Homework 7: Intro to Parsing and Context-Free Grammars

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October 13, 2025

Homework Problems

Using the context-free grammar:

$$\text{Exp} \rightarrow \text{Exp '+' Exp1} \tag{1}$$
$$\text{Exp1} \rightarrow \text{Exp1 '*' Exp2} \tag{2}$$
$$\text{Exp2} \rightarrow \text{Integer} \tag{3}$$
$$\text{Exp2} \rightarrow \text{'(' Exp ')'} \tag{4}$$
$$\text{Exp} \rightarrow \text{Exp1} \tag{5}$$
$$\text{Exp1} \rightarrow \text{Exp2} \tag{6}$$

Problem 1: Derivation Trees

Write out the derivation trees (also called parse trees or concrete syntax trees) for the following strings:

(a) $2 + 1$

(b) $1 + 2 * 3$

(c) $1 + (2 * 3)$

(d) $(1 + 2) * 3$

(e) $1 + 2 * 3 + 4 * 5 + 6$

Problem 2: Unparsable Strings

Why do the following strings not have parse trees (given the context-free grammar above)?

(a) $2 - 1$

(b) $1.0 + 2$

(c) $6/3$

(d) $8 \bmod 6$

Problem 3: Parse Tree Uniqueness

With the simplified grammar without precedence levels:

$$\text{Exp} \rightarrow \text{Exp '+' Exp} \tag{7}$$
$$\text{Exp} \rightarrow \text{Exp '*' Exp} \tag{8}$$
$$\text{Exp} \rightarrow \text{Integer} \tag{9}$$

How many parse trees can you find for the following expressions?

(a) $1 + 2 + 3$

(b) $1 * 2 * 3 * 4$

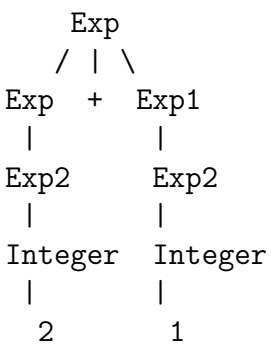
Answer the question above using instead the grammar:

$$\text{Exp} \rightarrow \text{Exp '+' Exp1} \tag{10}$$
$$\text{Exp} \rightarrow \text{Exp1} \tag{11}$$
$$\text{Exp1} \rightarrow \text{Exp1 '*' Exp2} \tag{12}$$
$$\text{Exp1} \rightarrow \text{Exp2} \tag{13}$$
$$\text{Exp2} \rightarrow \text{Integer} \tag{14}$$

Solutions

Solution 1: Derivation Trees

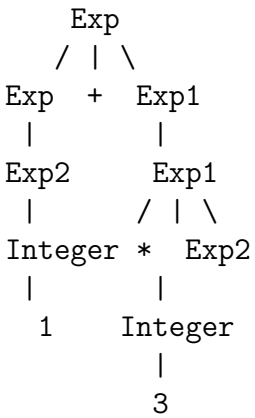
(a) Derivation tree for $2 + 1$:



Derivation steps:

$$\begin{aligned}
\text{Exp} &\rightarrow \text{Exp '+' Exp1} && (15)\\
&\rightarrow \text{Exp2 '+' Exp1} && (16)\\
&\rightarrow \text{Integer '+' Exp1} && (17)\\
&\rightarrow \text{Integer '+' Exp2} && (18)\\
&\rightarrow \text{Integer '+' Integer} && (19)\\
&\rightarrow \text{'2' '+' '1'} && (20)
\end{aligned}$$

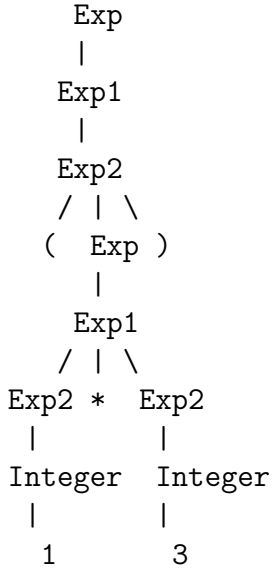
(b) Derivation tree for $1 + 2 * 3$:



Derivation steps:

$$\begin{aligned}
\text{Exp} &\rightarrow \text{Exp '+' Exp1} && (21)\\
&\rightarrow \text{Exp2 '+' Exp1} && (22)\\
&\rightarrow \text{Integer '+' Exp1} && (23)\\
&\rightarrow \text{Integer '+' Exp1 '*' Exp2} && (24)\\
&\rightarrow \text{Integer '+' Exp2 '*' Exp2} && (25)\\
&\rightarrow \text{Integer '+' Integer '*' Exp2} && (26)\\
&\rightarrow \text{Integer '+' Integer '*' Integer} && (27)\\
&\rightarrow \text{'1' '+' '2' '*' '3'} && (28)
\end{aligned}$$

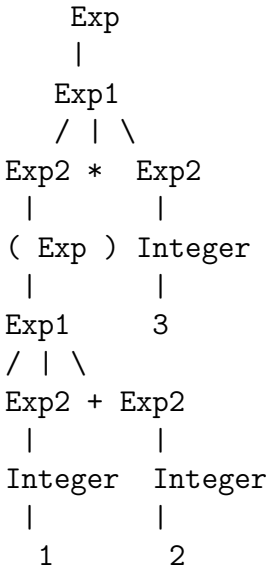
(c) Derivation tree for $1 + (2 * 3)$**:**



