



AGUDA is an imperative programming language composed of expressions alone.

A program in AGUDA is a sequence of declarations, each announced with keyword **let**. Value declarations introduce an identifier and a type. For example:

```
let maxSpeed : Int = 120
```

Declarations may also introduce functions, as in

```
let succ (n) : Int -> Int = n + 1
```

Functions can be recursive:

```
let add (n, m) : (Int, Int) -> Int =  
  if n == 0 then m else succ(add (n - 1, m))
```

or even mutually recursive:

```
let even (x) : Int -> Bool =  
  x == 0 || odd(x - 1)
```

```
let odd (x) : Int -> Bool =  
  x != 0 && even(x - 1)
```

Imperative means state changing. How do we change state? By means of (imperative) variables and assignment. Here's addition in imperative style:

```
let AddI (n, m) : (Int, Int) -> Int =  
  let sum : Int = m;  
  while n > 0 do (  
    set sum = sum + 1;  
    set n = n - 1  
  );  
  sum
```

A few points to notice:

- The body of the function is composed of three *expressions*, separated by a semicolon;
- The value of the function is given by the last expression in the semicolon separated sequence, namely, `sum`;

- The body of the **while** expression is a sequence of two expressions; we enclose them in parenthesis.
- Assignments are announced by the **set** keyword

Each expression has a value, including the **while** loop. The value of **while** is **unit**, the only value of type **Unit**. This means that we can “store” a while loop in a variable, as in

```
let x : Unit = while i > 0 do set i = i - 1
```

or “return” a while loop from a function, as in

```
let f (n) : Int -> Unit = while n > 0 do set n = n - 1
```

An assignment stores a value in a variable and returns the value. Given the below declaration, a call to `f (unit)` yields value 3.

```
let f (x) : Unit -> Int = let x : Int = let y : Int = 3
```

Equipped with the **unit** value, an if-then expression **if** `exp1` **then** `exp2` is an abbreviation for an if-then-else **if** `exp1` **then** `exp2` **else** **unit**.

Let us now look at the support provided for arrays. We start with an example: creating an  $n \times n$  matrix, where all entries are 0, except for the diagonal which is filled with 1.

```
let diagonal (n) : Int -> Int[][] =
  let a : Int[][] = new Int[] [n | new Int [n | 0]] ;
  let i : Int = 0 ;
  while i < length(a) do (
    set a[i][i] = 1 ;
    set i = i + 1
  ) ;
  a
```

Expression `new Int [n | 0]` creates an integer array of size `n`, each cell initialised to 0. Likewise, `new Int[] [n | new Int [n | 0]]` creates a  $n \times n$  matrix of zero values. Expression `set a[i][i] = 1` writes 1 in the `i`-th position of the `i`-th array of matrix `a`, that is, in line `i`, column `i` of `a`.

Below is a function that prints a matrix, line by line, each value terminated with a space. It includes two *primitive* functions: **length** returns the number of elements in an array; **print** sends to the prints `stdout` the textual representation of a given value. Both functions are overloaded: **length** works on any array; **print** accepts any value. More primitive functions shall be announced later.

```
let printMatrix (a) : Int[][] -> Unit =
  let i : Int = 0 ;
  while i < length(a) do (
```

```

while j < length(a[0]) do (
  print(a[i][j]) ; print(" ") ;
  set j = j + 1
);
print("\n") ;
set i = i + 1
)

```

A “little” main function exercises the two array-related functions.

```

let main : Unit =
  printMatrix(diagonal(10))

```

## Expressions

- Variable: `id`
- Literals: `...`, `-1`, `0`, `1`, `...`, `true`, `false`, `unit`, in addition to quote-enclosed string literals
- Binary operators: `;` `+` `-` `*` `/` `%` `^` `==` `!=` `<` `<=` `>` `>=` `!` `||` `&&`
- Unary operators: `-` `!`
- Function call: `id(exp1, ..., expn)` with  $n \geq 1$
- Assignment: `set LHS = exp`
- Variable declarations: `let id : type = exp`
- Conditionals: `if exp1 then exp2 else exp3` and `if exp1 then exp2`
- While loop: `while exp1 do exp2`
- Array creation: `new type [ exp1 | exp2 ]`
- Array access: `exp1[exp2]`
- Parenthetical expression: `(exp)`

The left-hand-side (LHS) of an assignment, that is, the part at the left of `=`, can be:

- A variable
- An array location: `LHS[exp]`

## Operator precedence and associativity

- For `+` `-` `*` `/` `%` `==` `!=` `<` `<=` `>` `>=` `!` `||` `&&` take the precedence and associativity of the Java programming language
- The unary minus binds tighter than any other arithmetic operator
- The power operator, `^`, is right associative and binds tighter than any other arithmetic operator (different from unary minus). For example, `2 ^ - 3 ^ 4 * 5` is to be understood as `(2 ^ ((-3) ^ 4)) * 5`
- The sequencing operator, `;`, associates to the right and binds loser than all other operators. For example, `1 ; 2 || 3 ; 4` is to be understood as `1 ; ((2 || 3) ; 4)`
- The precedence of **while** loops and conditionals (both **if-then-else** and **if-then**) seats between that of `;` and `||`. For example, **while** `b` **do** `false || true` should be understood as **while** `b` **do** `(false || true)`, but **while** `b` **do** `false ; true` should be understood as `(while b do false) ; true`. Similarly for conditionals
- Keywords **then** and **else** associate to the right, so that **if** `a` **then** **if** `b` **then** `c` **else** `d` is to be understood as **if** `a` **then** `(if b then c else d)`
- Array access, `[` binds tighter than unary minus, so that `-a[0]` denotes `-(a[0])`
- The arrow type operator, `->`, associates to the right

## Top-level declarations

- Variables: as in expressions
- Functions: **let** `id` (`id1, ..., idn`) : `type` = `exp` with  $n \geq 1$

**Programs** A non-empty sequence of declarations.

## Types

- Basic: **Int**, **Bool**, **Unit**, **String**
- Array: `type[]`
- Function: `type -> type` or `(type1, ..., typen) -> type` with  $n \geq 1$ . There are no zero-ary functions. If needed, use `f` `(_)` : **Unit** `-> type` = `exp` and call as `f(unit)`.

