Algoritmos e Estruturas de Dados III

Aula 6.1 - Hashing Dinâmico

Prof. Felipe Lara



Hash Dinâmico Definições



- Quando o arquivo de dados cresce ou diminui com frequência (muitas inclusões e exclusões), o índice também precisará ser ajustado.
- Uma tabela hash estática, para crescer, precisa reposicionar todos os registros.
- Basicamente, refazer todo o índice!

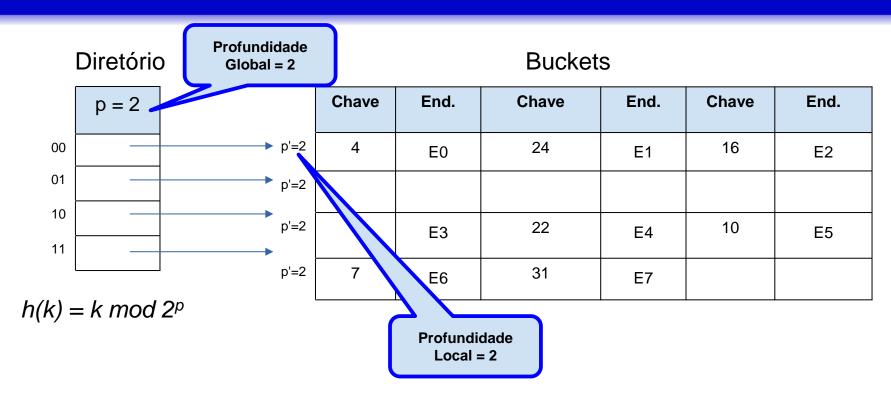
- Uma tabela hash extensível é uma tabela hash em que apenas alguns registros afetados (aqueles do bucket) precisam ser reposicionados.
- O diretório sempre tem tamanho equivalente a uma potência de dois

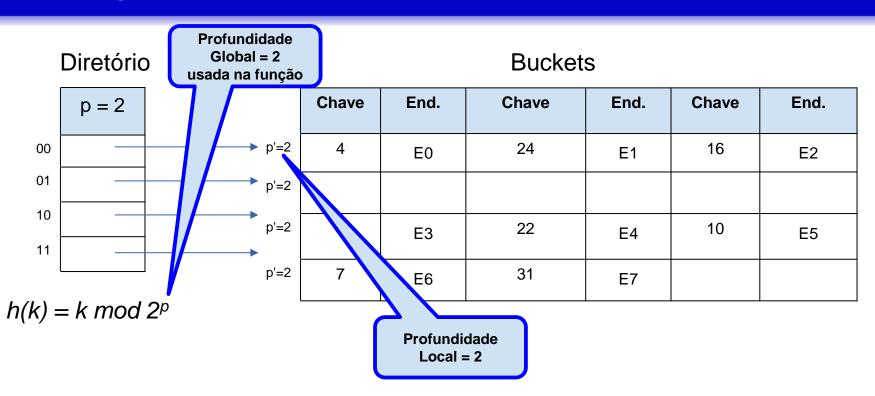
Diretório

Buckets

	Chave	End.	Chave	End.	Chave	End.
2	4	E0	24	E1	16	E2
2						
2	6	E3	22	E4	10	E5
2	7	E6	31	E7		

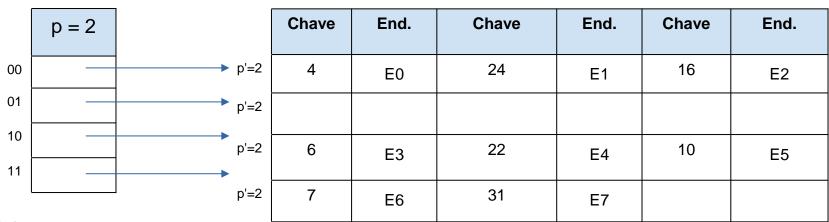
 $h(k) = k \mod 2^p$







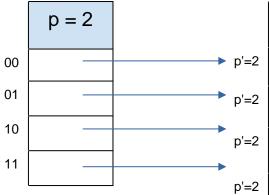
Buckets



 $h(k) = k \mod 2^p$

Função Hash depende de p

Diretório



 $h(k) = k \mod 2^2$

Buckets

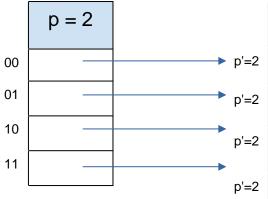
	Chave	End.	Chave	End.	Chave	End.
2	4	E0	24	E1	16	E2
2						
2	6	E3	22	E4	10	E5
2	7	E6	31	E7		

