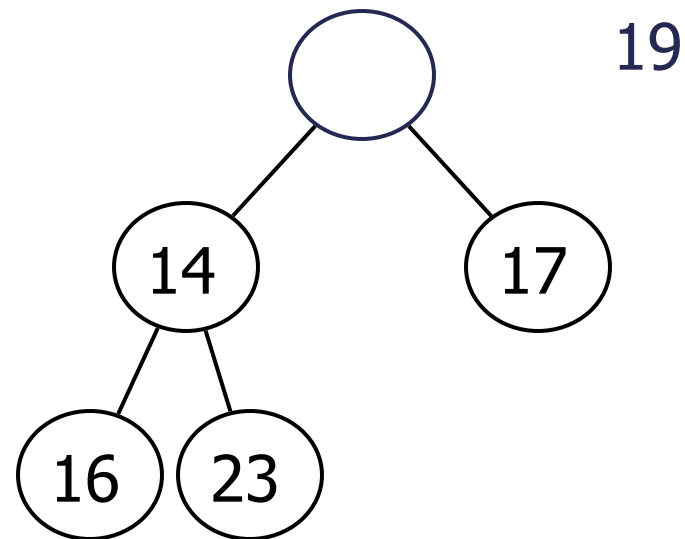
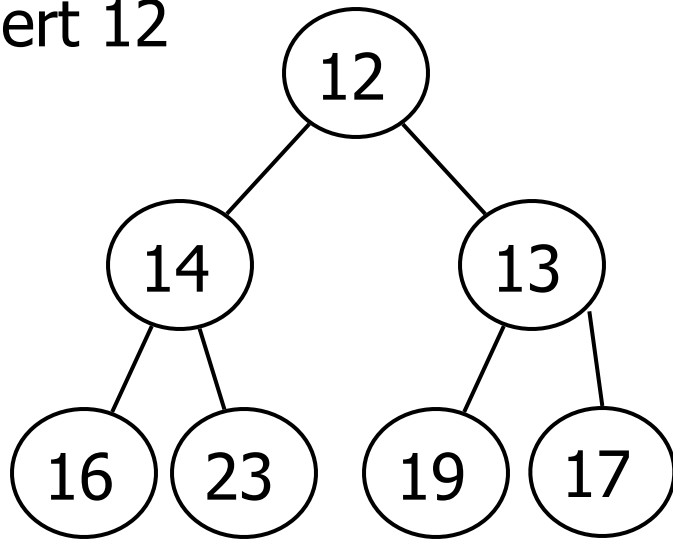
Remove 10



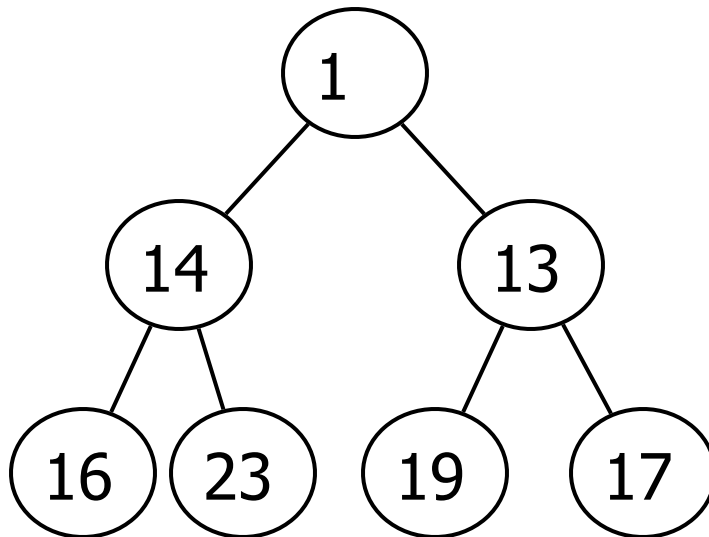
# Exam7 Q5

## Min Heap – Insert

Insert 12



# Min Heap - Remove



1,14,13,16,23,19,17  
13,14,17,16,23,19  
14,16,17,19,23  
16,19,17,23  
17,19,23  
19,23  
23

# Building a Heap

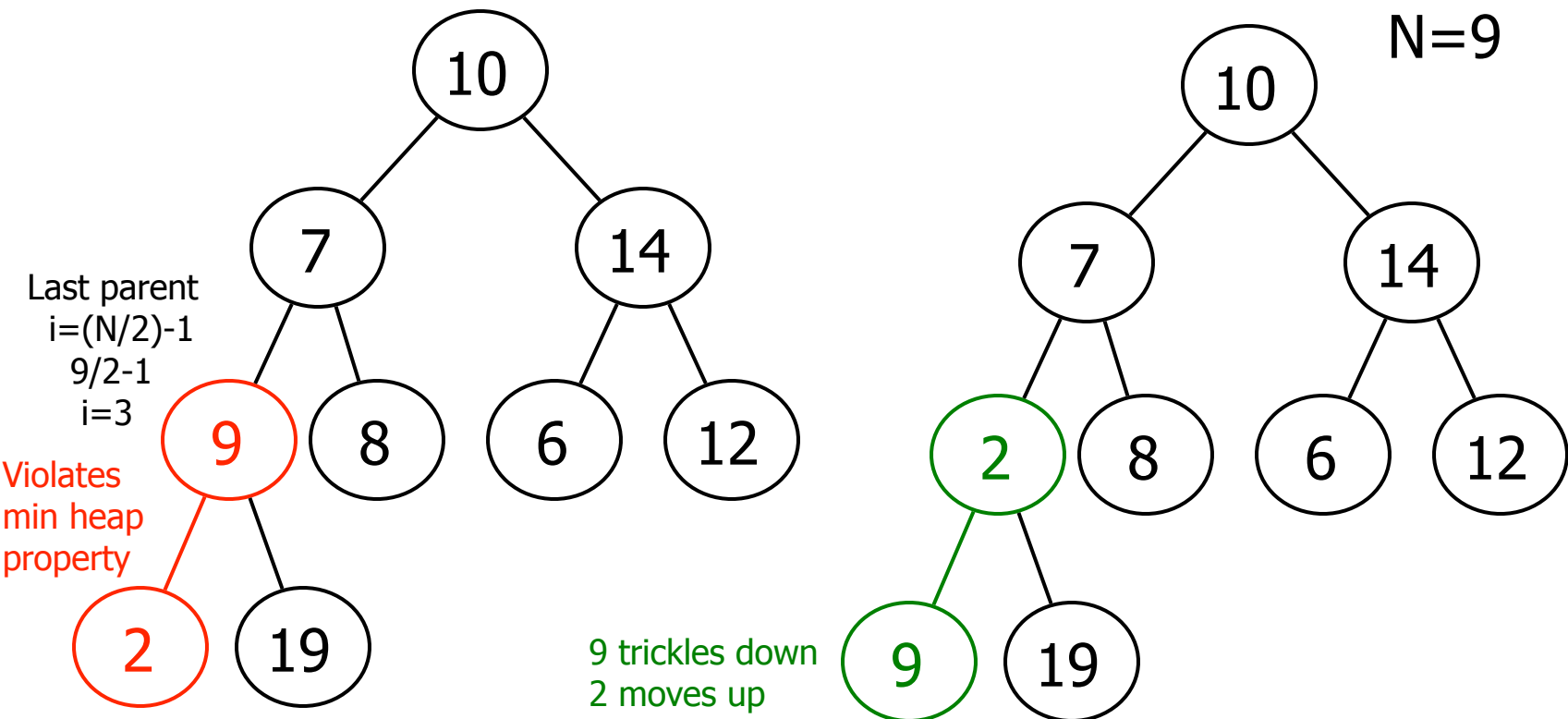
- Naïve method
  - Perform  $N$  insert operations -  $O(n \log n)$
- Better solution
  - Given an array of unordered items:

```
for (int i=(N/2)-1; i >= 0; i--)  
    Trickle down
```



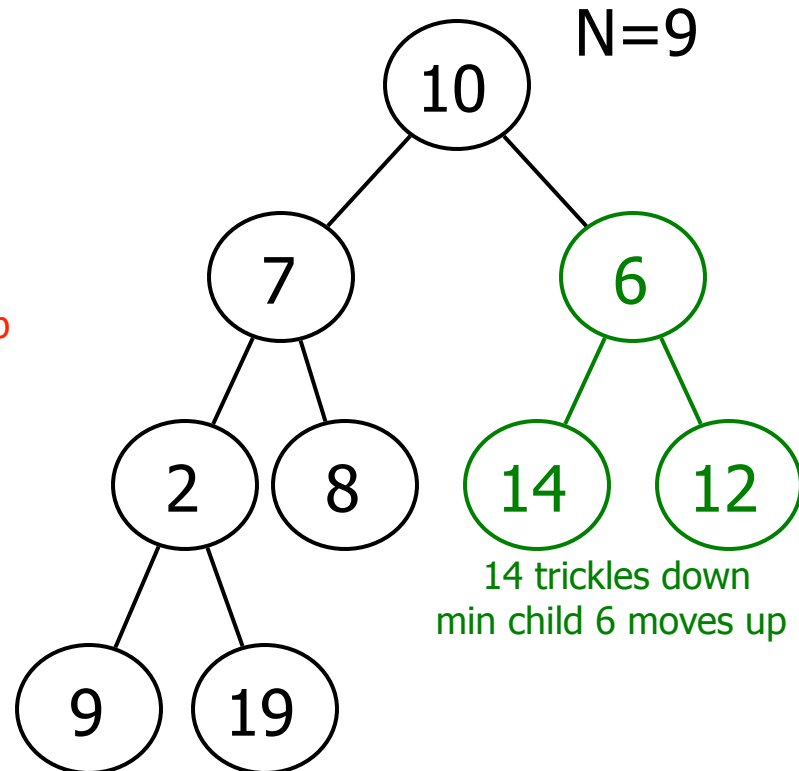
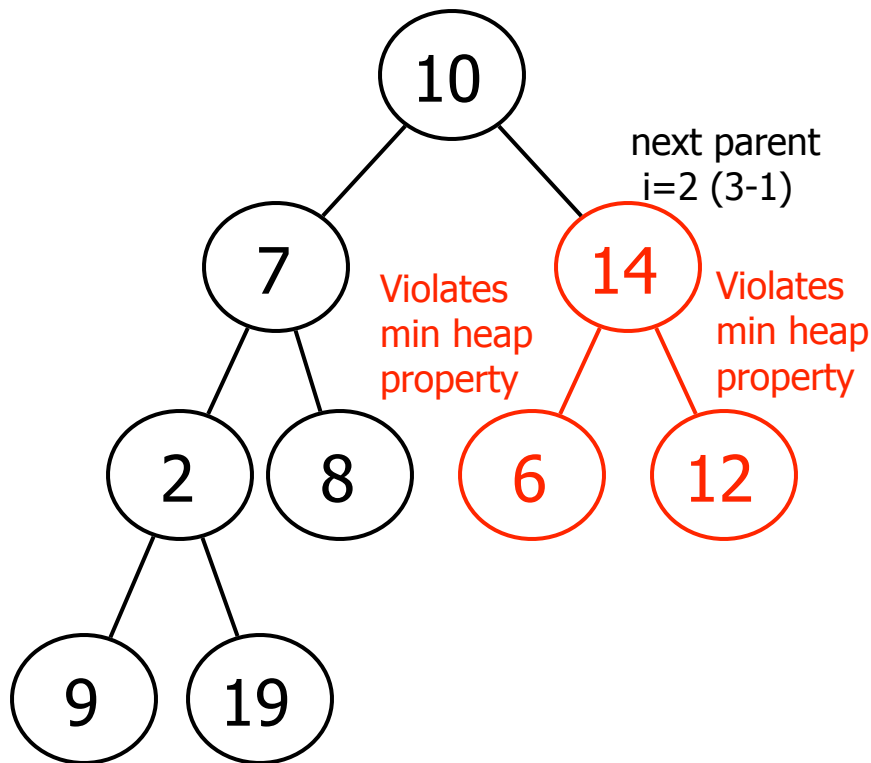
# Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
10	7	14	9	8	6	12	2	19



