

COLLABORATIVE INSURANCE SUSTAINABILITY AND NETWORK STRUCTURE

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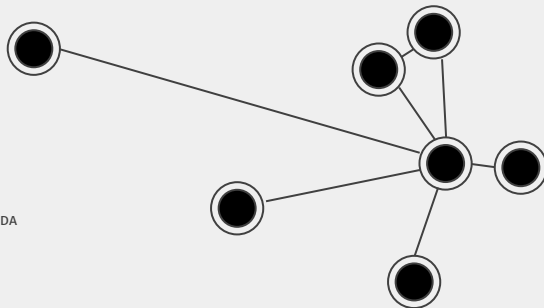
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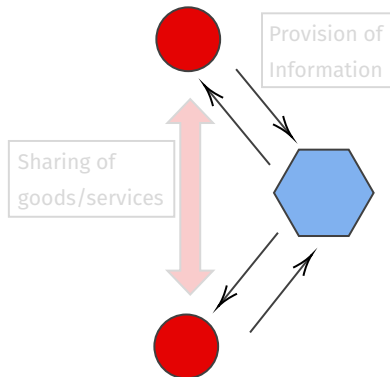
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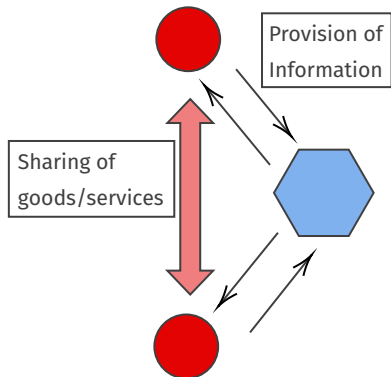
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THE MAIN IDEA: SHARING DEDUCTIBLES WITH FRIENDS



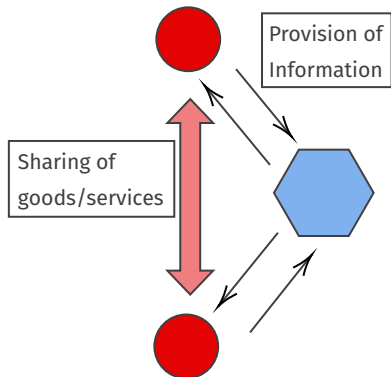
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 - ▶ Degree distribution and random graphs

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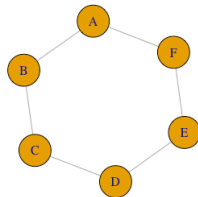
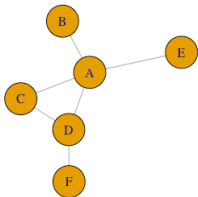


- In a nutshell - we will consider different risk sharing models between peers in a network and analyze the robustness of the mechanisms
- A (very) short primer on graph-terminology
 - ▶ Degree distribution and random graphs
- Risk sharing on a network
 - ▶ Reciprocal engagements & a simple example
 - ▶ Extensions for perceived fairness, solutions for graphs with high degree variance and extensions to friends-of-my-friends

Degree distribution

The degree of a vertex is the number of edges that are incident to it. The average degree of a graph is denoted as \bar{d} . The degree vector for each vertex can be calculated by taking $A\mathbf{1} = d$. Note that the average degree of a network is:

$$\mathbb{E}(D) = \frac{1}{|\mathcal{V}|} \sum_{v \in \mathcal{V}} d(v) = \frac{2|\mathcal{E}|}{|\mathcal{V}|} = \frac{1}{|\mathcal{V}|} \|d\|_1$$