Buckets com total de elementos fixo

Inserção



Diretório



Buckets

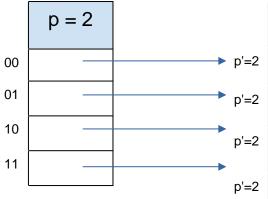
	Chave	End.	Chave	End.	Chave	End.
2	4	E0	24	E1	16	E2
2						
2	6	E3	22	E4	10	E5
₌2	7	E6	31	E7		

 $h(k) = k \mod 2^p$ $h(9) = 9 \mod 2^2$

 $h(9) = 9 \mod 4$

h(9) = 1

Diretório

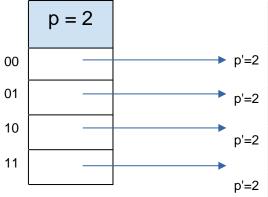


Buckets

	Chave	End.	Chave	End.	Chave	End.
2	4	E0	24	E1	16	E2
2	9	E8				
2	6	E3	22	E4	10	E5
2	7	E6	31	E7		

 $h(k) = k \mod 2^p$ $h(9) = 9 \mod 2^2$ $h(9) = 9 \mod 4$ h(9) = 1

Diretório

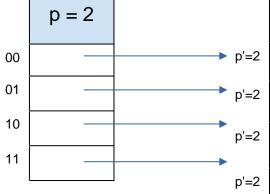


$h(k) = k \mod 2^p$
$h(20) = 20 \mod 2^2$
$h(20) = 20 \mod 4$
h(20) = 0

Buckets

	Chave	End.	Chave	End.	Chave	End.
	4	E0	24	E1	16	E2
	9	E8				
2	6	E3	22	E4	10	E5
: [7	E6	31	E7		

