

TABLE 1 The Truth Table for the Negation of a Proposition.

p	$\neg p$
T	F
F	T

TABLE 2 The Truth Table for the Conjunction of Two Propositions.

p	q	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

TABLE 3 The Truth Table for the Disjunction of Two Propositions.

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

TABLE 4 The Truth Table for the Exclusive Or of Two Propositions.

p	q	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

TABLE 5 The Truth Table for the Conditional Statement $p \rightarrow q$.

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

TABLE 2 De Morgan's Laws.

$$\begin{aligned}\neg(p \wedge q) &\equiv \neg p \vee \neg q \\ \neg(p \vee q) &\equiv \neg p \wedge \neg q\end{aligned}$$

TABLE 6 Logical Equivalences.

Equivalence	Name
$p \wedge T \equiv p$ $p \vee F \equiv p$	Identity laws
$p \vee T \equiv T$ $p \wedge F \equiv F$	Domination laws
$p \vee p \equiv p$ $p \wedge p \equiv p$	Idempotent laws
$\neg(\neg p) \equiv p$	Double negation law
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative laws
$(p \vee q) \vee r \equiv p \vee (q \vee r)$ $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	Associative laws
$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \wedge q) \equiv \neg p \vee \neg q$ $\neg(p \vee q) \equiv \neg p \wedge \neg q$	De Morgan's laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Absorption laws
$p \vee \neg p \equiv T$ $p \wedge \neg p \equiv F$	Negation laws

Proof by Contraposition

Making use of the fact that the conditional statement $p \rightarrow q$ is equivalent to its contrapositive, $\neg q \rightarrow \neg p$

e.g. in $p \rightarrow q$,

We get the contrapositive $\neg q \rightarrow \neg p$

Take $\neg q$ as premise

we show that $\neg p$ must follow

Proof of Equivalence

To prove a theorem that is a biconditional statement $p \leftrightarrow q$

we show that $p \rightarrow q$ and $q \rightarrow p$ are both true.

Based on the tautology $(p \leftrightarrow q) \leftrightarrow (p \rightarrow q) \wedge (q \rightarrow p)$

TABLE 5 The Truth Table for the Conditional Statement $p \rightarrow q$.**TABLE 6** The Truth Table for the Biconditional $p \leftrightarrow q$.

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

TABLE 8 Precedence of Logical Operators.

Operator	Precedence
\neg	1
\wedge	2
\vee	3
\rightarrow	4
\leftrightarrow	5

Soal:

Pernyataan:

"Jika n^2 habis dibagi 3, maka n habis dibagi 9."

Counterexample Proof:

- Ambil $n = 3$
- $n^2 = 9 \rightarrow$ habis dibagi 3 ✓
- Tapi $n = 3 \rightarrow$ tidak habis dibagi 9 ✗
- Jadi pernyataan salah, dan $n = 3$ adalah counterexample ✓

TABLE 7 Logical Equivalences Involving Conditional Statements.

$p \rightarrow q \equiv \neg p \vee q$
$p \rightarrow q \equiv \neg q \rightarrow \neg p$
$p \vee q \equiv \neg p \rightarrow q$
$p \wedge q \equiv \neg(p \rightarrow \neg q)$
$\neg(p \rightarrow q) \equiv p \wedge \neg q$
$(p \rightarrow q) \wedge (p \rightarrow r) \equiv p \rightarrow (q \wedge r)$
$(p \rightarrow r) \wedge (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$
$(p \rightarrow q) \vee (p \rightarrow r) \equiv p \rightarrow (q \vee r)$
$(p \rightarrow r) \vee (q \rightarrow r) \equiv (p \vee q) \rightarrow r$

TABLE 8 Logical Equivalences Involving Biconditional Statements.

$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$
$p \leftrightarrow q \equiv \neg p \leftrightarrow \neg q$
$p \leftrightarrow q \equiv (p \wedge q) \vee (\neg p \wedge \neg q)$
$\neg(p \leftrightarrow q) \equiv p \leftrightarrow \neg q$

TABLE 4 The Truth Table for the Exclusive Or of Two Propositions.

p	q	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

Definition

Another way to create a proposition from propositional function

In English words

all, some, many, none, and few

Types

- Universal quantifier → For all values... → $\forall x P(x)$
- Existential quantifier → For some values in... → $\exists x P(x)$
- Unique quantifier → For one values in... → $\exists! x P(x)$

Conclusion as $\neg \forall x P(x) \equiv \exists x \neg P(x)$

Negation of Quantifiers

Every student in your class has taken a course in calculus

True

Pernyataan: Ada tepat satu bilangan bulat x sehingga $2x + 3 = 11$.