Erdős-Rényi

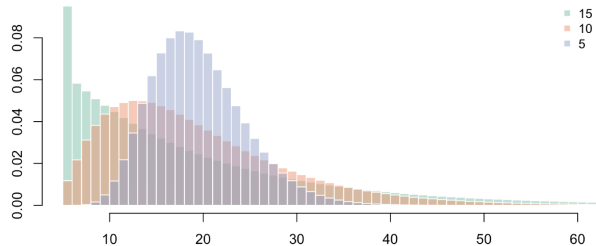
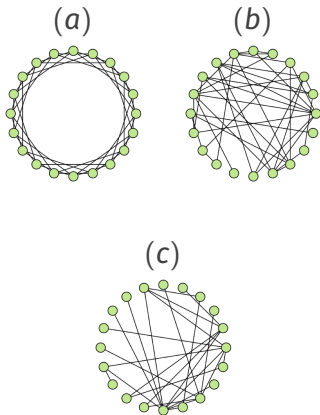
Each edge $i, j \in \mathcal{V} \times \mathcal{V}$ is included in the network with probability p independent from other edges. The Network then has a degree distribution that follows a binomial distribution $\mathcal{B}(n-1, p)$. Also, if p not too large but n_V :
 $p \sim \lambda/n_V, d(v) \sim \mathcal{P}(\lambda)$ See [5] for details.

Preferential attachment

To start with, consider a small network, and create a new node at every time-step. During every time-step, the new node is connected to the existing nodes with a probability p that is proportional to the degree of an existing node d_i . This way, every new node is more likely to connect to existing "popular" nodes. See [1] or [6] for details. Here the degree distribution follows a power law.

For a discussion on the nature of networks see eg. [3]

RANDOM NETWORK GENERATION



SETUP OF STUDY

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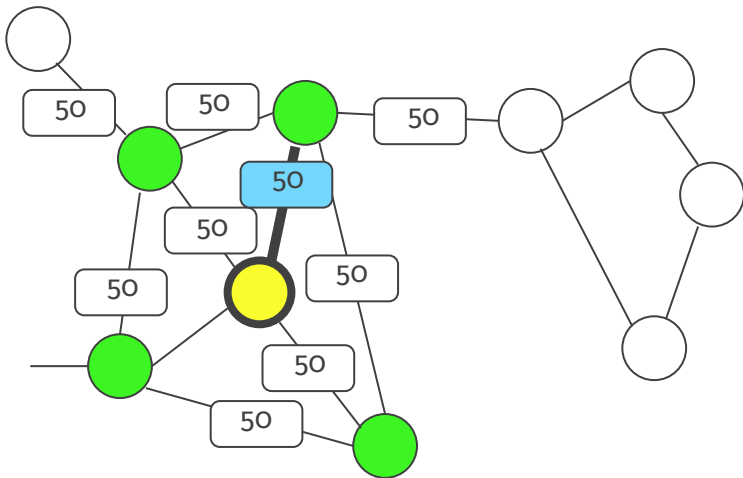
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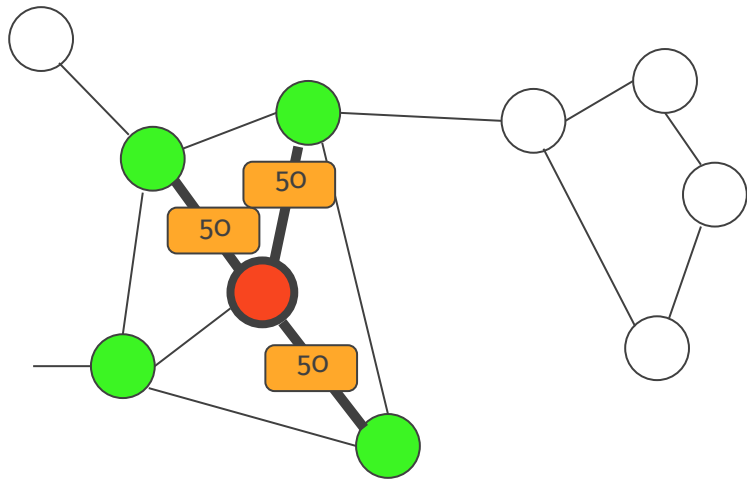
The random cost for a member after one period can then be summarized as:

$$X_i = Z_i \cdot \min\{s, Y_i\} = \begin{cases} 0 & \text{if no claim occurred } (Z_i = 0) \\ \min\{s, Y_i\} & \text{if a claim occurred } (Z_i = 1) \end{cases}$$