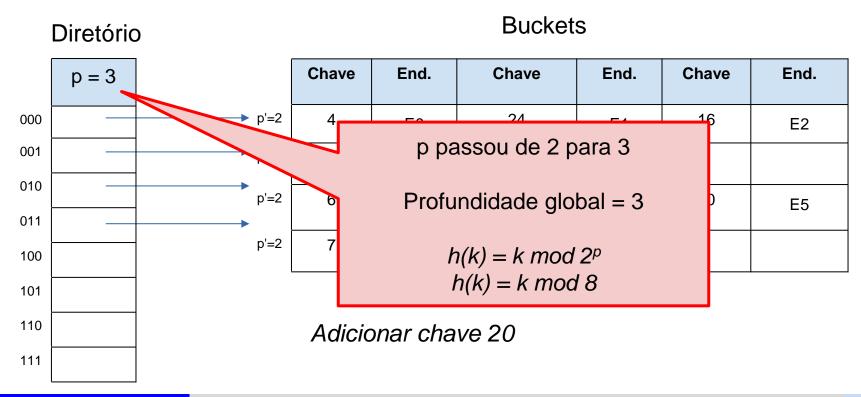
Diretório

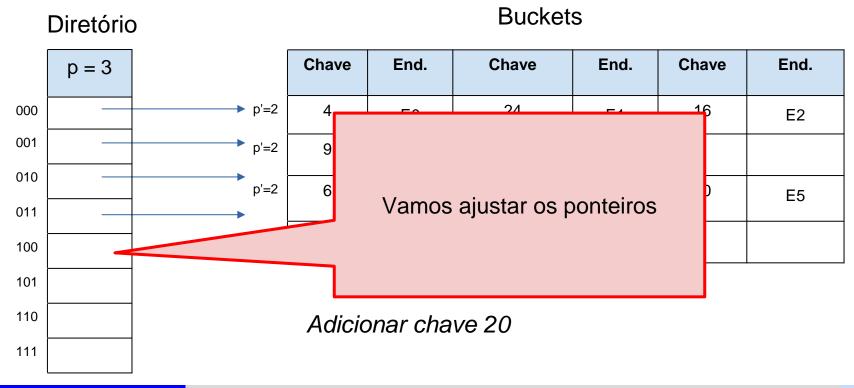


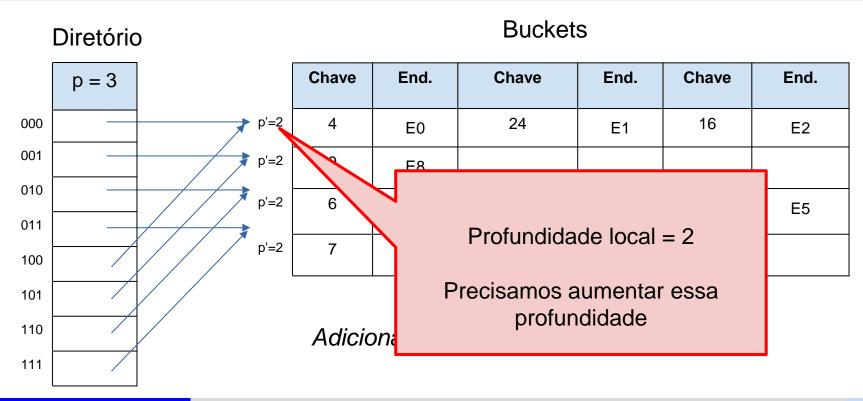
 $h(k) = k \mod 2^p$ $h(20) = 20 \mod 2^2$ $h(20) = 20 \mod 4$ h(20) = 0

Buckets

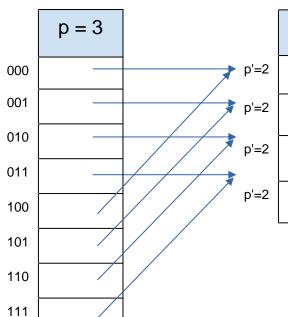
Chave	End.	Chave	End.	Chave	End.				
4	E0	24	E1	16	E2				
9	E8	~	1						
6 E3 22 E5									
Não cabe no bucket 0!									
Adicio	Adicionar chave 20								









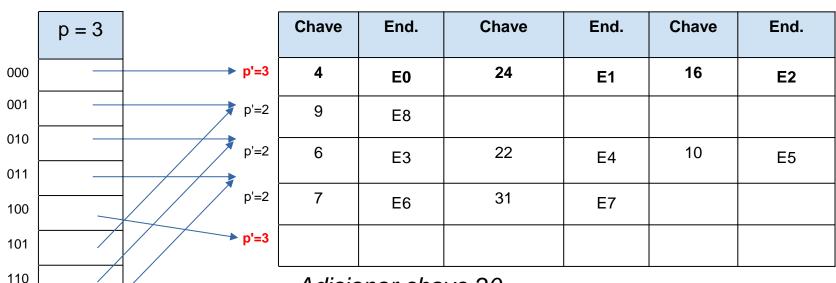


Buckets

Chave	End.	Chave	End.	Chave	End.
4	E0	24	E1	16	E2
9	E8				
6	E3	22	E4	10	E5
7	E6	31	E7		

