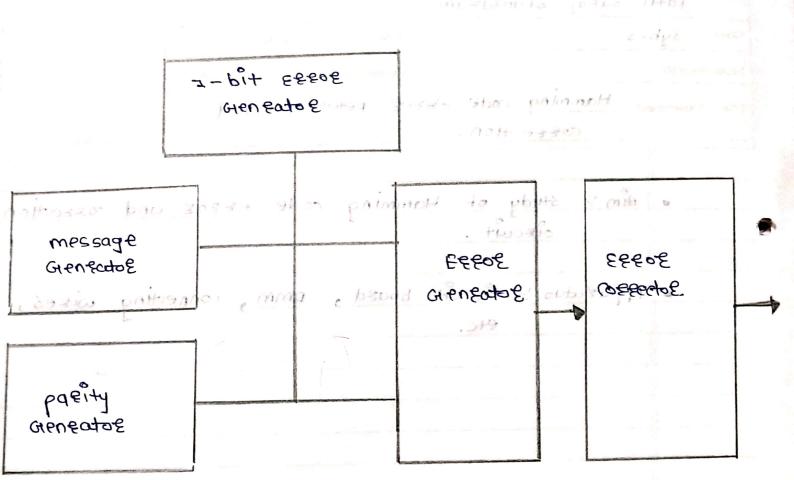
Dr. D. Y. Patil Unitech Soceity Dr. D. Y. Patil Science & Computer Science College, Akurdi, Pune - 411 044

DEPARTMENT OF ELECTRONICS

Name: Pa+1	sufaj shantafam.	
Class: ゴタb	Roll No.: Batch:	
Experiment No. :	Performed Date: 120	
Title of Experime	nt: Hamming code serbe petection and	
	cossection.	REMARKS
•	Am :- study of Hamming rode effor	
Transport of the Control of the Cont	ciecuit.	11023409
	1044 j	John ver
	Apparadus !- cherunt board, DMM, connec	Tho Wifes
	et c.	
		· · · · · · · · · · · · · · · · · · ·
A TO		zerto zerro
	No.	
		· ·
		200
		<u>A</u>

ciemit plageam !-



The second of th	
*	peocedure :-
7.	rake any 4 bit binney number information bits.
2,	And out the Hamming code too that I am
	either even parioty or odd parioty.
3	
	Honet se
4.	construct the H.C. as input to the circuit by using
4-30	Switches provided on Hamming road genrage
ಶ.	set pathy even as odd in it id
1	Daming Code 10110110
100	apply the discent the ds input to the dealt by
	neind enitches bearided of feeds menting
8.	set the posioty eventos odd, and not seen
g	Set the alb of 68808 getting and 68808 we 8860408
	notes down the old introbsequation table.
(0	cone to thit deepe in 7-bit hamming code by switches
	The peepe denente in this is found on the olp of
100 Mg	neene datente in binary rode.
12.	THE old of 65505 geterge 12 diney as 116 of
	lande in binney ode.
13	see the old of cossection ckt. note these sending
	a called dable
14	roon take new 4 pit intosmotion signed top add
	parioty and eppear the step 3 to 10.
	Note:
	The subject assaugement can detect and rossect the
	one bit effor in 4 bit information in even road
	parioty Hamming code.

rode are transferred from one point to a digital systems error can another point due to component maifunction of noise can sesult in I changing to 0 ०५ ८६६० although the probability of even a single bit error ounsauce in giality shatew is nesh ewall a basith bit may be the means of single bit essos detection a parioty bit is attached information bits (word) to make the total mose than one bit position. Hamming rode not only provides for dection of a bit essos put aso identities which but is in essos so that it can be rossected. This rade uses a number bits (depends on the no or information at restain positions in the group. essos detection and cossection are provided bi+6.

