String Algorithms and Data Structures Approximate Pattern Matching

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Approximate Pattern Matching

Input: A text *T*, a pattern *P*, and a distance *d*

Output: All positions in T where P has at most d mismatches or edits

P: word

T: There would have been a time for such a word Alignment 1: word Alignment 2: word

Not a match!

Distance 2 match!

Match!

Distance 0 match!

Hamming Distance

The number of **substitutions** required to turn one string into another

```
T: GGAAAAAGAGGTAGCGGCGTTTAACAGTAG

| | | | | | | | |

P: GTAACGGCG

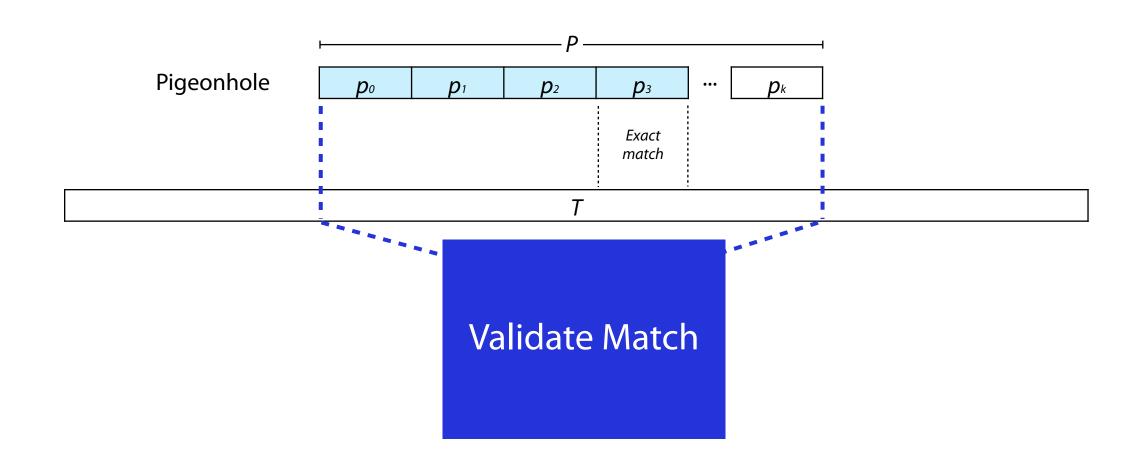
Mismatch

(Substitution)
```

```
Score = 248 bits (129), Expect = 1e-63
Identities = 213/263 (80%), Gaps = 34/263 (12%)
Strand = Plus / Plus
                                    Substitution
Query: 161 atatcaccacgtcaaaggtgactccaactcca---ccactccattttgttcagataatgc 217
         1111111111111111111111111111111
Sbjct: 481 atatcaccacgtcaaaggtgactccaact-tattgatagtgttttatgttcagataatgc 539
Query: 218 ccgatgatcatgtcatgcagctccaccgattgtgagaacgacagcgacttccgtcccagc 277
                Sbjct: 540 ccgatgactttgtcatgcagctccaccgattttg-g--
                                            ----ttccgtcccagc 586
                                      Deletion
Query: 278 c-gtgcc--aggtgctgcctcagattcaggttatgccgctcaattcgctgcgtatatcgc 334
          Sbjct: 587 caatgacgta-gtgctgcctcagattcaggttatgccgctcaattcgctgggtatatcgc 645
Query: 335 ttgctgattacgtgcagctttcccttcaggcggga-----ccagccatccgtc 382
         Insertion
Query: 383 ctccatatc-accacgtcaaagg 404
Sbjct: 706 atccatatcaaccacgtcaaagg 728
```

Approximate Pattern Matching

Find an exact match partition and validate the overall alignment



Learning Objectives



Review approximate pattern matching

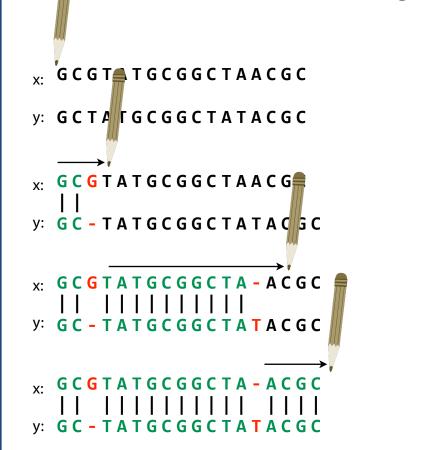
Formalize edit distance storage as an 'edit string'

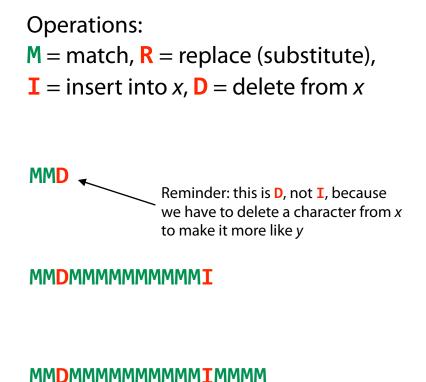
Discuss strategies for efficient APM with edits

Introduce dynamic programming

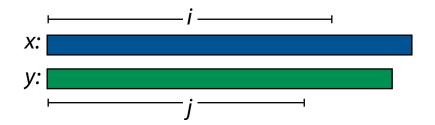
Imagine edits are introduced by an optimal editor working left-to-right:

Edit string summarizes how editor turns *x* into *y*:





D[i, j]: edit distance between length-i prefix of x and length-j prefix of y



Optimal edit string for D[i, j] is built by **extending a shorter optimal string by 1 operation**. 3 options:

Append D to transcript for D[i-1, j]

Append I to transcript for D[i, j-1]

Append M or R to transcript for D[i-1, j-1]

We choose based on whichever option has the fewest edits

X: GTTTAA Y: GGTTTA

D[5, 6] GTTTA

D[5, 5] GTTTA GGTTT

D[6, 5] GTTTAA

GTTTAA D[6, 6]

X: GTTTAA Y: GGTTTA **GTTTA** D[5, 6] **GGTTTA** GTTTAA **GTTTA** D[6, 6] D[5, 5] **GGTTT GGTTTA GTTTAA** D[6, 5] **GGTTT**

X: GTTTAA Y: GGTTTA

MIMMMM

D[5, 6]

G	-	Т	Т	Т	Α
G	G	Т	Т	Т	Α

MRMMR

D[5, 5]

MRMMRD

D[6, 5]

G	Т	Т	Т	Α	Α
G	G	Т	Т	Т	_

X: GTTTAA

Y: GGTTTA

MIMMMM

D[5, 6]

G	-	Т	Т	Т	Α
G	G	Т	T	T	Α

D[6, 6]¹

MRMMR

D[*5*, *5*]

 $D[6, 6]^2$

MRMMRD

D[6, 5]

 $D[6, 6]^3$

X: GTTTAA

Y: GGTTTA

MIMMMM

D[5, 6]

G	-	Т	Т	Т	Α
G	G	Т	Т	Т	Α

MIMMMMD

MRMMR

D[5, 5]

