the certifying compiler for properties like well typedness or safe memory access. As the certifying compiler is designed to be completely automatic, it will not be able to deal with rich functional or security properties.

We propose a veri cation framework with the following features:

compiler from source program annotation into bytecode annotation. Thus, bytecode

with

theorem prover

3 Bytecode Speci cation Language (BCSL)

In this section, we introduce a bytecode speci cation language which we call BCSL (short for ByteCode Speci cation Language). BCSL is based on the design principles of JML (Java Modeling Language), which is a behaviorial interface speci cation language following the design by contract approach [2].

Before going farther, we give a avour of what JML speci cations look like.

Fig. 2 shows an example of a Java class and its JML (based)4j 20.3897 2 Td (of)T5894.3882 0 Td (an)4j 15.23

Thus the \JML compiler" 2 compiles the JML source specil cation into user delined attributes. The compilation process has three stages:

1. compile the Java source le. This can be done by any Java compiler that supplies for every method in the generated class

```
JMLLoop_speci cation_attribute {
...
{ u2 index;
 u2 modi es_count;
 formula modi es[modi es_count];
 formula invariant;
 expression decreases;
}
```

substitution lemmas.

Returning back to the example, the expression c and st(c) stand respbac

resulting predicate is quanti ed over the expressions that may be modi ed by the called method. We also assume that if the invoked method terminates

related to the fact that the result type is boolean but the JVM encodes boolean expressions as integers (which is trivially true). This means that the proof obligations have also the same shape.

Another imp