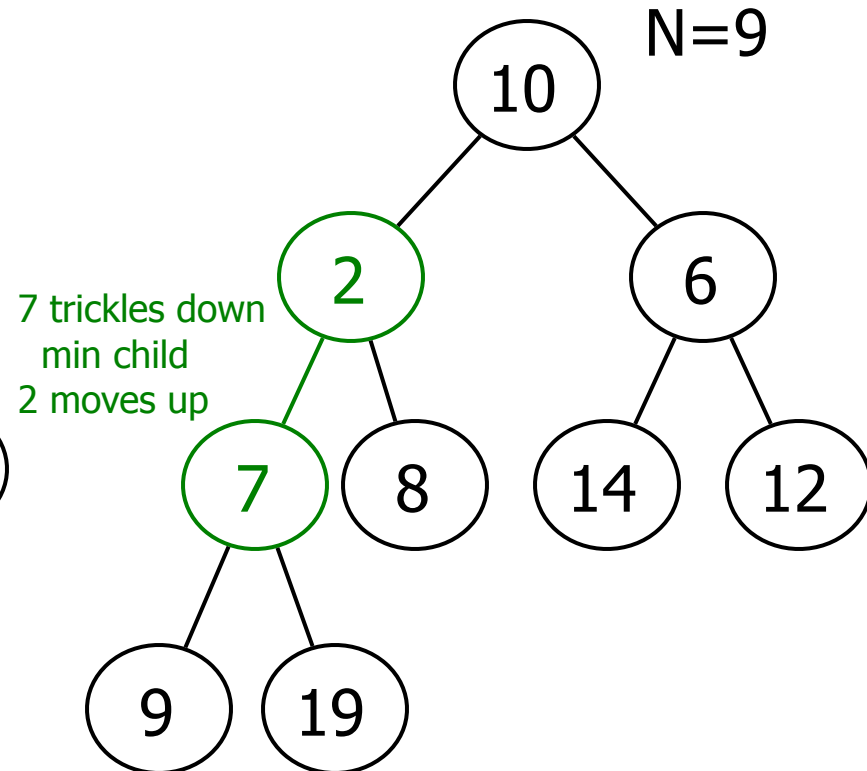
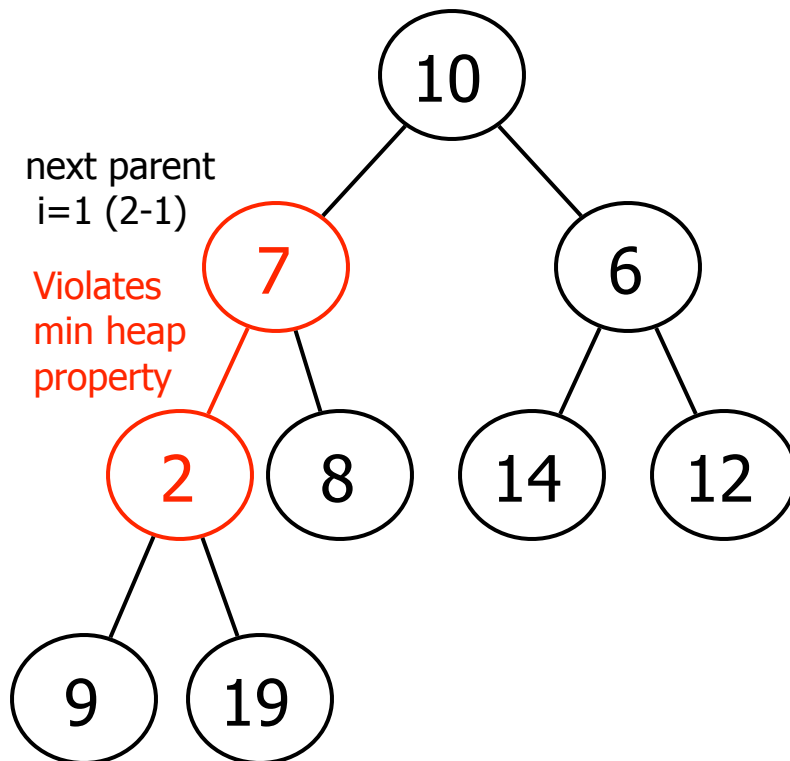
# Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
10	7	14	2	8	6	12	9	19



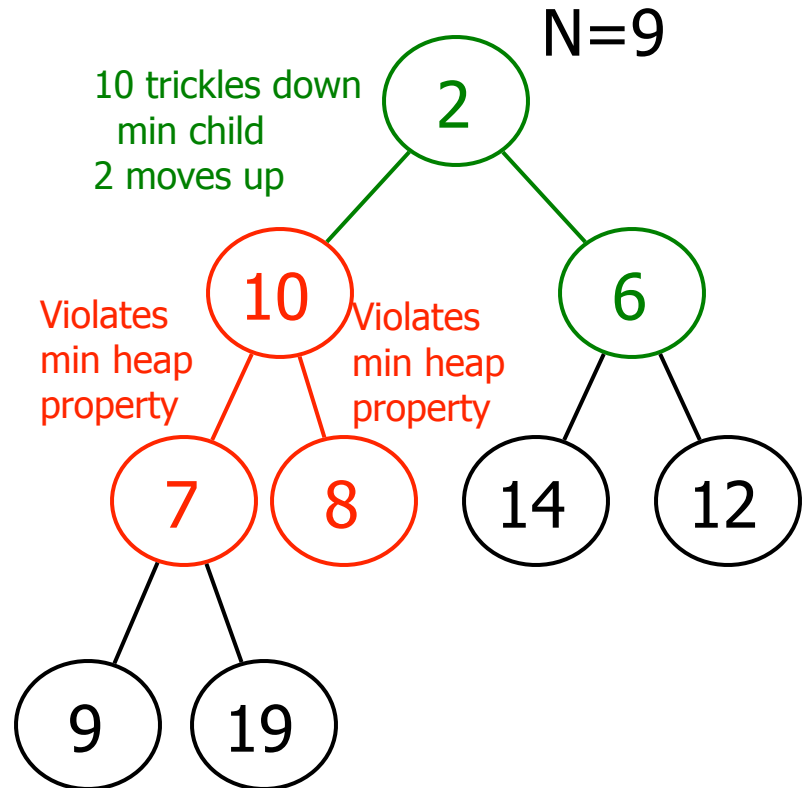
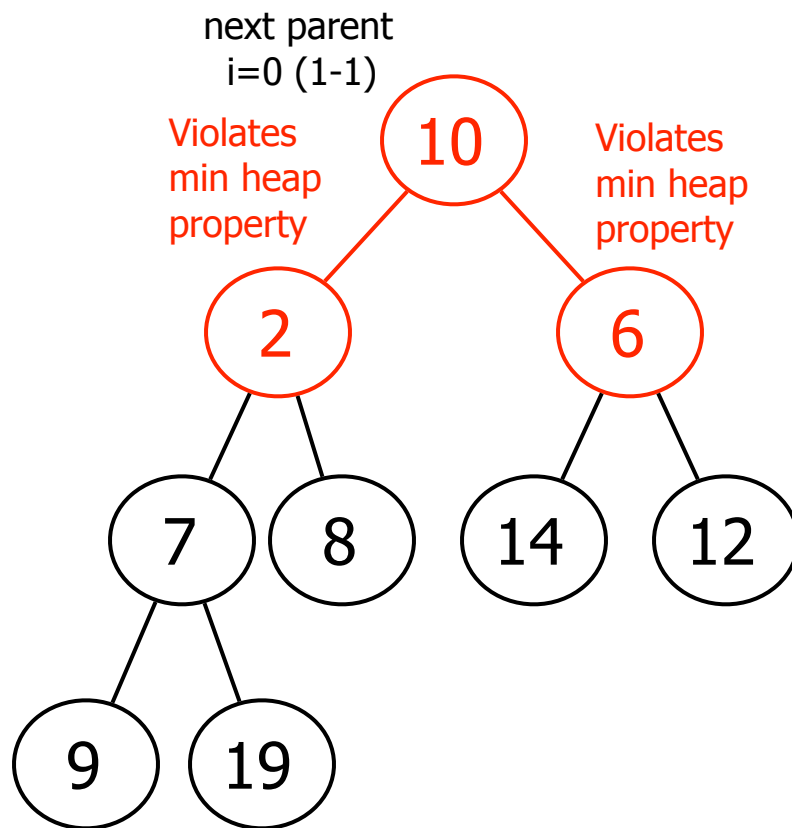
# Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
10	7	6	2	8	14	12	9	19



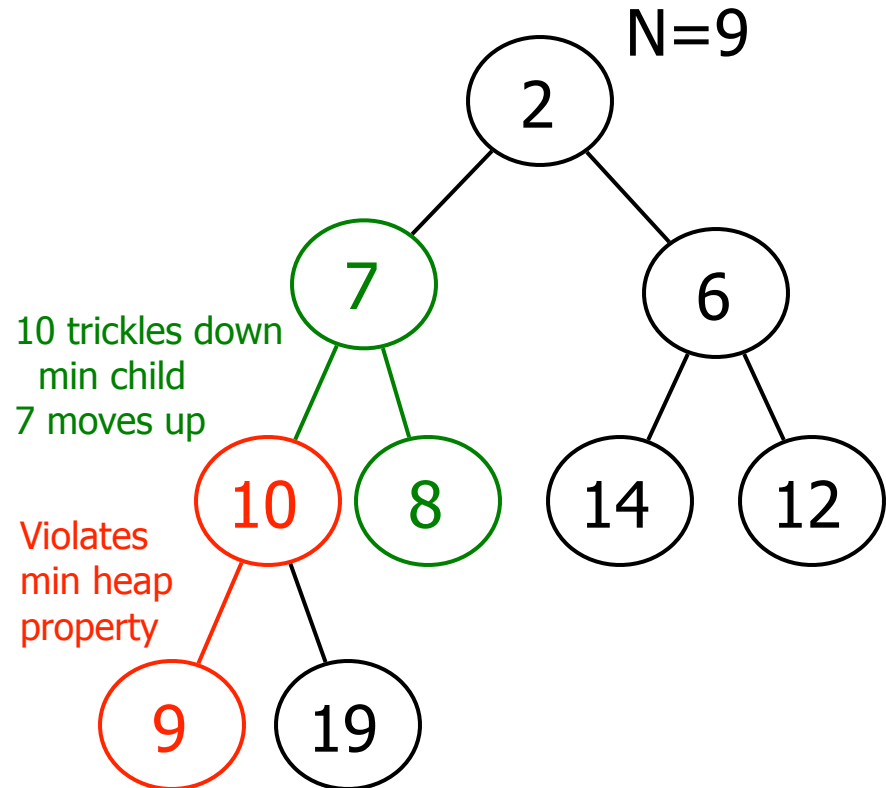
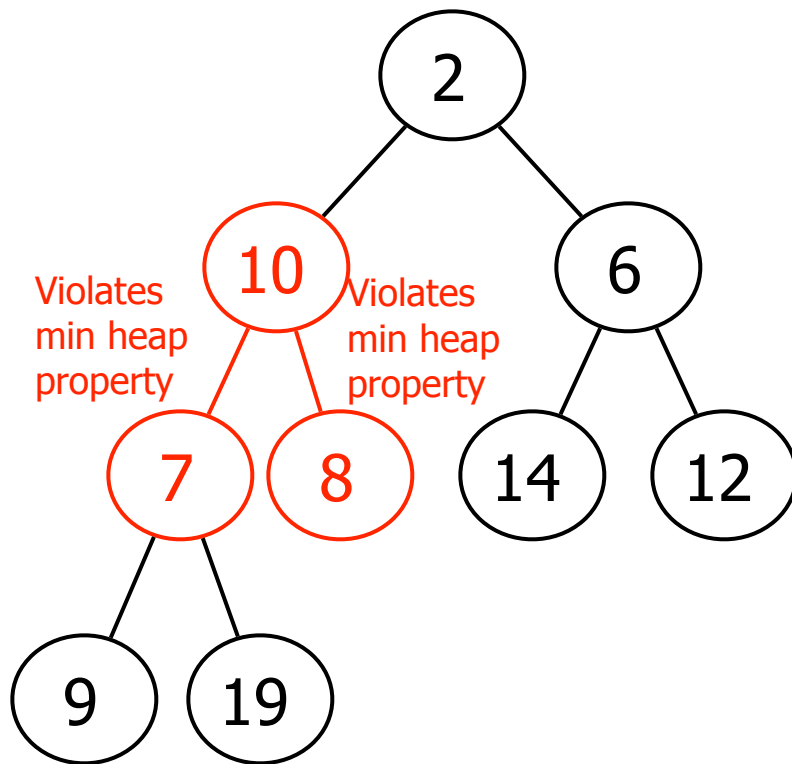
# Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
10	2	6	7	8	14	12	9	19