G	Т	Т	Т	Α
G	G	Т	Т	Т

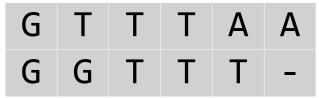
MRMMRM

G	Т	Т	Т	Α	A
G	G	Т	Т	Т	Α

$$D[6, 6]^2$$

MRMMRD

D[6, 5]

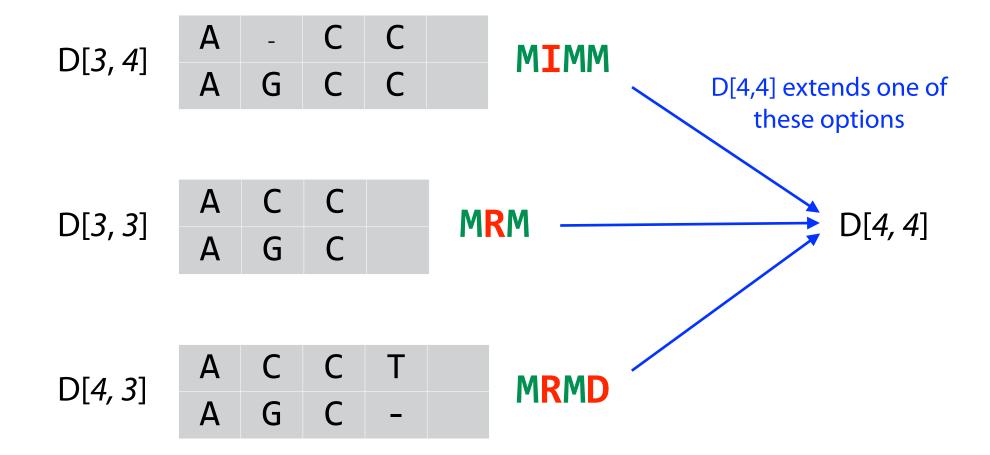


MRMMRDI

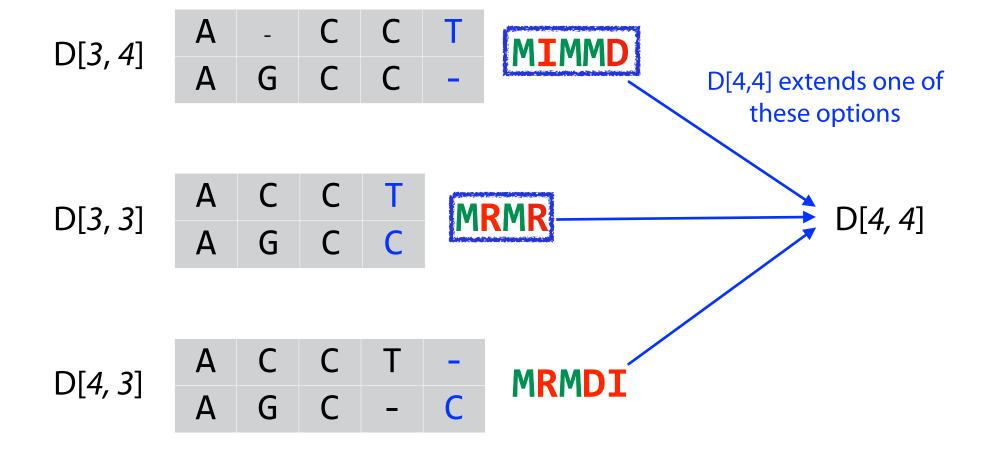
G	Т	Т	Τ	Α	A	-
G	G	Т	Т	Т	_	A

 $D[6, 6]^3$

X: ACCT Y: AGCC

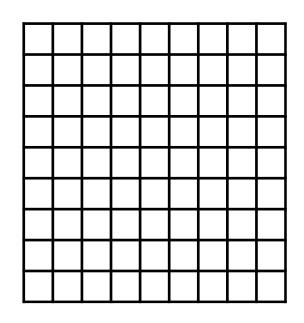


X: ACCT Y: AGCC



We can store D as a 2D matrix:

```
Let D[0,j]=j, and let D[i,0]=i
```

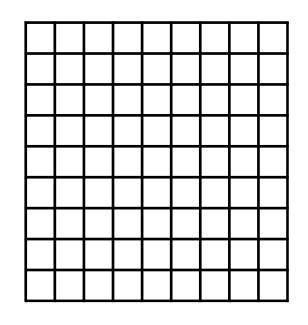


```
the longest----
|||||||
----longest day
```

We can store D as a 2D matrix:

Is at beginning Ds at beginning

Let
$$D[0, j] = j$$
, and let $D[i, 0] = i$



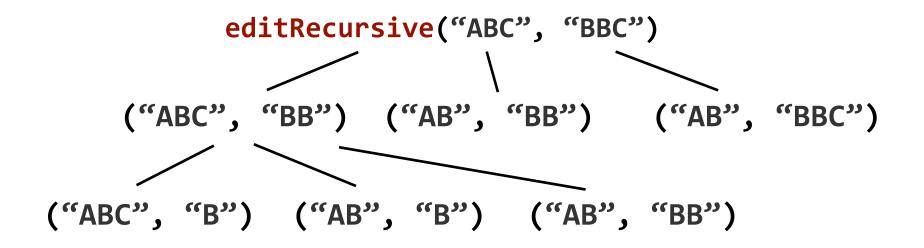


Otherwise, let
$$D[i,j] = \min$$

$$\begin{cases} D[i-1,j] + 1 & \text{horizontal (I)} \\ D[i,j-1] + 1 & \\ D[i-1,j-1] + \delta(x[i-1],y[j-1]) \end{cases}$$
 diagonal (M or R)

diagonal (M or R)

$$\delta(a,b)$$
 is 0 if $a=b$, 1 otherwise



(only part of recursion tree shown)

```
editRecursive("ABC", "BBC")
    ("ABC", "BB") ("AB", "BB") ("AB", "BBC")
("ABC", "B") ("AB", "B") ("AB", "BB")
                 (only part of recursion tree shown)
     editRecursive("Shakespeare", "shake spear")
     Calculate ("Shake", "shake") 8989 times
           How can we address this problem?
```

```
editRecursive("ABC", "BBC")

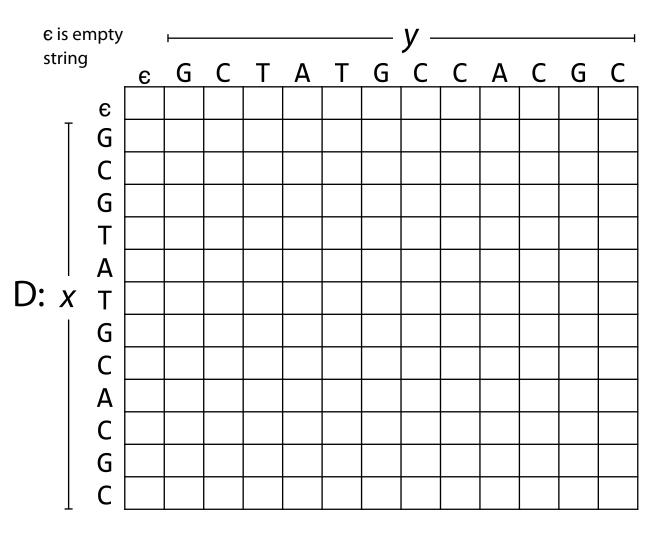
("ABC", "BB") ("AB", "BB") ("AB", "BBC")

("ABC", "B") ("AB", "B") ("AB", "BB")
```

