

Задачи по дискретной математике для контрольных и самостоятельных работ

О.-д. функции. Теория кодирования. Графы

Учебный практикум

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Практикум предназначен для студентов, изучающих курс "Дискретная математика", а также для преподавателей, ведущих практические занятия по данному курсу.

1. Построить диаграмму Мура для о.д.-функции $\varphi(x(1)x(2)\dots x(t)\dots) = y(1)y(1)\dots y(t)\dots$, где

1.1.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t-1) \sim x(t), & t \ge 2 \end{cases}$$

1.2.
$$y(t) = \begin{cases} 1, & t = 1 \\ x(t-1) \oplus x(t), & t \ge 2 \end{cases}$$

1.3.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t-1)x(t), & t \ge 2 \end{cases}$$

1.4.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t) \to x(t-1), & t \ge 2 \end{cases}$$

1.5.
$$y(t) = \begin{cases} 1, & t = 1 \\ x(t) \mid x(t-1), & t \ge 2 \end{cases}$$

1.6.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t-1) \lor x(t), & t \ge 2 \end{cases}$$

1.7.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t-1) \to x(t), & t \ge 2 \end{cases}$$

1.8.
$$y(t) = \begin{cases} 1, & t = 1 \\ x(t-1) \downarrow x(t), & t \ge 2 \end{cases}$$

1.9.
$$y(t) = \begin{cases} 0, & t = 1 \\ \overline{x(t-1)} \to x(t), & t \ge 2 \end{cases}$$

1.10.
$$y(t) = \begin{cases} \frac{1, & t = 1}{x(t-1) \lor x(t)}, & t \ge 2 \end{cases}$$

1.11.
$$y(t) = \begin{cases} 1, & t = 1 \\ x(t-1) \oplus \overline{x(t)}, & t \ge 2 \end{cases}$$

1.12.
$$y(t) = \begin{cases} \frac{0, & t = 1}{x(t-1)} & x(t), & t \ge 2 \end{cases}$$

1.13.
$$y(t) = \begin{cases} 1, & t = 1 \\ x(t-1) \to \overline{x(t)}, & t \ge 2 \end{cases}$$

1.14.
$$y(t) = \begin{cases} \frac{1}{x(t)} & t = 1 \\ \frac{1}{x(t)} & t = 1 \end{cases}, \quad t \ge 2$$

1.15.
$$y(t) = \begin{cases} \frac{0, & t = 1}{x(t)x(t-1)}, & t \ge 2 \end{cases}$$

1.16.
$$y(t) = \begin{cases} \frac{1, & t = 1}{x(t-1) \to x(t)}, & t \ge 2 \end{cases}$$

1.17.
$$y(t) = \begin{cases} \frac{1}{x(t)} & t = 1 \\ \frac{1}{x(t-1)}, & t \ge 2 \end{cases}$$

1.18.
$$y(t) = \begin{cases} 0, & t = 1 \\ \frac{1}{x(t-1) \downarrow x(t)}, & t \ge 2 \end{cases}$$

1.19.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t-1)\overline{x(t)}, & t \ge 2 \end{cases}$$

1.20.
$$y(t) = \begin{cases} \frac{1, & t = 1 \\ x(t-1) \lor x(t), & t \ge 2 \end{cases}$$

1.21.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t) \oplus \overline{x(t-1)}, & t \ge 2 \end{cases}$$

1.22.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t) \mid \overline{x(t-1)}, & t \ge 2 \end{cases}$$

1.23.
$$y(t) = \begin{cases} 1, & t = 1 \\ x(t-1) \lor \overline{x(t)}, & t \ge 2 \end{cases}$$

1.24.
$$y(t) = \begin{cases} \frac{0, & t = 1}{x(t) \mid x(t-1), & t \ge 2 \end{cases}$$

1.25.
$$y(t) = \begin{cases} \frac{0}{x(t)} & t = 1 \\ \frac{1}{x(t)} & t = 1 \end{cases}$$

1.26.
$$y(t) = \begin{cases} 1, & t = 1 \\ x(t-1)\overline{x(t)}, & t \ge 2 \end{cases}$$

1.27.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t) \downarrow \overline{x(t-1)}, & t \ge 2 \end{cases}$$

1.28.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t-1) \sim \overline{x(t)}, & t \ge 2 \end{cases}$$

1.29.
$$y(t) = \begin{cases} \frac{0, & t = 1 \\ x(t-1) \oplus x(t), & t \ge 2 \end{cases}$$

1.30.
$$y(t) = \begin{cases} \frac{1, & t = 1 \\ x(t-1) \to x(t), & t \ge 2 \end{cases}$$

2. Построить систему канонических уравнений для о.д.-функции $\varphi(x(1)x(2)\dots x(t)\dots)=y(1)y(1)\dots y(t)\dots$, где

2.1.
$$y(t) = \begin{cases} 0, & t = 1 \\ y(t-1) \sim x(t), & t \ge 2 \end{cases}$$

2.2.
$$y(t) = \begin{cases} 1, & t = 1 \\ y(t-1) \oplus x(t), & t \ge 2 \end{cases}$$

2.3.
$$y(t) = \begin{cases} 1, & t = 1 \\ y(t-1)x(t), & t \ge 2 \end{cases}$$

2.4.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t) \to y(t-1), & t \ge 2 \end{cases}$$

2.5.
$$y(t) = \begin{cases} 1, & t = 1 \\ x(t) \mid y(t-1), & t \ge 2 \end{cases}$$

2.6.
$$y(t) = \begin{cases} 0, & t = 1 \\ y(t-1) \lor x(t), & t \ge 2 \end{cases}$$