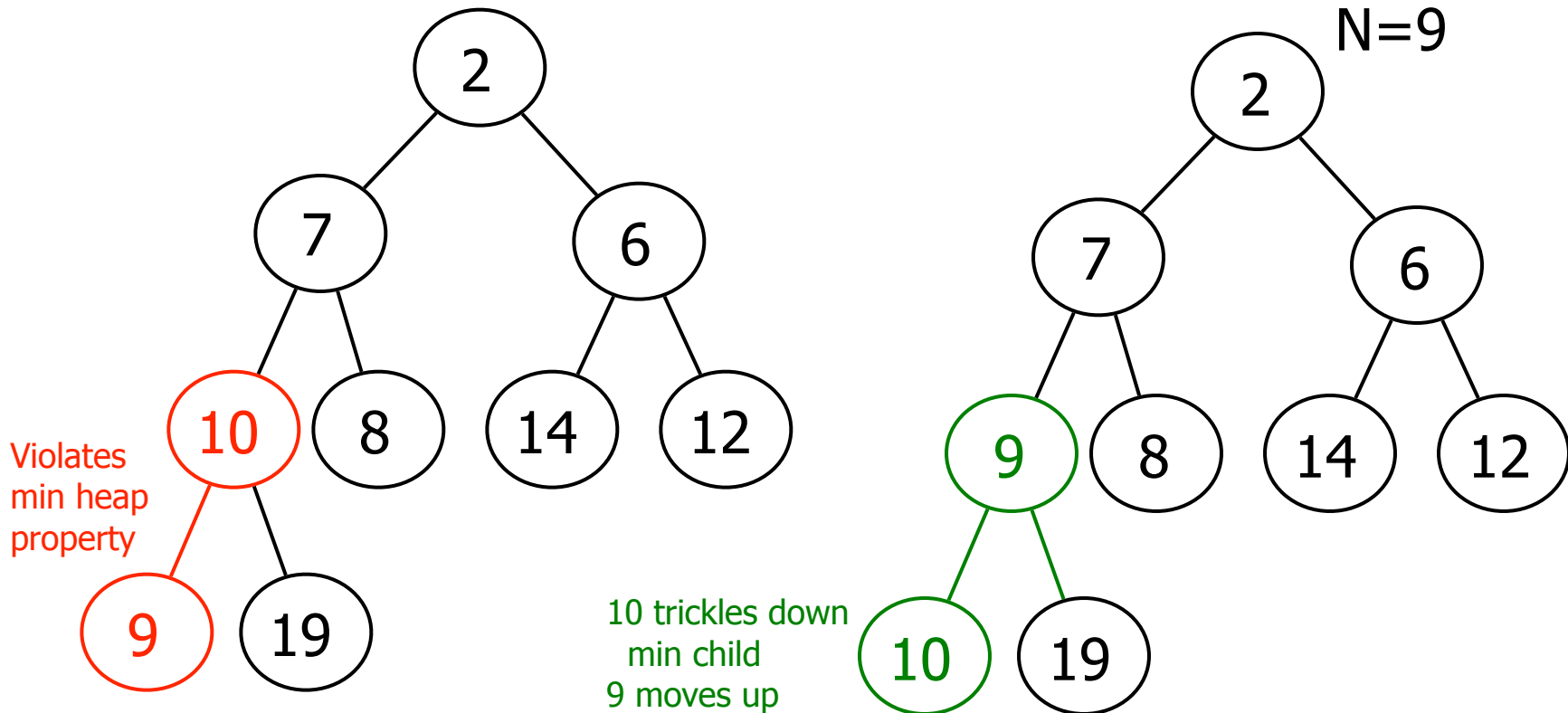
# Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
2	10	6	7	8	14	12	9	19



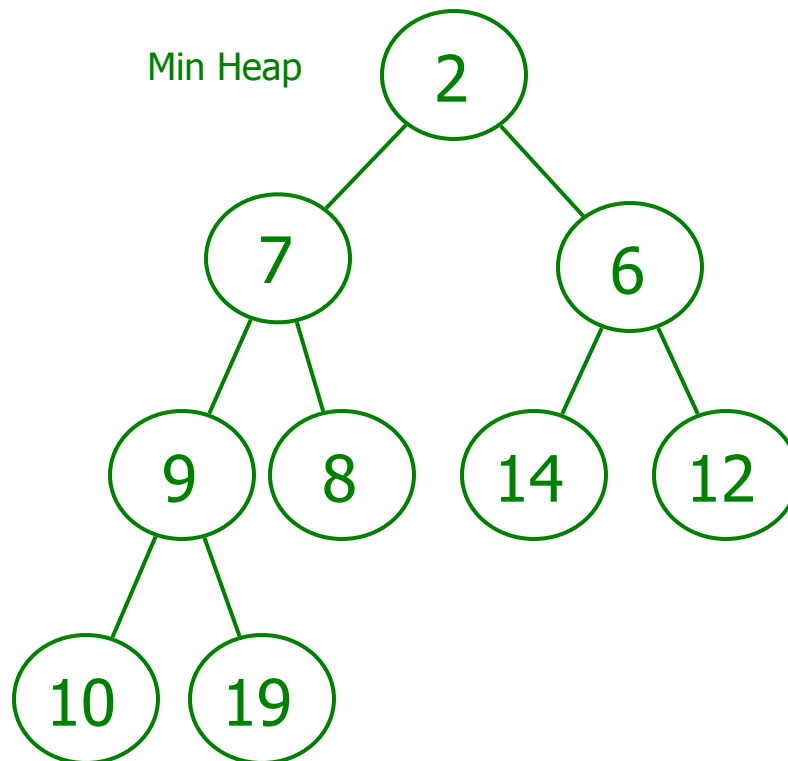
# Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
2	7	6	10	8	14	12	9	19



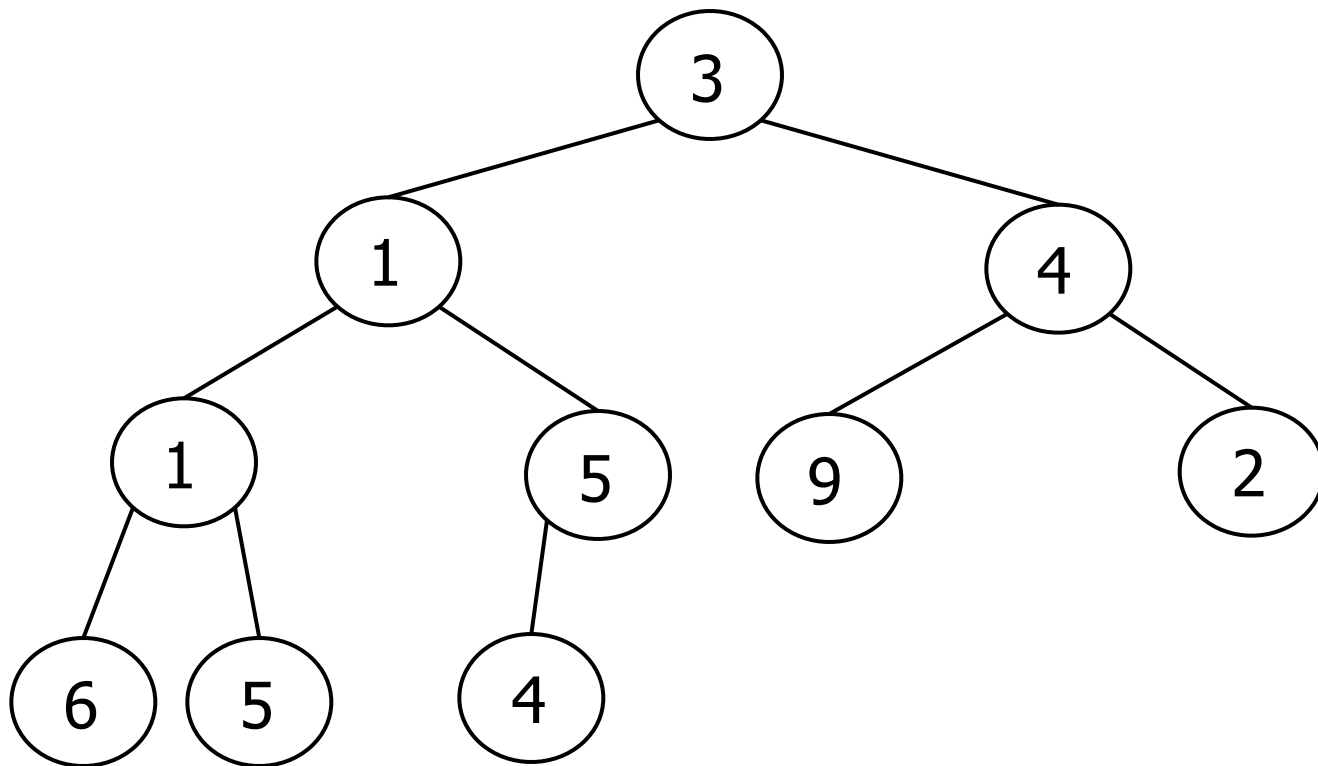
# Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
2	7	6	9	8	14	12	10	19



# ICE - Max Heap Build

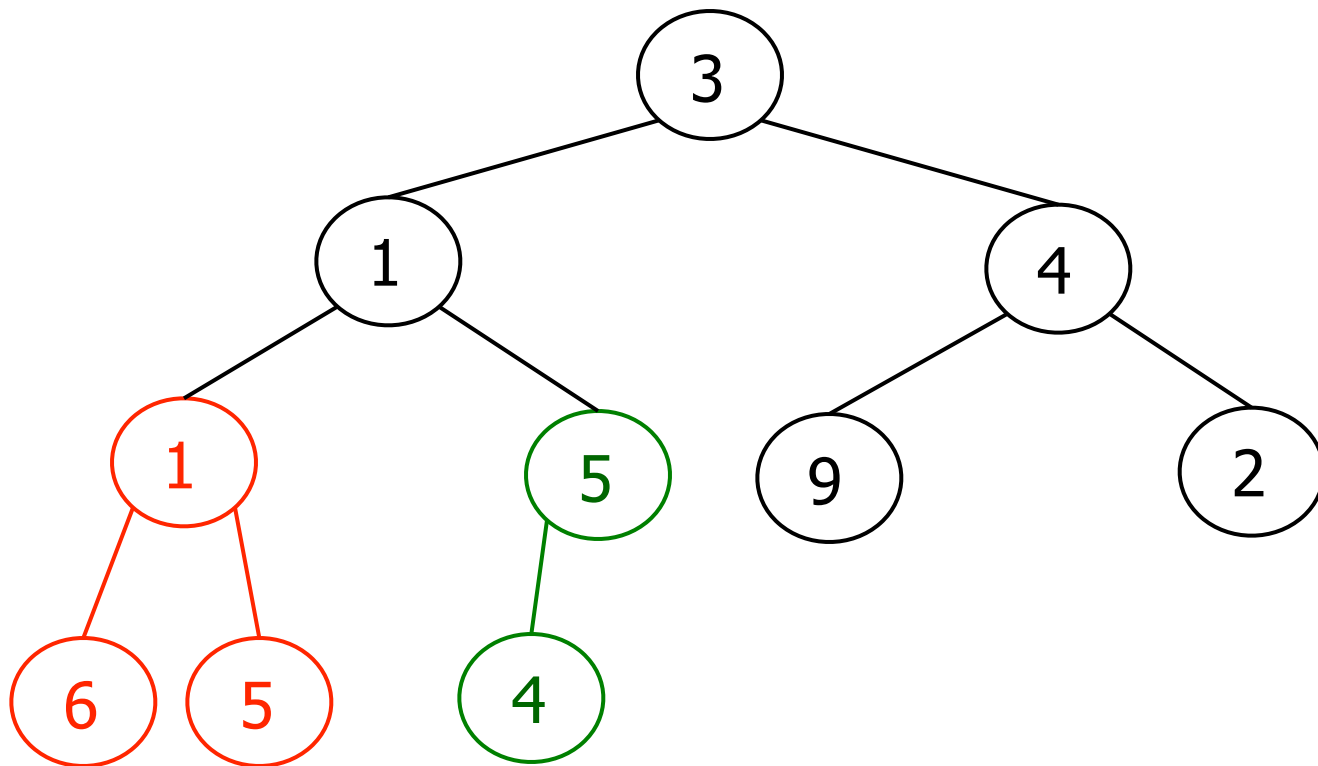
3	1	4	1	5	9	2	6	5	4
---	---	---	---	---	---	---	---	---	---





