

link-cut tree

```
namespace lct {
    struct Node {
        int size, fa, son[2];
        bool rev;
    } tree[N];

    bool isRoot(int x) {
        return tree[tree[x].fa].son[0] != x && tree[tree[x].fa].son[1] != x;
    }

    int which(int x) {
        return tree[tree[x].fa].son[1] == x;
    }

    void apply(int x) {
        swap(tree[x].son[0], tree[x].son[1]);
        tree[x].rev ^= 1;
    }

    void pushDown(int x) {
        if (tree[x].rev) {
            apply(tree[x].son[0]), apply(tree[x].son[1]);
            tree[x].rev = 0;
        }
    }

    void pushUp(int x) {
        tree[x].size = 1 + tree[tree[x].son[0]].size + tree[tree[x].son[1]].size;
    }

    void rotate(int x) {
        int y = tree[x].fa, z = tree[y].fa;
        int id = which(x), p = tree[x].son[1 - id];
        if (!isRoot(y)) tree[z].son[which(y)] = x;
        tree[x].fa = z, tree[y].son[id] = p, tree[p].fa = y;
        tree[x].son[1 - id] = y, tree[y].fa = x;
        pushUp(y), pushUp(x);
    }

    void dfs(int root) {
        if (!isRoot(root)) dfs(tree[root].fa);
        pushDown(root);
    }

    void Splay(int x) {
        dfs(x);
        while (!isRoot(x)) {
            int y = tree[x].fa;
            if (!isRoot(y)) {
                rotate((which(x) == which(y)) ? y : x);
            }
            rotate(x);
        }
    }
```

```

}

// access u -> root
int access(int u) {
    int tmp = 0;
    while (u != 0) {
        Splay(u), tree[u].son[1] = tmp, pushUp(u);
        tmp = u, u = tree[u].fa;
    }
    return tmp;
}

// find the root of this tree.
int findRoot(int root) {
    access(root), Splay(root);
    int tmp = root;
    while (tree[tmp].son[0]) {
        pushDown(tmp);
        tmp = tree[tmp].son[0];
    }
    Splay(tmp);
    return tmp;
}

// make the node become root.
void makeRoot(int root) {
    access(root), Splay(root), apply(root);
}

// link edge u -> v
void link(int u, int v) {
    makeRoot(u); tree[u].fa = v;
}

// cut edge u -> v
void cut(int u, int v) {
    makeRoot(u), access(v), Splay(v);
    tree[v].son[0] = tree[u].fa = 0, pushUp(v);
}
}

```

Steiner Tree

$O(n3^k)$

```

// maximum node num, key node num, infinity.
const int N = 1010, K = 5, inf = 0x3f3f3f3f;

struct edge { int v, w; };

vector<edge> G[N];
int dp[1 << K][N], a[N];
bool inq[N];
queue<int> Q;

void init() {
    memset(dp, 0x3f, sizeof dp);
    for (int i = 0; i < N; i++) {

```

```

        G[i].clear();
    }
}

void addEdge(int u, int v, int w);

void bellmanFord(int *dp) {
    while (!Q.empty()) {
        int u = Q.front(); Q.pop(), inq[u] = 0;
        for (auto& e : G[u]) {
            if (dp[e.v] > dp[u] + e.w) {
                dp[e.v] = dp[u] + e.w;
                if (!inq[e.v]) {
                    Q.push(e.v), inq[e.v] = 1;
                }
            }
        }
    }
}

// k -> key node num, n -> node num.
void solve(int k, int n) {
    for (int S = 1; S < (1 << k); S++) {
        for (int i = 0; i < n; i++) {
            for (int s = (S - 1) & S; s; s = (s - 1) & S) {
                dp[S][i] = min(dp[S][i], dp[s][i] + dp[s ^ S][i]);
            }
            if (dp[S][i] < inf) Q.push(i), inq[i] = 1;
        }
        bellmanFord(dp[S]);
    }
}

```

ZhuLiu's Algo

$O(nm)(?)$

```

const int N = 1010, M = 1010101, inf = 0x3f3f3f3f;

// directed edges.
// r -> root, x_i -> u, y_i -> v, z_i -> w
int r, dfn, x[M], y[M], z[M], last[N], weight[N], id[N];
bool vis[N], instk[N];

void dfs(int u) {
    if (u == r) return;
    instk[u] = vis[u] = 1;
    if (!vis[last[u]]) {
        dfs(last[u]);
    } else if (instk[last[u]]) {
        id[u] = ++dfn;
        for (int v = last[u]; v != u; v = last[v]) {
            id[v] = dfn;
        }
    }
    instk[u] = 0;
    if (!id[u]) id[u] = ++dfn;
}