2.7.
$$y(t) = \begin{cases} 0, & t = 1 \\ y(t-1) \to x(t), & t \ge 2 \end{cases}$$

2.8.
$$y(t) = \begin{cases} 1, & t = 1 \\ y(t-1) \downarrow x(t), & t \ge 2 \end{cases}$$

2.9.
$$y(t) = \begin{cases} 0, & t = 1 \\ y(t-1) \to x(t), & t \ge 2 \end{cases}$$

2.10.
$$y(t) = \begin{cases} \frac{1, & t = 1}{y(t-1) \lor x(t)}, & t \ge 2 \end{cases}$$

2.11.
$$y(t) = \begin{cases} 1, & t = 1 \\ y(t-1) \oplus \overline{x(t)}, & t \ge 2 \end{cases}$$

2.12.
$$y(t) = \begin{cases} \frac{0, & t = 1 \\ y(t-1) \sim x(t), & t \ge 2 \end{cases}$$

2.13.
$$y(t) = \begin{cases} 1, & t = 1 \\ y(t-1) \to \overline{x(t)}, & t \ge 2 \end{cases}$$

2.14.
$$y(t) = \begin{cases} \frac{1}{x(t)} & t = 1 \\ \frac{1}{y(t-1)}, & t \ge 2 \end{cases}$$

2.15.
$$y(t) = \begin{cases} 0, & t = 1 \\ \overline{x(t)y(t-1)}, & t \ge 2 \end{cases}$$

2.16.
$$y(t) = \begin{cases} \frac{1, & t = 1}{y(t-1) \to x(t)}, & t \ge 2 \end{cases}$$

2.17.
$$y(t) = \begin{cases} \frac{1}{x(t)} & t = 1 \\ \frac{1}{y(t-1)}, & t \ge 2 \end{cases}$$

2.18.
$$y(t) = \begin{cases} \frac{0, & t = 1}{y(t-1) \downarrow x(t)}, & t \ge 2 \end{cases}$$

2.19.
$$y(t) = \begin{cases} 1, & t = 1 \\ y(t-1)\overline{x(t)}, & t \ge 2 \end{cases}$$

2.20.
$$y(t) = \begin{cases} \frac{1, & t = 1 \\ y(t-1) \lor x(t), & t \ge 2 \end{cases}$$

2.21.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t) \oplus y(t-1), & t \ge 2 \end{cases}$$

2.22.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t) \mid y(t-1), & t \ge 2 \end{cases}$$

2.23.
$$y(t) = \begin{cases} 0, & t = 1 \\ y(t-1) \lor \overline{x(t)}, & t \ge 2 \end{cases}$$

2.24.
$$y(t) = \begin{cases} \frac{0, & t = 1}{x(t) \mid y(t-1)}, & t \ge 2 \end{cases}$$

2.25.
$$y(t) = \begin{cases} \frac{0}{x(t)} & t = 1 \\ \frac{1}{y(t-1)}, & t \ge 2 \end{cases}$$

2.26.
$$y(t) = \begin{cases} 1, & t = 1 \\ y(t-1)\overline{x(t)}, & t \ge 2 \end{cases}$$

2.27.
$$y(t) = \begin{cases} 0, & t = 1 \\ x(t) \downarrow \frac{1}{y(t-1)}, & t \ge 2 \end{cases}$$

2.28.
$$y(t) = \begin{cases} 0, & t = 1 \\ y(t-1) \sim \overline{x(t)}, & t \ge 2 \end{cases}$$

2.29.
$$y(t) = \begin{cases} 0, & t = 1 \\ y(t-1) \oplus x(t), & t \ge 2 \end{cases}$$

2.30.
$$y(t) = \begin{cases} 0, & t = 1\\ \frac{1}{y(t-1)} \to x(t), & t \ge 2 \end{cases}$$

3. Построить диаграмму Мура для о.д.-функции $\varphi(x(1)x(2)\dots x(t)\dots)=y(1)y(1)\dots y(t)\dots$, где

3.1.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ \hline x(t-1) \oplus x(t-2), & t \ge 3 \end{cases}$$

3.2.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ x(t-1) \lor \overline{x(t-2)}, & t \ge 3 \end{cases}$$

3.3.
$$y(t) = \begin{cases} 0, & t = 1\\ \frac{1, & t = 2}{x(t-1)} \sim x(t-2), & t \ge 3 \end{cases}$$

3.4.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ x(t-1) \to \overline{x(t-2)}, & t \ge 3 \end{cases}$$

3.5.
$$y(t) = \begin{cases} 1, & t = 1\\ 0, & t = 2\\ \overline{x(t-2)} \sim \overline{x(t-1)}, & t \ge 3 \end{cases}$$

3.6.
$$y(t) = \begin{cases} 0, & t = 1\\ \frac{1, & t = 2}{x(t-2)x(t-1)}, & t \ge 3 \end{cases}$$

3.7.
$$y(t) = \begin{cases} 1, & t = 1\\ 0, & t = 2\\ \hline x(t-1) \to x(t-2), & t \ge 3 \end{cases}$$

3.8.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ \hline x(t-2) \lor \overline{x(t-1)}, & t \ge 3 \end{cases}$$

3.9.
$$y(t) = \begin{cases} 0, & t = 1\\ \frac{1, & t = 2}{x(t-1) \mid x(t-2)}, & t \ge 3 \end{cases}$$

3.10.
$$y(t) = \begin{cases} 0, & t = 1\\ 1, & t = 2\\ x(t-1)\overline{x(t-2)}, & t \ge 3 \end{cases}$$