RISK SHARING ON A NETWORK



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Risk and reciprocal engagements

Let \mathcal{V}_i denote the set of friends (the set of nodes that are connected to node i). Also assume fixed contributions γ . The cost for node i then becomes:

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which can be shown to be a risk sharing principle (for a definition of a risk-sharing principle see eg. [4])

A SIMPLE EXAMPLE

Generate network with $n = 5000$ nodes, with degree distribution

$$D \stackrel{\mathcal{L}}{=} \min\{5 + [\Delta], n - 1\},$$

where $[\Delta]$ has a rounded Gamma distribution with mean $\bar{d} - 5$ and variance σ^2 .
Let σ vary between 0 and $4\bar{d}$ to take into account different network structures

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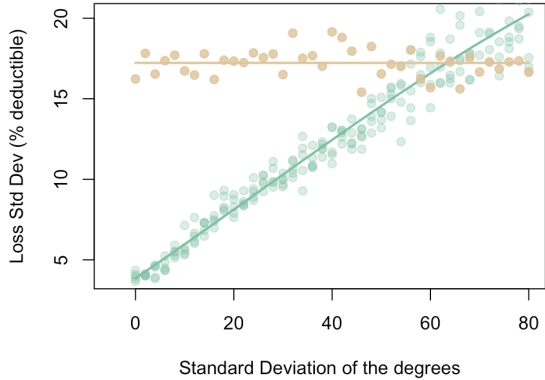
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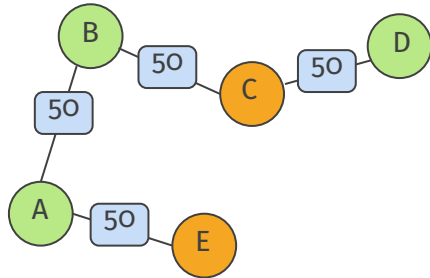
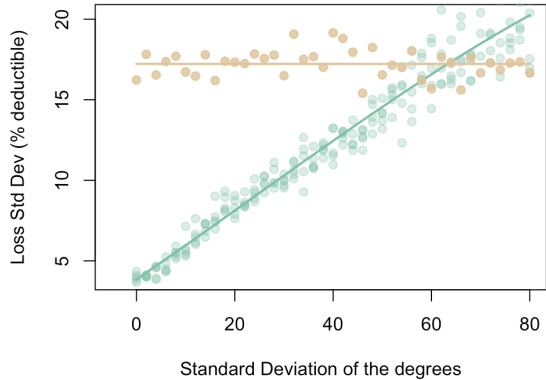
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Consider a risk sharing model where the reciprocal engagements weigh $\gamma = 50$ and $\mathbb{E}(X) = 4.5\%$ of s . The losses are distributed over the network via the Bernoulli mechanism

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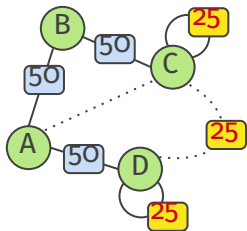
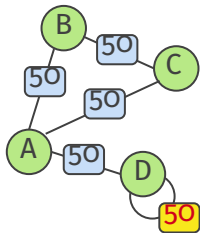
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EXTENSIONS OF THE SIMPLE FRAMEWORK