Memoization: Top-down

Dynamic Programming: Bottom-up

Both: Solve individual sub-problems once

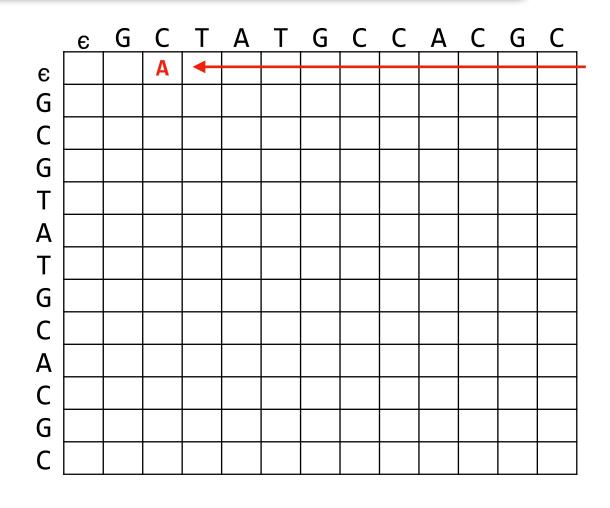


Let
$$n = |x|, m = |y|$$

D:
$$(n+1) \times (m+1)$$
 matrix

D[i, j] = edit distance b/t length-iprefix of x and length-j prefix of y

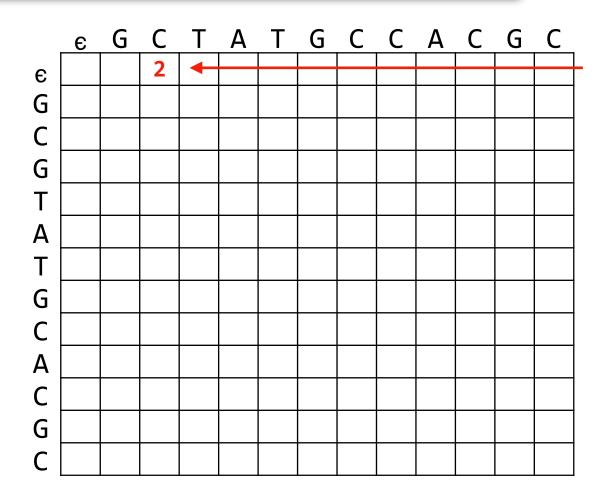
Let D[0, j] = j, and let D[i, 0] = i



What is A?

D[i, j] = edit distance b/t length-iprefix of x and length-j prefix of y

Let D[0, j] = j, and let D[i, 0] = i



What is A?

Let D[0,j] = j, and let D[i,0] = i

	<u> </u>	G	С	Τ	Α	Τ	G	C	C	Α	C	G	<u>C</u>
ϵ	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1												
C	2												
G	3												
Т	4												
Α	5												
Т	6												
G	7												
C	8												
Α	9												
C	10												
G	11												
C	12												

Let D[0, j] = j, and let D[i, 0] = i

	ϵ	G	C	T	Α	Τ	G	C	C	Α	C	G	C
ϵ	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1												
C	2												
G	3												
Τ	4												
Α	5												
Τ	6												
G	7												
C	8												
Α	9												
C	10												
G	11												
C	12												



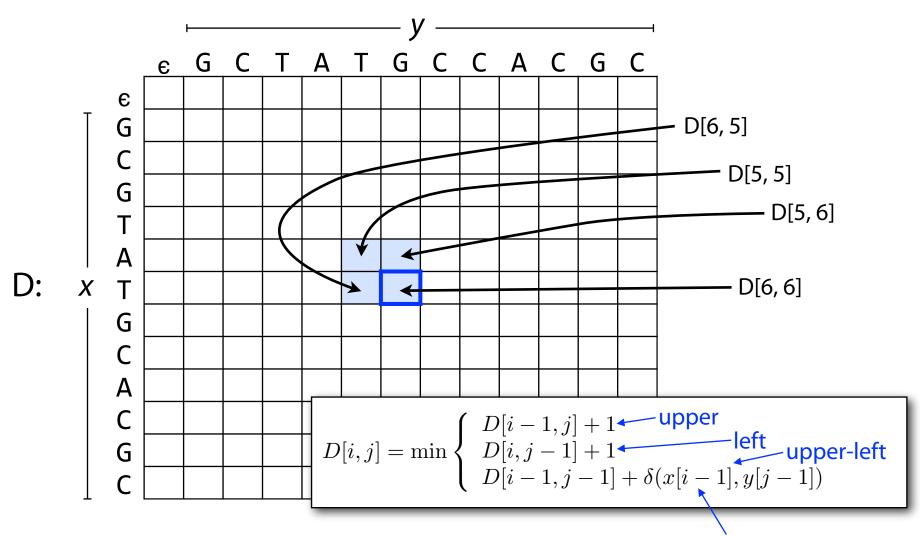
Let D[0,j] = j, and let D[i,0] = i

	ϵ	G	C	T	Α	Τ	G	C	C	Α	C	G	C
ϵ	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1												
C	2												
G	3												
Τ	4												
Α	5												
Τ	6												
G	7												
C	8												
Α	9												
C	10												
G	11												
C	12												

GCGTAT

Let D[0, j] = j, and let D[i, 0] = i

	ε	G	C	Т	Α	Т	G	C	C	Α	C	G	C						
ε	0	1	2	3	4	5	6	7	8	9	10	11	12						
G	1																		
C	2																		
G	3																		
T	4																		
Α	5																		
T	6																		
G	7																		
C	8																		
А	9																		
C	10																		
G	11					D i	S C	ne	rc)\/\	<i>/ c</i>	ဂါ၊	ım	n la	arge	r tء	har	า x /	/ \
C	12																		



Cell depends upon its upper, left, and upper-left neighbors

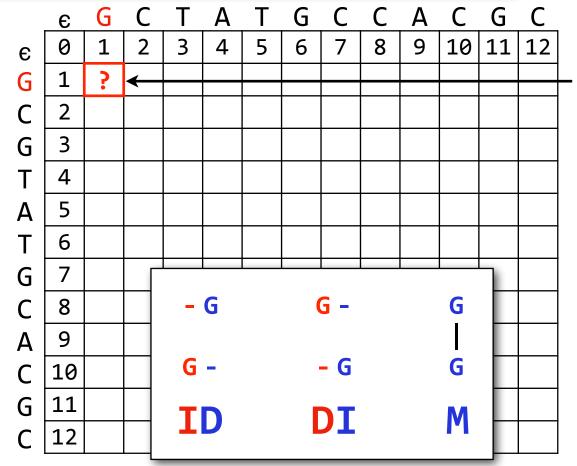
i, j are D-basedD[0] is 'empty string'

$$D[i,j] = \min \left\{ \begin{array}{l} D[i-1,j]+1 \\ D[i,j-1]+1 \\ D[i-1,j-1]+\delta(x[i-1],y[j-1]) \end{array} \right.$$

