

MACHINE m2

REFINES m1

SEES c0

VARIABLES

wait
process
clk
t1
t2
t3
position
index
qsize

INVARIANTS

inv1: $\forall p. (p \in process \wedge p \in dom(t2)) \Rightarrow clk - t2(p) \leq ddl2$
inv2: $\forall p. (p \in dom(t2) \wedge p \in dom(t3) \wedge t3(p) \geq t2(p)) \Rightarrow t3(p) - t2(p) \leq ddl2$
deadline(t2,t3,ddl1)
inv3: $\forall t, p. (p \in process \wedge p \in dom(t1) \wedge p \in dom(t2) \wedge t = t2(p) \wedge t2(p) \geq t1(p)) \Rightarrow t2(p) - t1(p) \leq ddl1$
inv4: $position \in 1 .. qsize \mapsto wait$
inv5: $index \in PROCESS \mapsto \mathbb{N}$
inv6: $qsize \in \mathbb{N}$

EVENTS

Initialisation $\langle \text{extended} \rangle$

begin

act1: $wait := \emptyset$
act2: $process := \emptyset$
act3: $clk := 0$
act4: $t1 := \emptyset$
act5: $t2 := \emptyset$
act6: $t3 := \emptyset$
act7: $position := \emptyset$
act8: $index := \emptyset$
act9: $qsize := 0$

end

Event wish $\langle \text{ordinary} \rangle \hat{=}$

extends wish

any

pro

where

grd1: $pro \in PROCESS \setminus wait$
grd2: $pro \in PROCESS \setminus process$

then

act1: $wait := wait \cup \{pro\}$
act2: $t1(pro) := clk$
act4: $index(pro) := qsize$
act3: $position(qsize + 1) := pro$
act5: $qsize := qsize + 1$

end

Event enter $\langle \text{ordinary} \rangle \hat{=}$

extends enter

any

pro

where

grd1: $pro \in wait$
grd2: $card(process) = 0$

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    grd3:  $qsize > 0$ 
    grd4:  $pro = position(1)$ 
  then
    act1:  $wait := wait \setminus \{pro\}$ 
    act2:  $process := process \cup \{pro\}$ 
    act3:  $t2(pro) := clk$ 
    act5:  $position :| position' \in 1..qsize-1 \mapsto wait \setminus \{position(1)\} \wedge (\forall i \cdot i \in 1..qsize-1 \Rightarrow position'(i) = position(i+1))$ 
    act6:  $qsize := qsize - 1$ 
  end
Event leave  $\langle ordinary \rangle \hat{=}$ 
extends leave
  any
     $pro$ 
  where
    grd1:  $pro \in process$ 
  then
    act1:  $process := process \setminus \{pro\}$ 
    act3:  $t3(pro) := clk$ 
  end
Event tick  $\langle ordinary \rangle \hat{=}$ 
refines tick
  when
    grd2:  $\forall p \cdot (p \in wait \wedge p \in dom(t1)) \Rightarrow clk + 1 - t1(p) \leq ddl1$ 
    grd4:  $\forall p \cdot (p \in process \wedge p \in dom(t2)) \Rightarrow clk + 1 - t2(p) \leq ddl2$ 
  then
    act1:  $clk := clk + 1$ 
  end
END

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