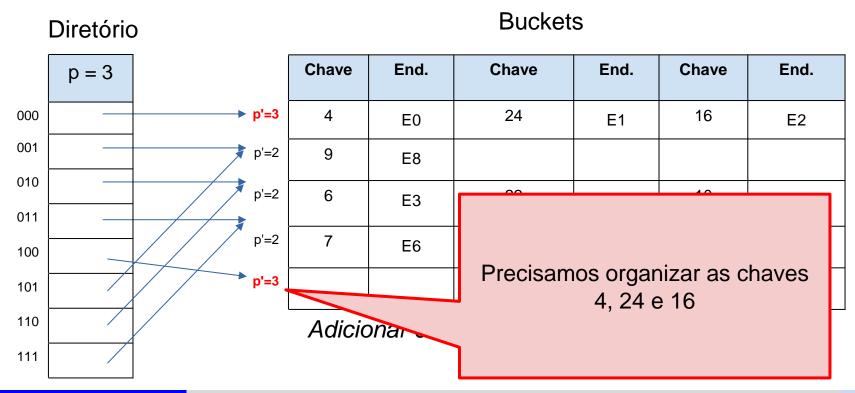


Buckets



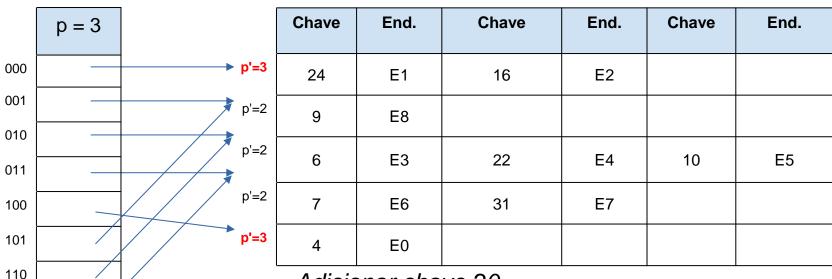
Adicionar chave 20

111



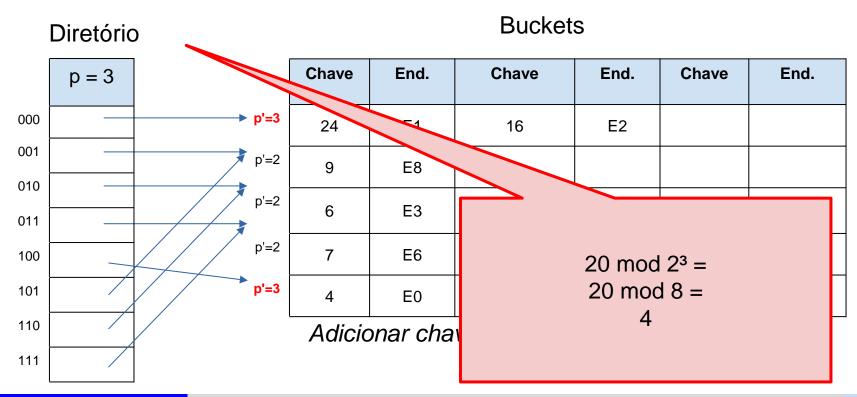


Buckets



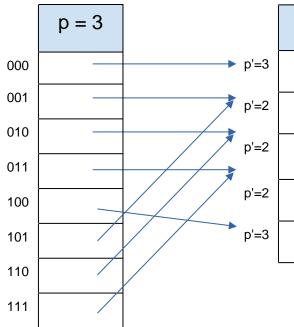
Adicionar chave 20

111





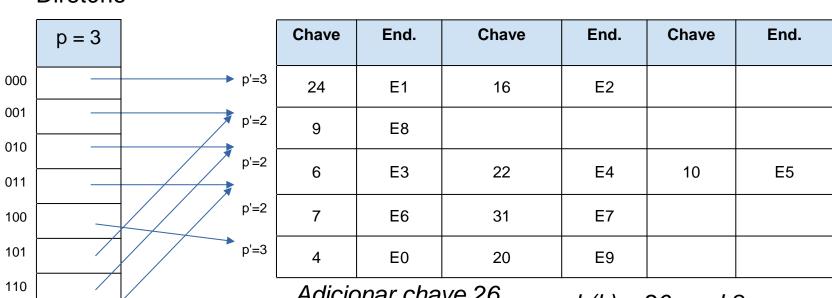




Chave	End.	Chave	End.	Chave	End.
24	E1	16	E2		
9	E8				
6	E3	22	E4	10	E5
7	E6	31	E7		
4	E0	20	E9		