	€	G	C	Τ	Α	Τ	G	С	С	Α	C	G	C
ϵ	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1												→
C	2	4											→
G	3	4											\rightarrow
Т	4	4											→
Α	5						etc						
Т	6												
G	7												
C	8												
Α	9												
C	10												
G	11												
C	12												

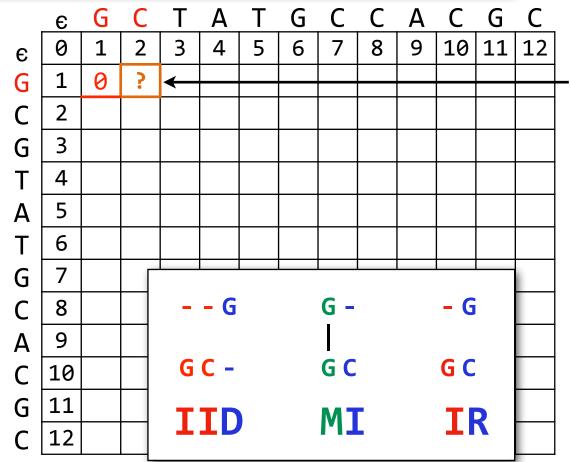
Fill remaining cells from top row to bottom and from left to right

$$D[i,j] = \min \left\{ \begin{array}{l} D[i-1,j] + 1 \\ D[i,j-1] + 1 \\ D[i-1,j-1] + \delta(x[i-1],y[j-1]) \end{array} \right.$$



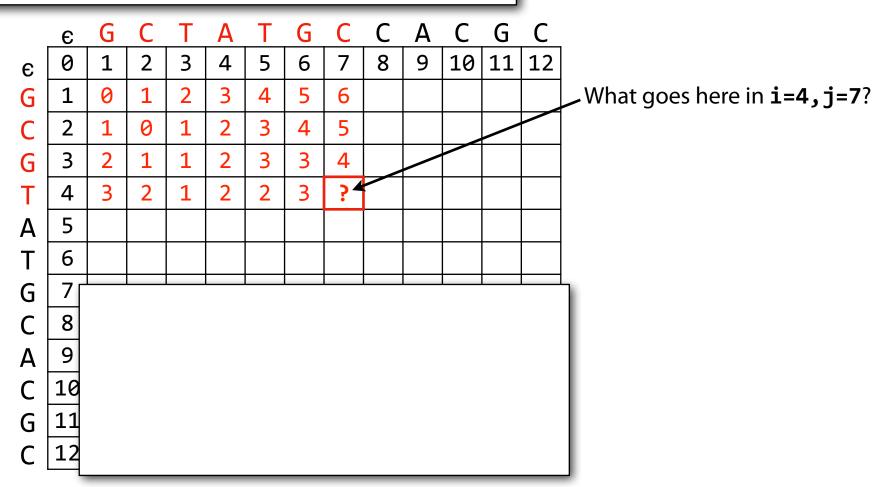
What goes here in i=1, j=1?

$$D[i,j] = \min \left\{ \begin{array}{l} D[i-1,j]+1 \\ D[i,j-1]+1 \\ D[i-1,j-1]+\delta(x[i-1],y[j-1]) \end{array} \right.$$

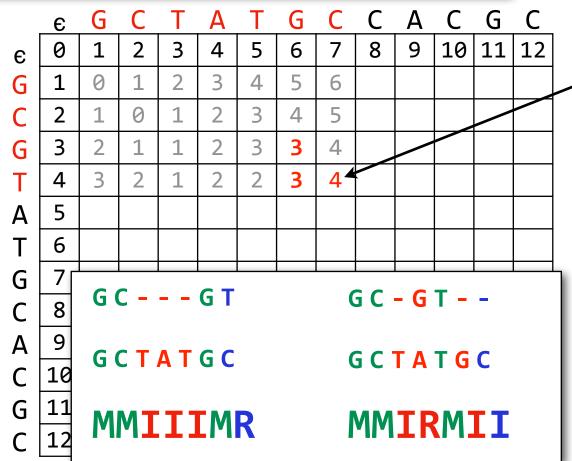


What goes here in i=1, j=2?

$$D[i,j] = \min \left\{ \begin{array}{l} D[i-1,j] + 1 \\ D[i,j-1] + 1 \\ D[i-1,j-1] + \delta(x[i-1],y[j-1]) \end{array} \right.$$



$$D[i,j] = \min \begin{cases} D[i-1,j] + 1 \\ D[i,j-1] + 1 \\ D[i-1,j-1] + \delta(x[i-1],y[j-1]) \end{cases}$$



What goes here in i=4, j=7?

```
x[i-1] = 'T',
y[j-1] = 'C',
so delt = 1
```

$$D[i,j] = \min \left\{ \begin{array}{l} D[i-1,j]+1 \\ D[i,j-1]+1 \\ D[i-1,j-1]+\delta(x[i-1],y[j-1]) \end{array} \right.$$

	€	G			<u> </u>		G			<u> </u>		G	<u> </u>
ε	0	1	2	ო	4	5	6	7	8	9	10	11	12
G	1	0	1	2	3	4	5	6	7	8	9	10	11
C	2	1	0	1	2	3	4	5	6	7	8	9	10
G	3	2	1	1	2	3	3	4	5	6	7	8	9
Т	4	3	2	1	2	2	3	4	5	6	7	8	9
Α	5	4	3	2	1	2	3	4	5	5	6	7	8
Т	6	5	4	3	2	1	2	3	4	5	6	7	8
G	7	6	5	4	3	2	1	2	3	4	5	6	7
C	8	7	6	5	4	3	2	1	2	3	4	5	6
Α	9	8	7	6	5	4	3	2	2	2	3	4	5
C	10	9	8	7	6	5	4	3	2	3	2	3	4
G	11	10	9	8	7	6	5	4	3	3	3	2	3
C	12	11	10	9	8	7	6	5	4	4	3	3	2 <

Fill remaining cells from top row to bottom and from left to right

-Edit distance for x, y

		Y						
		ε	C	А	А	T		
X	E	0	1	2	3	4		
	C	1	0	A	2	3		
	Α	2	1	0	В			
	Τ	3	2	С				

		Y						
		ϵ	C	А	А	Τ		
X	E	0	1	2	3	4		
	C	1	0	1	2	3		
	Α	2	1	0	В			
	Τ	3	2	C				

				Y		
		ϵ	C	Α	Α	T
	E	0	1	2	3	4
V	C	1	0	1	2	3
X	Α	2	1	0	1	
	Τ	3	2	С		

				Y		
		ε	C	Α	Α	T
	E	0	1	2	3	4
V	C	1	0	1	2	3
X	Α	2	1	0	1	
	Τ	3	2	1		

				Y		
		ϵ	C	Α	Α	Т
	E	0	1	2	3	4
X	C	1	0	1	2	3
A	A	2	1	0	1	2
	Τ	3	2	1	1	1

eMatrix buildEditMatrix(X, Y)