```

```

// n -> nodes num, m -> edges num.
int solve(int n, int m) {
    int ans = 0;
    while (true) {
        for (int i = 1; i <= n; i++) {
            weight[i] = inf;
        }
        for (int i = 1; i <= m; i++) {
            if (z[i] < weight[y[i]]) {
                last[y[i]] = x[i], weight[y[i]] = z[i];
            }
        }
        memset(id, 0, sizeof id);
        memset(vis, 0, sizeof vis);
        id[r] = dfn = 1;
        for (int i = 1; i <= n; i++) {
            if (i == r) continue;
            if (weight[i] == inf) return -1;
            ans += weight[i];
            if (!vis[i]) dfs(i);
        }
        if (dfn == n) return ans;
        int cnt = 0;
        for (int i = 1; i <= m; i++) {
            if (id[x[i]] != id[y[i]]) {
                z[++cnt] = z[i] - weight[y[i]];
                x[cnt] = id[x[i]], y[cnt] = id[y[i]];
            }
        }
        m = cnt, n = dfn, r = id[r];
    }
    return -1;
}

```

CartesianTree

```

// left son, right son, father.
int l[N], r[N], fa[N];

void build(int n, int arr[]) {
    static int top, stack[N];
    top = 0, stack[top++] = 1;
    for (int i = 2; i <= n; i++) {
        while (top && arr[stack[top - 1]] > arr[i]) --top;
        if (top > 0) {
            int x = i, y = stack[top - 1];
            l[x] = r[y], fa[l[x]] = x, fa[x] = y, r[y] = x;
            stack[top++] = x;
        } else {
            fa[stack[0]] = i, l[i] = stack[0];
            stack[top++] = i;
        }
    }
}

```

AC Automaton

```
template <int N, int charset> class acam {
    int tot;
    int fail[N], endpos[N];
    int son[N][charset];

    int encode(int c) { return c - 'A'; }
public:

    int fa[N], length[101010];
    bool vis[N];

    void initNode(int i) {
        fail[i] = 0, endpos[i] = -1;
        memset(son[i], 0, sizeof son[i]);
    }

    void insert(char *s, int index = 0) {
        int cur = 0;
        for (int i = 0; s[i]; i++) {
            int c = encode(s[i]);
            if (!son[cur][c]) {
                son[cur][c] = ++tot;
                initNode(tot);
            }
            fa[son[cur][c]] = cur;
            cur = son[cur][c];
            length[index] = i + 1;
        }
        endpos[index] = cur;
    }

    void build() {
        queue<int> Q;
        for (int i = 0; i < charset; i++) {
            if (son[0][i]) Q.push(son[0][i]);
        }
        while (!Q.empty()) {
            int u = Q.front(); Q.pop();
            for (int i = 0; i < charset; i++) {
                if (son[u][i]) {
                    fail[son[u][i]] = son[fail[u]][i];
                    Q.push(son[u][i]);
                } else {
                    son[u][i] = son[fail[u]][i];
                }
            }
        }
    }
};
```

pb_ds

```
#include<bits/extc++.h>
// 下面是set的例子 map的话将__gnu_pbds::null_type改成想要的结构即可
```

```

typedef __gnu_pbds::tree<int, __gnu_pbds::null_type, less<int>,
__gnu_pbds::rb_tree_tag, __gnu_pbds::tree_order_statistics_node_update> ordered_set;
/*
    iterator find_by_order(size_type order)
        找到第order+1小的迭代器，如果order太大会返回end()
    size_type order_of_key(const_key_reference r_key)
        询问这个tree中有多少个比r_key小的元素
    void join(tree &other)
        把other中所有元素移动到*this中（要求原来other和*this的key不能相交，否则会抛出异常）
*/
typedef __gnu_pbds::priority_queue<int, less<int>> pq;
/*
    void join(priority_queue &other)
        把other合并到*this，然后other会被清空
*/

```

树哈希

$$f_{now} = 1 + \sum f_{son_{now,i}} \times prime(siz[son_{now,i}])$$

Notes 其中 f_x 为以节点 x 为根的子树对应的哈希值。 $siz[x]$ 表示以节点 x 为根的子树大小。 $son_{x,i}$ 表示 x 的所有子节点之一。 $prime(i)$ 表示第 i 个质数。 选树的重心作根可保证一棵树仅有两个哈希值