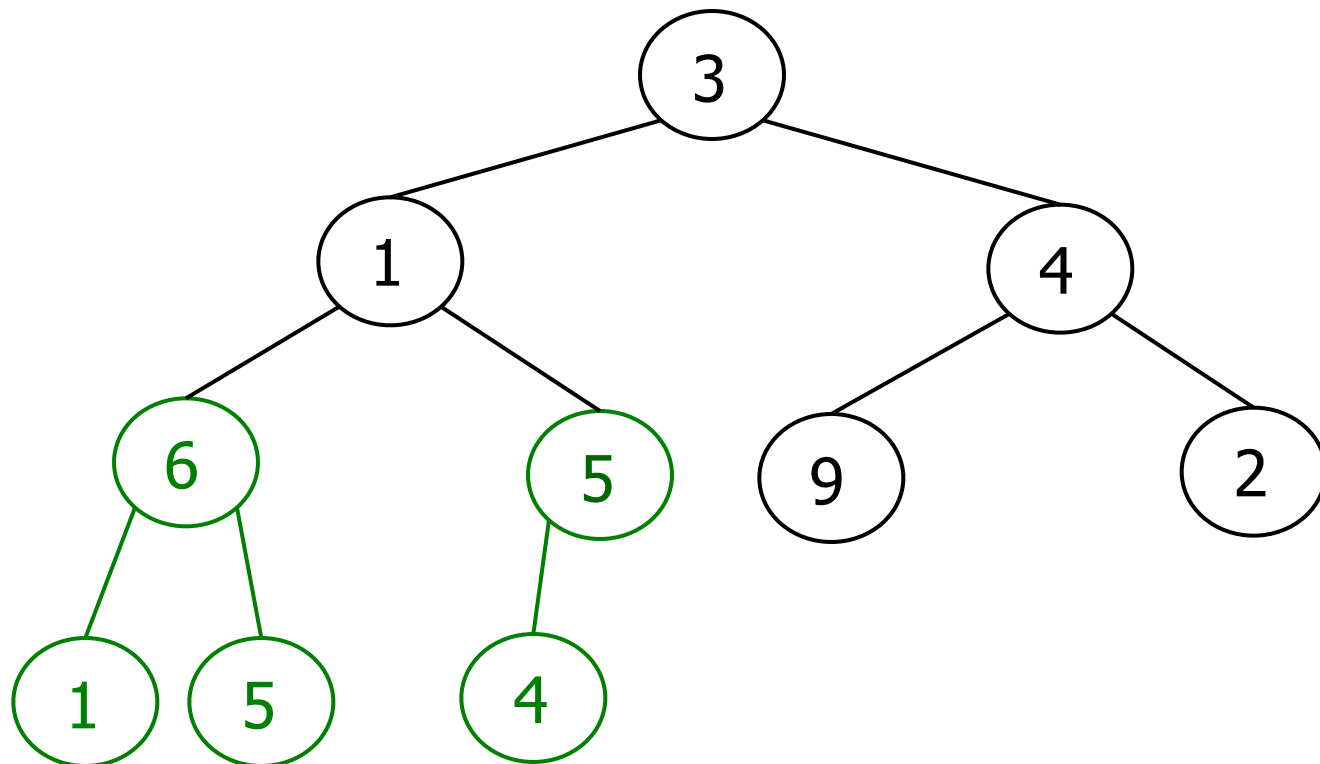
# ICE - Max Heap Build

3	1	4	1	5	9	2	6	5	4
---	---	---	---	---	---	---	---	---	---



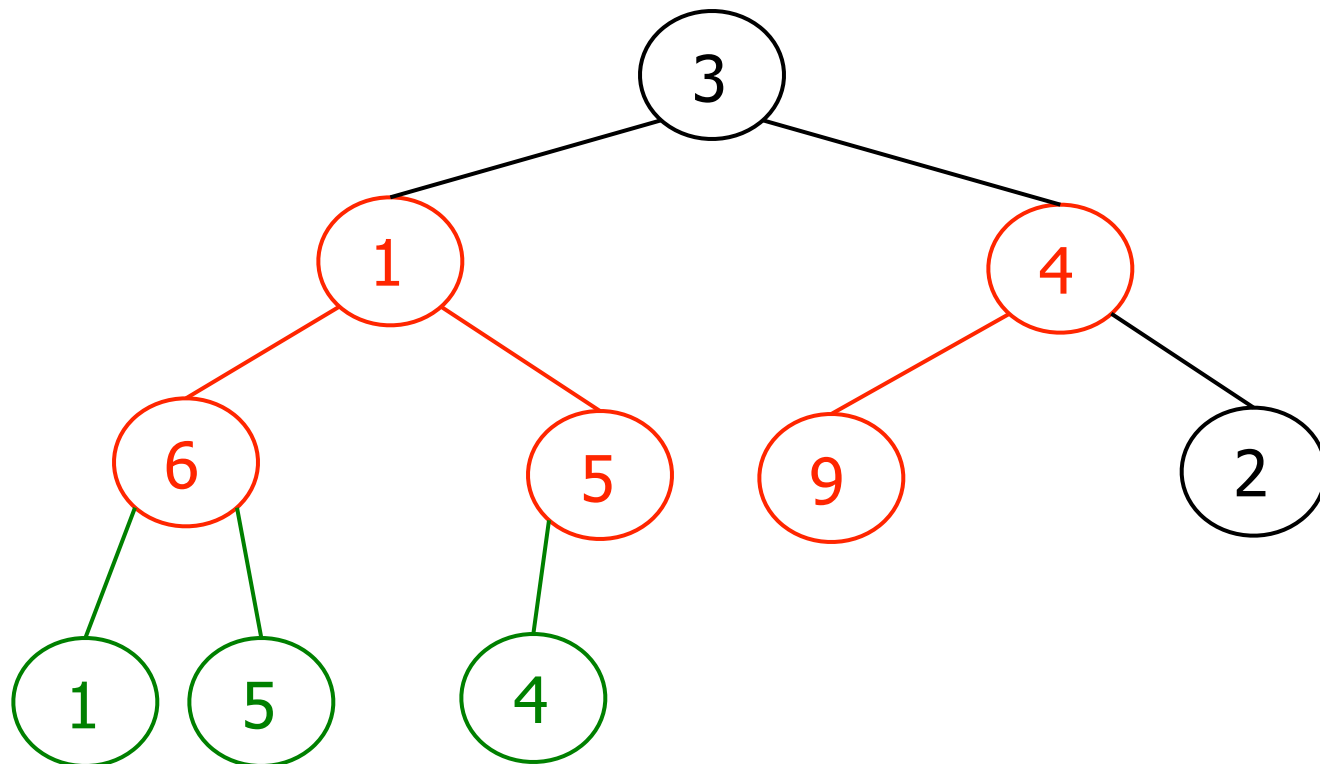
# ICE - Max Heap Build

3	1	4	6	5	9	2	1	5	4
---	---	---	---	---	---	---	---	---	---



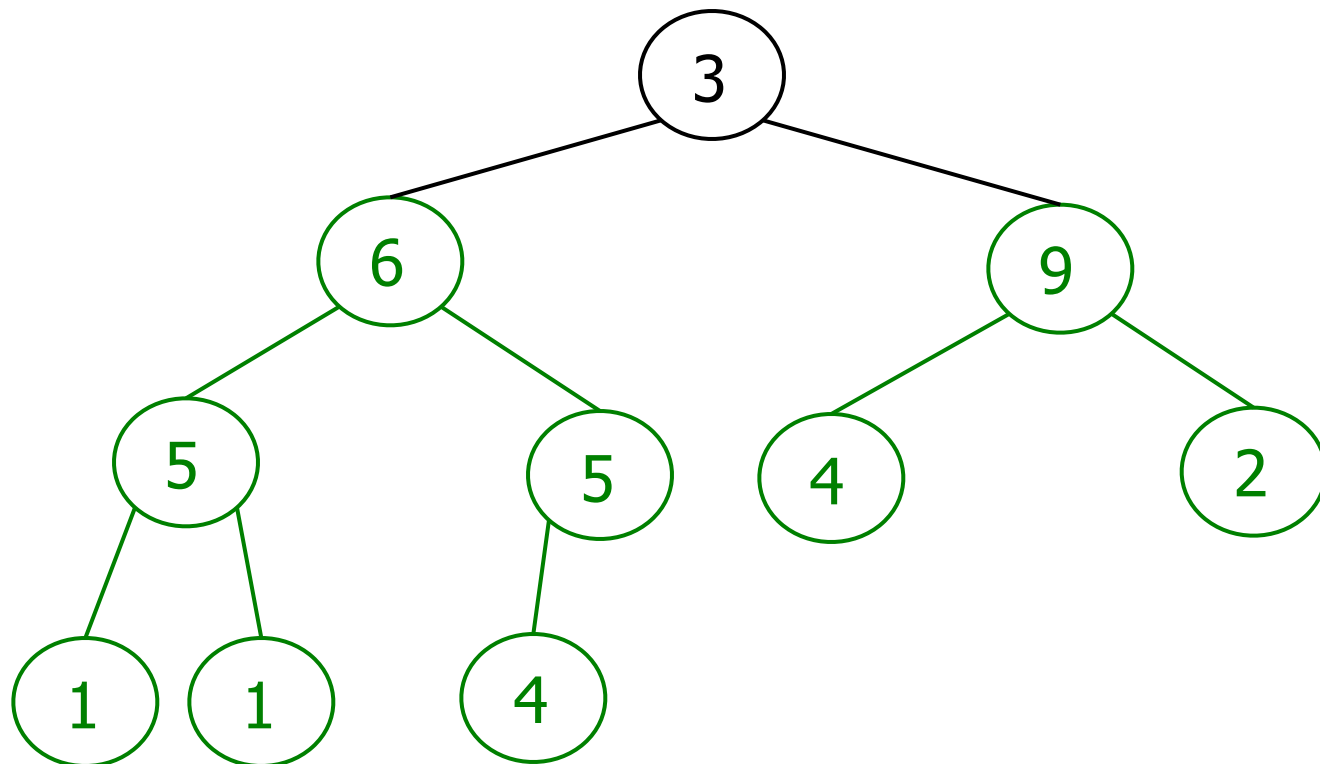
# ICE - Max Heap Build

3	1	4	6	5	9	2	1	5	4
---	---	---	---	---	---	---	---	---	---



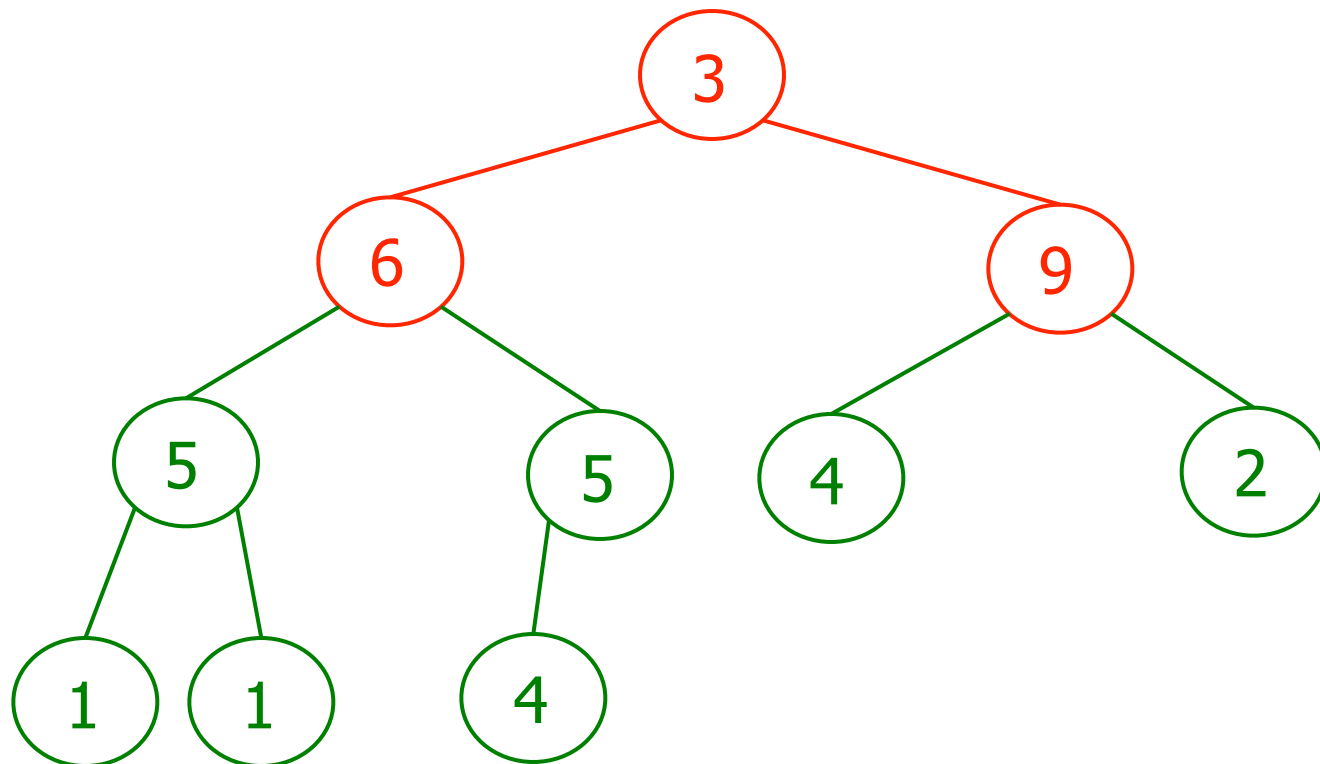
# ICE - Max Heap Build

3	6	9	5	5	4	2	1	1	4
---	---	---	---	---	---	---	---	---	---



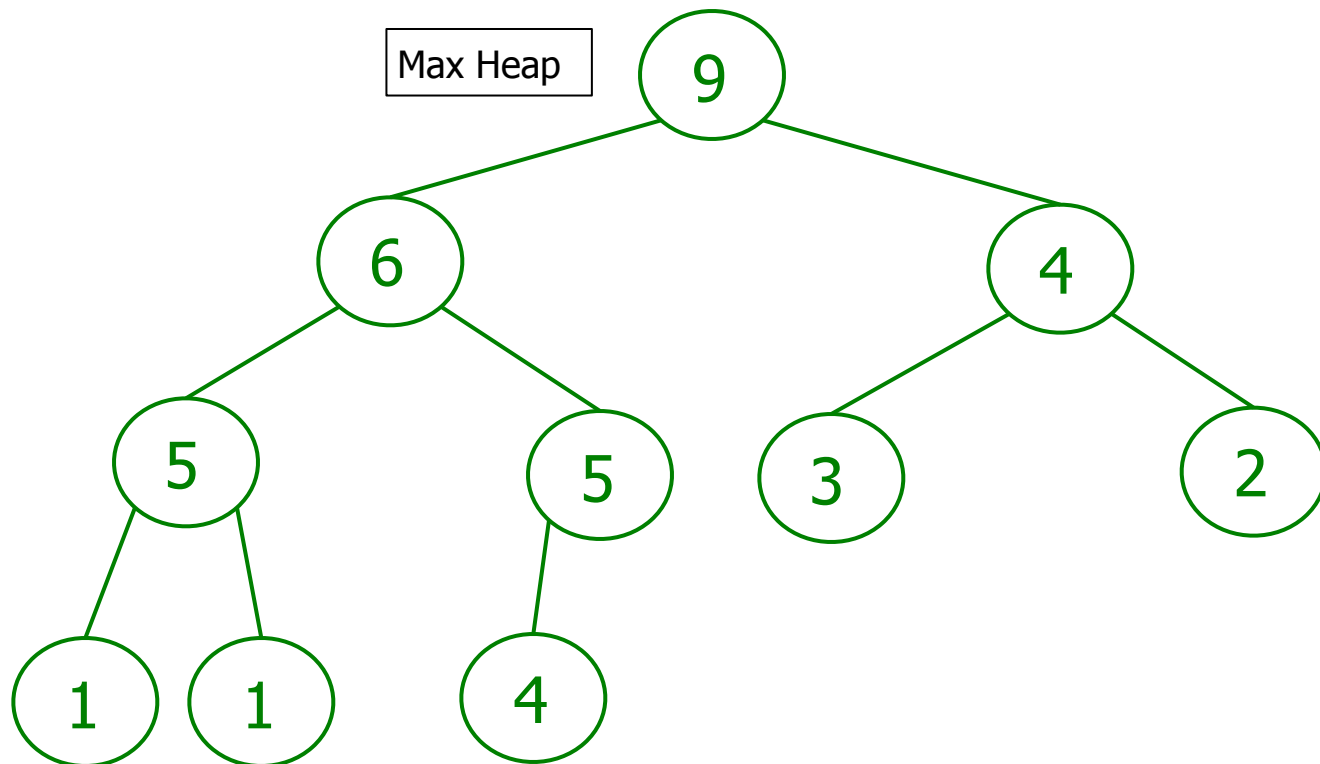