Input:

```
string X:Input string X (edits with respect to X)
```

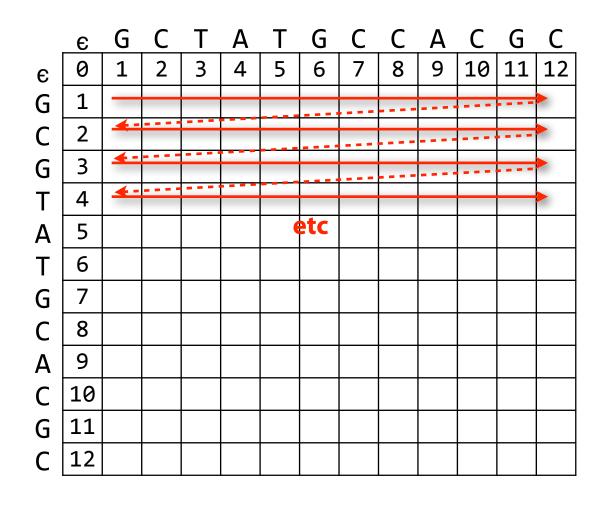
string Y:Input string Y (edits turn X into Y)

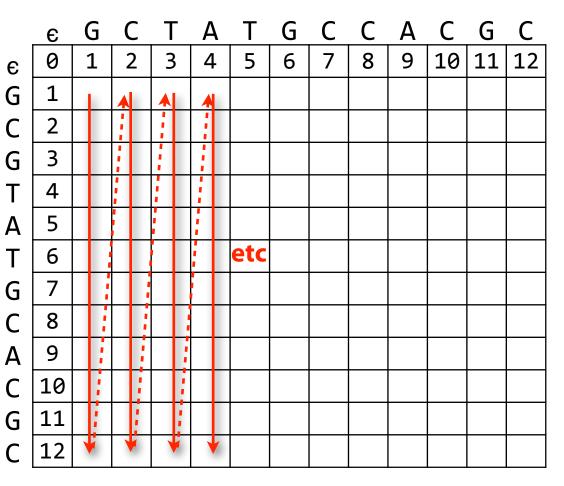
Output:

eMatrix: vector<vector<int>> storing all optimal edit distances

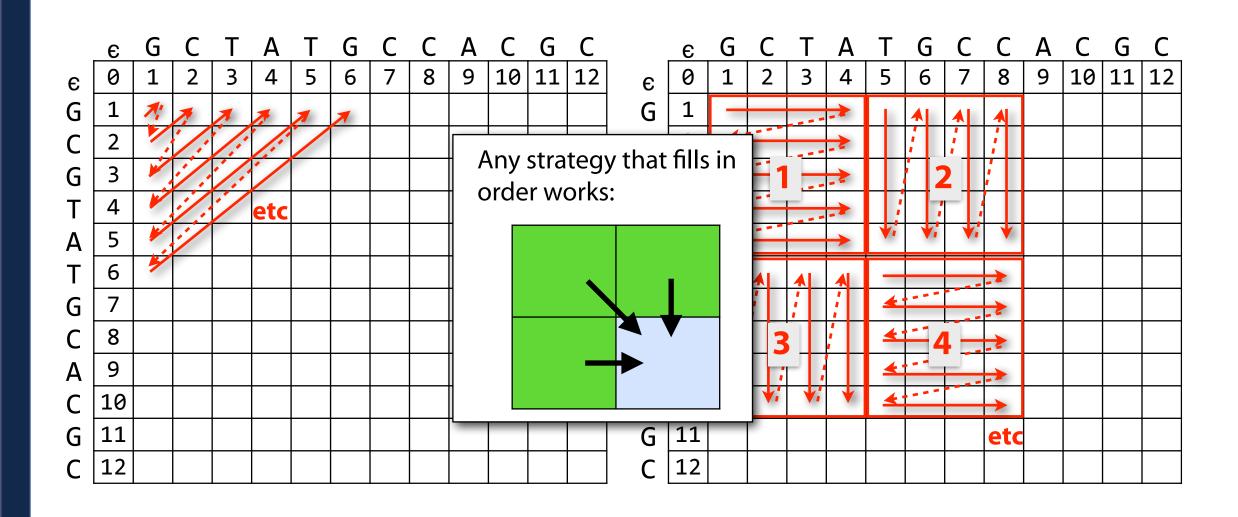
	ϵ	C	A	A	Т
ϵ	0	1	2	3	4
C	1	0	1	2	3
Α	2	1	0	1	2
Т	3	2	1	1	1

Edit Distance: dynamic programming





Edit Distance: dynamic programming



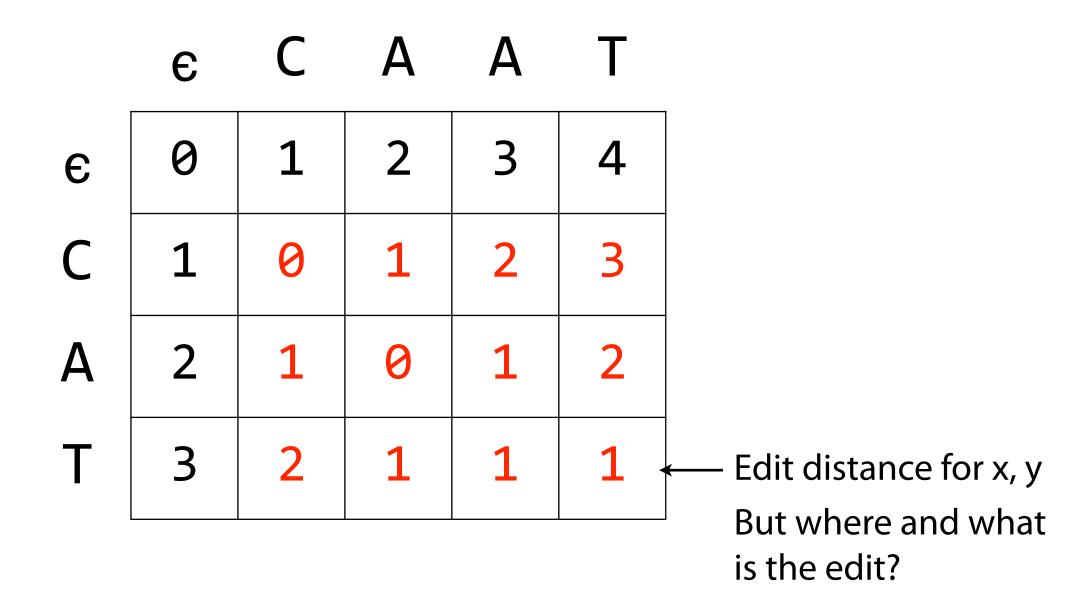
Assignment 11: a_edist

Learning Objective:

Use dynamic programming to build an edit distance matrix

Construct an optimal edit string from the edit matrix

Consider: Does substitution, insertion, and deletion need to have the same 'weight' as a penalty? How could you modify the code to take a user-specified input for each?