线性递推

```

namespace LinearRecurrence {
    VL ReedsSloane(const VL &s, ll Mod) {
        function<void(VL &, size_t)> extend = [](VL &v, size_t d) {
            if(d <= v.size()) return;
            v.resize(d, 0);
        };
        function<ll(ll, ll)> inverse = [](ll a, ll m) {
            pll ret = crt::exgcd(a, m);
            ll g = a * ret.first + m * ret.second;
            if(g != 1) return -1ll;
            return ret.first >= 0 ? ret.first : ret.first + m;
        };
        function<int(const VL&, const VL&> L = [](const VL &a, const VL &b) {
            int da = (a.size() > 1 || (a.size() == 1 && a[0])) ? a.size() - 1 : -1000;
            int db = (b.size() > 1 || (b.size() == 1 && b[0])) ? b.size() - 1 : -1000;
            return max(da, db + 1);
        });
        function<pair<VL, VL>(const VL&, ll, ll, ll)> prime_power = [&](const VL &s, ll
Mod, ll p, ll e) {
            vector<VL> a(e), b(e), an(e), bn(e), ao(e), bo(e);
            VL t(e), u(e), r(e), to(e, 1), uo(e), pw(e + 1);
            pw[0] = 1;
            for(int i = pw[0] = 1; i <= e; i++) pw[i] = pw[i - 1] * p;
            for(ll i = 0; i < e; i++) {
                a[i] = {pw[i]}, an[i] = {pw[i]};
                b[i] = {0}, bn[i] = {s[0] * pw[i] % Mod};
                t[i] = s[0] * pw[i] % Mod;
                if(t[i] == 0) t[i] = 1, u[i] = e;
                else {
                    for(u[i] = 0; t[i] % p == 0; t[i] /= p, u[i]++);
                }
            }
        };
    }
}

```

```

    }
}
for(size_t k = 1; k < s.size(); k++) {
    for(int g = 0; g < e; g++) {
        if(L(an[g], bn[g]) > L(a[g], b[g])) {
            ao[g] = a[e - 1 - u[g]];
            bo[g] = b[e - 1 - u[g]];
            to[g] = t[e - 1 - u[g]];
            uo[g] = u[e - 1 - u[g]];
            r[g] = k - 1;
        }
    }
}
a = an, b = bn;
for(int o = 0; o < e; o++) {
    ll d = 0;
    for(size_t i = 0; i < a[o].size() && i <= k; i++) {
        d = (d + a[o][i] * s[k - i]) % Mod;
    }
    if(d == 0) t[o] = 1, u[o] = e;
    else {
        for(u[o] = 0, t[o] = d; t[o] % p == 0; t[o] /= p, u[o]++);
        int g = e - 1 - u[o];
        if(L(a[g], b[g]) == 0) {
            extend(bn[o], k + 1);
            bn[o][k] = (bn[o][k] + d) % Mod;
        } else {
            ll coef = t[o] * inverse(to[g], Mod) % Mod * pw[u[o] - uo[g]] % Mod;
            int m = k - r[g];
            extend(an[o], ao[g].size() + m);
            extend(bn[o], bo[g].size() + m);
            for(size_t i = 0; i < ao[g].size(); i++) {
                an[o][i + m] -= coef * ao[g][i] % Mod;
                if(an[o][i + m] < 0) an[o][i + m] += Mod;
            }
            while(an[o].size() && an[o].back() == 0) an[o].pop_back();
            for(size_t i = 0; i < bo[g].size(); i++) {
                bn[o][i + m] -= coef * bo[g][i] % Mod;
                if(bn[o][i + m] < 0) bn[o][i + m] += Mod;
            }
            while(bn[o].size() && bn[o].back() == 0) bn[o].pop_back();
        }
    }
}
}
}
return make_pair(an[0], bn[0]);
};

vector<tuple<ll, ll, int>> fac;
for(ll i = 2; i * i <= Mod; i++) {
    if(Mod % i == 0) {
        ll cnt = 0, pw = 1;
        while(Mod % i == 0) Mod /= i, ++cnt, pw *= i;
        fac.emplace_back(pw, i, cnt);
    }
}
if(Mod > 1) fac.emplace_back(Mod, Mod, 1);
vector<VL> as;
size_t n = 0;
for(auto &&x: fac) {
    ll mod, p, e;
    VL a, b;