3.11.
$$y(t) = \begin{cases} 1, & t = 1\\ 0, & t = 2\\ \hline x(t-1) \to x(t-2), & t \ge 3 \end{cases}$$

3.12.
$$y(t) = \begin{cases} 0, & t = 1\\ 1, & t = 2\\ x(t-1) \mid \overline{x(t-2)}, & t \ge 3 \end{cases}$$

3.13.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ x(t-1) \mid x(t-2), & t \ge 3 \end{cases}$$

3.14.
$$y(t) = \begin{cases} 0, & t = 1\\ \frac{1}{x(t-1)} \to \overline{x(t-2)}, & t \ge 3 \end{cases}$$

3.15.
$$y(t) = \begin{cases} 0, & t = 1\\ \frac{1}{x(t-2)} \downarrow x(t-1), & t \ge 3 \end{cases}$$

3.16.
$$y(t) = \begin{cases} 0, & t = 1\\ \frac{1, & t = 2}{x(t-2)} \downarrow x(t-1), & t \ge 3 \end{cases}$$

3.17.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ \overline{x(t-1)} \overline{x(t-2)}, & t \ge 3 \end{cases}$$

3.18.
$$y(t) = \begin{cases} 0, & t = 1 \\ 1, & t = 2 \\ x(t-2) \downarrow \overline{x(t-1)}, & t \ge 3 \end{cases}$$

3.19.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ x(t-1) \oplus \overline{x(t-2)}, & t \ge 3 \end{cases}$$

3.20.
$$y(t) = \begin{cases} 0, & t = 1\\ 1, & t = 2\\ x(t-1) \sim x(t-2), & t \ge 3 \end{cases}$$

3.21.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ x(t-1) \lor x(t-2), & t \ge 3 \end{cases}$$

3.22.
$$y(t) = \begin{cases} 0, & t = 1\\ 1, & t = 2\\ x(t-1)x(t-2), & t \ge 3 \end{cases}$$

3.23.
$$y(t) = \begin{cases} 0, & t = 1\\ 1, & t = 2\\ x(t-2) \to x(t-1), & t \ge 3 \end{cases}$$

3.24.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ x(t-2) \downarrow x(t-1), & t \ge 3 \end{cases}$$

3.25.
$$y(t) = \begin{cases} 0, & t = 1 \\ 1, & t = 2 \\ x(t-1) \oplus x(t-2), & t \ge 3 \end{cases}$$

3.26.
$$y(t) = \begin{cases} 0, & t = 1\\ \frac{1}{x(t-1)} \downarrow \frac{x(t-2)}{x(t-2)}, & t \ge 3 \end{cases}$$

3.27.
$$y(t) = \begin{cases} 0, & t = 1\\ 1, & t = 2\\ x(t-2) \mid \overline{x(t-1)}, & t \ge 3 \end{cases}$$

3.28.
$$y(t) = \begin{cases} 0, & t = 1 \\ 1, & t = 2 \\ x(t-1) \sim \overline{x(t-2)}, & t \ge 3 \end{cases}$$

3.29.
$$y(t) = \begin{cases} 0, & t = 1\\ \frac{1, & t = 2}{x(t-1)} \lor x(t-2), & t \ge 3 \end{cases}$$

3.30.
$$y(t) = \begin{cases} 1, & t = 1 \\ 0, & t = 2 \\ \hline x(t-1) \to x(t-2), & t \ge 3 \end{cases}$$

4. Для о.д.-функции φ , заданной системой канонических уравнений, построить усеченное дерево.

$$4.1. \ \varphi : \begin{cases} y(t) = x_1(t) \sim x_2(t) \sim q(t-1) \\ q(t) = x_1(t) \oplus x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.2. \ \varphi : \begin{cases} y(t) = x_1(t) \vee x_2(t) \vee q(t-1) \\ q(t) = x_1(t) \oplus x_2(t) \\ q(0) = 1 \end{cases}$$

$$4.3. \ \varphi : \begin{cases} y(t) = x_1(t) \oplus x_2(t) \oplus q(t-1) \\ q(t) = x_1(t)x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.4. \ \varphi : \begin{cases} y(t) = (x_1(t) \rightarrow x_2(t)) \rightarrow q(t-1) \\ q(t) = x_1(t) \oplus x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.5. \ \varphi : \begin{cases} y(t) = x_1(t)x_2(t)q(t-1) \\ q(t) = x_1(t) \sim x_2(t) \\ q(0) = 1 \end{cases}$$

$$4.6. \ \varphi : \begin{cases} y(t) = x_1(t) \sim x_2(t) \sim q(t-1) \\ q(t) = x_1(t) \vee x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.7. \ \varphi : \begin{cases} y(t) = x_1(t) \rightarrow (x_2(t) \rightarrow q(t-1)) \\ q(t) = x_1(t) \sim x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.8. \ \varphi : \begin{cases} y(t) = x_1(t) \vee x_2(t) \vee q(t-1) \\ q(t) = x_1(t)x_2(t) \\ q(0) = 1 \end{cases}$$

$$4.9. \ \varphi : \begin{cases} y(t) = (x_1(t) \rightarrow x_2(t)) \rightarrow q(t-1) \\ q(t) = x_1(t)x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.10. \ \varphi : \begin{cases} y(t) = x_1(t) \vee x_2(t) \vee q(t-1) \\ q(t) = x_1(t)x_2(t) \rightarrow q(t-1) \\ q(t) = x_1(t) \rightarrow x_2(t) \\ q(0) = 1 \end{cases}$$

