

MACHINE m2

REFINES m1

SEES c0

VARIABLES

wait
process
clk
t1
t2
t3
position
index

INVARIANTS

inv1: $\forall t, p. (p \in process \wedge p \in dom(t2) \wedge t = t2(p)) \Rightarrow clk - t \leq ddl2$
inv2: $\forall p. (p \in dom(t2) \wedge p \in dom(t3) \wedge t3(p) \geq t2(p)) \Rightarrow t3(p) - t2(p) \leq ddl2$
 $deadline(t2, t3, ddl1)$
inv3: $\forall t, p. (p \in process \wedge p \in dom(t1) \wedge p \in dom(t2) \wedge t = t2(p) \wedge t2(p) \geq t1(p)) \Rightarrow t2(p) - t1(p) \leq ddl1$
inv4: $position \in wait \mapsto POSITION$
inv5: $index \in PROCESS \mapsto \mathbb{N}$

EVENTS

Initialisation $\langle \text{extended} \rangle$

begin

act1: $wait := \emptyset$
act2: $process := \emptyset$
act3: $clk := 0$
act4: $t1 := \emptyset$
act5: $t2 := \emptyset$
act6: $t3 := \emptyset$
act7: $position := \emptyset$
act8: $index := \emptyset$

end

Event wish $\langle \text{ordinary} \rangle \hat{=}$

extends wish

any

pro
pos

where

grd1: $pro \in PROCESS \setminus wait$
grd2: $pro \in PROCESS \setminus process$
grd3: $pos \in POSITION$
grd4: $pos \notin ran(position)$
grd7: $finite(ran(position))$
grd6: $position \neq \emptyset \Rightarrow pos = max(ran(position)) + 1$

then

act1: $wait := wait \cup \{pro\}$
act2: $t1(pro) := clk$
act3: $position := position \cup \{pro \mapsto pos\}$
act4: $index(pro) := pos$

end

Event enter $\langle \text{ordinary} \rangle \hat{=}$

extends enter

any

pro

where

grd1: $pro \in wait$
grd2: $card(process) = 0$

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    grd3:  $pro \in dom(position)$ 
    grd4:  $position(pro) = min(ran(position))$ 
  then
    act1:  $wait := wait \setminus \{pro\}$ 
    act2:  $process := process \cup \{pro\}$ 
    act3:  $t2(pro) := clk$ 
    act5:  $position := \lambda p. p \in wait \setminus \{pro\} | position(p) - 1$ 
  end
Event leave ⟨ordinary⟩  $\hat{=}$ 
extends leave
  any
     $pro$ 
  where
    grd1:  $pro \in process$ 
  then
    act1:  $process := process \setminus \{pro\}$ 
    act3:  $t3(pro) := clk$ 
  end
Event tick ⟨ordinary⟩  $\hat{=}$ 
refines tick
  when
    grd2:  $\forall t, p. (p \in wait \wedge p \in dom(t1) \wedge t = t1(p)) \Rightarrow clk + 1 - t \leq ddl1$ 
    grd4:  $\forall t, p. (p \in process \wedge p \in dom(t2) \wedge t = t2(p)) \Rightarrow clk + 1 - t \leq ddl2$ 
  then
    act1:  $clk := clk + 1$ 
  end
END

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