```

```

    tie(mod, p, e) = x;
    auto ss = s;
    for(auto &&x: ss) x %= mod;
    tie(a, b) = prime_power(ss, mod, p, e);
    as.emplace_back(a);
    n = max(n, a.size());
}
VL a(n);
vector<pll> c(as.size());
for(size_t i = 0; i < n; i++) {
    for(size_t j = 0; j < as.size(); j++) {
        c[j].first = get<0>(fac[j]);
        c[j].second = i < as[j].size() ? as[j][i] : 0;
    }
    a[i] = crt::crt(c).second;
}
return a;
}

VL BM(const VL &s, ll Mod) {
    VL C(1, 1), B(1, 1);
    int L = 0, m = 1, b = 1;
    for(size_t n = 0; n < s.size(); n++) {
        ll d = 0;
        for(int i = 0; i <= L; i++) {
            d = (d + C[i] * s[n - i]) % Mod;
        }
        if(d == 0) ++m;
        else if(2 * L <= int(n)) {
            VL T = C;
            ll inv = crt::exgcd(b, Mod).first;
            if(inv < 0) inv += Mod;
            ll c = Mod - d * inv % Mod;
            while(C.size() < B.size() + m) C.push_back(0);
            for(size_t i = 0; i < B.size(); i++) C[i + m] = (C[i + m] + c * B[i]) % Mod;
            L = n + 1 - L, B = T, b = d, m = 1;
        } else {
            ll inv = crt::exgcd(b, Mod).first;
            if(inv < 0) inv += Mod;
            ll c = Mod - d * inv % Mod;
            while(C.size() < B.size() + m) C.push_back(0);
            for(size_t i = 0; i < B.size(); i++) C[i + m] = (C[i + m] + c * B[i]) % Mod;
            ++m;
        }
    }
    return C;
}

int m;
VL ini, trans;
ll Mod;
void init(const VL &s, ll md, VL A={}) {
    Mod = md;
    if(A.empty()) {
        A = ReedsSloane(s, Mod);
        // A = BM(s, Mod);
    }
    if(A.empty()) A = {0};
    m = A.size() - 1;
    trans.resize(m);
    for(int i = 0; i < m; i++) {

```



```

        trans[i] = (Mod - A[i + 1]) % Mod;
    }
    reverse(trans.begin(), trans.end());
    ini = {s.begin(), s.begin() + m};
}

ll calc(ll n) {
    if(Mod == 1) return n;
    if(n < m) return ini[n];
    VL v(m), u(m << 1);
    int msk = !!n;
    for(ll m = n; m > 1; m >>= 1) msk <<= 1;
    v[0] = 1 % Mod;
    for(int x = 0; msk; msk >>= 1, x <<= 1) {
        fill_n(u.begin(), m * 2, 0);
        x |= !(n & msk);
        if(x < m) u[x] = 1 % Mod;
        else {
            for(int i = 0; i < m; i++) {
                for(int j = 0, t = i + (x & 1); j < m; j++, t++) {
                    u[t] = (u[t] + v[i] * v[j]) % Mod; // to better
                }
            }
            for(int i = m * 2 - 1; i >= m; i--) {
                for(int j = 0, t = i - m; j < m; j++, t++) {
                    u[t] = (u[t] + trans[j] * u[i]) % Mod; // to better
                }
            }
        }
    }
    v = {u.begin(), u.begin() + m};
}

ll ret = 0;
for(int i = 0; i < m; i++) {
    ret = (ret + v[i] * ini[i]) % Mod; // to better
}
return ret;
}
}

```

线性基

```

typedef unsigned long long ull;
struct LB {
    const static int L = 64; // insert(x) 0 <= x < (1 << L)
    ull a[L];
    LB() { this->init(); }
    void init() { memset(a, 0, sizeof a); }
    ull &operator[](const size_t &id) { return a[id]; }
    const ull &operator[](const size_t &id) const { return a[id]; }
    // 询问x是否在线性基中可以仿造下面的函数来写
    // 即将`return true;`上面三行删去 然后把返回值取反
    // 插入一个数x ==> obj(x) 一边插入一边高斯消元 O(L)
    bool operator()(ull x) {
        for(int i = L - 1; ~i; i--) {
            if((x >> i) & 1) {
                if(!a[i]) {
                    for(int j = 0; j < i; j++) if((x >> j) & 1) x ^= a[j];

```

```

        for(int j = i + 1; j < L; j++) if((a[j] >> i) & 1) a[j] ^= x;
        a[i] = x;
        return true;
    } else {
        x ^= a[i];
    }
}
if(!x) {
    return false;
}
}
return true;
}
// 线性基求交  $O(L^2)$ 
friend LB operator&(const LB &A, const LB &B) {
    LB C, D, E;
    for(int i = L - 1; ~i; i--) {
        if(A[i]) {
            C(A[i]);
            D[i] = 1ull << i;
        }
    }
    for(int i = 0; i < L; i++) {
        if(!B[i]) {
            continue;
        }
        bool can = true;
        ull v = 0, x = B[i];
        for(int j = L - 1; ~j; j--) {
            if((x >> j) & 1) {
                if(C[j]) {
                    x ^= C[j], v ^= D[j];
                } else {
                    can = false, C[j] = x, D[j] = v;
                    break;
                }
            }
        }
    }
    if(can) {
        ll m = 0;
        for(int j = L - 1; ~j; j--) {
            if((v >> j) & 1) {
                m ^= A[j];
            }
        }
        E(m);
    }
    return E;
}
// 线性基求并  $O(L^2)$ 
friend LB operator|(const LB &x, const LB &y) {
    LB z;
    for(int i = 0; i < L; i++) if(x[i]) z(x[i]);
    for(int i = 0; i < L; i++) if(y[i]) z(y[i]);
    return z;
}
};