Penyelesaian (Unique Proof):

De Morgan's Law

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

p	q	$\neg p$	$\neg q$	$\neg(p \wedge q)$	$\neg p \vee \neg q$
0	0	1	1	0	1
0	1	1	0	1	0
1	0	0	1	1	0
1	1	0	0	1	0

Soal:
Carilah x yang memenuhi: $2x + 3 = 11 \Rightarrow 2x = 8 \Rightarrow x = 4$
Jadi ada $x = 4$ yang memenuhi ✓

2. Uniqueness:
Misalkan ada y lain yang juga memenuhi: $2y + 3 = 11 \Rightarrow 2y = 8 \Rightarrow y = 4$
Jadi $y = x = 4$
Tidak mungkin ada bilangan lain yang memenuhi ✓

Kesimpulan: Ada tepat satu bilangan bulat $x = 4$ ✓

Proofs**Vacuous Proof**

$p \rightarrow q$ is true if p is false

p	q	$p \rightarrow q$
0	0	1
0	1	1
1	0	0
1	1	1

Symbol \rightarrow

Symbol \neg

Symbol \wedge

Symbol \vee

Symbol \rightarrow

TABLE 1 Rules of Inference.

Rule of Inference	Tautology	Name
$\frac{p}{p \rightarrow q}$ p $\therefore q$	$(p \wedge (p \rightarrow q)) \rightarrow q$	Modus ponens
$\frac{\neg q}{\neg p \rightarrow q}$ $\neg q$ $p \rightarrow q$ $\therefore \neg p$	$(\neg q \wedge (p \rightarrow q)) \rightarrow \neg p$	Modus tollens
$\frac{p \rightarrow q}{p \rightarrow q}$ $q \rightarrow r$ $\therefore p \rightarrow r$	$((p \rightarrow q) \wedge (q \rightarrow r)) \rightarrow (p \rightarrow r)$	Hypothetical syllogism
$\frac{p \vee q}{\neg p}$ $\therefore q$	$((p \vee q) \wedge \neg p) \rightarrow q$	Disjunctive syllogism
$\frac{p}{p}$ $\therefore p \vee q$	$p \rightarrow (p \vee q)$	Addition
$\frac{p \wedge q}{p}$ $\therefore p$	$(p \wedge q) \rightarrow p$	Simplification
$\frac{p}{p}$ q $\therefore p \wedge q$	$((p \wedge q) \wedge (p \wedge q)) \rightarrow (p \wedge q)$	Conjunction
$\frac{p \vee q}{p \vee q}$ $\neg p \vee r$ $\therefore q \vee r$	$((p \vee q) \wedge (\neg p \vee r)) \rightarrow (q \vee r)$	Resolution

Rule of Inference	Name
$\frac{}{\forall x P(x)}$ $\therefore P(c)$	Universal instantiation
$P(c)$ for an arbitrary c $\therefore \forall x P(x)$	Universal generalization
$\frac{}{\exists x P(x)}$ $\therefore P(c)$ for some element c	Existential instantiation
$P(c)$ for some element c $\therefore \exists x P(x)$	Existential generalization

Soal:

Buktikan bahwa untuk setiap bilangan bulat n dengan $1 \leq n \leq 4$, $n^2 \leq 16$.

Penyelesaian (exhaustive proof):

Periksa semua nilai $n = 1, 2, 3, 4$:

1. $n = 1: 1^2 = 1 \leq 16 \checkmark$
2. $n = 2: 2^2 = 4 \leq 16 \checkmark$
3. $n = 3: 3^2 = 9 \leq 16 \checkmark$
4. $n = 4: 4^2 = 16 \leq 16 \checkmark$

"Jika n adalah bilangan bulat, maka $n^2 + n$ adalah genap."

Penyelesaian (Proof by Cases):

1. Kasus 1: n genap, misal $n = 2k$
 - $n^2 + n = (2k)^2 + 2k = 4k^2 + 2k = 2(2k^2 + k) \checkmark$ genap
2. Kasus 2: n ganjil, misal $n = 2k + 1$
 - $n^2 + n = (2k + 1)^2 + (2k + 1) = 4k^2 + 4k + 1 + 2k + 1 = 4k^2 + 6k + 2 = 2(2k^2 + 3k + 1) \checkmark$ genap

Karena kedua kasus menunjukkan hasil genap, pernyataan benar untuk semua bilangan bulat n .

