"You don't have to be a genius to code, you just have to be persistent."





# Graphs 1



## Agenda :

~~~~~~

Searching

in Gaph

Introduction to Graph

2. Types of Graphs

(3. How to store data in Graph

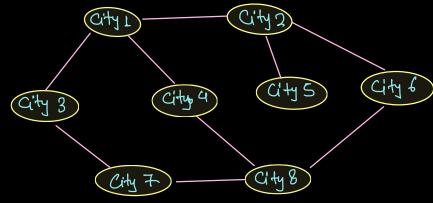
traversal { 4. BFS (Breadth First Search)

5. Is Path Available from Source to Destination

## Introduction to Graph

.

Want to store Enformation of cities and their connectivity.



With the help of graph we com
Store this kind of information.

cities = verter links = Edge

No. of Edges = 9

No. of vertex = 8

Definition: Collection of vortex and edgs is known as Graph.

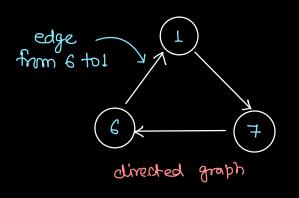
Neighbour of  $L \rightarrow \text{City 2}$ , city 3 and City 4

neighbour of & -> city 4, city 6, and city 7 (nbx)

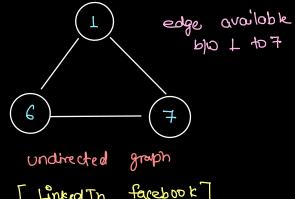
Direct connected vertex are known as Neighbours.

### Types of Graphs

1. Based on type of edges:

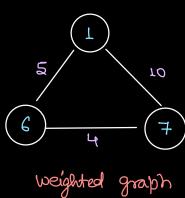


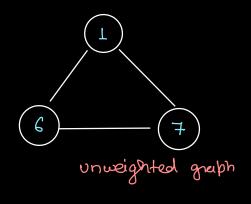
[ instagram, youtube scriber]



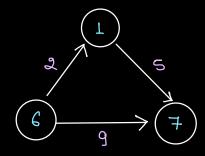
[Linked In, facebook] connection

2. Based on Edge wt. present or not:





also possible: types  $\mathcal{O}_{\mathcal{M}}$ 2. Combination above of



directed weighted graph

#### ~~~~~~~~~~~~~~~~~~~~~~

### How to store data in Graph

There is two famous implementation of graph is available:

- 1 Adjacency matrix
- a Adjacency List

### L. Adjacency Matrix:

int[][] graph = new int[v+x][vtx];

|   | 0          | T | 2 | 3 | 4          | 5          | 6             |
|---|------------|---|---|---|------------|------------|---------------|
| Q | O          | 1 | D | 1 | 0          | 0          | O             |
| Τ | L          | 0 | 1 | O | $\bigcirc$ | O          | $\Diamond$    |
| 2 |            |   | 0 |   |            |            |               |
| 3 | 1          | O | 1 | 0 | L          | 0          | $\bigcirc$    |
| 4 | $\bigcirc$ | O | O | L | D          | T          | 上             |
| S | $\bigcirc$ | 0 | 0 | O | 1          | $\Diamond$ | 上             |
| 6 | 0          | 0 | 0 | 0 | 7          | 1          | $\mathcal{D}$ |

| vertex: 7, | Edges = 8 |
|------------|-----------|
|------------|-----------|

| ,                  |   |   |                 |
|--------------------|---|---|-----------------|
|                    | Ð | T | 1               |
| Ŷ <del>-&gt;</del> | 2 | 3 | <u></u>         |
|                    | 3 | 4 | /               |
|                    | T | 2 | レ               |
| nt u= edge[i][o];  | Ч | 5 | <b>/</b>        |
|                    | Ч | 6 | <u></u>         |
| nt to edge (i)[i]; | 5 | 6 | ~               |
| / Edge blu U&V     |   |   | <b>*</b> *>==== |

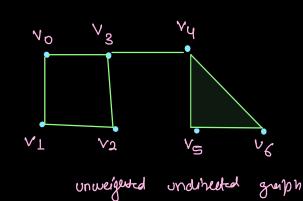
Edge-

undirected graph

// Edge blo u f
graph[u][v] =1
graph [v][u] =1

Jup[1][1]=1; - Jupp[3][3]=1

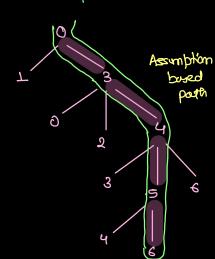
|   | 0 | T | 2 | 3             | Ч          | 5          | 6             |
|---|---|---|---|---------------|------------|------------|---------------|
| O | O | L | O | 1             | 0          | 0          | 0             |
| 1 | T | 0 | 1 | $\mathcal{O}$ | $\Diamond$ | 0          | Q             |
| 2 | 0 | L | 0 | L             | 0          | 0          | 0             |
| 3 | 1 | D | 1 | 0             | T          | 0          | Ø             |
| 4 | C | 0 | 0 | L             | O          | 1          | 上             |
| S | Q | 0 | 0 | 0             | 1          | $\bigcirc$ | 上             |
| 6 | 0 | 0 | 0 | O             | 1          | 1          | $\mathcal{Q}$ |



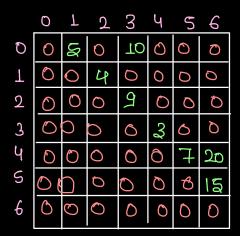
# undirected graph.

|   | 0 | 7 | 2 | 3 | Ч | 5 | 6 |
|---|---|---|---|---|---|---|---|
| O | O | L | O | 1 | ٥ | ٥ | ٥ |
| 1 | L | 0 | L | ٥ | 0 | D | D |
| 2 | 0 | 1 | 0 | 1 | 0 | ŋ | 0 |
| 3 | 1 | Ö | T | Q | L | ٥ | D |
| 4 | 0 | 0 | 0 | T | 0 | T | T |
| S | 0 | O | 0 | ٥ | 1 | 0 | L |
| 6 | O | Ð | ೦ | 0 | 1 | 1 | O |

| Source = 0 . | deet'n ad'on =6 |
|--------------|-----------------|
|--------------|-----------------|

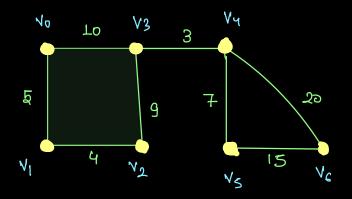


### climated kleighted Grapho



int u: edge[i][o];
int v = edge[i][i];
int wt = edge[i][g];
// Edge from u to v
with welget with
graph[u][v-]= wt;

|   | 0 | T          | 2 | 3  | Ч      | 5 | 6  |
|---|---|------------|---|----|--------|---|----|
| O | 0 | S          | 0 | TO | 0      | 0 | 0  |
| Τ | 0 | 0          | 4 | 0  | 0      | 0 | 0  |
| 2 | 0 | 0          | 0 | 9  | 0      | 0 | Q  |
| 3 | 0 | 0          | 0 | 0  | ς<br>γ | 0 | Q  |
| Ч |   | $\bigcirc$ |   | 0  | Ô      | 7 | 20 |
| S | 0 |            | 0 | 0  | O      | В | เล |
| 6 | 0 | Q          | 0 | 0  | 0      | 0 | 0  |



Major disadvantge of Adjacency matrix:

major disadvantge is wastaye of Space, that is why.

most of the tim we will dead with adjacency

list problem