```

线性基区间最大值

```
namespace LBRMQ {
    const int N = 1e6 + 10, L = 32;
    int b[N][L], pre[N][L];
    void init() {
        memset(b[0], 0, sizeof b[0]);
        memset(pre[0], 0, sizeof pre[0]);
    }
    // index start from 1
    void add(int x, int r) {
        int maxl = r;
        memcpy(b[r], b[r - 1], sizeof(int) * L);
        memcpy(pre[r], pre[r - 1], sizeof(int) * L);
        for(int i = L - 1; ~i; i--) {
            if((x >> i) & 1) {
                if(!b[r][i]) {
                    b[r][i] = x;
                    pre[r][i] = maxl;
                    return;
                }
                if(pre[r][i] < maxl) {
                    swap(pre[r][i], maxl);
                    swap(b[r][i], x);
                }
                x ^= b[r][i];
            }
        }
    }
    int query(int l, int r) {
        int ans = 0;
        for(int i = L - 1; ~i; i--) {
            if(pre[r][i] >= l) {
                ans = max(ans, ans ^ b[r][i]);
            }
        }
        return ans;
    }
}
```

多项式全家桶

```
// 引用 exgcd 和 fft::multiply_mod
// 确保所有输入的数在[0, MOD)的区间中
// MOD < 1073741823 以及 最好是质数
// 传入的vector为{a_0, a_1, a_2, ..., a_n} 即认定为  $y = \sum_{i=0}^n a_i \cdot x^i$ 
namespace poly {
    int MOD = 998244353ll;
    vector<int> inv(const vector<int> &a) {
        if(a.size() == 1) {
            const int inv = exgcd(a[0], MOD).first;
            return vector<int>(1, inv < 0 ? inv + MOD : inv);
        }
        const int na = a.size(), nb = (na + 1) >> 1;
        vector<int> b(a.begin(), a.begin() + nb);
        b = inv(b);
    }
}
```

```

vector<int> c = fft::multiply_mod(b, b, MOD);
c.resize(na);
c = fft::multiply_mod(a, c, MOD);
b.resize(na), c.resize(na);
for(int i = 0; i < na; i++) {
    c[i] = (((211 * b[i] - c[i]) % MOD) + MOD) % MOD;
}
return c;
}

// A = B * C + D (mod x^n) (n = A.size())
// always use with the next function mod
// make sure A.size() >= B.size() or else it will return an empty vector
vector<int> divide(const vector<int> &a, const vector<int> &b) {
    const int n = a.size(), m = b.size();
    if(n < m) return {};
    vector<int> A(a), B(b);
    reverse(A.begin(), A.end()), reverse(B.begin(), B.end());
    A.resize(n - m + 1), B.resize(n - m + 1);
    B = inv(B);
    vector<int> C = fft::multiply_mod(A, B, MOD);
    C.resize(n - m + 1), reverse(C.begin(), C.end());
    return C;
}

vector<int> mod(const vector<int> &a, const vector<int> &b, const vector<int> &c) {
    const int n = a.size(), m = b.size();
    if(n < m) return a;
    vector<int> e = fft::multiply_mod(b, c, MOD);
    e.resize(m - 1);
    for(int i = 0; i < m - 1; i++) {
        e[i] = a[i] - e[i];
        if(e[i] < 0) {
            e[i] += MOD;
        }
    }
    return e;
}

// 构造一个多项式  $\prod_{i=left}^{right} (x - \text{vec}_i)$ 
vector<int> buildPoly(const vector<int> &vec, const int left, const int right) {
    if(left == right) {
        vector<int> ret;
        ret.push_back(MOD - vec[left]);
        ret.push_back(1);
        return ret;
    }
    const int mid = (left + right) >> 1;
    return fft::multiply_mod(buildPoly(vec, left, mid), buildPoly(vec, mid + 1,
right), MOD);
}

void multipointCalc(const vector<int> &poly, const vector<int> &vec, const int left,
const int right, vector<int> &ret) {
    const int n = poly.size(), mid = (left + right) >> 1;
    if(n == 1) {
        for(int i = left; i <= right; i++) {
            ret[i] = poly[0];
        }
        return;
    }
}