$$4.11. \ \varphi : \begin{cases} y(t) = x_1(t)x_2(t) \rightarrow q(t-1) \\ q(t) = x_1(t) \oplus x_2(t) \\ q(0) = 1 \end{cases}$$

$$4.12. \ \varphi : \begin{cases} y(t) = (x_1(t) \to x_2(t)) \oplus q(t-1) \\ q(t) = x_1(t) \sim x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.13. \ \varphi : \begin{cases} y(t) = (x_1(t) \oplus x_2(t)) \vee q(t-1) \\ q(t) = x_1(t) \to x_2(t) \\ q(0) = 1 \end{cases}$$

$$4.14. \ \varphi : \begin{cases} y(t) = x_1(t)(x_2(t) \sim q(t-1)) \\ q(t) = x_1(t) \vee x_2(t) \\ q(0) = 1 \end{cases}$$

$$4.15. \ \varphi : \begin{cases} y(t) = x_1(t) \to \overline{x_2(t)q(t-1)} \\ q(t) = x_1(t) \times x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.16. \ \varphi : \begin{cases} y(t) = \overline{(x_1(t) \to x_2(t))} \downarrow q(t-1) \\ q(t) = x_1(t) \to (x_2(t) \to x_1(t)) \\ q(0) = 1 \end{cases}$$

$$4.17. \ \varphi : \begin{cases} y(t) = x_1(t) \times x_2(t) \oplus q(t-1) \\ q(t) = x_1(t) \vee \overline{x_2(t)} \oplus q(t-1) \\ q(0) = 1 \end{cases}$$

$$4.18. \ \varphi : \begin{cases} y(t) = \overline{x_1(t)} \downarrow x_2(t) \oplus q(t-1) \\ q(t) = x_1(t) \oplus x_2(t) \oplus 1 \\ q(0) = 0 \end{cases}$$

$$4.19. \ \varphi : \begin{cases} y(t) = (x_1(t) \vee \overline{x_2(t)}) q(t-1) \\ q(t) = (x_1(t) \oplus x_2(t)) \vee x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.20. \ \varphi : \begin{cases} y(t) = \overline{x_1(t)} \overline{(x_2(t) \vee q(t-1))} \\ q(t) = x_1(t) \to x_2(t) \\ q(0) = 1 \end{cases}$$

$$4.21. \ \varphi : \begin{cases} y(t) = \overline{x_1(t)} \downarrow \overline{(x_2(t) \vee q(t-1))} \\ q(t) = x_1(t) \times x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.22. \ \varphi : \begin{cases} y(t) = \overline{x_1(t)} \downarrow \overline{(x_2(t) \vee q(t-1))} \\ q(t) = x_1(t) \oplus x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.23. \ \varphi : \begin{cases} y(t) = (x_1(t) \to x_2(t))(x_2(t) \to q(t-1)) \\ q(t) = x_1(t) \to x_2(t) \\ q(0) = 0 \end{cases}$$

$$4.23. \ \varphi : \begin{cases} y(t) = (x_1(t) \to x_2(t))(x_2(t) \to q(t-1)) \\ q(t) = x_1(t) \vee x_2(t) \\ q(0) = 0 \end{cases}$$

4.24.
$$\varphi$$
:
$$\begin{cases} y(t) = x_1(t)\overline{(x_2(t) \mid q(t-1))} \\ q(t) = x_1(t)x_2(t) \\ q(0) = 0 \end{cases}$$
4.25. φ :
$$\begin{cases} y(t) = (x_1(t) \oplus \overline{x_2(t)}) \mid q(t-1) \\ q(t) = x_1(t) \sim x_2(t) \\ q(0) = 0 \end{cases}$$
4.26. φ :
$$\begin{cases} y(t) = x_1(t) \vee x_2(t) \vee q(t-1) \\ q(t) = x_1(t)\overline{x_2(t)} \\ q(0) = 1 \end{cases}$$
4.27. φ :
$$\begin{cases} y(t) = (x_1(t) \vee x_2(t)) \downarrow q(t-1) \\ q(t) = x_1(t) \oplus x_2(t) \\ q(0) = 0 \end{cases}$$
4.28. φ :
$$\begin{cases} y(t) = (x_1(t)x_2(t)) \mid q(t-1) \\ q(t) = x_1(t) \sim x_2(t) \\ q(0) = 0 \end{cases}$$
4.29. φ :
$$\begin{cases} y(t) = x_1(t) \oplus (x_2(t) \rightarrow q(t-1)) \\ q(t) = x_1(t)(x_1(t) \vee x_2(t)) \\ q(0) = 0 \end{cases}$$
4.30. φ :
$$\begin{cases} y(t) = x_1(t) \oplus (x_2(t) \mid q(t-1)) \\ q(t) = x_1(t) \rightarrow x_2(t) \\ q(0) = 1 \end{cases}$$

5. Для о.д.-функции φ , заданной системой канонических уравнений, построить диаграмму Мура.

