S: $r_1(x)\, r_1(y)\, r_2(y)\, r_3(y)\, w_2(x)\, r_1(z)\, w_2(z)\, w_1(y)\, r_1(z)\, w_3(y)$

VSR ?

Legge da $\{ r_1(z) \text{ ld } w_2(z) \}$

Scr finale $\{ w_3(y), w_2(z), w_2(x) \}$

T_1: $r_1(x)\, r_1(y)\, r_1(z)\, w_1(y)\, r_1(z)$

T_2: $r_2(y)\, w_2(x)\, w_2(z)$

T_3: $r_3(y)\, w_3(y)$

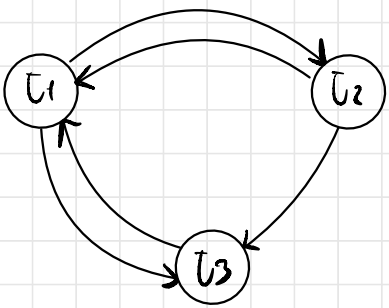
Com. sole $T_2\, T_1\, T_3$

$r_2(y)\ \underline{w_2(x)}\ w_2(z)\ \underline{r_1(x)}\ r_1(y)\ r_1(z)\ \underline{w_1(y)}\ r_1(z)\ \underline{r_3(y)}\ w_3(y)$

ha generato $r_1(x)$ legge da $w_2(x)$ che non era presente $\Rightarrow$ non è in VSR

c) Dato lo Schedule r1(x)r1(y)r2(y)r3(y)w2(x)r1(z)w2(z)w1(y)r1(z)w3(y)

- CSR ?
- VSR ?
- 2PL produce attese o stalli ?



=> NON è CSR

Legge-da { $r_1(z)$ legge da $w_2(z)$ }

Scritti finali { $w_3(y)$, $w_2(z)$, $w_2(x)$ }

T_1 : $r_1(x)\ r_1(y)\ r_1(z)\ w_1(y)\ r_1(z)$

T_2 : $r_2(y)\ w_2(x)\ w_2(z)$

T_3 : $r_3(y)\ w_3(y)$

=> NON è VSR

r1(x)r1(y)r2(y)r3(y)w2(x)r1(z)~~w2(z)~~w1(y)~~r1(z)~~w3(y)

A A A

	r. lock	w. lock
X	T_1	
Y	T_1 T_2 T_3	
Z	T_1	

T_2 attende T_1

T_1 attende T_2 e T_3

T_3 attende T_1 e T_2

⇒ c'è Deadlock