

Bidirectional Search Algorithm

Lam

Breadth first Search (BFS)

- It is the most common Search strategy to traversing a tree/graph.
- This algorithm searches breadthwise in a tree & graph, so it is called breadth-first search.
- BFS algorithm starts searching from the root node of the tree and expands all success node at the current level before moving to the next level.
- BFS algorithm is an example of general-graph search algorithm.
- BFS implemented using FIFO queue data structure.

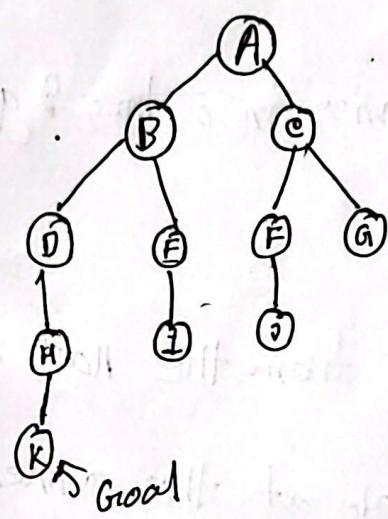
Advantages:

- BFS will provide a solution if any solution exists.
- If there are more than one solution in a given problem, the BFS will provide minimal solution which requires the least no of steps.

Disadvantages:

- It requires lot of memory since each level of the tree must be saved into memory to expand the next level.
- BFS needs lots of time if the soln is far away from the root nodes.

Example BFS



Time complexity

$d = \text{depth of shallow}$

$b = \text{node of every state}$

$$T(b) > 1 + b^2 + b^3 + \dots + b^d \\ = O(b^d)$$

$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H \rightarrow I \rightarrow J \rightarrow K \quad \text{Space} = O(b^d)$

* BFS is complete.

optimality: optimal

Bidirectional Search Algorithm

→ Two different searches are run simultaneously. Searching starts from both ends. (forward search and backward search)

$S \rightarrow G$

$G \rightarrow S$

→ Here single search graph is replaced with two small graphs.

→ Any search technique can be used (BFS, DFS, ...)

→ When graphs intersects with each other it stops searching

Advantages

→ Fast

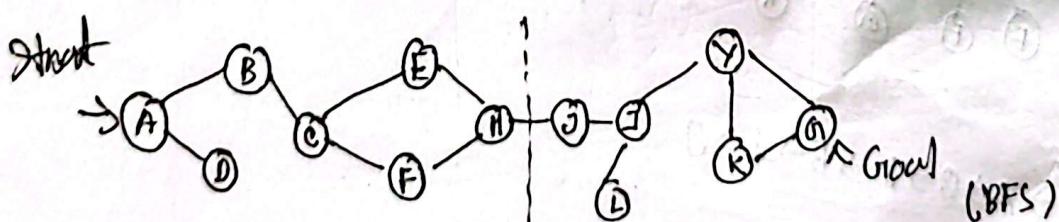
→ Less memory

Disadvantages:

→ Implementation is difficult

→ Goal state should be known in advance

Example:



Forward

ABDC~~E~~ FH

Backward

GIKXILJH

Path: $A \rightarrow B \rightarrow D \rightarrow C \rightarrow F \rightarrow H \rightarrow J \rightarrow L \rightarrow I \rightarrow X \rightarrow K \rightarrow G$

Depth-Limited Search Algorithm

→ Working is similar to DFS but with predetermined limit.

→ Helps in solving the problem of DFS (Infinite Path)

Termination Conditions:

{ Time complexity $O(b^d)$
Space complexity $O(bd)$

i) Failure Value! There is no solution

ii) Cutoff failure: terminates on reaching predetermined depth.