$$5.1. \ \varphi : \begin{cases} y(t) = (x(t) \to q_1(t-1)) \sim q_2(t-1) \\ q_1(t) = x(t) \oplus q_1(t-1) \\ q_2(t) = x(t) \to q_2(t-1) \\ q_1(0) = q_2(0) = 0 \end{cases}$$

$$5.2. \ \varphi : \begin{cases} y(t) = (x(t) \sim q_1(t-1)) \vee q_2(t-1) \\ q_1(t) = x(t) \to q_1(t-1) \\ q_2(t) = x(t)(x(t) \vee q_2(t-1)) \\ q_1(0) = q_2(0) = 1 \end{cases}$$

$$5.3. \ \varphi : \begin{cases} y(t) = x(t)(q_1(t-1) \oplus q_2(t-1)) \\ q_1(t) = x(t) \vee q_1(t-1) \\ q_2(t) = x(t) \oplus q_2(t-1) \\ q_1(0) = q_2(0) = 1 \end{cases}$$

$$5.4. \ \varphi : \begin{cases} y(t) = x(t) \to q_1(t-1)q_2(t-1) \\ q_1(t) = x(t) \to q_1(t-1)q_2(t-1) \\ q_1(t) = x(t) \to q_1(t-1) \to x(t) \end{cases}$$

$$1.5. \ \varphi : \begin{cases} y(t) = (x(t) \to q_1(t-1)) \mid q_2(t-1) \\ q_1(t) = x(t) \to q_1(t-1) \to x(t) \end{pmatrix} = (x(t) \to q_1(t-1) \to x(t)) \\ q_2(t) = x(t) \to q_1(t-1) \to x(t) \to x(t) \to x(t) \end{cases}$$

$$1.5. \ \varphi : \begin{cases} y(t) = x(t) \downarrow q_1(t-1) \to x(t) \\ q_1(t) = x(t) \lor q_1(t-1) \to x(t) \to x(t) \\ q_1(t) = x(t) \lor q_1(t-1) \to x(t) \to x(t) \to x(t) \\ q_1(t) = x(t) \to x(t) \to x(t) \to x(t) \to x(t) \to x(t) \end{cases}$$

$$1.5. \ \varphi : \begin{cases} y(t) = x(t) \downarrow q_1(t-1) \to x(t) \\ q_1(t) = x(t) \lor q_1(t-1) \\ q_2(t) = x(t) \downarrow q_2(t-1) \\ q_1(t) = x(t) \to x(t) \downarrow x(t) \to x(t) \to x(t) \to x(t) \\ x(t) = x(t) \downarrow x(t) \to x(t) \to x(t) \to x(t) \to x(t) \end{cases}$$

$$1.5. \ \varphi : \begin{cases} y(t) = (x(t) \lor x(t) \lor x(t) \to x(t) \to x(t) \to x(t) \to x(t) \\ y_1(t) = (x(t) \to x(t) \\ y_1(t) = (x(t) \to x(t) \to$$

$$5.9. \ \varphi : \begin{cases} y(t) = \overline{x(t)} \overline{(q_1(t-1) \vee q_2(t-1))} \\ q_1(t) = x(t) \rightarrow q_1(t-1) \\ q_2(t) = x(t) \sim q_2(t-1) \\ q_1(0) = q_2(0) = 1 \end{cases}$$

$$5.10. \ \varphi : \begin{cases} y(t) = (x(t) \sim \overline{q_1(t-1)})(q_1(t-1) \sim \overline{q_2(t-1)}) \\ q_1(t) = x(t)q_1(t-1) \\ q_2(t) = x(t) \sim q_2(t-1) \\ q_1(0) = q_2(0) = 0 \end{cases}$$

$$5.11. \ \varphi : \begin{cases} y(t) = \overline{x(t) \downarrow (q_1(t-1) \mid \overline{q_2(t-1)})} \\ q_1(t) = x(t) \sim q_1(t-1) \\ q_2(t) = x(t)q_2(t-1) \\ q_1(0) = q_2(0) = 0 \end{cases}$$

$$5.12. \ \varphi : \begin{cases} y(t) = (x(t) \rightarrow q_1(t-1))(q_1(t-1) \rightarrow q_2(t-1)) \\ q_1(t) = x(t) \vee q_1(t-1) \\ q_2(t) = x(t) \sim q_2(t-1) \\ q_1(0) = q_2(0) = 1 \end{cases}$$

$$5.13. \ \varphi : \begin{cases} y(t) = x(t)q_1(t-1)q_2(t-1) \\ q_1(t) = x(t) \oplus q_1(t-1) \\ q_2(t) = x(t) \rightarrow q_2(t-1) \\ q_1(0) = q_2(0) = 1 \end{cases}$$

$$5.14. \ \varphi : \begin{cases} y(t) = x(t) \oplus q_1(t-1) \oplus q_2(t-1) \\ q_1(t) = x(t) \vee q_1(t-1) \oplus q_2(t-1) \\ q_1(t) = x(t) \vee q_1(t-1) \rightarrow q_2(t-1) \\ q_1(t) = x(t) \oplus q_1(t-1) \rightarrow q_2(t-1) \end{cases}$$

$$5.15. \ \varphi : \begin{cases} y(t) = x(t) \oplus q_1(t-1) \rightarrow q_2(t-1) \\ q_1(t) = x(t) \oplus q_1(t-1) \oplus q_2(t-1) \\ q_1(t) = x(t) \oplus q_1(t-1) \oplus q_2(t-1) \\ q_1(t) = x(t) \oplus q_1(t-1) \rightarrow q_2(t-1) \\ q_1(t) = x(t) \oplus q_1(t-1) \oplus q_2(t-1) \\ q_1(t) = x(t) \oplus q_1(t-1) \rightarrow q_2(t-1) \\ q_1(t) = x(t) \oplus q_2(t-1) \\$$

$$5.18. \ \varphi : \begin{cases} y(t) = x(t) \lor q_1(t-1) \lor q_2(t-1) \\ q_1(t) = x(t) \to q_1(t-1) \\ q_2(t) = x(t) \lor q_2(t-1) \\ q_1(0) = q_2(0) = 1 \end{cases}$$

