

We can then simplify the summations on either side to

$$\sum_{i=0}^{m-1} \left(\sum_{c=0}^{s_i-1} f(s_{i,c}) \right) < \sum_{i=0}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right) \quad (16)$$

which is the def of I_f so $I_f(s) < I_f(s')$ \square

*bottom

top

Note, this does not hold for "entire" positions.

It works for sequences RBRB etc from left to right * $R=0$ $B=1$ as binary number

$\boxed{R} = 0$ $S = \{R\} \mapsto$ doesn't follow rules of game.

$\boxed{B} = 1$ $S = \{B\}$ completed = (BRBR, BBRR, BR₀RB)

10 = $\boxed{B} \boxed{R}$ $S = \{B, R\}$ completed = (BBRR or BRRB)

$$\underbrace{f(s_j, s_j)}_{f(s_{j+1})} \leq \sum_{c=s_j}^{s'_j-1} f(s_{j,c}) \text{ less than } f(s'_{j+1}) \quad (3)$$

Remember $s_j = s'_j$ b/c agree until j
 so then $\forall c, s_{j,c} = s'_{j,c} \quad (9)$

$$\sum_{c=s_j}^{s'_j-1} f(s_{j,c}) = \sum_{c=s_j}^{s'_j-1} f(s'_{j,c}) = (10) \text{ just (9) by zero or more}$$

and then

$$\sum_{c=s_j}^{s'_j-1} f(s'_{j,c}) \leq \sum_{c=s_j}^{s'_j-1} f(s'_{j,c}) + \sum_{i=j+1}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right) = (11)$$

we use (7), (8), (10), (11) to put together (12):

$$\sum_{i=j+1}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right) < \sum_{c=s_j}^{s'_j-1} f(s'_{j,c}) + \sum_{i=j+1}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right)$$

again from (9)

$$\sum_{c=0}^{s'_j-1} f(s'_{j,c}) = \sum_{c=0}^{s_j-1} f(s_{j,c}) \quad (13)$$

so

$$\sum_{i=j}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right) < \sum_{c=0}^{s'_j-1} f(s'_{j,c}) + \sum_{i=j+1}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right)$$

w/o sec term

so

$$\sum_{i=j}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right) < \sum_{i=j+1}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right) \quad (14) \text{ b/c remove } j^{\text{th}} \text{ term}$$

Because $i < j \Rightarrow s_i = s'_i$ and $i \leq j \Rightarrow s_i = s'_i \Rightarrow s_{i,c} = s'_{i,c}$