# ICE - Max Heap Build

3	6	9	5	5	4	2	1	1	4
---	---	---	---	---	---	---	---	---	---



# ICE - Max Heap Build

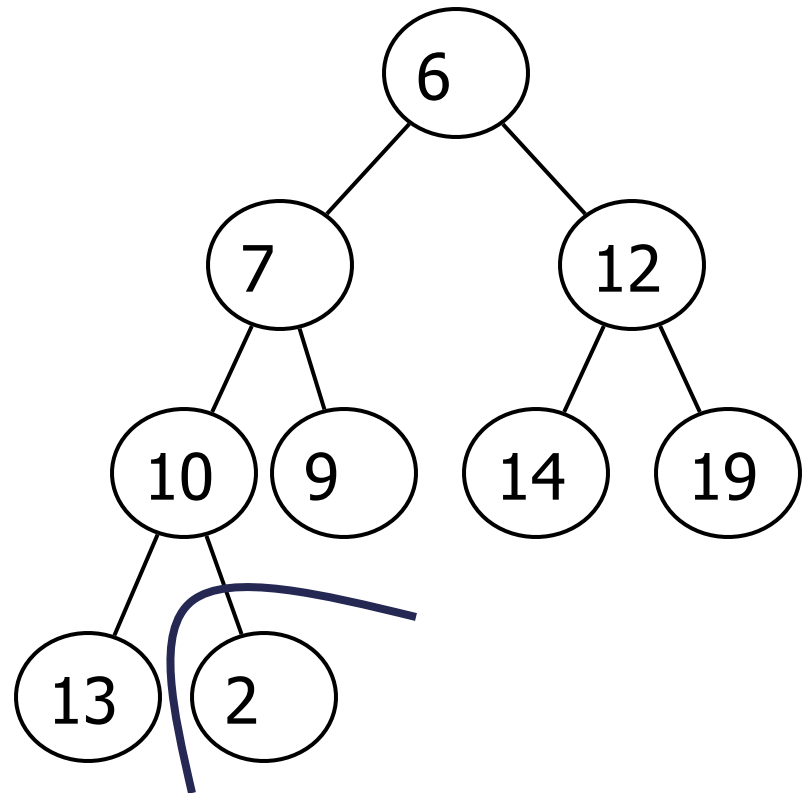
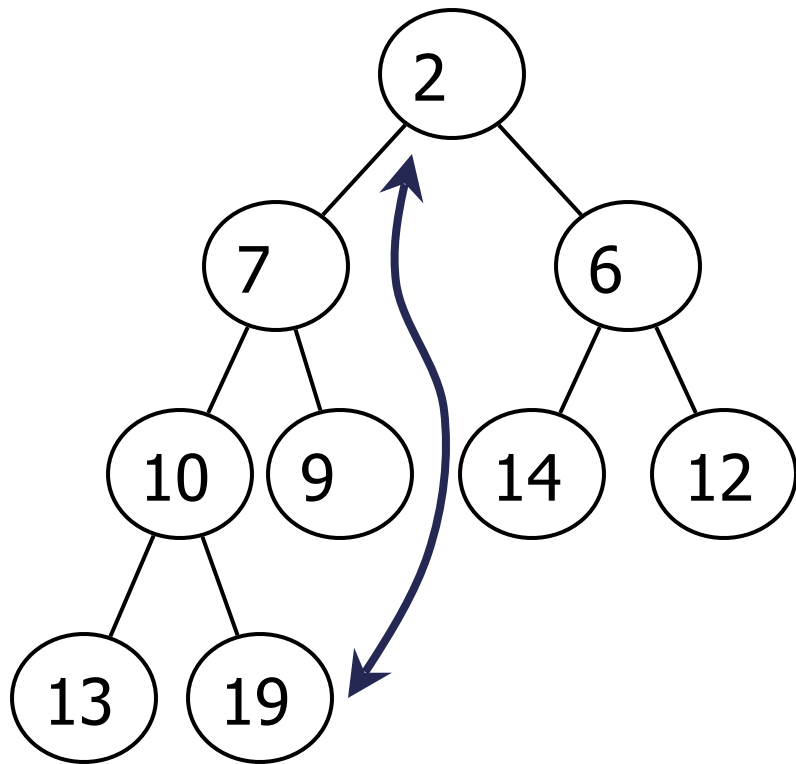
9	6	4	5	5	3	2	1	1	4
---	---	---	---	---	---	---	---	---	---



# Heapsort

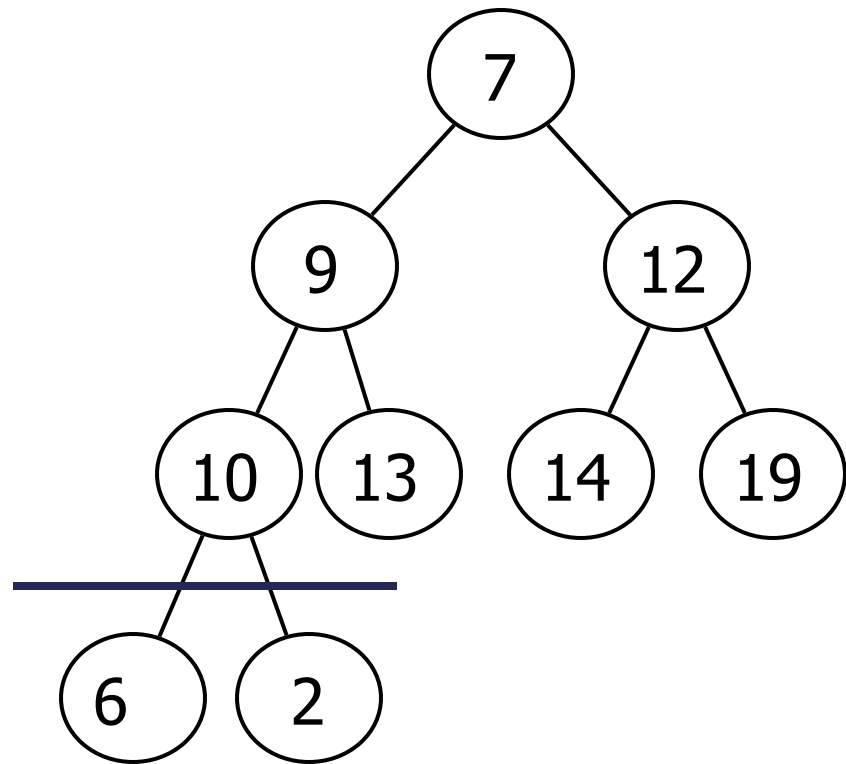
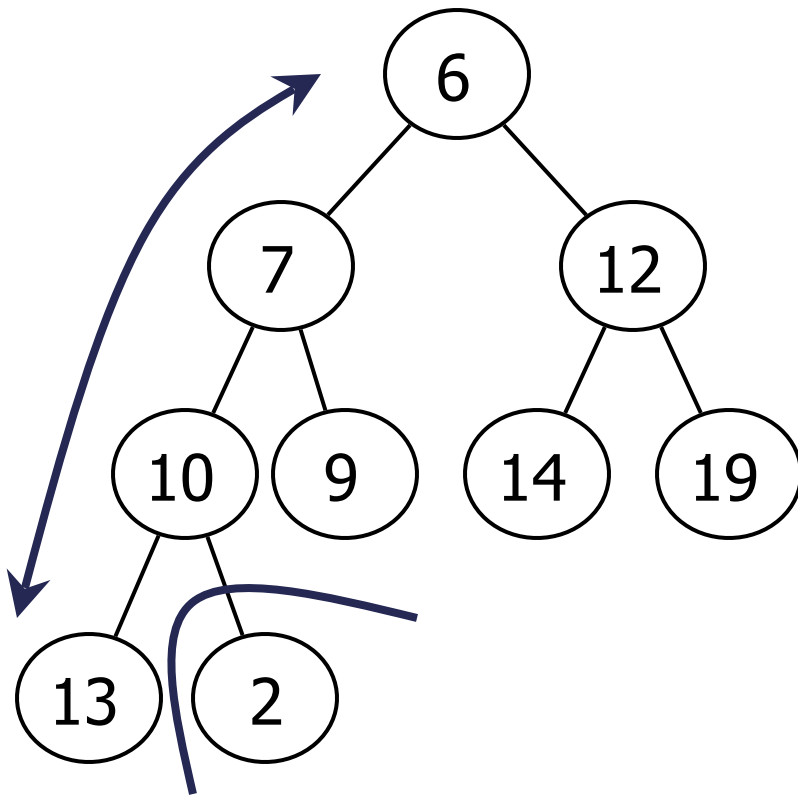
- Build heap for the opposite of what you want
  - Max heap for ascending order
  - Min heap for descending order
- Take root and place in last array position, then think of array as 1 smaller
- Trickle down from root to rebuild heap