Buckets

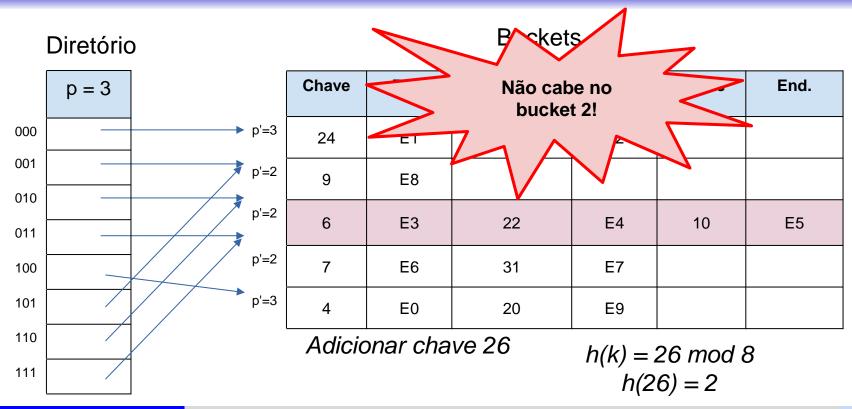


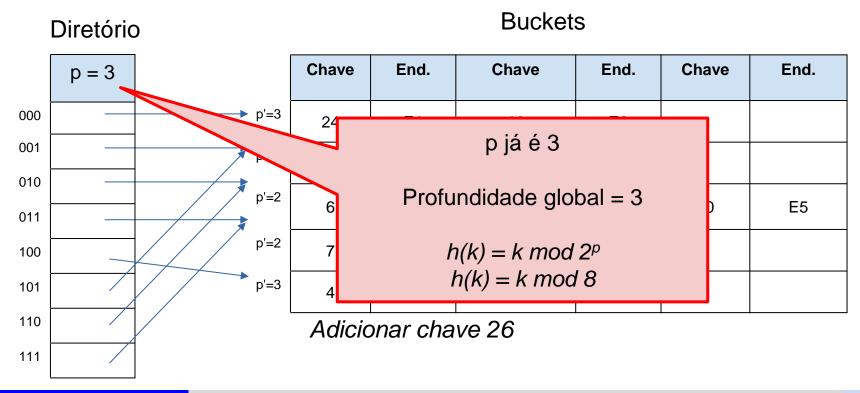
Adicionar chave 26

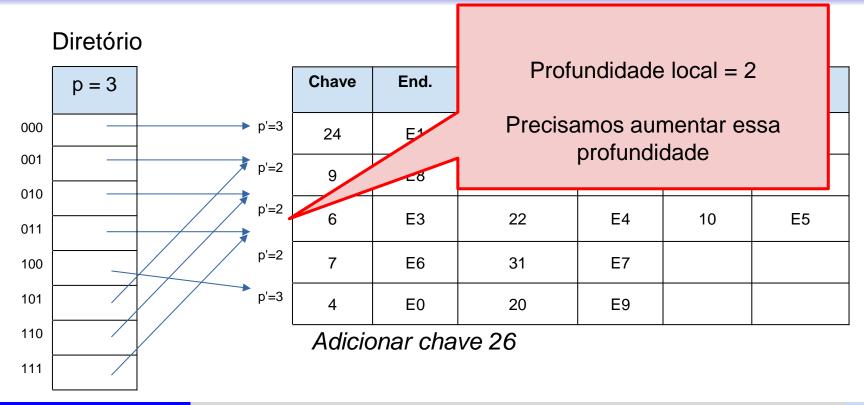
$$h(k) = 26 \mod 8$$

 $h(26) = 2$

111

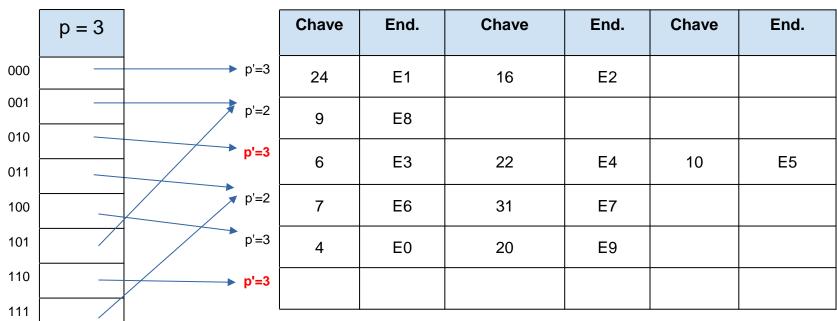


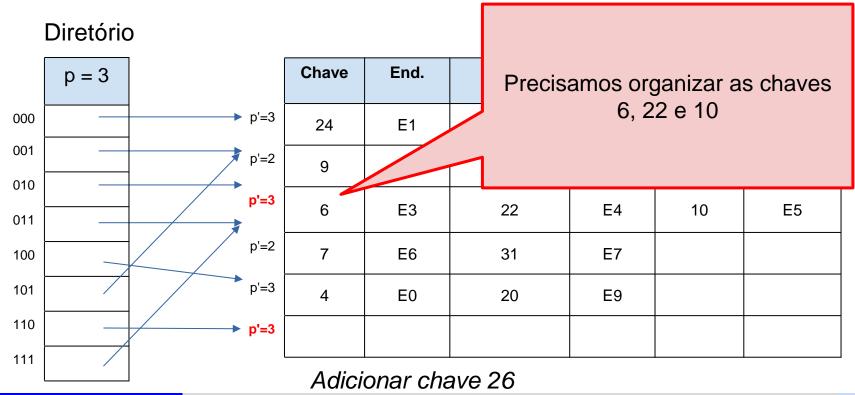






