

## **Programming languages and compiler**

### **Phase2: Parser Generator**

Mahmoud Mohamed	<b>19016578</b>
Gamal Abdel Hameed Nasef Nowesar	<b>19015550</b>
Moamen Mahmoud Gharib	<b>19016208</b>
Salah Eldin Eltenihy	19015854

## 1- A description of the used data structures

Data Structure	Usage
<code>vector&lt;pair&lt;string,string&gt;&gt;Grammar_rules;</code>	Keep track of Grammar Rules LHS = RHS
<code>map&lt;string,vector&lt;string&gt;&gt; rules_map;</code>	Grammar rules in other form A -> B   C   D A -> {B, C, D}
<code>map&lt;string,set&lt;string&gt;&gt; first_sets;</code>	Keep track of first sets LHS -> first(LHS)
<code>map&lt;string,map&lt;string,string&gt;&gt; table;</code>	Parse table LHS -> Map between T -> (Production Rule)
<code>map&lt;string,set&lt;string&gt;&gt; follow_sets;</code>	Keep track of follow sets LHS -> follow(LHS)

## 2- Explanation of all algorithms and techniques used

Algorithm	Steps
<b>Eliminate Left Recursion</b>	<pre>For LHS:   For prod in RHS:     Alpha = ""     Beta = ""     If LHS exists in first of prod:       Alpha += prod     Else       Beta += prod   If there is Left Recursion:     Modify LHS in Grammar_rules:       LHS -&gt; Beta LHS_DASH       LHS_DASH -&gt; Alpha LHS_DASH Epsilon</pre>
<b>Left Factor</b>	<pre>//Left factor RHS factors of same predecessor Curr_group = For LHS:   For prod in RHS:     map&lt;string, vector&lt;string&gt;&gt; groups     If Curr_group substr of prod:       Add prod to group     Else:       Wrap up old group [substr[A1, A2, ...]] and add it to groups   Curr_group = prod</pre>
<b>First Set</b>	<pre>For LHS:   For prod in RHS:     If prod is terminal:       Add to first set     Else:</pre>

getFollowSets	<p>Make a map to store the rightmost non-terminals of each LHS.</p> <p>Then apply topological sort.</p> <p>Then get the follow set for each non-terminal 'str':</p> <p>loop over all the productions and for each production, loop over the production's parts:</p> <p>If the 'str' is followed directly by a terminal, add this terminal to the follow set of 'str'.</p> <p>If the 'str' is followed directly by a non-terminal, then add the first set of this non-terminal to the follow set of 'str'. Then remove 'Epsilon' from the follow set if it exists.</p> <p>If 'str' is a rightmost nonterminal of the current LHS, then add the follow set of the LHS to the follow set of 'str'.</p>
createTable	<p>Iterate through each non-terminal in the grammar. Iterate through each production of the current non-terminal. Split the production into individual parts based on spaces. Check if the first symbol of the production is a terminal:</p> <p>If the entry in the parsing table is already occupied, print an error.</p> <p>Else, Set the entry in the parsing table for the current non-terminal and</p>

	<p>terminal symbol.</p> <p>Else If the first symbol of the production is a non-terminal:</p> <p>Iterate through the First set of the first non-terminal in the production:</p> <p>If the entry in the parsing table is already occupied, print an error.</p> <p>Else, Set the entry in the parsing table for the current non-terminal and terminal symbol.</p> <p>Iterate through the Follow set of the current non-terminal:</p> <p>If the entry in the parsing table is not present or is set to "Sync": Set the entry in the parsing table to "Epsilon" or "Sync" based on the presence of "Epsilon" in the First set.</p> <p>If "Epsilon" is in the First set but the entry in the parsing table is not "Epsilon," print an error.</p>
LL1_parse	<p>Initially, flag is true.  Loop until flag is false:  Get the top of the stack.</p> <p>If top is the end-of-input marker "\$":  If input is "\$", then the accepted,  else rejected.</p> <p>If top is a terminal symbol:</p> <p>If top is a single-character terminal</p>

	<p>(e.g., 'a'), then match the input with the terminal symbol, and pop it from the stack.</p> <p>Else, Error (the input is missing).</p> <p>Else, if the top of the stack is a non-terminal:</p> <p>If the entry in the parsing table for the current non-terminal and terminal symbol of input is empty, then error.</p> <p>Else, if the production is "Sync", then pop from the stack.</p> <p>If the production is "Epsilon", then pop from the stack.</p> <p>Else, Expand the stack with the production parts. Pop from the stack and push the production parts to the stack in reverse order.</p>
--	---

### 3- Screenshots:

```
1 |
2 | After eliminating Left Recursion:
3 | METHOD_BODY, STATEMENT_LIST
4 | STATEMENT_LIST, STATEMENT STATEMENT_LIST_dashLR
5 | STATEMENT, DECLARATION | IF | WHILE | ASSIGNMENT
6 | DECLARATION, PRIMITIVE_TYPE 'id' ';'
7 | PRIMITIVE_TYPE, 'int' | 'float'
8 | IF, 'if' '(' EXPRESSION ')' '{' STATEMENT '}' 'else' '{' STATEMENT '}'
9 | WHILE, 'while' '(' EXPRESSION ')' '{' STATEMENT '}'
10 | ASSIGNMENT, 'id' 'assign' EXPRESSION ';'
11 | EXPRESSION, SIMPLE_EXPRESSION | SIMPLE_EXPRESSION 'relop' SIMPLE_EXPRESSION
12 | SIMPLE_EXPRESSION, TERM SIMPLE_EXPRESSION_dashLR | SIGN TERM SIMPLE_EXPRESSION_dashLR
13 | TERM, FACTOR TERM_dashLR
14 | FACTOR, 'id' | 'num' | '(' EXPRESSION ')'
15 | SIGN, '+' | '-'
16 | STATEMENT_LIST_dashLR, 'int' 'id' ';' STATEMENT_LIST_dashLR | 'float' 'id' ';' STATEMENT_LIST_dashLR
17 | SIMPLE_EXPRESSION_dashLR, 'addop' TERM SIMPLE_EXPRESSION_dashLR | Epsilon
18 | TERM_dashLR, 'mulop' FACTOR TERM_dashLR | Epsilon
19 |
20 |
21 | After Left Factoring:
22 | METHOD_BODY, STATEMENT_LIST
23 | STATEMENT_LIST, STATEMENT STATEMENT_LIST_dashLR
24 | STATEMENT, ASSIGNMENT | DECLARATION | IF | WHILE
25 | DECLARATION, PRIMITIVE_TYPE 'id' ';'
26 | PRIMITIVE_TYPE, 'float' | 'int'
27 | IF, 'if' '(' EXPRESSION ')' '{' STATEMENT '}' 'else' '{' STATEMENT '}'
28 | WHILE, 'while' '(' EXPRESSION ')' '{' STATEMENT '}'
29 | ASSIGNMENT, 'id' 'assign' EXPRESSION ';'
30 | EXPRESSION, SIMPLE_EXPRESSION SIMPLE_EXPRESSION_dashLF
```