## Lexing

- Identifiers (variable or function names) start with a letter and are followed by zero or more letters, digits, underscore symbols (`_`) and single quotes (`'`)
- Integer values start with an optional sign, followed by a non-zero digit and then followed by zero or more digits
- Strings are sequences of characters (not including new line), enclosed in quotes (`"`)
- Comments: Line comments only, starting with `--`.

## Wildcards

- Wildcards may appear only in binding positions:  
`let _ : type = exp`
- Hence, cannot appear in expressions: `2 * _`
- Wildcards may appear more than once, even if with different types:  
`let f (_, _) : (Unit, Int) -> Int = 5`

**Typing** The types for expressions are as follows.

- The type of a literal is the corresponding type. For example, the type of 5 is `Int`
- The type of array creation `new type [ exp1 | exp2 ]` is `type[]`. Furthermore, `exp1` should be of type `Int` and `exp2` should be of type `type`
- The type of array access `exp1 [exp2]` is `type` if `exp1` is of type `type[]` and `exp2` of type `Int`
- The type of call `print (exp1, ..., expn)` is `Unit` if `n = 1` and `exp1` has a type (any type)
- The type of call `length (exp1, ..., expn)` is `Int` if `n = 1` and `exp1` has a type `type[]`
- The type of call `id (exp1, ..., expn)` (with `id ≠ print, length`) is `type` if `id` is of type `(type1, ..., typem) -> type` and `n = m` and each `expi` has type `typei`
- The type of variable declaration `let id : type = exp` is `Unit` if `exp` is of type `type`. Furthermore, if the declaration appears at the left of a semicolon `let id : type = exp1 ; exp2`, then the type of `id` is used to validate `exp2`

- The type of an identifier is that more recently introduced by a **let** (either expression or top-level **let** declaration)
- The type of conditional **if**  $\text{exp1}$  **then**  $\text{exp2}$  **else**  $\text{exp3}$  is **type** if  $\text{exp1}$  is of type **Bool** and both  $\text{exp2}$  and  $\text{exp3}$  are of type **type**. In the case of **if**  $\text{exp1}$  **then**  $\text{exp2}$ , expression  $\text{exp2}$  must be of type **Unit**
- The type of while loop **while**  $\text{exp1}$  **do**  $\text{exp2}$  is **Unit** if  $\text{exp1}$  is of type **Bool** and  $\text{exp2}$  has a type (any type)
- The type of assignment **let**  $\text{LHS} = e$  is **Unit** if  $\text{LHS}$  and  $e$  share the same type
- The type of sequential composition  $\text{exp1}; \text{exp2}$  is the type of  $\text{exp2}$  if  $\text{exp1}$  has a type (any type)
- The type of the remaining operators (binary and unary) are as customary in programming languages

The types for LHS are as follows.

- Variables, as in expressions
- The type of  $\text{LHS}[\text{exp}]$  is **type** if  $\text{LHS}$  is of type **type[]** and  $\text{exp}$  is of type **Int**

Top-level declarations have no types. The validation rules are as follows.

- Variables, as in expressions
- Function declaration **let**  $\text{id}(\text{id1}, \dots, \text{idn}) : \text{type} = \text{exp}$  is valid if  $\text{type} = (\text{type1}, \dots, \text{typen}) \rightarrow \text{type}'$  and  $n = m$  and  $\text{exp}$  is of type **type'** in a context augmented with  $\text{id1}:\text{type1}, \dots, \text{idn}:\text{typen}$
- No variable of function may be named **print** or **length**
- Function declarations may be (mutually) recursive
- Functions and top-level variables cannot be declared twice, even if with different signatures

The validation rule for programs is as follows. A program is valid if:

- The identifiers of all its (top level) declarations are pairwise distinct and
- All its (top level) declaration are valid and
- There is a function with signature **let**  $\text{main}(x) : \text{Unit} \rightarrow \text{Unit}$

**Semantics** A further restriction for the semantics:

- An expression `exp` in a program (top level) variable declaration `let id : type = exp` should be a literal (a compile time constant)

An AGUDA program is executed starting from a function call `main (unit)`. Here's a few pointers.

- Function arguments are evaluated left-to-right until they all become values. The corresponding function is then called (AGUDA is call-by-value)
- Binary expressions `exp1 op exp2` evaluate `exp1` and `exp2` (in this order) and then apply the corresponding operator
- The exception are boolean expressions. These are evaluated in a *short-circuit* manner: they are evaluated left-to-right until the truth value of the expression may be determined. In particular, in expression `exp1 && exp2`, if `exp1` turns out to be false, then `exp2` is *not* evaluated. Dually for disjunction
- Variable declaration (global or local) and function parameters introduce storage to hold values of the corresponding type
- For scalar (non-array) types, the value in the store may be read via the identifier expression `id` and updated via expression `set id = exp`
- The value of `set id = exp` is `unit`
- A conditional expression `if exp1 then exp2 else exp3` evaluates to `exp2` or to `exp3` according to the truth value of expression `exp1`. The value of the conditional is the value of `exp1` or `exp2`
- A `while exp1 do exp2` loop evaluates `exp2` while `exp1` remains true. The value of the loop is `unit`

We refrain from describing the semantics of arrays for we shall not implement them.