$$5.19. \ \varphi : \begin{cases} y(t) = x(t)q_1(t-1) \to q_2(t-1) \\ q_1(t) = x(t) \sim q_1(t-1) \\ q_2(t) = x(t) \lor q_2(t-1) \\ q_1(0) = q_2(0) = 1 \end{cases}$$

$$5.20. \ \varphi : \begin{cases} y(t) = x(t)\overline{(q_1(t-1) \mid q_2(t-1))} \\ q_1(t) = x(t) \to q_1(t-1) \\ q_2(t) = x(t)q_2(t-1) \\ q_1(0) = q_2(0) = 0 \end{cases}$$

$$5.21. \ \varphi : \begin{cases} y(t) = x(t) \oplus q_1(t-1) \oplus q_2(t-1) \\ q_1(t) = x(t) \sim q_1(t-1) \\ q_2(t) = x(t) \to q_2(t-1) \\ q_1(0) = q_2(0) = 0 \end{cases}$$

$$5.22. \ \varphi : \begin{cases} y(t) = x(t) \lor q_1(t-1) \lor q_2(t-1) \\ q_1(t) = x(t) \sim q_1(t-1) \\ q_2(t) = x(t) \sim q_2(t-1)) \lor q_2(t-1) \\ q_1(0) = q_2(0) = 1 \end{cases}$$

$$5.23. \ \varphi : \begin{cases} y(t) = x(t) \sim q_1(t-1) \sim q_2(t-1) \\ q_1(t) = x(t) \sim q_1(t-1) \sim q_2(t-1) \\ q_1(t) = x(t) \sim q_1(t-1) \\ q_2(t) = x(t) \sim q_2(t-1) \\ q_1(0) = q_2(0) = 0 \end{cases}$$

$$5.24. \ \varphi : \begin{cases} y(t) = (x(t) \to q_1(t-1)) \to q_2(t-1) \\ q_1(t) = x(t) \lor q_2(t-1) \\ q_1(t) = x(t) \lor q_2(t-1) \\ q_1(t) = x(t) \to q_1(t-1) \\ q_2(t) = x(t) \lor q_2(t-1) \\ q_1(t) = x(t) \to q_1(t-1) \\ q_1(t) = x(t) \to$$

5.27.
$$\varphi : \begin{cases} y(t) = (x(t) \lor q_1(t-1)) \downarrow q_2(t-1) \\ q_1(t) = x(t) \sim q_1(t-1) \\ q_2(t) = x(t) \lor q_2(t-1) \\ q_1(0) = q_2(0) = 0 \end{cases}$$
5.28.
$$\varphi : \begin{cases} y(t) = (x(t)q_1(t-1)) \mid q_2(t-1) \\ q_1(t) = x(t) \oplus q_1(t-1) \\ q_2(t) = x(t) \rightarrow q_2(t-1) \\ q_1(0) = q_2(0) = 0 \end{cases}$$
5.29.
$$\varphi : \begin{cases} y(t) = x(t) \sim (q_1(t-1) \rightarrow q_2(t-1)) \\ q_1(t) = x(t)(x(t) \lor q_1(t-1)) \\ q_2(t) = x(t) \oplus q_2(t-1) \\ q_1(0) = q_2(0) = 0 \end{cases}$$
5.30.
$$\varphi : \begin{cases} y(t) = x(t) \sim (q_1(t-1) \mid q_2(t-1)) \\ q_1(t) = x(t) \rightarrow q_1(t-1) \\ q_2(t) = q_2(t-1) \rightarrow x(t) \\ q_1(0) = q_2(0) = 1 \end{cases}$$

- 6. Проверить однозначную декодируемость кода C.
 - 6.1. $C = \{1001, 012, 0120, 112, 12, 10, 02\}$
 - 6.2. $C = \{1000, 022, 221, 12, 2012, 00, 20\}$
 - 6.3. $C = \{12, 2100, 0212, 2010, 01, 11, 010\}$
 - 6.4. $C = \{011, 01, 22, 122, 21, 0220, 1011\}$
 - 6.5. $C = \{200, 12, 202, 112, 210, 11, 20, 212\}$
 - 6.6. $C = \{02, 10, 121, 1022, 102, 11, 001, 0211\}$
 - 6.7. $C = \{02, 0222, 22, 12, 0100, 11, 221, 0022\}$
 - 6.8. $C = \{0210, 22, 010, 20, 01, 001, 0021\}$
 - 6.9. $C = \{2221,000,00,01,22,10,0012,011\}$
- 6.10. $C = \{00, 22, 221, 02, 220, 121, 222, 100\}$
- 6.11. $C = \{21, 10, 002, 2110, 1020, 0012, 200\}$
- 6.12. $C = \{20, 01, 1000, 0002, 12, 22, 11, 1112\}$
- 6.13. $C = \{01, 122, 10, 0122, 121, 22, 200, 012\}$
- 6.14. $C = \{112, 0210, 02, 21, 11, 221, 212, 1211\}$
- 6.15. $C = \{1000, 222, 001, 02, 2021, 22, 210\}$
- 6.16. $C = \{0220, 20, 1021, 12, 02, 00, 111, 1110\}$
- 6.17. $C = \{10, 00, 11, 0021, 2002, 002, 121\}$
- 6.18. $C = \{1122, 11, 1100, 122, 001, 2211\}$
- 6.19. $C = \{21, 10, 22, 0111, 11, 2201, 020, 2011\}$
- 6.20. $C = \{21, 2210, 02, 122, 0221, 012, 211\}$
- 6.21. $C = \{200, 22, 1122, 10, 2110, 1022, 0222\}$
- 6.22. $C = \{112, 210, 100, 11, 10, 001, 20, 1111\}$
- 6.23. $C = \{112, 022, 12, 221, 2220, 0012, 00\}$
- 6.24. $C = \{201, 0001, 01, 12, 10, 21, 220, 22\}$
- 6.25. $C = \{022, 22, 0000, 00, 0220, 210, 112\}$
- 6.26. $C = \{11, 20, 2110, 10, 00, 011, 202, 1222\}$
- 6.27. $C = \{2210, 21, 222, 0222, 0220, 221\}$
- 6.28. $C = \{011, 0000, 00, 0122, 21, 22, 022\}$
- 6.29. $C = \{00, 11, 0022, 21, 20, 002, 02, 001\}$
- 6.30. $C = \{1201, 20, 201, 0201, 120, 01, 200\}$