# Heapsort - Sort in descending order

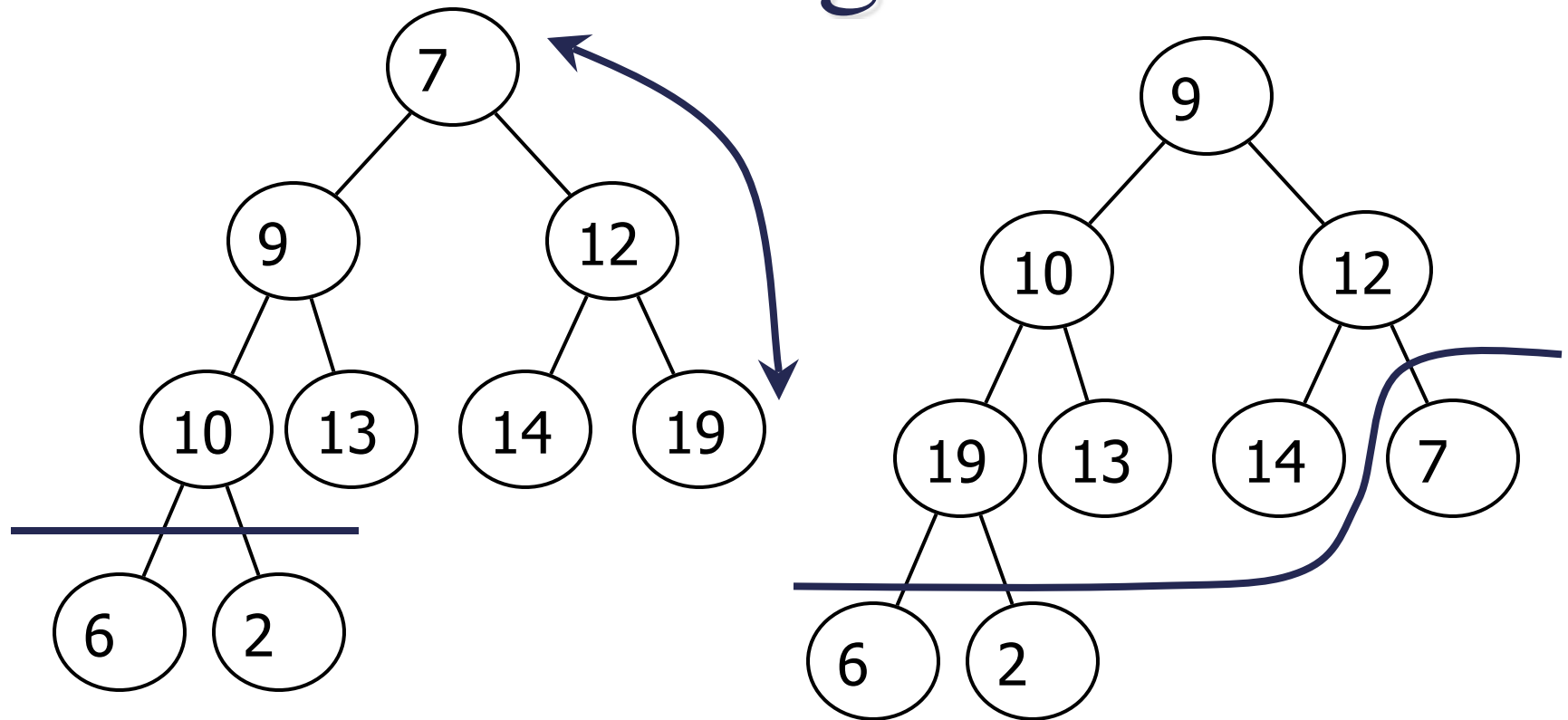




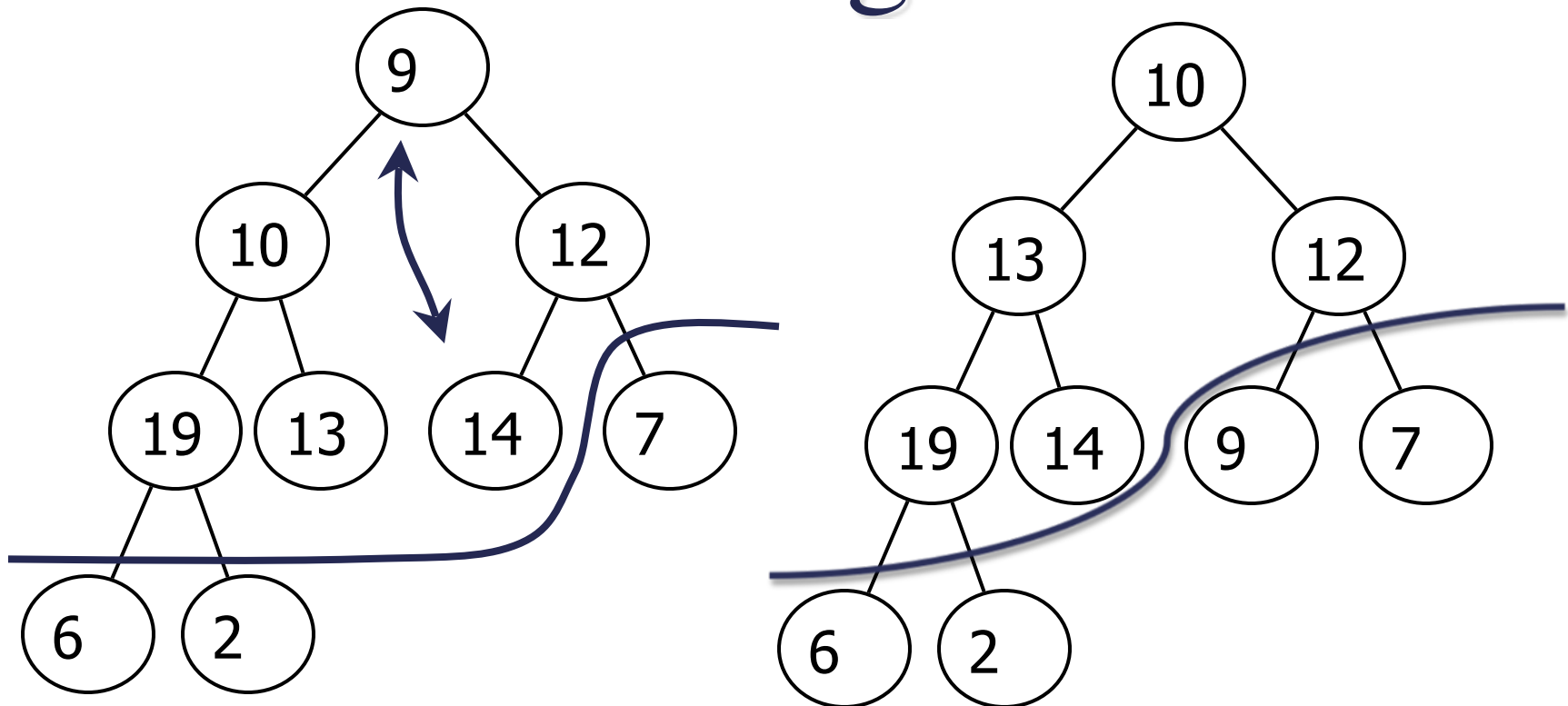
# Heapsort - Sort in descending order



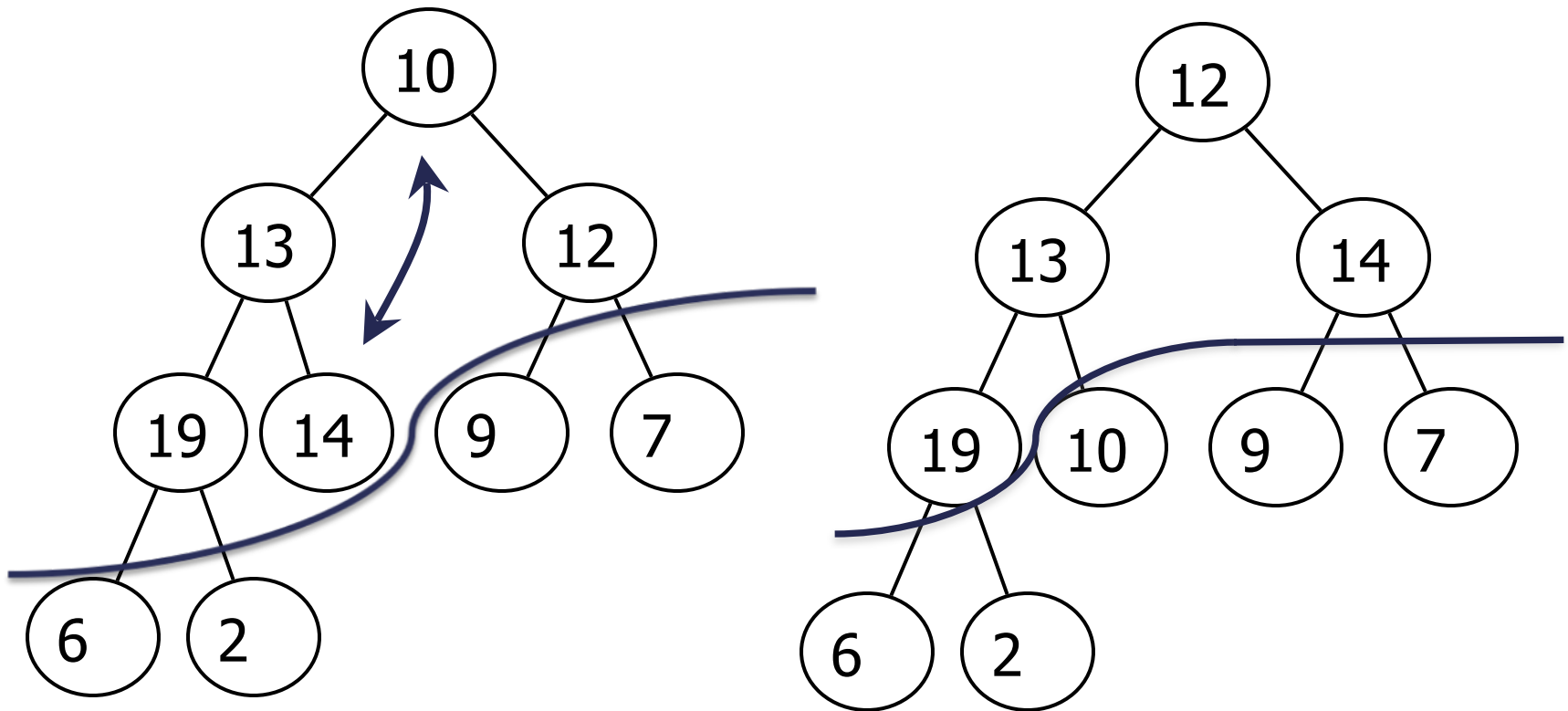
# Heapsort - Sort in descending order



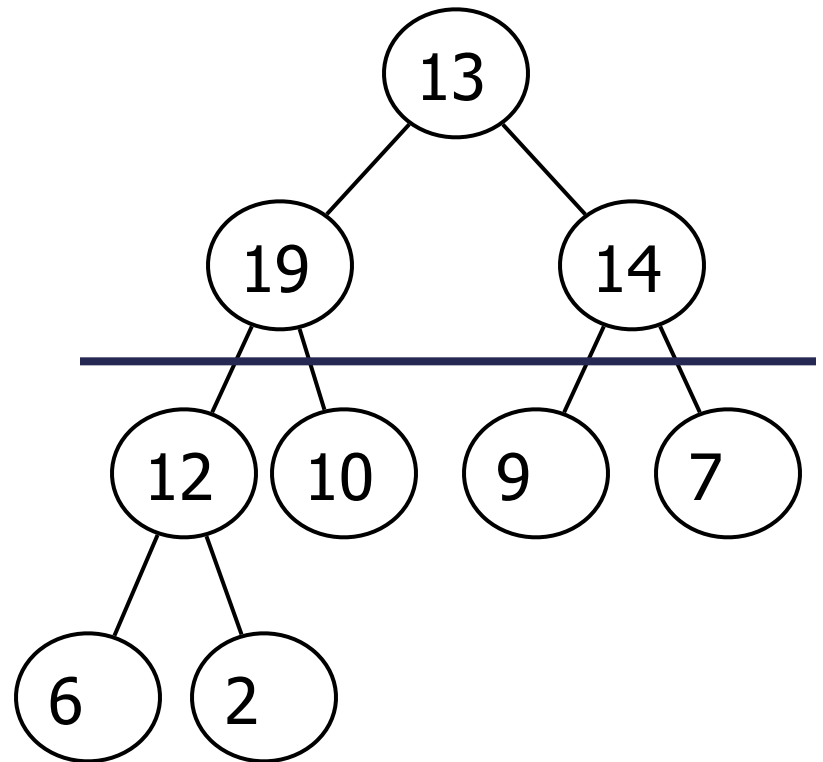
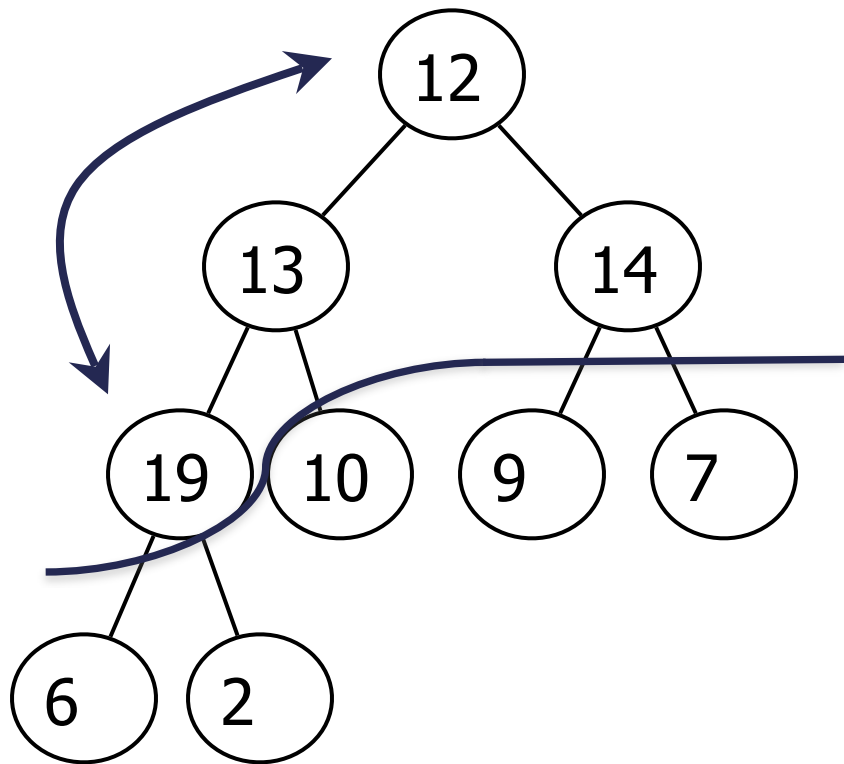
# Heapsort - Sort in descending order