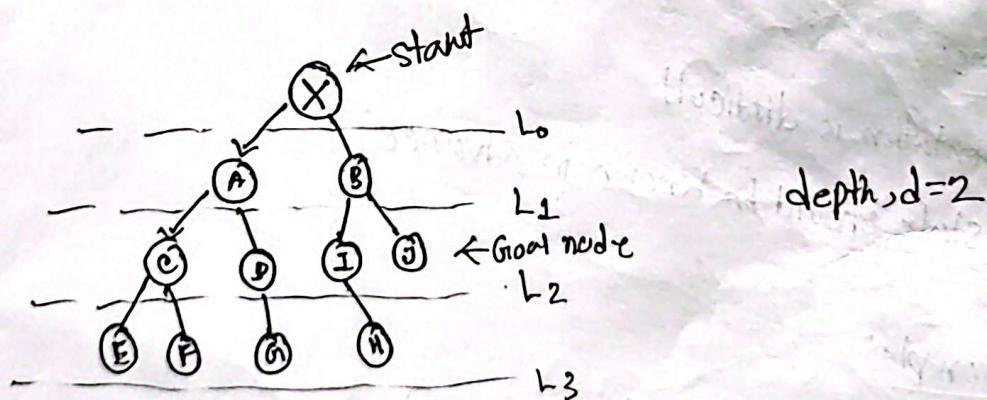
Advantages:

→ memory efficient

Disadvantages:

→ can be terminated without finding solution (incompleteness)

→ not optimal.



Path $\rightarrow X \rightarrow A \rightarrow C \rightarrow D \rightarrow B \rightarrow I \rightarrow J$

Uniform Cost Search (UCS)

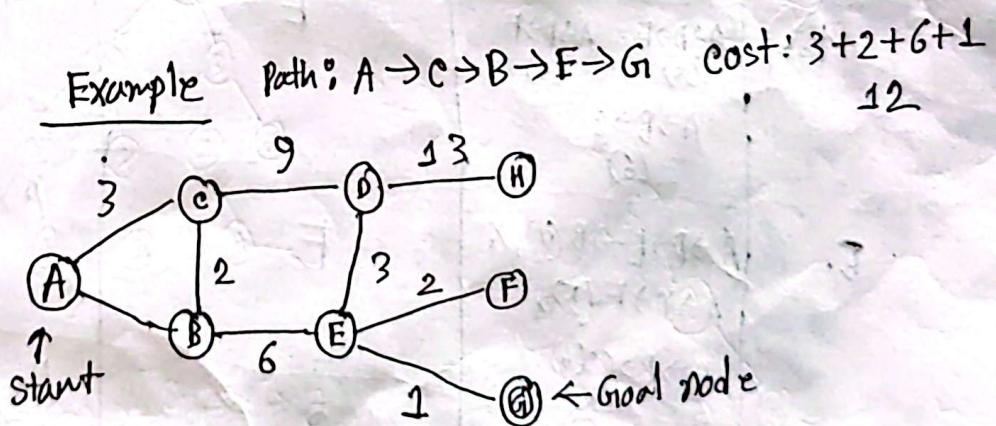
- It is used for weighted Tree/Graph traversal
- Goal is to path finding to goal node with lowest cumulative cost
- Node expansion is based on path costs.
- Priority Queue is used for implementation. (Higher Priority to minimum cost)
- Supports backtracking, complete.

Advantages

- optimal Solution

Disadvantages

- Stuck in infinite loop.

