

Question 2

Sparce matrix and vector products

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Objective : Accept matrix indices and values from an input file ,store them using linked lists and pointers and use it to compute both Ax and $x^T A$.

Storing A

Two linked lists have been used here : sa1 and sa2.

Note that only non-zero elements have been stored even if input file consists of zero elements.

```

/*****START*****/
typedef struct S1    //for storing the non-zero diagonal
elements
{
    float value;
    int i;
    struct S1 *p;
}SA1;

typedef struct S2    //for storing the non-zero non-diagonal
elements
{
    float value;
    int i,j;
    struct S2 *p;
}SA2;

struct S1 *sa1=NULL;    //linked list of non-zero
diagonal elements in order
struct S2 *sa2=NULL;    //linked list of non-zero
non-diagonal elements in order, from starting row and
```

starting column and traverse row-wise. ie, for a particular i, j from 0 to N-1

******END******

sa1 holds the non-zero diagonal elements. Each time a non-zero diagonal element is read from the input file, it is inserted into the linked list such that the linked list is sorted in ascending order with respect to row number i. This is achieved by means of the function :

```
void createlist1(int r,float val)    //create the linked
    list for non-zero diagonal elements
{
    .....
}
```

The new element details are stored in temp and sa1 points to the first element of the list. The pointer top is used for navigating through the list and inserting the new element in the correct position in the list.

```
*****START*****
if(sa1==NULL)        //if list is empty
{
    sa1=temp;    //temp is the first element in the
                list and hence sa1 points to temp
}
else if(sa1 != NULL )    //if list is not empty
{
    top=sa1;        //to navigate through the list

    if((temp->i)<(top->i))    //if the new item if
        to be added in front of sa1(first element in
        the linked list with lowest row index up till
        now)
    {
        temp->p=top;    //temp points to top
    }
}
```

```

        sa1 = temp;          //sa1 points to temp as the
                             element with the lowest row index
    }

    else if((temp->i)>(top->i))    //if the new
        item's index is greater than the row-index of
        top
    {
        while(flag==1)    //till new element is
            inserted
        {
            if(top->p==NULL)    //if temp is to be
                appended to the end of the list
            {
                top->p=temp;    //last element of
                                sa1 points to temp
                flag=0;        //denotes that new
                                element has been inserted
            }
            else if((temp->i)>(top->i) &&
                (temp->i)<(top->p->i))    //if the
                element is to be inserted in between
                the list
            {
                temp->p=top->p;    //temp is linked
                                to the element to which top is
                                linked to
                top->p=temp;        //temp is linked
                                to top
                flag=0;            //denotes that new
                                element has been inserted
            }
            top=top->p;    //incrementing top to
                next element in the linked list
        }
    }
}
}
}

```

/******END******/

sa2 holds the non-zero non-diagonal elements. Each time a non-zero non-diagonal element is read from the input file, it is inserted into the linked list such that the linked list is sorted in ascending order with respect the quantity : $i*N+j$, where i is row number , j is column number of the element and N

is the dimension of matrix A. This is achieved by means of the function :

```
void createlist2(int r, int c, float val)    //create the
    linked list for non-zero non-diagonal elements
{
    .....
}
```

The new element details are stored in Temp and sa2 points to the first element of the list. The pointer Top is used for navigating through the list and inserting the new element in the correct position in the list.

```
*****START*****

if(sa2==NULL)        //if list is empty
{
    sa2=Temp;        //Temp is the first element in the
                    list and hence sa2 points to Temp
}
else if(sa2 != NULL)    //if list is not empty
{
    Top=sa2;          //to navigate through the list

    index1=(Temp->i)*N + (Temp->j);    //index
                                    value for Temp
    index2=(Top->i)*N + (Top->j);      //index
                                    value for Top

    if(index1<index2)        //if the new item is to
                            be added in front of sa2(first element in the
                            linked list with lowest index up till now)
    {
        Temp->p=Top;        //Temp points to Top
        sa2 = Temp;        //sa2 points to Temp as
                            the element with the lowest index
    }

    else if(index1>index2)    //if the new item's
                            index is greater than the index of Top
    {
        while(flag==1)        //till new element is
```

```

        inserted
    {
        if(Top->p==NULL)    //if Temp is to be
            appended to the end of the list
        {
            Top->p=Temp;    //last element of
                sa2 points to Temp
            flag=0;        //denotes that new
                element has been inserted
        }
        else if(index1>index2 &&
            index1<((Top->p->i)*N + (Top->p->j)))
            //if the element is to be inserted in
            between the list
        {
            Temp->p=Top->p;    //Temp is linked
                to the element to which Top is
                linked to
            Top->p=Temp;        //Temp is linked
                to Top
            flag=0;            //denotes that
                new element has been inserted
        }
        Top=Top->p;            //incrementing
            top to next element in the linked list
        index2=(Top->i)*N + (Top->j);
            //re-initialising the value of index2
    }
}

/*****END*****/