```

```

const vector<int> b0 = buildPoly(vec, left, mid);
multipointCalc(mod(poly, b0, divide(poly, b0)), vec, left, mid, ret);
if(left != right) {
    const vector<int> b1 = buildPoly(vec, mid + 1, right);
    multipointCalc(mod(poly, b1, divide(poly, b1)), vec, mid + 1, right, ret);
}
}

// 多点求值
vector<int> multipointCalc(const vector<int> &poly, const vector<int> &vec) {
    const int n = vec.size();
    vector<int> ret(n);
    multipointCalc(poly, vec, 0, n - 1, ret);
    return ret;
}

vector<int> multiInv(const vector<int> &vec) {
    const int n = vec.size();
    vector<int> a(n + 1), ret(n); a[0] = 1;
    for(int i = 1; i <= n; i++) {
        a[i] = 111 * a[i - 1] * vec[i - 1] % MOD;
    }
    int cur = (exgcd(a[n], MOD).first + MOD) % MOD;
    for(int i = n - 1; i >= 0; i--) {
        ret[i] = 111 * cur * a[i] % MOD;
        cur = 111 * cur * vec[i] % MOD;
    }
    return ret;
}

// 快速插值 {{x0, y0}, {x1, y1}, {x2, y2}, ...}
vector<int> interpolate(const vector<pair<int, int>> &p) {
    const int n = p.size(), n0 = (n + 1) >> 1, n1 = n - n0;
    if(n == 1) {
        return {p[0].second};
    }
    vector<pair<int, int>> p0(p.begin() + n1, p.end());
    vector<int> f0 = interpolate(p0);
    vector<int> x(n);
    for(int i = 0; i < n; i++) {
        x[i] = p[i].first;
    }
    vector<int> g0 = buildPoly(x, n1, n - 1);
    x.resize(n1);
    vector<int> fx = multipointCalc(f0, x), gx = multipointCalc(g0, x);
    gx = multiInv(gx);
    p0.resize(n1);
    for(int i = 0; i < n1; i++) {
        p0[i].first = p[0].first;
        p0[i].second = (p[i].second - fx[i] + MOD) % MOD;
        p0[i].second = 111 * p0[i].second * gx[i] % MOD;
    }
    fx = interpolate(p0);
    fx = fft::multiply_mod(fx, g0, MOD);
    fx.resize(n), f0.resize(n);
    for(int i = 0; i < n; i++) {
        fx[i] = (fx[i] + f0[i]) % MOD;
    }
    return fx;
}
}

```

```

namespace SA {
    const size_t sz = 3e5 + 5;
    int bucket[sz], bucket1[sz], sa[sz], rk[sz], ht[sz];
    bool type[sz << 1];
    bool isLMS(const int i, const bool *type) {
        return i > 0 && type[i] && !type[i - 1];
    }

    template<class T>
    void inducedSort(const T &s, int *sa, const int len, const int sigma, const int
bucketSize, bool *type, int *bucket, int *cntbuf, int *p) {
        memset(bucket, 0, sizeof(int) * sigma);
        memset(sa, -1, sizeof(int) * len);
        for (int i = 0; i < len; i++) {
            bucket[s[i]]++;
        }
        cntbuf[0] = bucket[0];
        for (int i = 1; i < sigma; i++) {
            cntbuf[i] = cntbuf[i - 1] + bucket[i];
        }
        for (int i = bucketSize - 1; i >= 0; i--) {
            sa[--cntbuf[s[p[i]]]] = p[i];
        }
        for (int i = 1; i < sigma; i++) {
            cntbuf[i] = cntbuf[i - 1] + bucket[i - 1];
        }
        for (int i = 0; i < len; i++) {
            if (sa[i] > 0 && !type[sa[i] - 1]) {
                sa[cntbuf[s[sa[i] - 1]]++] = sa[i] - 1;
            }
        }
        cntbuf[0] = bucket[0];
        for (int i = 1; i < sigma; i++) {
            cntbuf[i] = cntbuf[i - 1] + bucket[i];
        }
        for (int i = len - 1; i >= 0; i--) {
            if (sa[i] > 0 && type[sa[i] - 1]) {
                sa[--cntbuf[s[sa[i] - 1]]] = sa[i] - 1;
            }
        }
    }

    template<class T>
    void sais(const T &s, int *sa, int len, bool *type, int *bucket, int *bucket1, int
sigma) {
        int i, j, bucketSize = 0, cnt = 0, p = -1, x, *cntbuf = bucket + sigma;
        type[len - 1] = 1;
        for (i = len - 2; i >= 0; i--) {
            type[i] = s[i] < s[i + 1] || (s[i] == s[i + 1] && type[i + 1]);
        }
        for (i = 1; i < len; i++) {
            if (type[i] && !type[i - 1]) {
                bucket1[bucketSize++] = i;
            }
        }
        inducedSort(s, sa, len, sigma, bucketSize, type, bucket, cntbuf, bucket1);
        for (i = bucketSize = 0; i < len; i++) {
            if (isLMS(sa[i], type)) {
                sa[bucketSize++] = sa[i];
            }
        }
    }
}