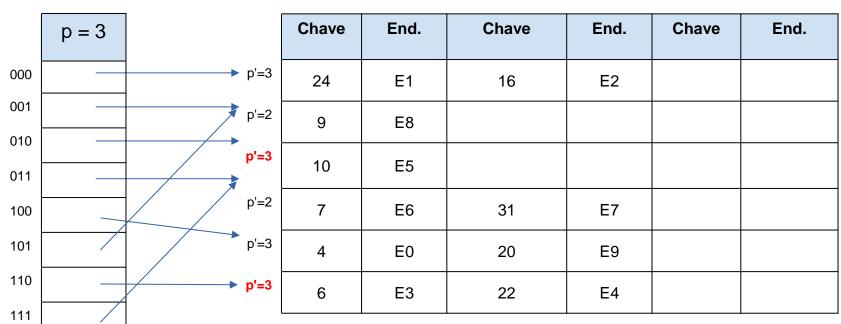
Buckets

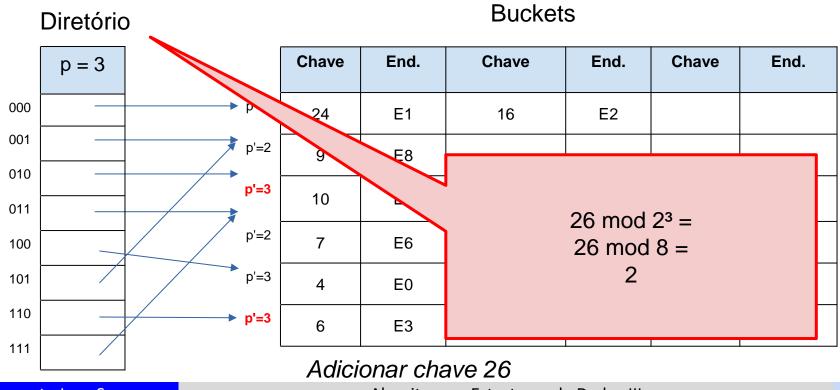






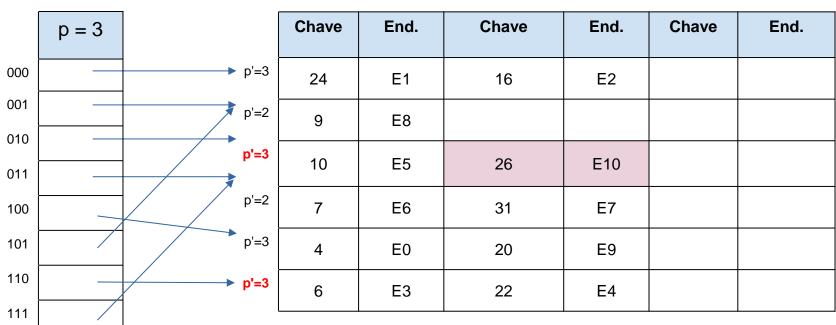
Buckets





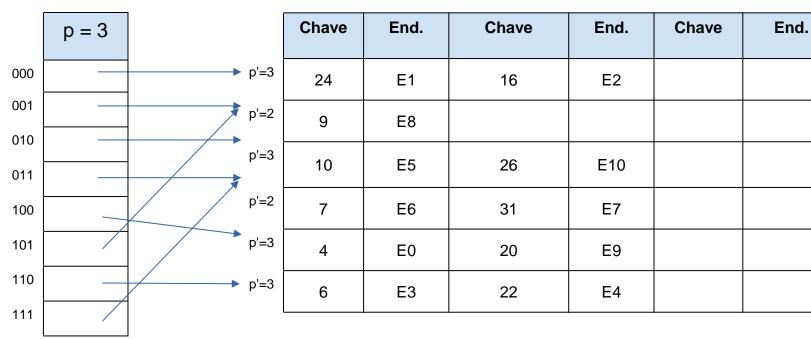


Buckets





Buckets

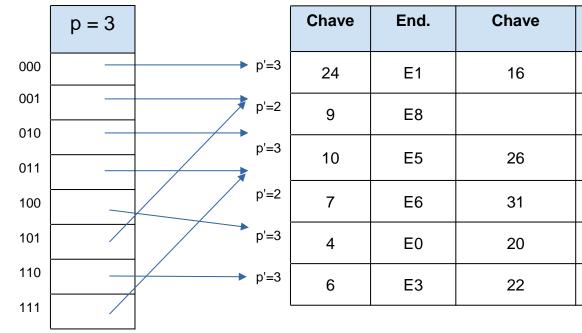