- 7. Проверить, является ли B кодом ровно одного сообщения в кодировании C.
 - 7.1. $C = \{1201, 0110, 110, 00, 212, 10, 21\}, B = 12012121102100212212$
 - 7.2. $C = \{022, 22, 120, 2011, 21, 00, 0000\}, B = 2220110000201122022022$
 - 7.3. $C = \{1000, 21, 22, 222, 12, 0101, 0202\}, B = 1221222202020222220202$
 - 7.4. $C = \{11, 2102, 01, 1221, 0122, 22, 020\}, B = 2201222102012201221221$
 - 7.5. $C = \{101, 21, 20, 102, 221, 00, 01, 010\}, B = 00010101010100011022001$
 - 7.6. $C = \{01, 20, 1120, 12, 012, 02, 21, 100\}, B = 100211120100011120012$
 - 7.7. $C = \{0202, 02, 210, 21, 2200, 00, 001\}, B = 00121022000202001022200$
 - 7.8. $C = \{10, 1212, 2022, 21, 100, 212, 201\}, B = 20121221121221100212$
 - 7.9. $C = \{20, 22, 0200, 1112, 0222, 2002\}, B = 20022002111202002020202$
- 7.10. $C = \{120, 20, 221, 122, 22, 2201, 02, 10\}, B = 1222220202022221122122$
- 7.11. $C = \{1210, 00, 200, 21, 111, 121, 021\}, B = 0212112121200111121200$
- 7.12. $C = \{20, 02, 2120, 010, 120, 0221, 102\}, B = 2021202120022102010102$
- 7.13. $C = \{02, 1102, 21, 0100, 020, 002, 0220\}, B = 0220211102110211021102$
- 7.14. $C = \{20, 21, 2020, 12, 01, 11, 210, 201\}, B = 2020202020212100112201$
- 7.15. $C = \{102, 112, 12, 2210, 00, 11, 121, 2120\}, B = 12111212111222102210$
- $7.16.\ C=\{22,022,212,20,1011,02,211,2200\}, B=201011101121222211212$
- 7.17. $C = \{12,0110,110,20,00,10,02,22,200\}, B = 11012110002002200110$
- 7.18. $C = \{2011, 201, 220, 01, 12, 0201, 0021\}, B = 220002120112012202011$
- 7.19. $C = \{1120, 22, 0100, 021, 100, 220, 2202\}, B = 22112022011200100021$
- 7.20. $C = \{211, 22, 011, 1212, 21, 1100, 121\}, B = 22121211211211211210111100$
- 7.21. $C = \{21, 22, 122, 02, 0011, 2111, 1220\}, B = 00111220220011220011$
- 7.22. $C = \{01, 001, 1022, 120, 12, 010, 211\}, B = 120102212121200101010$
- 7.23. $C = \{20, 121, 22, 0011, 01, 1100, 10, 001\}, B = 00120110012120012201$
- $7.24.\ C=\{00,000,12,20,22,0202,2110,01\}, B=20000000002110002000$
- 7.25. $C = \{00, 12, 22, 0021, 01, 100, 11, 222\}, B = 01100110021122222100$
- 7.26. $C = \{2101, 002, 12, 120, 21, 1002, 122\}, B = 122122120002122121002$
- 7.27. $C = \{2021, 00, 000, 0112, 11, 200, 21\}, B = 21202100000112202100$
- 7.28. $C = \{12,011,210,202,21,2112,112\}, B = 120111120112121001112$
- 7.29. $C = \{0012, 20, 210, 220, 2121, 2102\}, B = 21022102202102210202121$
- 7.30. $C = \{21,0011,20,2122,220,0220,211\}, B = 212221022021102200220$

- 8. Построить двоичный префиксный код с заданной последовательностью длин кодовых слов L.
 - 8.1. L = (3, 3, 3, 3)
 - 8.2. L = (1, 2, 5, 6)
 - 8.3. L = (3, 4, 4, 4)
 - 8.4. L = (2, 2, 2, 3)
 - 8.5. L = (1, 3, 3, 5)
 - 8.6. L = (2, 3, 3, 4)
 - 8.7. L = (3, 4, 4, 5)
 - 8.8. L = (1, 2, 5, 5)
 - 8.9. L = (3, 3, 4, 5)
- 8.10. L = (3, 4, 6, 6)
- 8.11. L = (3, 5, 5, 5)
- 8.12. L = (3, 3, 3, 4)
- 8.13. L = (1, 3, 4, 5)
- 8.14. L = (1, 2, 3, 6)
- 8.15. L = (3, 4, 5, 6)
- 8.16. L = (1, 4, 5, 6)
- 8.17. L = (2, 2, 3, 3)
- 8.18. L = (1, 2, 4, 4)
- 8.19. L = (2, 3, 4, 5)
- 8.20. L = (2, 2, 3, 4)
- 8.21. L = (1, 3, 3, 3)
- 8.22. L = (2, 2, 2, 5)
- 8.23. L = (1, 2, 3, 5)
- 8.24. L = (3, 4, 5, 5)
- 8.25. L = (1, 3, 3, 4)
- 8.26. L = (2, 3, 3, 3)
- 8.27. L = (1, 4, 4, 4)
- 8.28. L = (2, 2, 2, 4)
- 8.29. L = (2, 3, 4, 4)
- 8.30. L = (1, 2, 3, 4)