```

```

    }
}
for (i = bucketSize; i < len; i++) {
    sa[i] = -1;
}
for (i = 0; i < bucketSize; i++) {
    x = sa[i];
    for (j = 0; j < len; j++) {
        if (p == -1 || s[x + j] != s[p + j] || type[x + j] != type[p + j]) {
            cnt++, p = x;
            break;
        } else if (j > 0 && (isLMS(x + j, type) || isLMS(p + j, type))) {
            break;
        }
    }
    x = (~x & 1 ? x >> 1 : (x - 1) >> 1), sa[bucketSize + x] = cnt - 1;
}
for (i = j = len - 1; i >= bucketSize; i--) {
    if (sa[i] >= 0) {
        sa[j--] = sa[i];
    }
}
int *s1 = sa + len - bucketSize, *bucket2 = bucket1 + bucketSize;
if (cnt < bucketSize) {
    sais(s1, sa, bucketSize, type + len, bucket, bucket1 + bucketSize, cnt);
} else {
    for (i = 0; i < bucketSize; i++) {
        sa[s1[i]] = i;
    }
}
for (i = 0; i < bucketSize; i++) {
    bucket2[i] = bucket1[sa[i]];
}
inducedSort(s, sa, len, sigma, bucketSize, type, bucket, cntbuf, bucket2);
}

void getHeight(const vector<int> &s, int n) {
    int i, j, k = 0;
    for(i = 1; i <= n; i++) {
        rk[sa[i]] = i;
    }
    for(i = 0; i < n; i++) {
        if(k) k--;
        for(j = sa[rk[i] - 1]; s[i + k] == s[j + k]; k++);
        ht[rk[i]] = k;
    }
}

template<class T>
void solve(const T &s) {
    const int n = s.size();
    vector<int> v(n + 1);
    int sigma = 0;
    for(int i = 0; i < n; i++) {
        v[i] = int(s[i]) + 1;
        sigma = max(sigma, v[i] + 1);
    }
    v[n] = 0;
    sais(v, sa, n + 1, type, bucket, bucket1, sigma);
    getHeight(v, n);
}

```

```
}
```

旋转卡壳

```
const int N=1e5+7;
struct point{
    double x,y;
}S[N];
inline double Cross(point a,point b,point c){
    return (a.x-c.x)*(b.y-c.y)-(a.y-c.y)*(b.x-c.x);
}
inline double Dis(point a,point b){
    return sqrt((a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y));
}
//先求凸包再旋转卡壳
double FarthestPointPair(int top){
    if(top==1)return Dis(S[0],S[1]);
    S[++top]=S[0];int j=2;double ans=0;
    for(int i=0;i<top;++i){
        while(Cross(S[i],S[i+1],S[j])<Cross(S[i],S[i+1],S[j+1])) {
            j=(j+1)%top;
        }
        ans=max(ans,max(Dis(S[i],S[j]),Dis(S[i+1],S[j])));
    }
    return ans;
}
```

化行最简形

```
void Gauss(vector<vector<int>> &v) {
    const int m = v.size(), n = v[0].size();
    vector<int> id(n, -1);
    int i = 0, j = 0;
    while(i < m && j < n) {
        int mi = i;
        for(int k = i + 1; k < m; k++) {
            if(v[k][j] > v[mi][j]) {
                mi = k;
            }
        }
        if(v[mi][j] != 0) {
            if(i != mi) {
                for(int k = 0; k < n; k++) {
                    swap(v[i][k], v[mi][k]);
                }
            }
            const int inv = fpow(v[i][j], MOD - 2);
            for(int k = j + 1; k < n; k++) {
                v[i][k] = 1ll * v[i][k] * inv % MOD;
            }
            id[j] = i;
            v[i][j] = 1;
            for(int r = i + 1; r < m; r++) {
                for(int c = j + 1; c < n; c++) {
```



```

        v[r][c] -= 111 * v[r][j] * v[i][c] % MOD;
        if(v[r][c] < 0) v[r][c] += MOD;
    }
    v[r][j] = 0;
}
i++;
}
j++;
}
for(int i = 0; i < m; i++) {
    for(int j = 0; j < n; j++) {
        if(v[i][j] && id[j] > i) {
            for(int k = j + 1; k < n; k++) {
                v[i][k] -= 111 * v[i][j] * v[id[j]][k] % MOD;
                if(v[i][k] < 0) v[i][k] += MOD;
            }
            v[i][j] = 0;
        }
    }
}
}
}