Deletando





Buckets

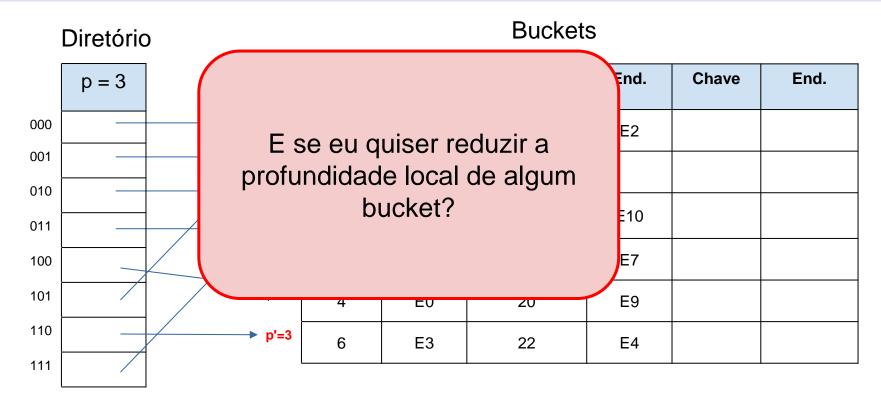


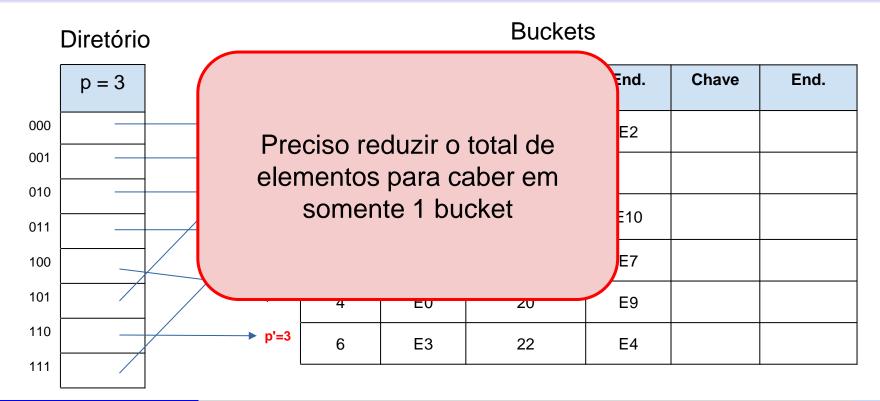
	0110110				
	24	E1	16	E2	
	9	E8			
}	10	E5	26	E10	
	7	E6	31	E7	
	4	E0	20	E9	
	6	E3	22	E4	

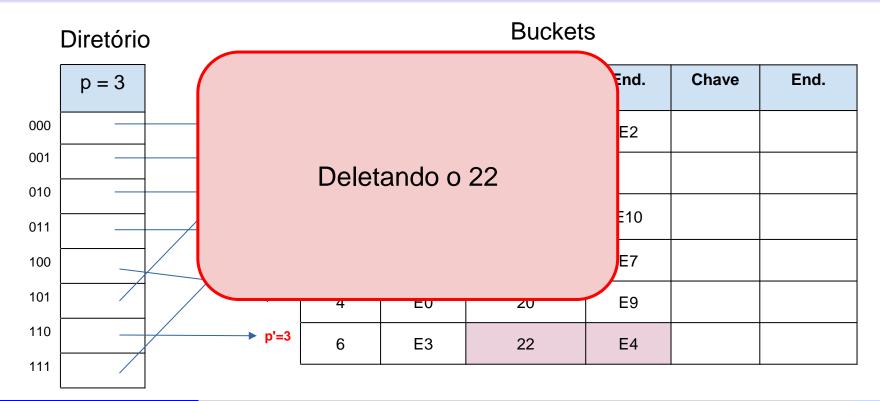
End.

Chave

End.