We can consider some extensions

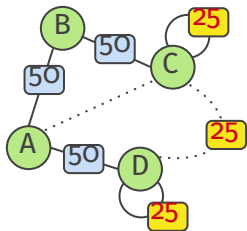
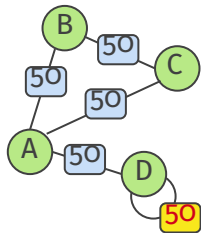
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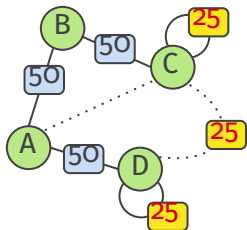
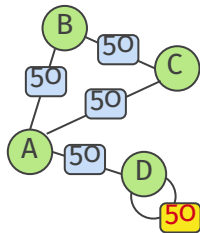
- Introducing self-contribution to ensure fairness in the mechanism and alleviate the variance issue
- Calculate optimal reciprocal engagements subject to some constraint - to take into account the structure of the network. This can be combined with engagements which are restricted to a subgraph (intuitively - we can cut some edges by setting the engagement to zero)



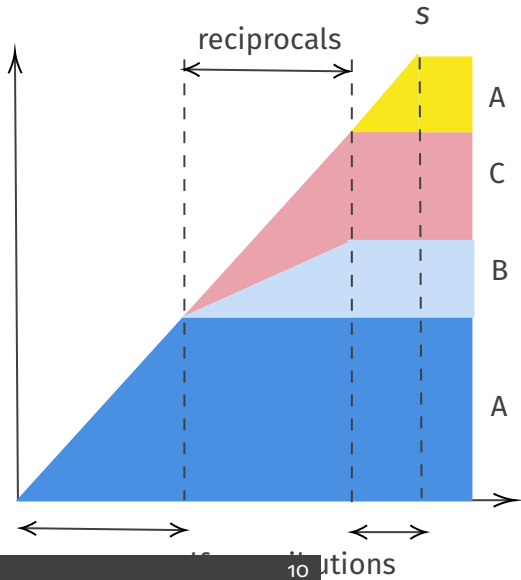
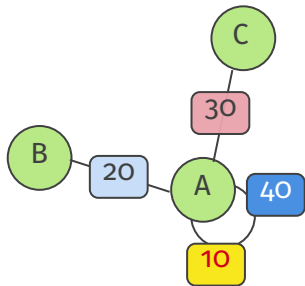
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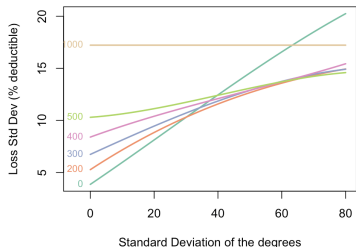
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- Extending the reciprocal engagements to friends-of-friends to expand the set of possible edges for nodes with only a few links



RECIPROCAL ENGAGEMENT WITH SELF-CONTRIBUTION



RISK SHARING WITH SELF-CONTRIBUTION



- It can be shown that in many cases, it is likely that the nodes who actually claim a loss will end up paying *less* than its adjacent vertices. The system seems fair ex-ante but ex-post (after the occurrence of the claim) it might seem unfair
- Self contribution makes a "fair" ex-post outcome more likely

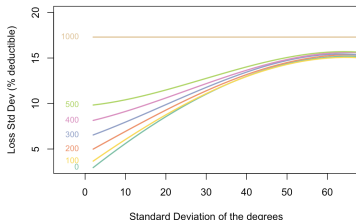
RISK SHARING WITH OPTIMAL RECIPROCAL ENGAGEMENTS

Instead of a fixed γ we can also formulate the γ -vector as a linear programming optimization problem. This can lead to smaller total contributions, especially for "popular" nodes:

$$\gamma_i^* = \operatorname{argmax} \left\{ \sum_{(i,j) \in \mathcal{E}} \gamma_{(i,j)} \right\}$$

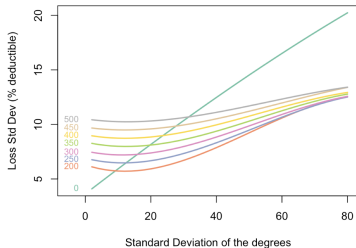
$$\text{s.t. } \gamma_{(i,j)} \in [0, \gamma], \forall (i,j) \in \mathcal{E}$$

$$\sum_{j \in \mathcal{V}_i} \gamma_{(i,j)} \leq s, \forall i \in \mathcal{V}$$



In words, this means we have a limited liability s and our objective is to maximize the overall coverage for our network. Sparsity of the solution can also be achieved by reformulating the problem into a mixed integer problem.