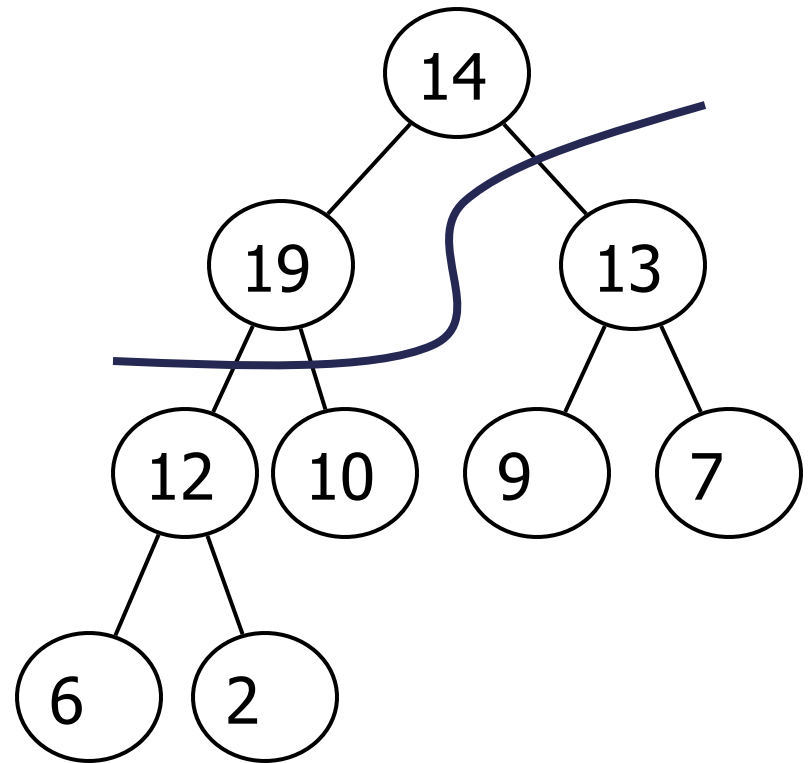
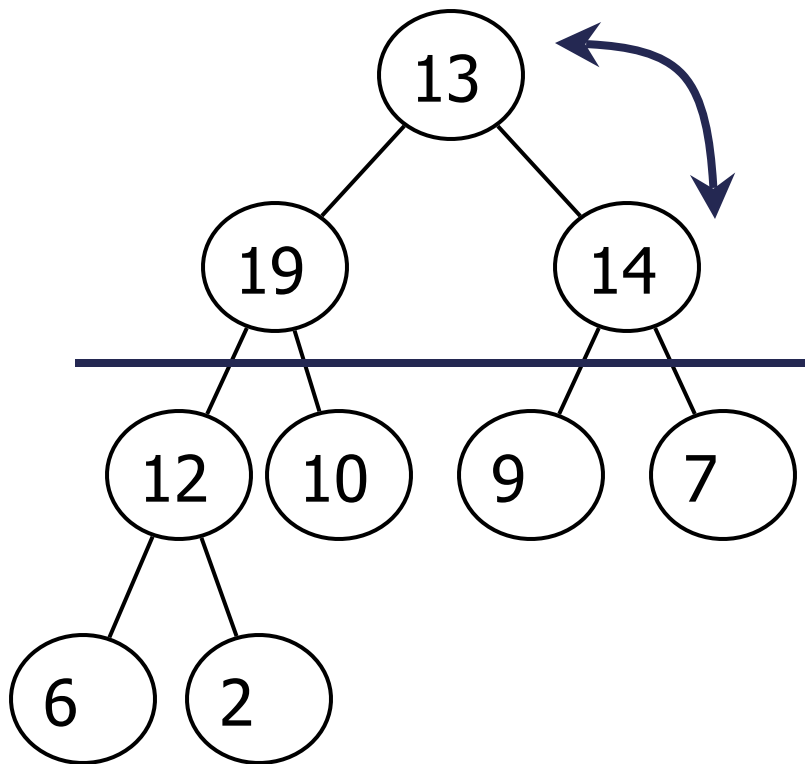
# Heapsort - Sort in descending order



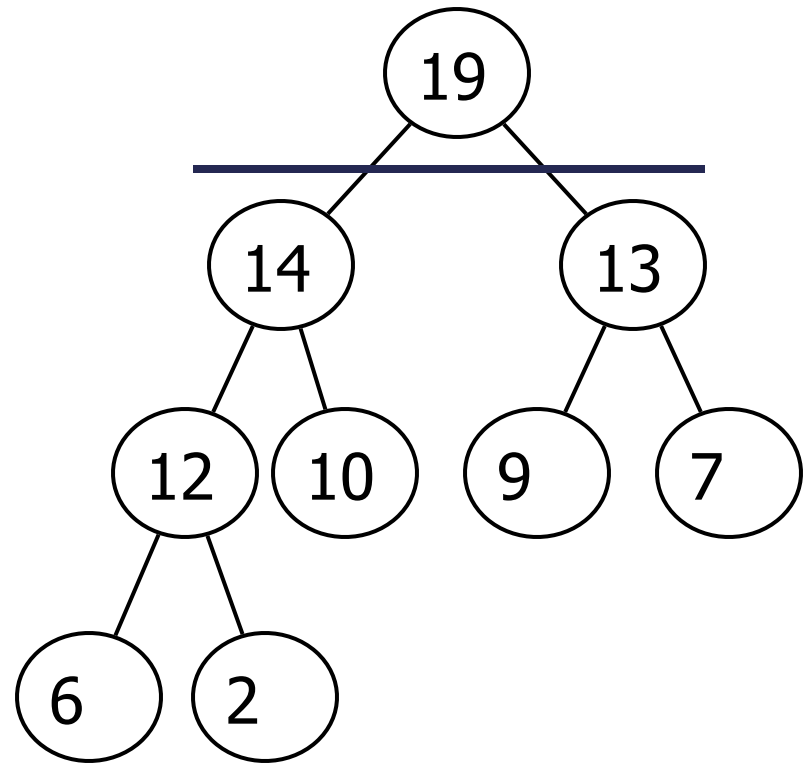
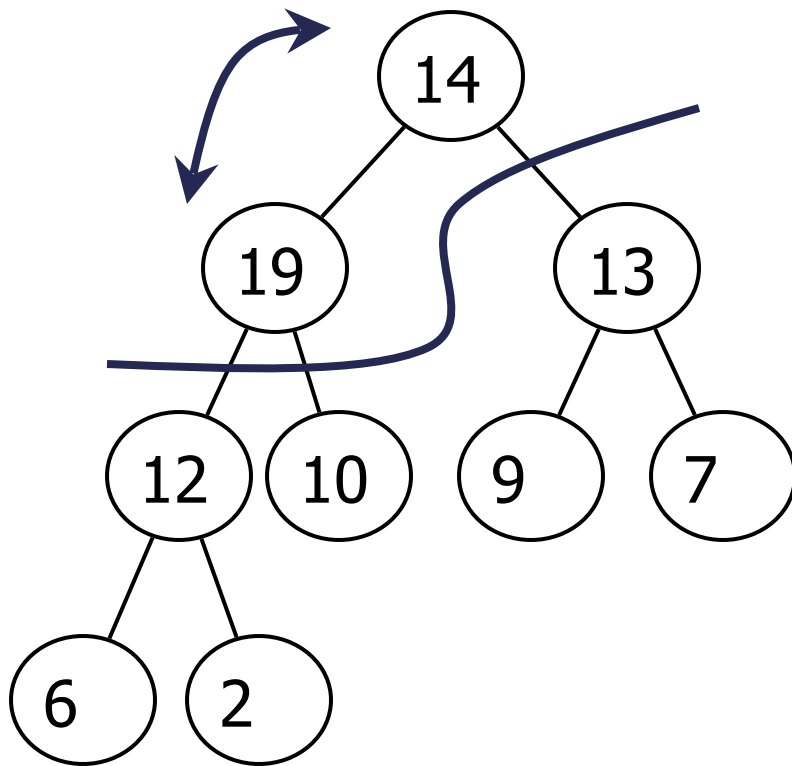
# Heapsort - Sort in descending order



# Heapsort - Sort in descending order



# Heapsort - Sort in descending order



# Heapsort Run Time

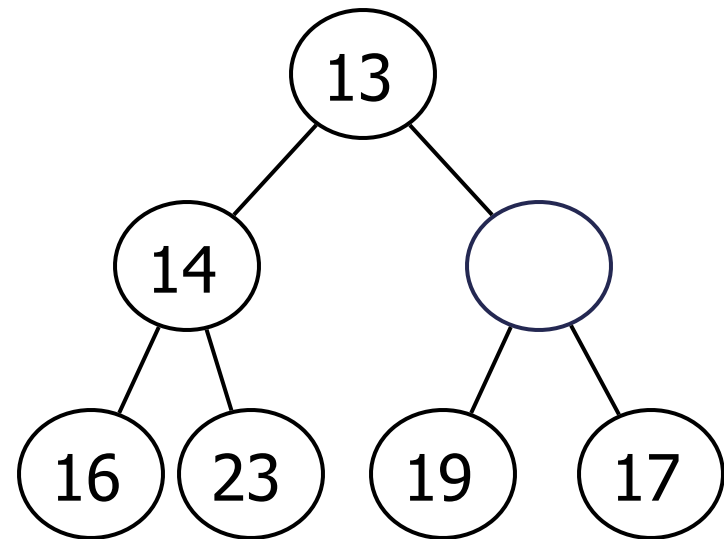
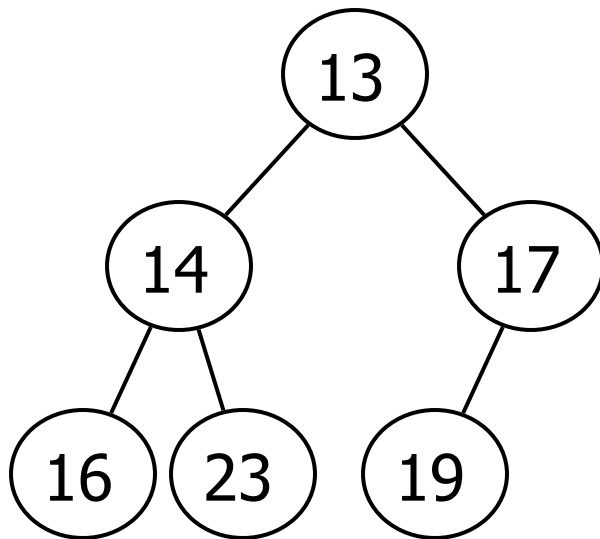
- Build step -  $O(n)$
- Sorting step -  $O(n \log n)$
- Heapsort -  $O(n) + O(n \log n) = O(n \log n)$



# Exam 7 Min Heap - Insert

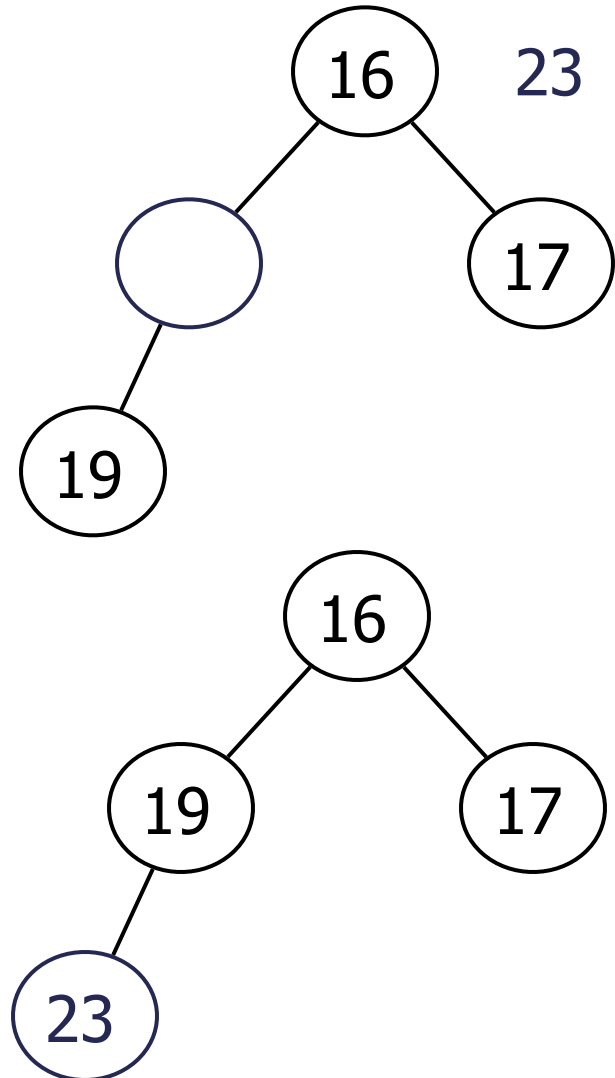
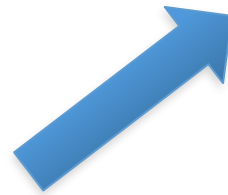
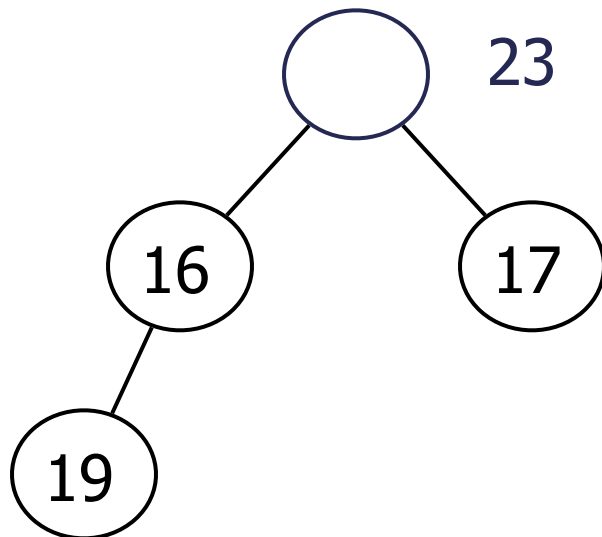
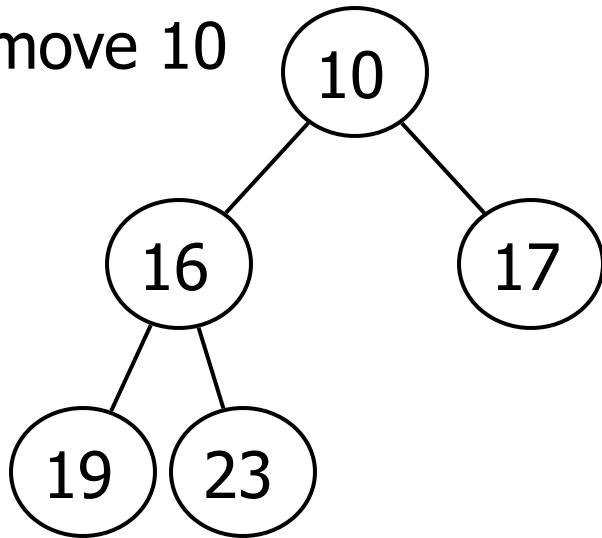
13, 16, 19, 14, 23, 17, 12