	4	obse	evation	, table			the state of the s	11561100 B
		1+					0	
	do	-0			0		Consulation and the state of th	
		_ N	Q	0			0	
-	<u> </u>	7	0	1-16	0		0	0
	Jag .	n	-	11-12/1	1 4/1/1/		D	O
	pavasau	2		0	0	in in the control of	6-	CO
		-	0	19/2-1482	11-1	MINOWH .	0 4 6	
	was g	20	-	0	_		150 -1	0
	203 Pr	02		-	0	O	EPF.	0
	ट्ड्यूट्ट वसल्पायह ट्राफ्स्य (व्यट	10	he some consultant consultant	0	٥	_	0	
		+	011	-		-	_	
	ming bit	<u>v</u>		-330	0	14 010	0	0
	5	N	0	0	+	011	3- /	
	2 70	7	0	0	1110	1 -0	0	011
	7-61+ har	19		-1	- 0			0
	4-1-00	2	0		01	0 /	00	TO THE PROPERTY OF THE PROPERTY CONTINUES AND RECOVER AND ADDRESS OF THE PROPERTY OF THE PROPE
	TF 6	+1	-	-				t- 640
	به	9		2 71	000	0	0	0
	+ 8	N	0	0.6				
	_		0	_	0		0	
	9 6	M					0	 0
	+ 6	0		0	1 0	-		0
	set fam	1 -	00	-114	9 -	101	0 000	-
		1	2	, bb)	1	1 0 bbd	
	4 45	10000	Jec	ppo	ه	odd	96.0	0 d d
	set Paerty		Ü	ō	1000 E	Ö	E S	0
	-	70	71 -	m- s	7 - E	-191	19-17	
	+	ρ ₃	_	2 _ 5	0	0	0	0
	4-19-4	29	٥	0	0		9 _ /	((191 <u>2</u>)
100	tas					*		
	3,5	Td		_			0	0
	07.72		戸	2]	F)	7	(5)	~
		and the state of t						

					Company of the compan				
calcul	ation!					1 j.			
		e code	= 10	(1			74		
11	-	m+ p+ 1							9
* 44		4+1+1		lammina	RW	5 Nat	Satis	Fied	Billionery promiser and consistency account of the service of the
		4+2+1							
		4+3+1		1		1.7	in the same of the	1.7	
***************************************		8		ing Ru	1e 50+	isfied			-
	2, L=					-	-	(2)	
i, i		443	j*						* *
76,		70		-	J				
LDP		rength	20	Hammin	0 000	2.0	3	*	
	11(19	5		4		_			
	91			63			7 O	-	
even.		P2	- (0	0	1.18	24		•
099 -		٥			0				1
O dd -		Yem	*			**	10	•	5 T
***	PI =	0 1						22	Maria de la companya della companya
Seci		0 (0	1 200	S au :				
			0			**			
	p2 =	2 3		7 .		4	_ 5	6	7
	५५०		1 9	1	406		0	- (-1 -
A second	odd	0 1			edo		0		
						-			
<u> </u>	\ \	01		-	•		had	. 7	
	P 1	P2	w1	P3 -	m_2	· W3	m	4	
9	\	02	3 .	4	5	6		F	300
4 n6U	→ 1	0	1	0	- 1	0			-
eqq	→ 0	1	1		1	0		1	1"
		,	- 11		•			Peru	
7			processor.	Pines			20 3000		

	10 1 2 1600 (0)
	P1= 1 3 5 7 P2= 2 3 6 7
	ENEU = 1 1 1 1 606U = 0 1 0 1
	odd = 0 1 1 1 0 dd = 1 1 0 1
	P3 = 4 5 6 7
	even = 0 1 0 1 and the
-	odd = 1 1 0 1
10 10	trough 30 to 3600 painment bailing food in 1
	3 0 1 0 1 to 1 to 1 to 1 to 1 to 1 to 1 t
D. Trailey	> P1 P2 m1 20 P3 1 m20 m3 m40
	even = 0 1 0 0 1. 1. 0 0 1. 1. 0
10 0 14.	odd = 01 00 0210 9d1 000 pho 6 painted 15
	his prizes pates enibbe es notasses
to Atpo	11 183 =1 15147th 500 2161 171717 10 m 917 (1)
	even = 0 1 0 . Toitement i etch
-/ 9-	0 old = 1 10 19010 900 -111 prisag to on 5
	41m 21 1242 0 600
•	Result: 1 tet + 100
	© pada = 1 0 11 8 = 3
	Hamming code
	b1 b5 w1 b3 w5 w8 wA
	Eneu = 0 1 1 0 0 1 1
	odd= 1 0 0 1 0 1
	(1) = 1 dob (1)
	Hamming rode
-	P1 P2 m1 P3 m2 m3 m4
	even = 1 0 1 0 1
	odd = 1 1 1 1 0 1
1	