Traceback corresponds to an optimal alignment / edit transcript

At each step, ask: which neighbor (\mathbb{Q} , \Leftarrow or \mathbb{Q}) gave the minimum?

	ϵ	C	A	A	Т
E	0	1	2	3	4
C	1	0	1	2	3
Α	2	1	0	1	2
Т	3	2	1	1	1 •

Traceback corresponds to an optimal alignment / edit transcript

At each step, ask: which neighbor ($\langle \rangle$, \leftarrow or $\uparrow \rangle$) gave the minimum?

	ϵ	C	Α	А	Т	
ϵ	0	1	2	3	4	D[2, 4] =
C	1	0	1	2	3	
А	2	1	0	1	2	
Т	3	2	1	1	1 •	— Q: How did I get here?

Traceback corresponds to an optimal alignment / edit transcript

At each step, ask: which neighbor (\mathbb{Q} , \Leftarrow or \Uparrow) gave the minimum?

	ϵ	C	Α	А	T	
ϵ	0	1	2	3	4	D[3, 3] =
C	1	0	1	2	3	
A	2	1	0	1	2	
T	3	2	1	1	1 •	— Q: How did I get here?

Traceback corresponds to an optimal alignment / edit transcript

At each step, ask: which neighbor (\triangleright , \leftarrow or $\uparrow \vdash$) gave the minimum?

	ϵ	C	Α	Α	Т	
ϵ	0	1	2	3	4	D[2, 3] =
C	1	0	1	2	3	
A	2	1	0	1	2	
T	3	2	1	1	1 •	— Q: How did I get here?

Traceback corresponds to an optimal alignment / edit transcript

At each step, ask: which neighbor (\triangleright , \Leftarrow or \Uparrow) gave the minimum?

	ϵ	C	A	A	T
ε	0	1	2	3	4
C	1	0	1	2	3
Α	2	1	0	1	2
Т	3	2	1	1	1

$$D[1, 3] =$$

$$D[1, 2] =$$

$$D[2, 2] =$$

Traceback corresponds to an optimal alignment / edit transcript

At each step, ask: which neighbor (\triangleright , \leftarrow or $\uparrow \vdash$) gave the minimum?

	ϵ	C	A	A	Т
ϵ	0	1	2	3	4
C	1	0	1	2	Tie!
Α	2	1	0	1	7
Т	3	2	1	1	1

$$D[1, 3] = 2 + 1$$

$$D[1, 2] = 1 + 0$$

$$D[2, 2] = 0 + 1$$

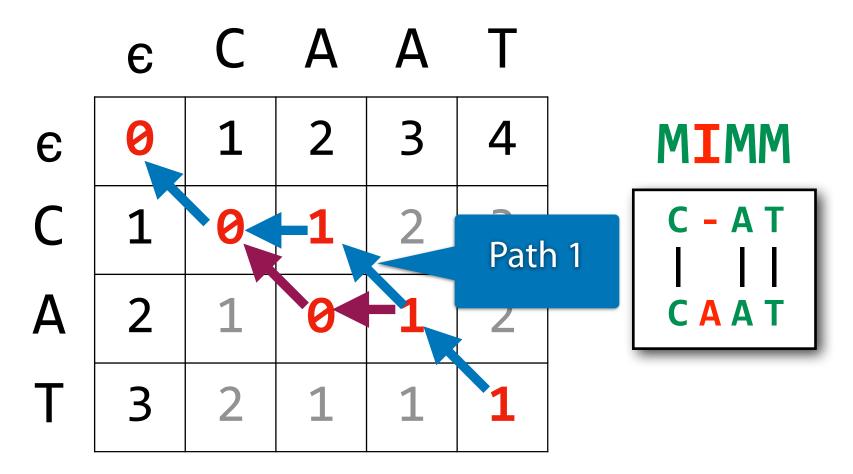
Traceback corresponds to an optimal alignment / edit transcript

At each step, ask: which neighbor (\triangleright , \Leftarrow or \Uparrow) gave the minimum?

	ϵ	C	Α	Α	T
ϵ	0	1	2	3	4
C	1	0	-1	2	3
Α	2	1	0	1	2
Т	3	2	1	1	1

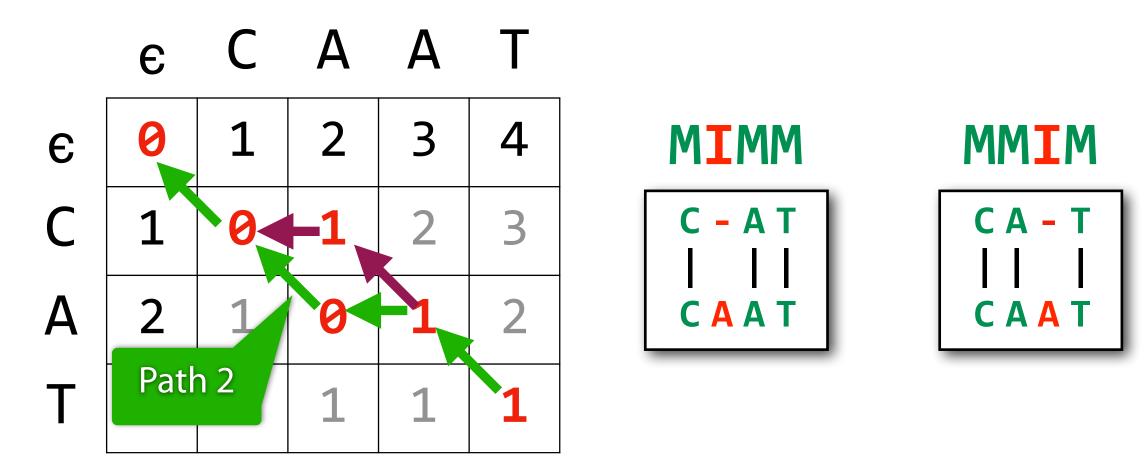
Traceback corresponds to an optimal alignment / edit transcript

At each step, ask: which neighbor (\triangleright , \leftarrow or $\uparrow \vdash$) gave the minimum?