Insert 12



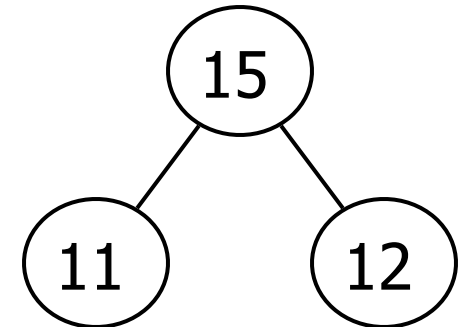
# Exam 7 Min Heap - Remove

Remove 10

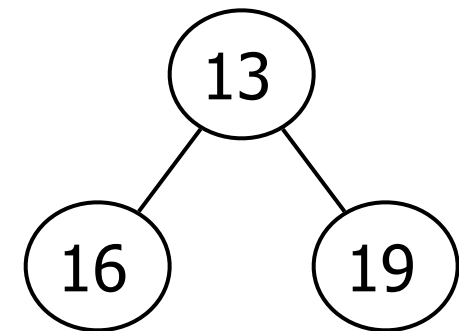


# Heap Order Property

- Every root or sub-root has a higher priority than all of its children nodes.
- **Min Heap** –
  - For each node  $X$ ,  $X.key > (X.parent).key$
  - Smallest number is at the root
- **Max Heap** –
  - For each node  $X$ ,  $X.key < (X.parent).key$
  - Largest number is at the root



Max-heap

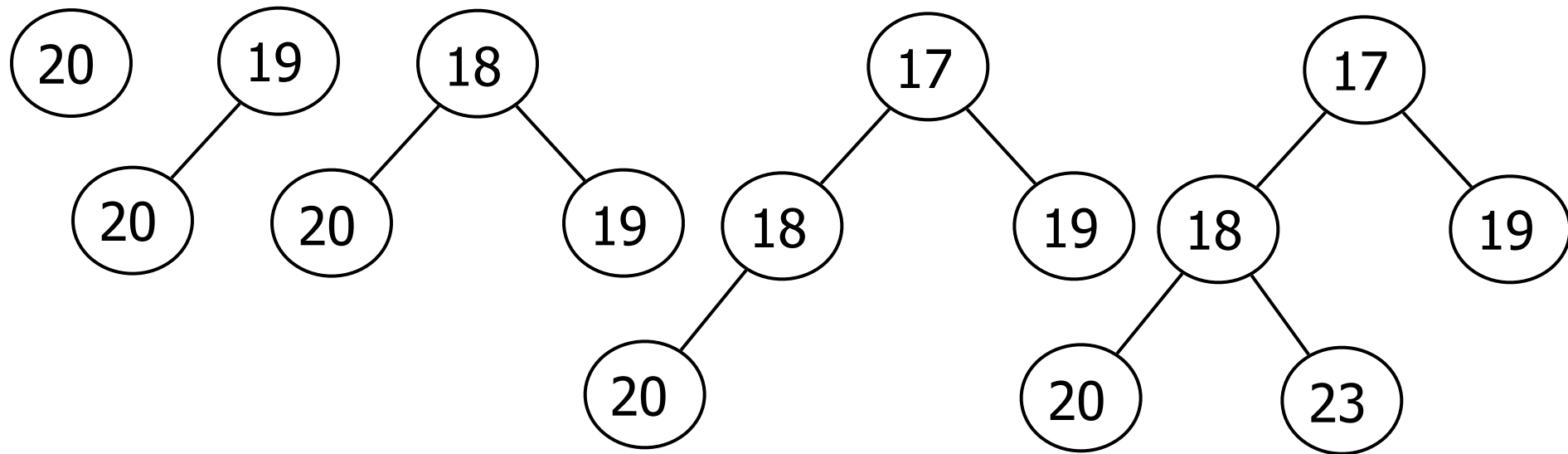


Min-heap

Exam 5 Q2:  
Summer 17

# Min-Heap Insert

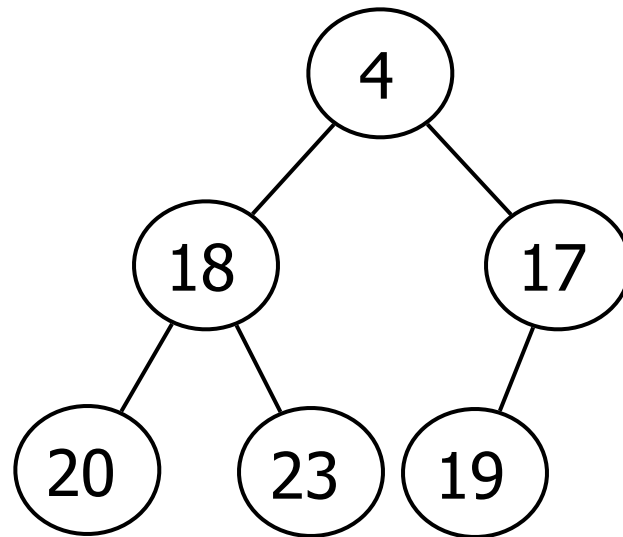
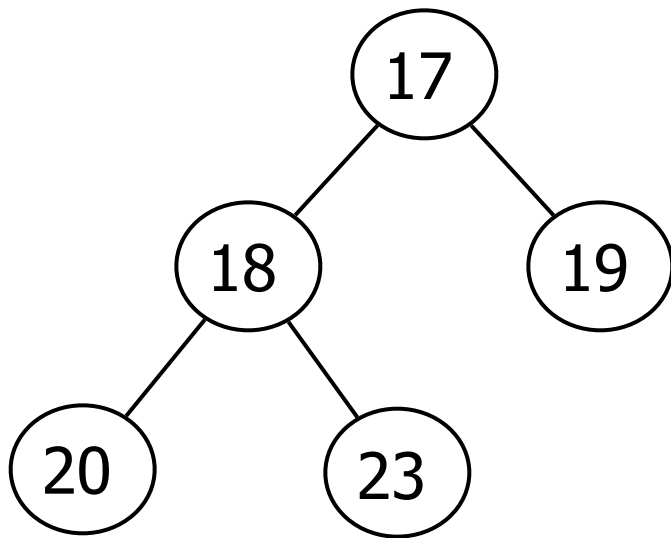
20, 19, 18, 17, 23, 4



Exam 5 Q3:  
Summer 17

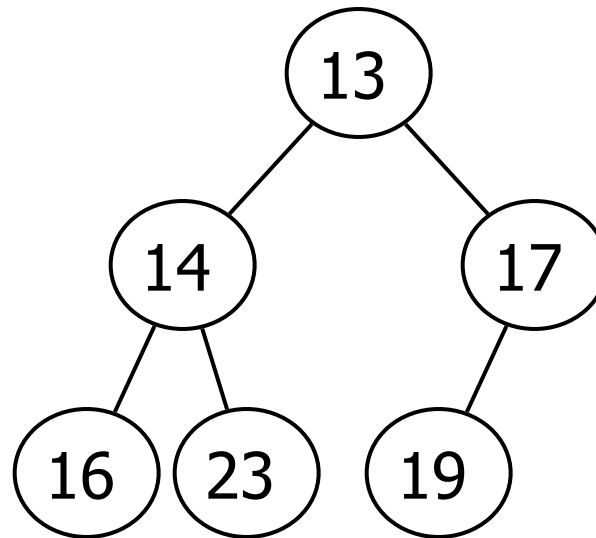
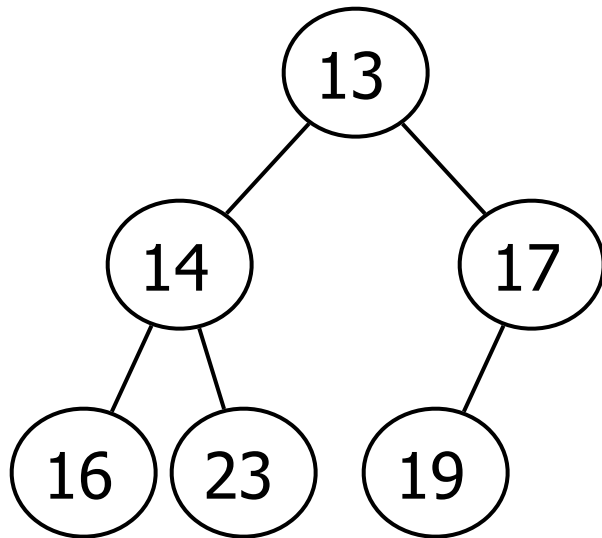
# Min-Heap Insert

20, 19, 18, 17, 23, 4



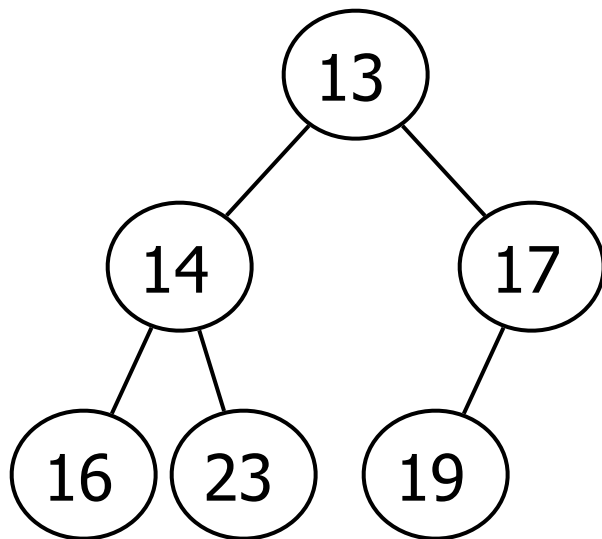
Exam 5 Q3:  
Summer 17

# Binary Heap Insert



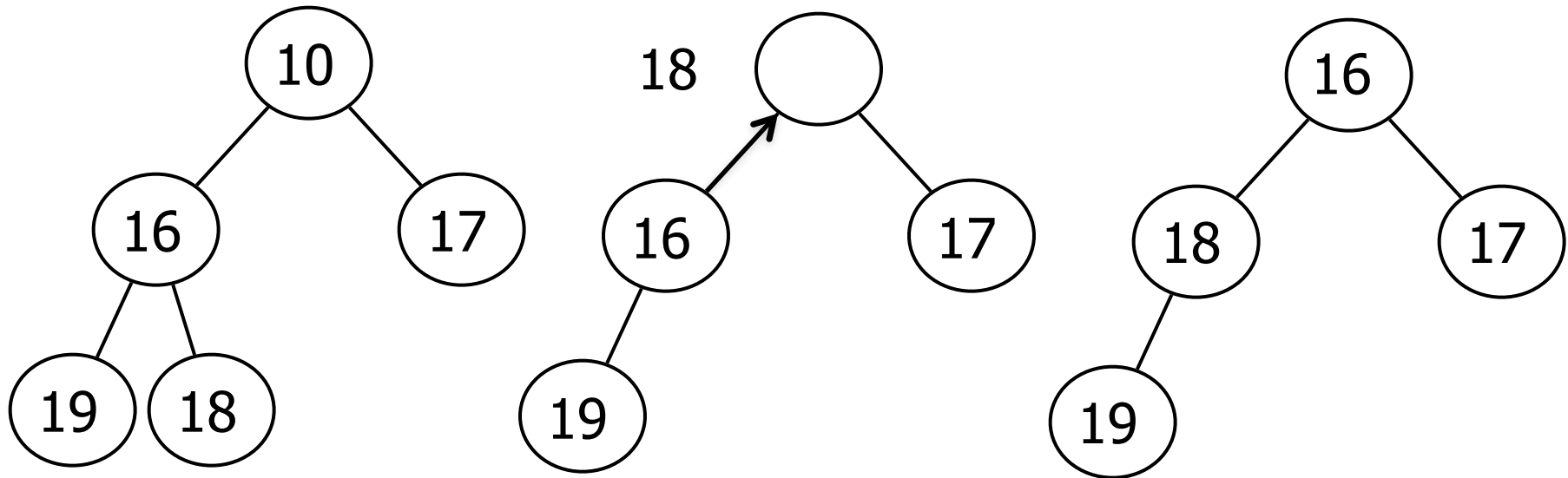
Exam 5 Q3:  
Summer 17

# Binary Heap Insert



Exam 5 Q3:  
Summer 17

# Binary Heap Remove



Exam 5 Q8:  
Summer 17