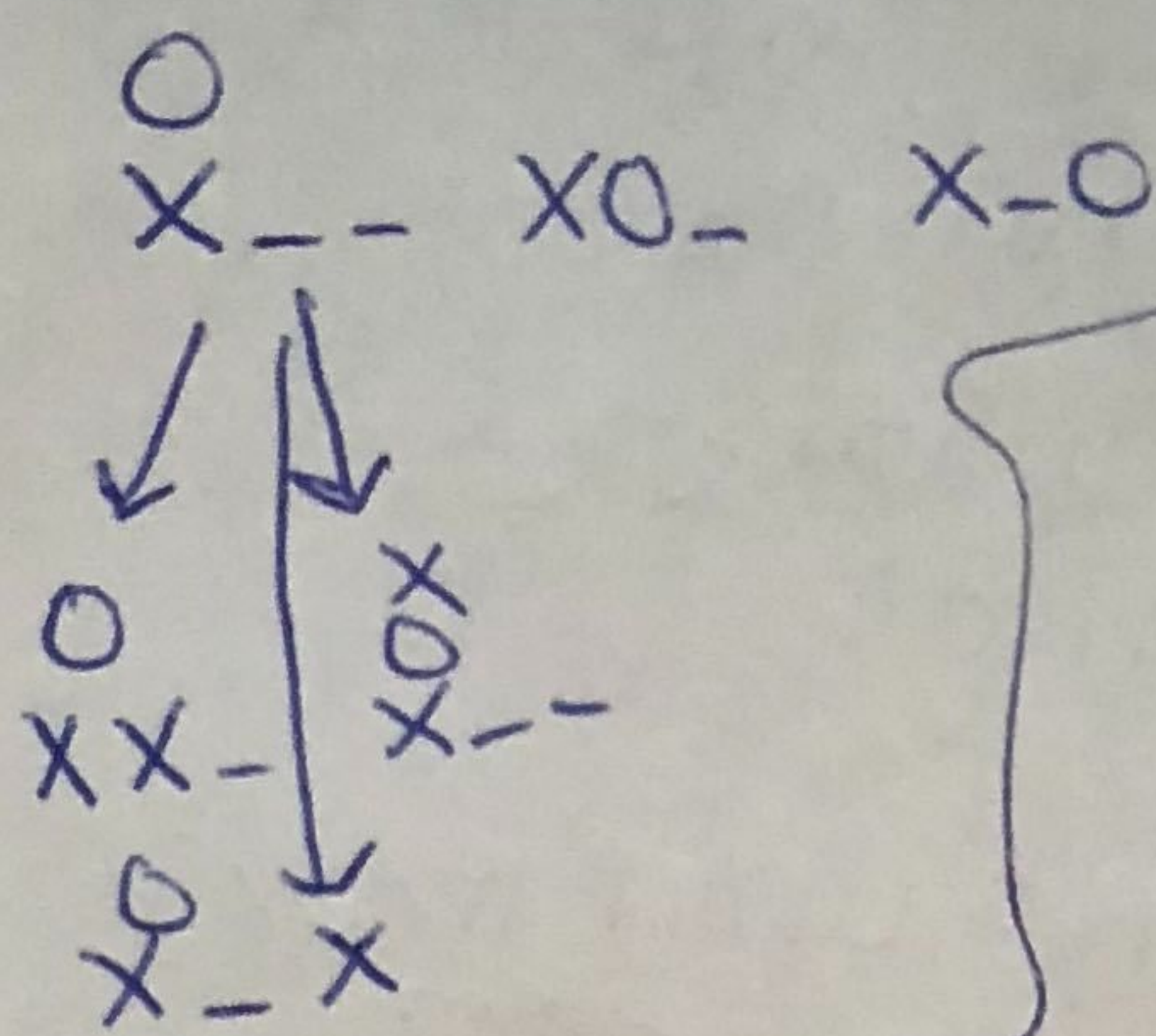
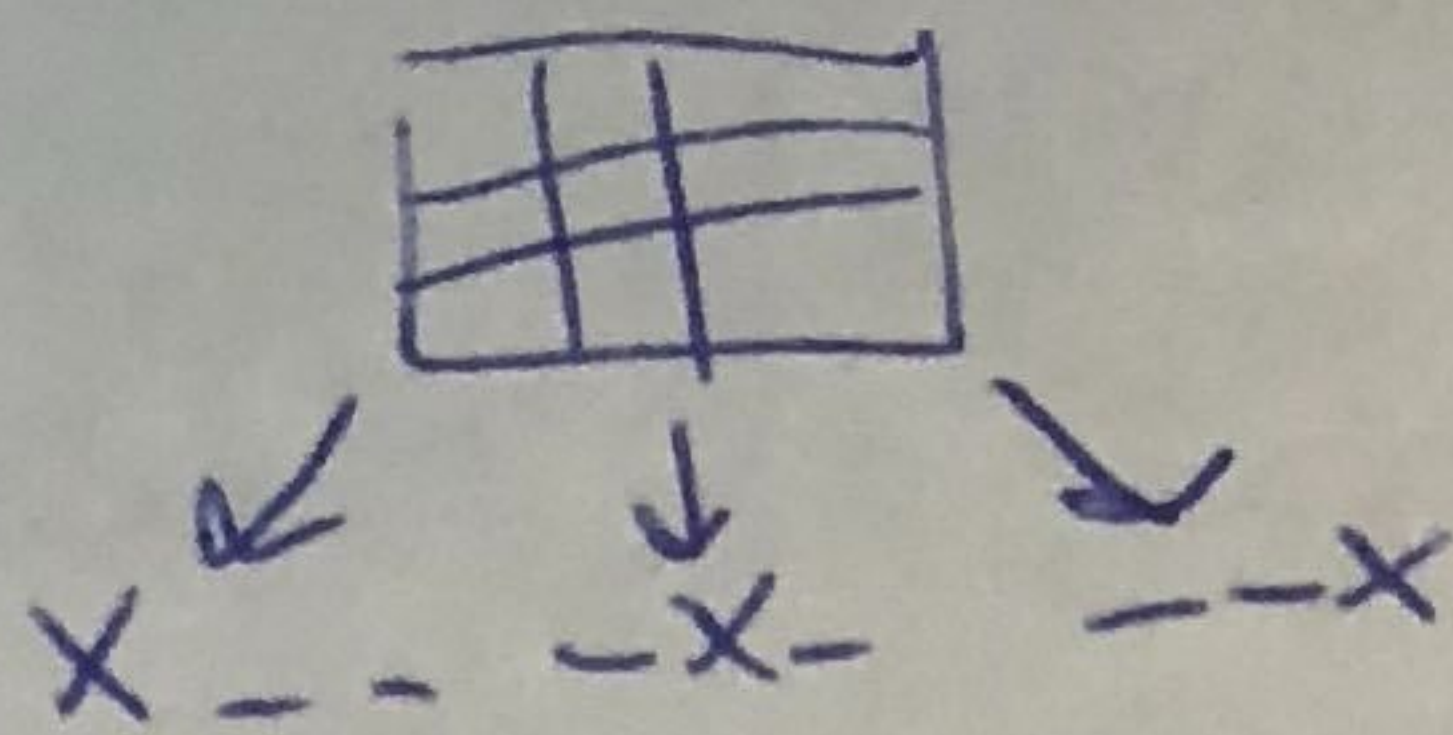
from (14) we can say