```

min_25

*min_25*用途

1. 筛出质数

2. 求出所有的 $\sum_{i=2}^{\lfloor \frac{n}{x} \rfloor} [i \text{ 是质数}] f(i)$

3. 求积性函数前缀和

前提

- 当 i 为质数时 $f(i)$ 需要是一个多项式。
- 对于求积性函数前缀和而言 $f(p^k)$ 需要快速求出，求多个值的时候一般不适用 `min_25` 筛。

计算

处理质数

$$g(a, b) = \sum_{i=2}^a [i \text{ 是质数 或 } p \min_i > \text{prime}_b] * i^k$$

需要求每一个

$$g(\lfloor \frac{n}{x} \rfloor, \infty) = \sum_{i=2}^{\lfloor \frac{n}{x} \rfloor} [i \text{ 是质数}] f(i)$$

那么有

$$g(a, b) = \begin{cases} g(a, b-1), & a < prime_b^2 \\ g(a, b-1) - prime_b^k \left(g(\lfloor \frac{a}{prime_b} \rfloor, b-1) - g(prime_{b-1}, b-1) \right), & a \geq prime_b^2 \end{cases}$$

滚动数组叠上去即可求出。

计算前缀和

$$S(a, b) = \sum_{i=2}^a [pmin_i \geq prime_b] f(i)$$

前缀和即

$$\sum_{i=1}^n f(i) = S(n, 1) + f(1)$$

那么有

$$S(a, b) = \begin{cases} 0, & a < prime_b \\ g(a, \infty) - g(prime_{b-1}, \infty) + \sum_{i=b}^{\infty} \sum_{t \geq 1, prime_i^{t+1} \leq a} \left(S(\lfloor \frac{a}{prime_i^t} \rfloor, i+1) * f(prime_i^t) + f(prime_i^{t+1}) \right), & a \geq prime_b \end{cases}$$

无需记忆化，递归求解。

防忘代码

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
void Main();
#ifdef ConanYu
#include "local.hpp"
#else
#define debug(...) do { } while(0)
int main() {
    ios::sync_with_stdio(false), cin.tie(0), cout.tie(0);
    Main();
}
#endif

const int MOD = 1e9 + 7;
int fpow(int a, int b) {
    int ans = 1;
    for(; b >= 1, a = 1ll * a * a % MOD) {
        if(b & 1) ans = 1ll * ans * a % MOD;
    }
    return ans;
}

const int N = 1e5 + 10, INV2 = fpow(2, MOD - 2);
ll n, a[N << 1];
int sn, cnt, id, prime[N], g[N << 1], h[N << 1];
int idx(ll x) {
    return x <= sn ? x : id - n / x + 1;
}

void sub(int &a, int b) {
    a -= b;
```

```

    if(a < 0) a += MOD;
}

int solve(long long a, int b) {
    if(a < prime[b]) return 0;
    int ans = g[idx(a)];
    sub(ans, g[idx(prime[b - 1])]);
    for(int i = b; i <= cnt && a / prime[i] >= prime[i]; i++) {
        ll k = prime[i], m = k - 1;
        while(a / k >= prime[i]) {
            ans += 1ll * solve(a / k, i + 1) * m % MOD;
            if(ans >= MOD) ans -= MOD;
            m = m * prime[i] % MOD;
            ans += m;
            if(ans >= MOD) ans -= MOD;
            k *= prime[i];
        }
    }
    return ans;
}

void Main() {
    cin >> n;
    sn = sqrt(n);
    cnt = id = 0;
    for(ll i = 1; i <= n; i = a[id] + 1) {
        a[++id] = n / (n / i);
        g[id] = a[id] % MOD * ((a[id] + 1) % MOD) % MOD * INV2 % MOD;
        sub(g[id], 1);
        h[id] = (a[id] - 1) % MOD;
    }
    for(int i = 2; i <= sn; i++) {
        if(h[i] != h[i - 1]) {
            prime[++cnt] = i;
            for(int j = id; a[j] / i >= i; j--) {
                const int ta = idx(a[j] / i), tb = idx(prime[cnt - 1]);
                int A = g[ta], B = h[ta];
                sub(A, g[tb]), sub(B, h[tb]);
                sub(g[j], 1ll * i * A % MOD);
                sub(h[j], B);
            }
        }
    }
    for(int i = 1; i <= id; i++) {
        g[i] -= h[i];
        if(g[i] < 0) g[i] += MOD;
    }
    cout << ((solve(n, 1) + 1) % MOD) << "\n";
}

```