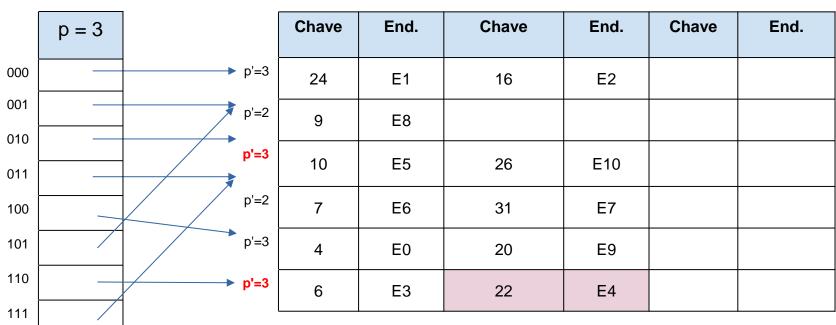






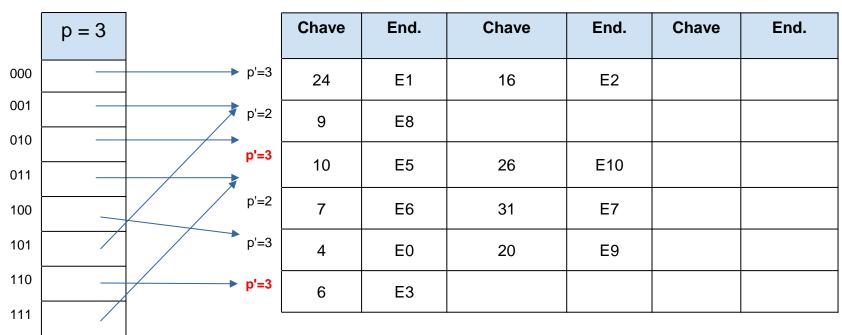


Buckets

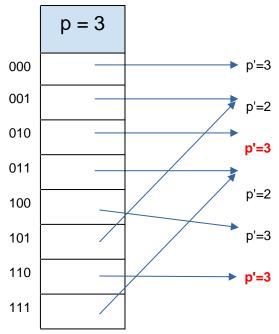




Buckets



Diretório



Chave

24

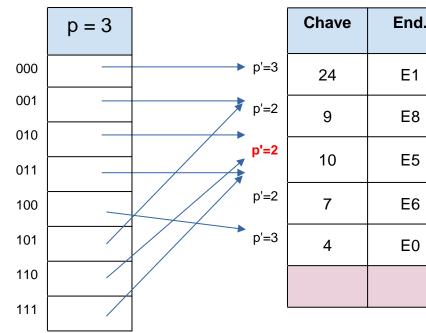
Juntando os buckets 2 e 6

Tenho 3 elementos e o bucket cabe 3 elementos

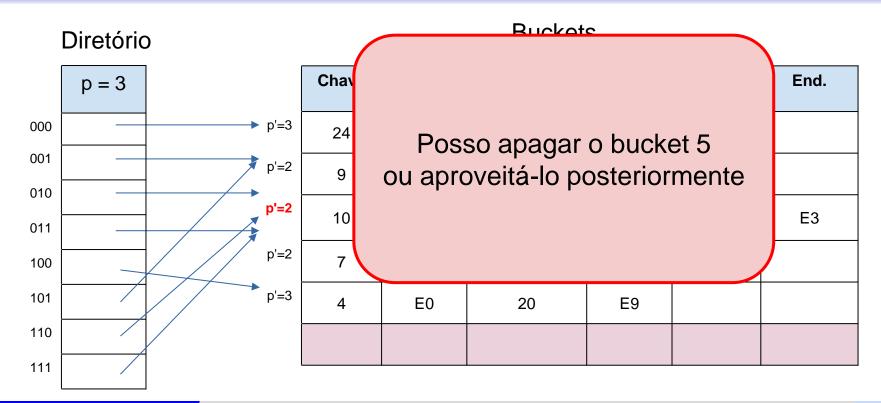
9 E8					
7 E6 31 E7 4 E0 20 E9	9	E8			
4 E0 20 E9	10	E5	26	E10	
	7	E6	31	E7	
6 E3	4	E0	20	E9	
	6	E3			



Buckets



Chave	End.	Chave	End.	Chave	End.
24	E1	16	E2		
9	E8				
10	E5	26	E10	6	E3
7	E6	31	E7		
4	E0	20	E9		



Conclusões



- O diretório cresce, sem precisarmos reposicionar todos os registros (do índice)
- O índice (lista de buckets) cresce de acordo com a necessidade
- Como não há encadeamento dos buckets, não há perda de eficiência

Qual a melhor estrutura auxiliar para índices:

Árvore B, B+, B* ou Hash Dinâmico?

Suponha que façamos a inserção das chaves 30, 21, 12, 42 e 31, nesta ordem. Que chave provocará o primeiro aumento de profundidade do diretório?

///_