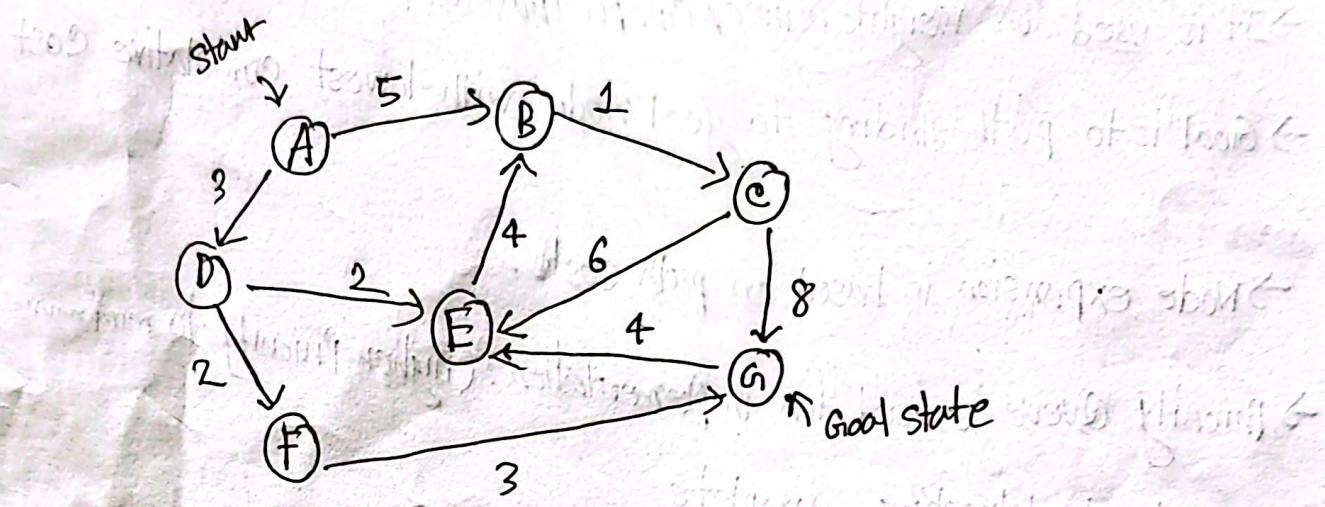


For Goal node 'H'

Path: A → C → B → F → G → E → F → D → H

$$\text{cost: } 3+2+6+1+1+2+2+3+13 = 33$$

Consider the following graph. Let the starting node be A and the Goal node be G



| | Frontier List | Expand List | Explained List |
|----|---|-------------|----------------------------------|
| 1. | A | A | - |
| 2. | $A \rightarrow D$, $A \rightarrow B$ (3) (5) | D | A |
| 3. | $A \rightarrow D \rightarrow E$, $A \rightarrow D \rightarrow F$, $A \rightarrow B$ (5) (5) (5) | B | A, D |
| 4. | $A \rightarrow D \rightarrow E$, $A \rightarrow D \rightarrow F$ (5) (5) $A \rightarrow B \rightarrow C$ (6) | E | A, D, B |
| 5. | $A \rightarrow D \rightarrow E \rightarrow B$ (8), X $A \rightarrow D \rightarrow F$ (5), $A \rightarrow B \rightarrow C$ (6), | F | A, D, B, E |
| 6. | $A \rightarrow D \rightarrow F \rightarrow G$ (8) $A \rightarrow B \rightarrow C$ (6) | C | A, D, B, E, F |
| 7. | $A \rightarrow D \rightarrow F \rightarrow G$ (8) $A \rightarrow B \rightarrow C \rightarrow E$ (12) X $A \rightarrow B \rightarrow C \rightarrow G$ (14) | G | A, D, B, E, F, C |
| 8 | $A \rightarrow D \rightarrow F \rightarrow G$ (8) (path) | NVLL | A, D, B, E, F, C, G traversed |

Iterative Deepening Depth-first search (IDFS/IDS)

- It is the combination of both DFS and BFS.
- Best Depth limit is found out by gradually increasing limit.

Initially $d=0$

→ every iteration increase by 1

→ Complete if the branching factor is finite.

[It's use stack]

Advantage:

→ Incorporates benefits of both DFS and BFS

Fast and less memory required

Disadvantage

→ Repeat the work / process.

{ Time complexity $O(b^d)$
Space complexity $O(bd)$

Working of IDFS / IDS

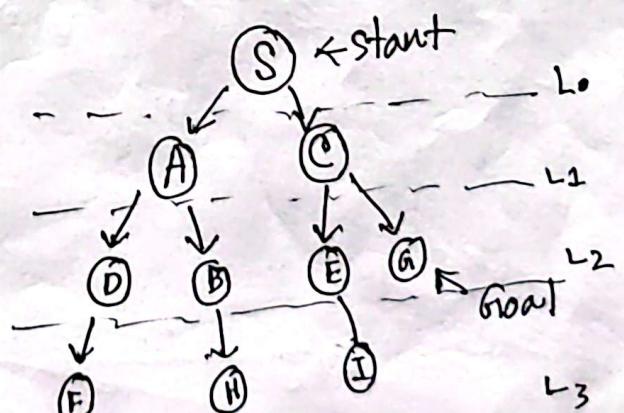
1st Iteration, $d=0$ [S]