Traceback corresponds to an optimal alignment / edit transcript

At each step, ask: which neighbor (\mathbb{Q} , \Leftarrow or \Uparrow) gave the minimum?



i	€	G	C	Τ	Α	Τ	G	C	C	Α	C	G	<u>C</u>
€	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1	0	1	2	3	4	5	6	7	8	9	10	11
C	2	1	0	1	2	3	4	5	6	7	8	9	10
G	3	2	1	1	2	3	3	4	5	6	7	8	9
T	4	3	2	1	2	2	3	4	5	6	7	8	9
Α	5	4	3	2	1	2	3	4	5	5	6	7	8
T	6	5	4	3	2	1	2	3	4	5	6	7	8
G	7	6	5	4	3	2	1	2	3	4	5	6	7
C	8	7	6	5	4	3	2	1	2	M	4	5	6
Α	9	8	7	6	5	4	3	2	2	2	3	4	5
C	10	9	8	7	6	5	4	3	2	M	2	3	4
G	11	10	9	8	7	6	5	4	3	3	3	2	3
C	12	11	10	9	8	7	6	5	4	4	3	3	2

<u> </u>	<u>G</u>	<u> </u>	<u>A</u>	<u> </u>	<u> </u>	<u>G</u>	T_{-}	<u>A</u>	T_{-}	<u> </u>	<u>G</u>	ε	
12	11	10	9	8	7	6	5	4	3	2	1	0	ϵ
11	10	9	8	7	6	5	4	3	2	1	0	1	G
10	9	8	7	6	5	4	3	2	1	0	1	2	C
9	8	7	6	5	4	3	3	2	1	1	2	3	G
9	8	7	6	5	4	3	2	2	1	2	3	4	T
8	7	6	5	5	4	3	2	1	2	3	4	5	Α
8	7	6	5	4	3	2	1	2	3	4	5	6	T
7	6	5	4	3	2	1	2	3	4	5	6	7	G
6	5	4	3	2	1	2	3	4	5	6	7	8	C
5	4	3	2	2	2	3	4	5	6	7	8	9	Α
4	3	2	3	2	3	4	5	6	7	8	9	10	C
3	2 ←	3	3	3	4	5	6	7	8	9	10	11	G
2	3	3	4	4	5	6	7	8	9	10	11	12	C
	6 5 4 3	5 4 3 2 3	4 3 2 3 3	3 2 2 2 3	2 1 2 3 4	1 2 3 4 5	2 3 4 5 6	3 4 5 6 7	4 5 6 7 8	5 6 7 8 9	6 7 8 9 10	7 8 9 10 11	G C A C

$$D[11, 12] = 3 + 1$$

$$D[11, 11] = 2 + 0$$

$$D[12, 11] = 3 + 1$$

A: From here

	E	G	C	Τ	Α	T	G	C	C	Α	C	G	C
ε	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1	0	1	2	3	4	5	6	7	8	9	10	11
C	2	1	0	1	2	3	4	5	6	7	8	9	10
G	3	2	1	1	2	3	3	4	5	6	7	8	9
T	4	3	2	1	2	2	3	4	5	6	7	8	9
Α	5	4	3	2	1	2	3	4	5	5	6	7	8
T	6	5	4	3	2	1	2	3	4	5	6	7	8
G	7	6	5	4	3	2	1	2	3	4	5	6	7
C	8	7	6	5	4	3	2	1	2	3	4	5	6
А	9	8	7	6	5	4	3	2	2	2	3	4	5
C	10	9	8	7	6	5	4	3	2	3	2	3	4
G	11	10	9	8	7	6	5	4	3	3	3	2	73
C	12	11	10	9	8	7	6	5	4	4	3	3	2

Ī	€	G	<u>C</u>	T	<u>A</u>	<u>T</u>	G	<u>C</u>	<u>C</u>	<u>A</u>	<u>C</u>	G	<u>C</u>	
ϵ	0	1	2	3	4	5	6	7	8	9	10	11	12	
G	1	0	1	2	3	4	5	6	7	8	9	10	11	
C	2	1	0	1	2	3	4	5	6	7	8	9	10	
G	3	2	1	1	2	3	3	4	5	6	7	8	9	
T	4	3	2	1	2	2	3	4	5	6	7	8	9	
Α	5	4	3	2	1	2	3	4	5	5	6	7	8	
Т	6	5	4	3	2	1	2	3	4	5	6	7	8	
G	7	6	5	4	3	2	1	2	3	4	5	6	7	
C	8	7	6	5	4	3	2	1	2	3	4	5	6	
Α	9	8	7	6	5	4	3	2	2	2	3	4	5	
C	10	9	8	7	6	5	4	3	2	3	2	3	4	
G	11	10	9	8	7	6	5	4	3	3	3	2	3	
C	12	11	10	9	8	7	6	5	4	4	3	3	2	

$$D[10, 11] = 3 + 1$$

$$D[10, 10] = 2 + 0$$

$$D[11, 10] = 3 + 1$$

- A: From here

	€	G	C	T	Α	Т	G	C	C	Α	C	G	<u>C</u>
ϵ	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1	0	1	2	3	4	5	6	7	8	9	10	11
C	2	1	0	1	2	3	4	5	6	7	8	9	10
G	3	2	1	1	2	3	3	4	5	6	7	8	9
T	4	3	2	1	2	2	3	4	5	6	7	8	9
Α	5	4	3	2	1	2	3	4	5	5	6	7	8
T	6	5	4	3	2	1	2	3	4	5	6	7	8
G	7	6	5	4	3	2	1	2	3	4	5	6	7
C	8	7	6	5	4	3	2	1	2	3	4	5	6
Α	9	8	7	6	5	4	3	2	2	2	3	4	5
C	10	9	8	7	6	5	4	3	2	M	N	7	4
G	11	10	9	8	7	6	5	4	3	3	3	2	3
C	12	11	10	9	8	7	6	5	4	4	3	3	2

i	€	<u>_G</u> _	<u> </u>		<u>A</u>		<u>G</u>	<u> </u>	<u> </u>	<u>A</u>	<u> </u>	<u>G</u>	<u> </u>
€	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1	0	1	2	3	4	5	6	7	8	9	10	11
C	2	1	0	1	2	3	4	5	6	7	8	9	10
G	3	2	1	1	2	3	3	4	5	6	7	8	9
T	4	3	2	1	2	2	3	4	5	6	7	8	9
Α	5	4	3	2	1	2	3	4	5	5	6	7	8
T	6	5	4	3	2	1	2	3	4	5	6	7	8
G	7	6	5	4	3	2	1	2	3	4	5	6	7
C	8	7	6	5	4	3	2	1	2	3	4	5	6
Α	9	8	7	6	5	4	3	2	2	2	m	4	5
C	10	9	8	7	6	5	4	ß	2	3	M	3	4
G	11	10	9	8	7	6	5	4	3	3	3	A	3
C	12	11	10	9	8	7	6	5	4	4	3	3	2

$$D[9, 10] = 3 + 1$$

$$D[9, 9] = 2 + 0$$

$$D[10, 9] = 3 + 1$$

- A: From here

	€	<u>G</u>	<u> </u>	<u> T </u>	<u>A</u>	<u>T</u>	<u>G</u>	<u> </u>	<u> </u>	<u>A</u>	<u> </u>	<u>G</u>	<u> </u>
E	a	1	2	3	4	5	6	7	8	9	10	11	12
G	1	Q	1	2	3	4	5	6	7	8	9	10	11
C	2	1	8	1	2	3	4	5	6	7	8	9	10
G	3	2	1	1	2	3	3	4	5	6	7	8	9
Т	4	3	2	A	2	2	3	4	5	6	7	8	9
Α	5	4	3	2	A	2	3	4	5	5	6	7	8
T	6	5	4	3	2	A	2	3	4	5	6	7	8
G	7	6	5	4	3	2	٠	4	3	4	5	6	7
C	8	7	6	5	4	3	2	1	2	3	4	5	6
Α	9	8	7	6	5	4	3	2	2	SA	3	4	5
C	10	9	8	7	6	5	4	3	2	3	See	3	4
G	11	10	9	8	7	6	5	4	3	3	3	2	3
C	12	11	10	9	8	7	6	5	4	4	3	3	2