$$\sum_{i=0}^{j-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right) + \sum_{i=j}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right) < \sum_{i=0}^{j-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right) + \sum_{i=j}^{m-1} \left(\sum_{c=0}^{s'_i-1} f(s'_{i,c}) \right)$$

b/c sec term RHS > LHS

→

②



$$\sum_{c=0}^{s_t-1} f(s_{t,c}) + f(s_{t+1}) \leq \sum_{c=0}^{s_{t+1}} f(s_{t,c}) + \sum_{c=s_t}^{d-1} f(s_{t,c}) \quad (6)$$

Substitute

$$(6) = \sum_{c=0}^{d-1} f(s_{t,c}) = f(s_t)$$

We substitute into the assumption from the step

$$\sum_{c=0}^{s_t-1} f(s_{t,c}) + \sum_{i=t+1}^{m-1} \left(\sum_{c=0}^{s_i-1} f(s_{i,c}) < f(s_t) \text{ by (2) and (6)} \right)$$

So we condense LHS to $\sum_{i=t}^{m-1} \left(\sum_{c=0}^{s_i-1} f(s_{i,c}) \right) < f(s_t)$
 which was our first goal to prove \square

We can now make the statement that

$$(7) = \sum_{i=j+1}^{m-1} \left(\sum_{c=0}^{s_i-1} f(s_{i,c}) \right) < f(s_{j+1}) = f(s_j, s_j) \text{ by (1) and (3)}$$

because we know that $s_j < s'_j$

$$f(s_j, s_j) \leq f(s_j, s_j) + \sum_{c=s_j+1}^{s'_j-1} f(s_{j,c}) = \sum_{c=s_j}^{s'_j-1} f(s_{j,c}) = (8)$$

~~scribbled out text~~

$$\sum_{i=j+1}^{m-1} \left(\sum_{c=0}^{s_i-1} f(s_{i,c}) \right) < f(s_j, s_j) + \sum_{c=s_j+1}^{s'_j-1} f(s_{j,c})$$

hash(s):
 \forall length i

$$s = [s_0 \dots s_{m-1}]$$

①

\forall num c less than or equal to s_i

total += ways to complete seq: ending in c

$I_f(s)$ is indexing function for some $f \equiv \text{hash } f$

$$I_f(s) = \sum_{i=0}^{m-1} \sum_{c=0}^{s_i-1} f(s_i, c) \mapsto \text{seq } s, \text{ terminating @ } s_{i-1} \text{ ending in } c$$

S and S' two seq which are valid and agree until index j

$s'_j \neq s_j$. We further constrain $s_j < s'_j$ s'_j 's j^{th} item is bigger

for all such t
 first, $0 \leq t \leq m$, show $\sum_{i=t}^{m-1} \left(\sum_{c=0}^{s_i-1} f(s_i, c) \right) < f(s_t)$

all the ways to finish sequences which are at least as long as t and end in a char $> s[t] < f(s_t)$

① let $t=m \mapsto 0 < 1$ left no terms $\because f(s_m) = 1$ by def.

② inductively, assume $\sum_{i=t+1}^{m-1} \left(\sum_{c=0}^{s_i-1} f(s_i, c) \right) < f(s_{t+1})$

note that $S_{t+1} = S_t, s_t \}$ seq up to t followed by t^{th} item

$$\text{so } (2) = f(s_t, s_t) \mapsto (4)$$

then (5) is:

$$\sum_{c=s_t}^d f(s_t, c) = f(s_t, s_t) + \sum_{c=s_t+1}^d f(s_t, c) = f(s_{t+1}) + \sum_{c=s_t+1}^d f(s_t, c)$$

$$\underline{f(s_{t+1})} \leq f(s_{t+1}) + \sum_{c=s_t+1}^d f(s_t, c) = \sum_{c=s_t}^d f(s_t, c) \text{ by (5) and (4)}$$