2nd Iteration, $d=0+1=1$ [S → A → C]

3rd Iteration $d=1+1=2$

[S → A → D → B → C → F → G]

4th Iteration $d=2+1=3$



Heuristic Search [Informed Search]

Heuristic Search: Tries to optimize a problem using heuristic function.
→ tries to solve problem in minimum steps/cost

Heuristic function: It is a function $H(n)$ that gives an estimation on the cost of getting from node 'n' to the Goal state.
→ (estimated value)

Types of Heuristic

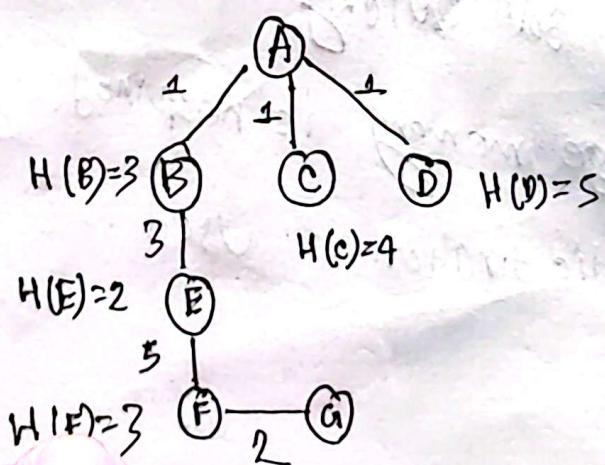
(1) Admissible: In this Heuristic function, never overestimates the cost of reaching the Goal.

$$H(n) \leq H'(n)$$

$H'(n)$ is always less than or equal to actual cost of lowest-cost path from node n to goal.

(2) Non-Admissible :- Overestimate

$$H(n) > H'(n)$$



Blind Search: It is also known as Unknown/Uninformed Search.

- There is no info about the searching
- No knowledge of where the goal.
- Eg:- Depth first, Breadth first Search
- Efficiency is low
- Slower than Heuristic
- Large memory is used.

Heuristic Search: It is a method of solving problem more easily and fast. They have knowledge of where goal or finish of the graph. (Informed Search)

Eg: Hill Climbing, A*, AO*

→ Highly efficient

$\begin{cases} \text{less time} \\ \text{less cost} \end{cases}$

→ finds soln quickly

→ no large memory is required.

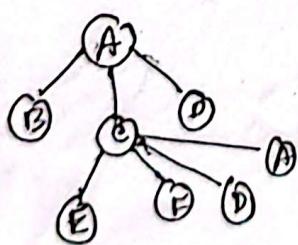
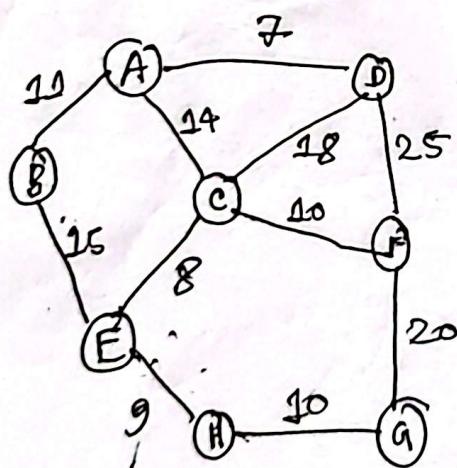
→ Heuristic function is used.

Breadth Search:

optimized version of Best first Search

- Heuristic Search Algorithm
- Explores a Graph by expanding the most promising node in a limited set
- Reduces Memory Requirement. (Greedy Algorithm)

Hence Only Predetermined no of Best Practical solutions are kept as Candidates.



$$\begin{aligned}
 A \rightarrow G_1 &= 40 \\
 B \rightarrow G_1 &= 32 \\
 C \rightarrow G_1 &= 25 \\
 D \rightarrow G_1 &= 35 \\
 E \rightarrow G_1 &= 19 \\
 F \rightarrow G_1 &= 17 \\
 H \rightarrow G_1 &= 10 \\
 G_1 \rightarrow G_1 &= 0
 \end{aligned}$$

Breadth Search
($B = 2$)