Derivation steps:

$$\begin{aligned}
\text{Exp} &\rightarrow \text{Exp1} && (29)\\
&\rightarrow \text{Exp2} && (30)\\
&\rightarrow \text{'(' Exp ')'} && (31)\\
&\rightarrow \text{'(' Exp1 ')'} && (32)\\
&\rightarrow \text{'(' Exp1 '*' Exp2 ')'} && (33)\\
&\rightarrow \text{'(' Exp2 '*' Exp2 ')'} && (34)\\
&\rightarrow \text{'(' Integer '*' Exp2 ')'} && (35)\\
&\rightarrow \text{'(' Integer '*' Integer ')'} && (36)\\
&\rightarrow \text{'(' '1' '*' '3' ')'} && (37)
\end{aligned}$$

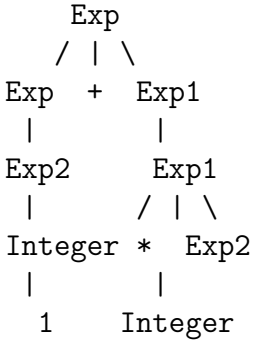
(d) Derivation tree for $(1 + 2) * 3$**:**



Derivation steps:

$$\begin{aligned}
\text{Exp} &\rightarrow \text{Exp1} && (38)\\
&\rightarrow \text{Exp1 '*' Exp2} && (39)\\
&\rightarrow \text{Exp2 '*' Exp2} && (40)\\
&\rightarrow \text{'(' Exp ')' '*' Exp2} && (41)\\
&\rightarrow \text{'(' Exp1 ')' '*' Exp2} && (42)\\
&\rightarrow \text{'(' Exp '+' Exp1 ')' '*' Exp2} && (43)\\
&\rightarrow \text{'(' Exp2 '+' Exp1 ')' '*' Exp2} && (44)\\
&\rightarrow \text{'(' Integer '+' Exp1 ')' '*' Exp2} && (45)\\
&\rightarrow \text{'(' Integer '+' Exp2 ')' '*' Exp2} && (46)\\
&\rightarrow \text{'(' Integer '+' Integer ')' '*' Exp2} && (47)\\
&\rightarrow \text{'(' Integer '+' Integer ')' '*' Integer} && (48)\\
&\rightarrow \text{'(' '1' '+' '2' ')' '*' '3'} && (49)
\end{aligned}$$

(e) Derivation tree for $1 + 2 * 3 + 4 * 5 + 6$**:**
This is a complex expression. The tree would be:



```
|
3
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Actually, let me be more careful. The full derivation would be quite large, but the key insight is that this parses as $1 + (2 * 3) + (4 * 5) + 6$ due to the precedence rules in the grammar.

Solution 2: Unparsable Strings

The following strings cannot be parsed because the grammar only defines rules for addition ($+$) and multiplication ($*$), but not for:

(a) $2 - 1$**:** The grammar has no rule for subtraction ($-$).

(b) $1.0 + 2$**:** The grammar only handles integers, not decimal numbers like 1.0.

(c) $6/3$**:** The grammar has no rule for division ($/$).

(d) 8 mod 6**:** The grammar has no rule for the modulo operation.

To make these strings parsable, we would need to add new rules to the grammar, such as:

$$\text{Exp} \rightarrow \text{Exp '-' Exp1} \tag{50}$$
$$\text{Exp} \rightarrow \text{Exp '/' Exp1} \tag{51}$$
$$\text{Exp} \rightarrow \text{Exp 'mod' Exp1} \tag{52}$$
$$\text{Integer} \rightarrow \text{Float} \tag{53}$$
$$\text{Float} \rightarrow \text{Integer '.' Integer} \tag{54}$$

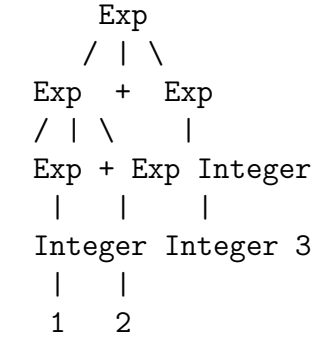
Solution 3: Parse Tree Uniqueness

With the simplified grammar:

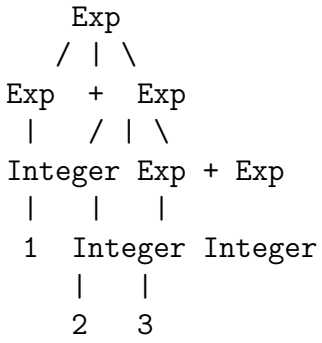
$$\text{Exp} \rightarrow \text{Exp '+' Exp} \tag{55}$$
$$\text{Exp} \rightarrow \text{Exp '*' Exp} \tag{56}$$
$$\text{Exp} \rightarrow \text{Integer} \tag{57}$$

(a) For $1 + 2 + 3$**:** This expression is ambiguous and has 2 different parse trees:

Tree 1: $(1 + 2) + 3$



Tree 2: $1 + (2 + 3)$



(b) For $1 * 2 * 3 * 4$**:** This expression is also ambiguous and has 5 different parse trees corresponding to the different ways of parenthesizing the multiplication.

With the precedence grammar:

$$\text{Exp} \rightarrow \text{Exp '+' Exp1} \tag{58}$$

$$\text{Exp} \rightarrow \text{Exp1} \tag{59}$$

$$\text{Exp1} \rightarrow \text{Exp1 '*' Exp2} \tag{60}$$

$$\text{Exp1} \rightarrow \text{Exp2} \tag{61}$$

$$\text{Exp2} \rightarrow \text{Integer} \tag{62}$$

(a) For $1 + 2 + 3$**:** This grammar forces left associativity, so there is only 1 parse tree: $((1 + 2) + 3)$.

(b) For $1 * 2 * 3 * 4$**:** This grammar also forces left associativity for multiplication, so there is only 1 parse tree: $((1 * 2) * 3) * 4$.

The key difference is that the precedence grammar eliminates ambiguity by using different nonterminals (Exp, Exp1, Exp2) to enforce operator precedence and associativity rules.