	(3) 00dd : 0 1 01
	P1 P2 m1 P3 m2 m3 m4
	euen = 0 1 0 0 1 0 1
	odd = 1 0 0 1 1 0 1
	condusion :-
	1311 10
	we have studied Hamming rade expos detection and
	rofermion rode.
2]	It is one of the most common essos detection and
	roferción rode.
75	Hamming code can be use the expos detection and
	resection by adding exten pasity bit.
4)	The no. of parity bits are dependent on length of
	dodg information.
5]	
5]	no of parity bits are given by formula are:
5]	no. of parity bits are given by formula are:
5]	no. of parity bits are given by formula are: $2^{n} > m + p + 1$ $2^{3} > 4 + 3 + 1$ $= 4 + 3$
5]	no. of parity bits are given by formula are: $2^{n} \times m + p + 1$ $2^{3} \times 4 + 3 + 1$ $3^{3} \times 4 + 3 + 1$ $3^{$
5]	no. of parity bits are given by formula are: $2^{n} > m + p + 1$ $2^{3} > 4 + 3 + 1$ $= 4 + 3$
5]	no. of parity bits are given by formula are: $2^{n} > m + p + 1 \qquad L = m + p$ $2^{3} > 4 + 3 + 1 \qquad = 4 + 3 \qquad \text{therefore}$ $\vdots 8 > 8 \qquad 1 \qquad 0 = 7 \qquad \text{therefore}$
5]	no.of parity bits are given by formula are: 2n > m + p + 1 2³ > 4 + 3 + 1 ∴ 8 > 8 10 = 7 1000 1000 1000 1000 1000 1000 1000 1
5]	10.0F parity bits are given by formula are: 2n x m + p + 1 23 x 4 + 3 + 1 - 4 + 3 - 8 x 8 - 9 box primary 1 m em em en
5]	no.of parity bits are given by formula are: 2n > m + p + 1 2³ > 4 + 3 + 1 ∴ 8 > 8 10 = 7 1000 1000 1000 1000 1000 1000 1000 1
5]	10.0F parity bits are given by formula are: 2n x m + p + 1 23 x 4 + 3 + 1 - 4 + 3 - 8 x 8 - 9 box primary 1 m em em en
5]	10.0F parity bits are given by formula are: 2n x m + p + 1 23 x 4 + 3 + 1 - 4 + 3 - 8 x 8 - 9 box primary 1 m em em en
5]	no. of parity bits are given by formula are: $2^{n} \times m + p + 1$ $2^{n} \times m + p + 1$ $3^{n} \times 4 + 3 + 1$ $3^{n} \times 4 + 3 + 1$ $3^{n} \times 4 + 3 + 1$ $3^{n} \times 4 \times $
	10.0F parity bits are given by formula are: 2n x m +p+1 1=m+p 23 x 4+3+1 38 x 8 10 = 7 1000 100
	no. of parity bits are given by formula are: $2^{n} > m + p + 1$ $1 = m + p$ 1
	no. of parity bits are given by formula are: $2^{n} \times m + p + 1$ $1 = m + p$ $1 \times 3^{n} \times 4 + 3 + 1$ $1 \times 8 \times 8$ $1 = 7$ $1 \times 3^{n} \times 3^{n} \times 3^{n}$ $1 \times 3^{n} \times 3^{n} \times 3^{n} \times 3^{n} \times 3^{n}$ $1 \times 3^{n} \times 3^{n} \times 3^{n} \times 3^{n} \times 3^{n}$ $1 \times 3^{n} \times 3^{n} \times 3^{n} \times 3^{n} \times 3^{n}$ $1 \times 3^{n} \times 3^$