Alignment:

MMDMMMMIMMMMM

Dynamic Programming fills our table with optimal distances

Traceback identifies the optimal *edit string* to convert *X* to *Y*

What is our efficiency? (|X| = m, |Y| = n)

	ε	G	C	Т	Α	Т	G	C	C	Α	C	G	C
ε	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1	9	1	2	3	4	5	6	7	8	9	10	11
C	2	1	9	1	2	3	4	5	6	7	8	9	10
G	3	2	4	1	2	3	3	4	5	6	7	80	9
Т	4	3	2	A	2	2	3	4	5	6	7	8	9
Α	5	4	3	2	×	2	3	4	5	5	6	7	8
Т	6	5	4	3	2	A	2	3	4	5	6	7	8
G	7	6	5	4	3	2	\mathbf{H}	4	w	4	5	6	7
C	8	7	6	5	4	3	2	1	24	3	4	5	6
Α	9	8	7	6	5	4	3	2	2	2	3	4	5
C	10	9	8	7	6	5	4	3	2	თ	A	3	4
G	11	10	9	8	7	6	5	4	3	3	3	Δ	3
C	12	11	10	9	8	7	6	5	4	4	3	3	2

Dynamic Programming fills our table with optimal distances

Traceback identifies the optimal *edit string* to convert *X* to *Y*

	ε	G	C	Т	Α	Т	G	C	C	Α	C	G	C
ε	0	1	2	3	4	5	6	7	8	9	10	11	12
G	1	9	1	2	3	4	5	6	7	8	9	10	11
C	2	1	9	1	2	3	4	5	6	7	8	9	10
G	3	2	4	1	2	3	ო	4	5	6	7	80	9
Т	4	3	2	A	2	2	M	4	5	6	7	8	9
Α	5	4	3	2	X	2	M	4	5	5	6	7	8
Т	6	5	4	3	2	A	2	3	4	5	6	7	8
G	7	6	5	4	3	2	H	7	3	4	5	6	7
C	8	7	6	5	4	3	2	1	24	3	4	5	6
Α	9	8	7	6	5	4	M	2	2	2	3	4	5
C	10	9	8	7	6	5	4	3	2	თ	A	3	4
G	11	10	9	8	7	6	5	4	3	3	3	Δ	3
C	12	11	10	9	8	7	6	5	4	4	3	3	2

What is our efficiency? (|X| = m, |Y| = n)

Table filling: Filling (m + 1) (n + 1) cells, each requiring constant work, so O(mn)

Traceback: Each step goes \triangleleft , \Leftarrow or

 \Uparrow . Worst case: traceback never moves diagonally, requiring $m \Uparrow$'s and $n \Leftarrow$'s, so O(m + n)

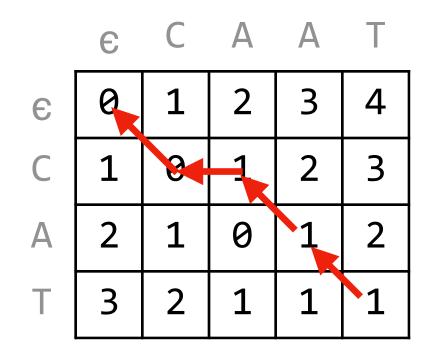
string buildEditString(X, Y)



Input: **string X:**Input string X (edits with respect to X)

string Y:Input string Y (edits turn X into Y)

Output: **string**: An optimal edit string produced by the matrix



On tie: prioritize diagonal, then vertical, then horizontal

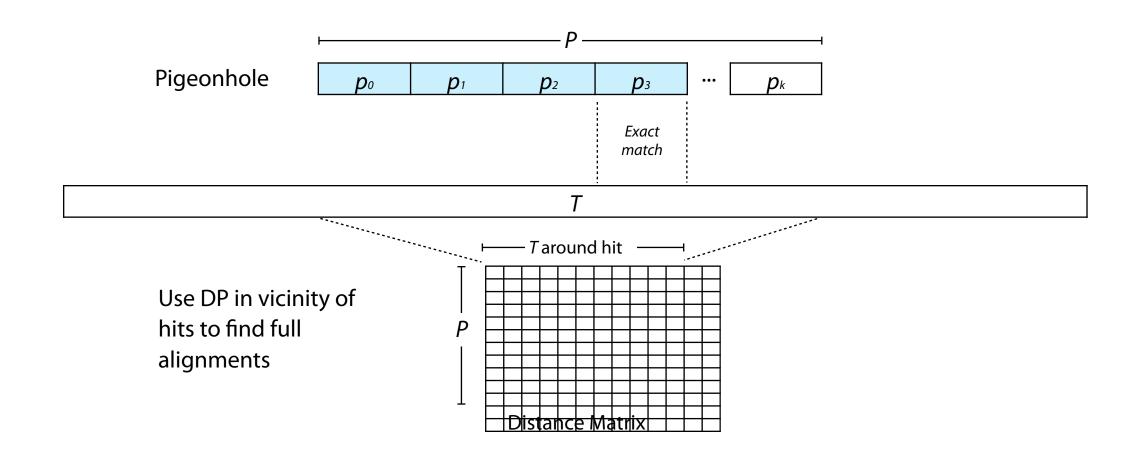


MIMM

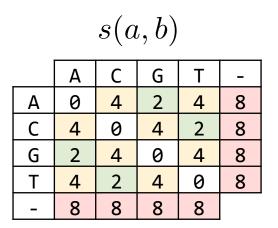


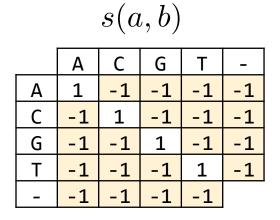
Approximate Pattern Matching

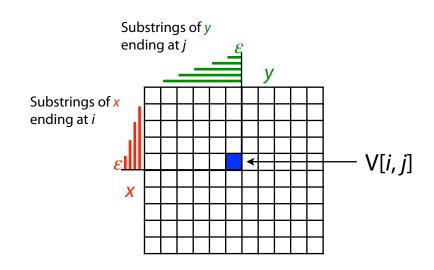
"Seed and extend" works for edit distance too!

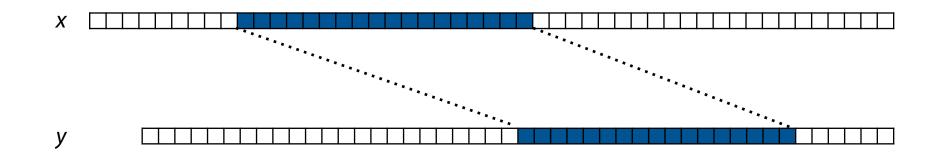


Bonus Slide









... and much more!