```

Compute Ax and $x^T A$

It has been carried out in the function :

```

int Ax_xTA(int x[])    //for calculating the required
    matrix products
{
    .....
}

```

The algorithm makes use of the fact that :

1. $b = Ax$:

For diagonal element : $b[i] = A[i][i] * x[i]$

For non-diagonal element : $b[i] = A[i][j] * x[j]$

top and Top are pointers used to traverse through the lists sa1 and sa2 respectively.

```

/*****START*****/

while(top!=NULL)    //condition for reaching end of
    the linked list
{
    b[top->i]+=(top->value)*x[top->i]; //diagonal
        elements multiplied by the corresponding
        element in vector x and stored in the
        corresponding location in b
    top=top->p;        //incrementing top to
        next element in the linked list
}

while(Top!=NULL)    //condition for reaching end of
    the linked list
{
    b[Top->i]+=(Top->value)*x[Top->j]; // the
        element pointed to by Top is multiplied by
        the corresponding element in x and added to
        corresponding element in b
    Top=Top->p; //incrementing top to next
        element in the linked list
}

/*****END*****/

```

2. $c = x^T A$:

$$(A^T x)^T = x^T A$$

Since $c[]$ is a 1-D array, the transpose doesn't matter.

For diagonal element : $c[i] = A[i][i] * x[i]$

For non-diagonal element : $c[j]=A[i][j]*x[i]$

top and Top are pointers used to traverse through the lists sa1 and sa2 respectively.

```

/*****START*****/

while(top!=NULL)    //condition for reaching end of
    the linked list
{
    c[top->i]=b[top->i];          //diagonal
    elements multiplied by the corresponding
    element in vector x and stored in the
    corresponding location in c
    top=top->p;                  //incrementing top to
    next element in the linked list
}

while(Top!=NULL)    //condition for reaching end of
    the linked list
{
    c[Top->j]+=(Top->value)*x[Top->i]; // the
    element pointed to by Top is multiplied by
    the corresponding element in x and added to
    corresponding element in c
    Top=Top->p; //incrementing top to next element
    in the linked list
}

/*****END*****/
```

The above two codes are edited versions of the ones used in the original program. They have been edited to increase readability and understanding of the code.

Note :

The above stated multiplications could have been achieved without two different linked lists sa1 and sa2 and without any sorting. That is, we could have achieved the above vectors $b[]$ and $c[]$ by inserting the read elements directly into a linked list of structure type SA2 and using the algorithm for

non-diagonal element multiplication. That would have made the code even more simpler. But the advantage of having 2 different linked lists comes into play when we want to calculate trace of the matrix A or if we want to access just the diagonal elements say, to find the determinant in case of a triangular matrix A. It also comes in handy while determining diagonalisability of the matrix. The sorting comes in handy as having an order is always a good thing. We will know what index to expect for next. It makes row-wise access easier too (Suppose while calculating the average no. of elements in a row).

Comparison of Algorithms

This algorithm would be more efficient than the one given in Numerical Recipes in C.

- The book algorithm converts a matrix A into sparse matrix and hence a lot of storage space is wasted while storing matrix A and then converting it to a sparse matrix. The above given algorithm directly accepts disordered matrix elements and puts them into proper place, just that only non-zero elements are stored, and that too as in the form of a matrix expanded into a 1-D array.
- The algorithm given in the book stores all the diagonal elements, regardless of whether they are non-zero or not, whereas the above specified algorithm stores only non-zero relevant elements.
- The row-indexed sparse storage mode given in the book requires storage of about two times the number of matrix elements, ie, number of $2 \times$ non-diagonal non-zero elements + $2 \times N$ (dimension of A) + 2, whereas the above specified method requires $3 \times$ number of non-zero non-diagonal elements + $2 \times$ number of non-zero diagonal elements + number of non-zero elements (links or pointers). Depending on the type of matrix, one of them would be more storage efficient.
- In Numerical Recipes, they don't store the indices for the diagonal elements nor the column indices for more than 2 elements in the same row, though all column indices of non-diagonal non-zero element are stored. In the above approach, both row and column indices of non-zero non-diagonal elements and row-indices of non-zero diagonal elements have to be stored.
- Access of any diagonal element in Numerical Recipes is $O(1)$ whereas in the above approach, we have to traverse the list in order to locate the diagonal element.

- The time complexity for storing the sparse matrix
 In Numerical Recipes : $O(n^2)$ where n = dimension of A
 In the above algorithm : $O(n)$ almost along with the overhead of function calls and if -else conditions.
- Time complexity for calculating matrix product
 In Numerical Recipes : $O(nm)$ where n = matrix dimension and m = no. of non-zero non-diagonal elements
 In the above algorithm : $O(p)$ where p = no. of non-zero elements in matrix A