Contoh Soal 3 (lebih abstrak)

Soal:

Buktikan:

"Untuk setiap bilangan bulat positif n , ada bilangan bulat positif m sehingga $m > n$."

Penyelesaian (Construction Proof):

- Cara konstruksi: ambil $m = n + 1$
- Jelas $m > n$ dan m positif \checkmark

Kita menunjukkan secara langsung cara membuat m yang memenuhi, jadi bukti selesai.

Soal:

Buktikan:

"Ada bilangan prima yang lebih besar dari 1000."

Penyelesaian (Non-Construction Proof):

- Kita tahu bilangan prima tak terbatas (Teorema Euclid).
- Maka pasti ada bilangan prima yang lebih besar dari 1000. \checkmark
- Kita tidak menulis bilangan primanya, cukup argumen eksistensi.

Soal:

Buktikan:

Teks soal. Use a proof by contraposition to show: if $x + y \geq 2$ (real x, y) then $x \geq 1$ or $y \geq 1$.
 $\neg x \geq 1 \quad \neg y \geq 1$
 $\neg x \rightarrow \neg P \quad \neg y \rightarrow \neg P$
 $\neg x \wedge \neg y \rightarrow \neg P$
 $\neg x \wedge \neg y \rightarrow x + y < 2$ $\neg P: x + y < 2$

Step 1	Kalimat	Simbol
It rains	R	
It is foggy	F	
The sailing race will be held	S	
The lifesaving demonstration will go on	L	
The trophy will be awarded	T	

Step 2	Kalimat	Simbol
"If it does not rain or if it is not foggy, then the sailing race will be held and the lifesaving demonstration will go on"	$(\neg R \vee \neg F) \rightarrow (S \wedge L)$	
"If the sailing race is held, then the trophy will be awarded"	$S \rightarrow T$	
"The trophy was not awarded"	$\neg T$	
"It rained"	R	

Step 1	Kalimat	Simbol
x is a student in this class	$C(x)$	
x has read the book	$B(x)$	
x passed the first exam	$P(x)$	

Step 2	Kalimat	Simbol
"A student in this class has not read the book"	$(\exists x) B(\neg x)$	
"Everyone in this class passed the first exam"	$\forall x (C(x) \rightarrow P(x))$	
"Someone who passed the first exam has not read the book"	$(\exists x) P(x) \wedge B(\neg x)$	

Step 1	Kalimat	Simbol
" $\neg x \vee \neg B(x)$ "	$\neg x \vee \neg B(x)$	Expansion (De Morgan's Law)
" $\neg B(x) \rightarrow C(x)$ "	$\neg B(x) \rightarrow C(x)$	Universal instantiation (Generalization)
" $\neg B(x) \rightarrow \neg B(y)$ "	$\neg B(x) \rightarrow \neg B(y)$	Conditional reasoning (6) and (7)

Step 2	Kalimat	Simbol
" $\neg B(x) \rightarrow \neg B(y)$ "	$\neg B(x) \rightarrow \neg B(y)$	Universal instantiation (Generalization)
" $\neg B(x) \rightarrow C(x)$ "	$\neg B(x) \rightarrow C(x)$	Universal instantiation (Generalization)
" $\neg B(x) \rightarrow \neg B(y) \wedge \neg B(x) \rightarrow C(x)$ "	$\neg B(x) \rightarrow \neg B(y) \wedge \neg B(x) \rightarrow C(x)$	Universal instantiation (Generalization)

Step 3	Alasan
1	Premis
2	Premis
3	Modus Tollens (1) dan (2)
4	Premis
5	Contrapositive
6	De Morgan's Law
7	Addition Rule dari $\neg S$ (3)
8	Modus Ponens (6) dan (7)
9	Simplification

Buktikan:

"Jika $x + y \geq 2$, maka $x \geq 1$ atau $y \geq 1$ ".

Penyelesaian (Backward Reasoning):

1. Tujuan: $Buktikan x + y \geq 2 \rightarrow x \geq 1$
2. Diketahui: $x + 1 \geq 2$
3. Maka $x = (x + 1) - 1 = 2m + 1$

"Jika $x + 1 \geq 2$, maka $x \geq 1$ ".

Penyelesaian (Forward Reasoning):

1. Diketahui: $x \geq 1$
2. Hitung x^2 :

$$x^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1$$

3. Karena $x^2 = 2$ (bilangan bulat) + 1, maka x^2 ganjil \checkmark

Penalaran maju ini langsung dari fakta ke kesimpulan tanpa mundur dari tujuan.