RISK SHARING WITH FRIENDS-OF-FRIENDS



A further simple extension is to consider not only the adjacent nodes but also their adjacent nodes (friends-of-my-friends). This will naturally increase the number of connections. After the first stage (optimal reciprocal engagements) nodes not yet fully covered can satisfy their demand with the friends of their friends (typically γ_2 can be smaller than γ_1)

$$\gamma_2^* = \operatorname{argmax} \left\{ \sum_{(i,j) \in \mathcal{E}^{(2)}} \gamma_{(i,j)} \right\}$$

$$\text{s.t. } \gamma_{(i,j)} \in [0, \gamma_2], \forall (i,j) \in \mathcal{E}^{(2)}, \mathcal{E}^{(2)} \text{ from } A^2$$

$$\sum_{j \in \mathcal{V}_i^{(1)}} \gamma_{1:(i,j)}^* + \sum_{j \in \mathcal{V}_i^{(2)}} \gamma_{(i,j)} \leq s, \forall i \in \mathcal{V}$$

CONCLUDING REMARKS

- Using Networks/Graphs allows us to tap into some hidden and interesting information (eg. [2]) and can be employed with novel digital techniques to enhance transparency
- An insurance product based on networks can be considered a risk sharing mechanism
- The proposed basic mechanism can be extended in many ways

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But...

- The proposed model is inherently static
- The model (as of yet) also neglects some locality properties such as higher-level groups

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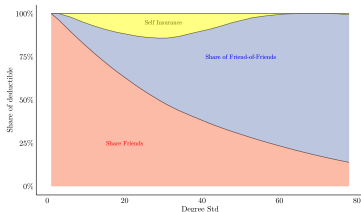
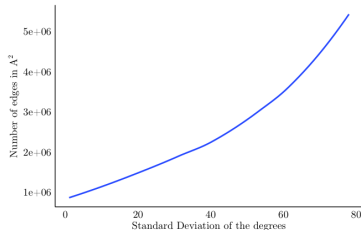
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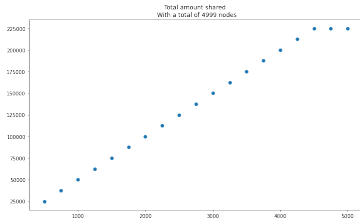
RISK SHARING WITH FRIENDS-OF-FRIENDS



In this setup, the computational cost increases drastically with the degree variance. First, because more is left over from the first part, but also because the number of edges in A^2 increases.

But, depending of the level of γ_2 , this property actually ensures that the network mechanism is still able to allocate most of the risk (as the number of edges A^2 runs in the opposite direction as the amount of risk "sharable" in A^1)

SPARSIFICATION OF RESULTS



By changing the problem from its linear form to its mixed-integer form, we can achieve the desired level of sparsity. M comes naturally here and actually makes the problem slightly simpler in terms of memory usage. The important part here part here is the m parameter which restricts the number of nonzero edges.

$$\begin{cases} \mathbf{z}^* = \operatorname{argmax}_{\mathbf{z} \in \mathbb{R}_+^m} \{ \mathbf{1}^\top \mathbf{z} \} \\ \text{s.t. } \mathbf{T}^\top \mathbf{z} \leq s \mathbf{1}_m \\ z_i \leq M y_i \\ \sum_i y_i \leq m \end{cases}$$

, where $y_i \in \{0, 1\}$