- 9. Построить двоичный префиксный код с заданной последовательностью длин кодовых слов L.
 - 9.1. L = (2, 2, 2, 5, 6, 7)
 - 9.2. L = (3, 3, 3, 4, 4, 4)
 - 9.3. L = (3, 4, 4, 5, 5, 5)
 - 9.4. L = (2, 2, 3, 4, 5, 6)
 - 9.5. L = (3, 4, 5, 6, 6, 7)
 - 9.6. L = (2, 2, 3, 3, 3, 4)
 - 9.7. L = (1, 2, 5, 6, 7, 7)
 - 9.8. L = (2, 3, 4, 4, 5, 5)
 - 9.9. L = (3, 4, 5, 5, 6, 6)
- 9.10. L = (2, 3, 4, 5, 6, 6)
- 9.11. L = (2, 3, 3, 4, 4, 4)
- 9.12. L = (1, 2, 3, 4, 5, 6)
- 9.13. L = (1, 2, 4, 4, 5, 6)
- 9.14. L = (3, 3, 4, 5, 5, 5)
- 9.15. L = (3, 4, 4, 4, 5, 5)
- 9.16. L = (2, 3, 4, 5, 6, 7)
- 9.17. L = (3, 4, 4, 4, 4, 5)
- 9.18. L = (1, 3, 4, 4, 5, 6)
- 9.19. L = (2, 3, 4, 4, 4, 4)
- 9.20. L = (2, 2, 3, 3, 4, 4)
- 9.21. L = (3, 3, 3, 4, 5, 5)
- 9.22. L = (3, 3, 4, 4, 4, 4)
- 9.23. L = (1, 3, 3, 4, 5, 6)
- 9.24. L = (3, 3, 4, 4, 5, 5)
- 9.25. L = (2, 3, 3, 4, 4, 5)
- 9.26. L = (3, 3, 3, 3, 3, 3, 4)
- 9.27. L = (3, 4, 5, 6, 7, 7)
- 9.28. L = (2, 3, 4, 4, 4, 5)
- 9.29. L = (1, 2, 3, 5, 5, 6)
- 9.30. L = (2, 2, 2, 5, 6, 6)

- 10. Построить q-ичный префиксный код (q=3) с заданной последовательностью длин кодовых слов L.
 - 10.1. L = (2, 3, 3, 3)
 - 10.2. L = (2, 2, 2, 3)
 - 10.3. L = (3, 3, 3, 5)
 - 10.4. L = (2, 2, 3, 3)
 - 10.5. L = (1, 3, 3, 4)
 - 10.6. L = (3, 4, 4, 5)
 - 10.7. L = (1, 1, 2, 4)
 - 10.8. L = (3, 3, 3, 3)
 - 10.9. L = (1, 1, 4, 4)
- 10.10. L = (1, 2, 2, 3)
- 10.11. L = (1, 1, 2, 5)
- 10.12. L = (3, 4, 4, 4)
- 10.13. L = (1, 1, 2, 2)
- 10.14. L = (3, 3, 4, 5)
- 10.15. L = (1, 1, 3, 3)
- 10.16. L = (3, 4, 5, 6)
- 10.17. L = (2, 2, 3, 4)
- 10.18. L = (3, 3, 4, 4)
- 10.19. L = (3, 3, 3, 4)
- 10.20. L = (1, 2, 3, 4)
- 10.21. L = (1, 2, 2, 2)
- 10.22. L = (2, 3, 4, 4)
- 10.23. L = (1, 1, 3, 4)
- 10.24. L = (1, 2, 3, 3)
- 10.25. L = (2, 2, 2, 4)
- 10.26. L = (2, 3, 3, 4)
- 10.27. L = (2, 2, 2, 5)
- 10.28. L = (2, 2, 2, 2)
- 10.29. L = (1, 1, 2, 3)
- 10.30. L = (3, 4, 5, 5)

- 11. Построить q-ичный префиксный код (q=3) с заданной последовательностью длин кодовых слов L.
 - 11.1. L = (3, 3, 3, 3, 3, 3, 3)
 - 11.2. L = (3, 3, 3, 3, 3, 5)
 - 11.3. L = (3, 4, 5, 6, 7, 8)
 - 11.4. L = (1, 2, 3, 3, 4, 4)
 - 11.5. L = (3, 3, 4, 4, 5, 5)
 - 11.6. L = (2, 3, 3, 4, 4, 4)
 - 11.7. L = (3, 4, 5, 5, 5, 6)
 - 11.8. L = (2, 3, 4, 4, 5, 6)
 - 11.9. L = (1, 2, 2, 3, 4, 4)
- 11.10. L = (2, 2, 2, 2, 2, 2)
- 11.11. L = (2, 2, 3, 3, 3, 4)
- 11.12. L = (1, 2, 2, 2, 3, 3)
- 11.13. L = (2, 2, 2, 2, 2, 3)
- 11.14. L = (2, 2, 3, 3, 3, 3)
- 11.15. L = (1, 1, 2, 2, 3, 4)
- 11.16. L = (3, 3, 3, 3, 4, 5)
- 11.17. L = (1, 2, 2, 2, 3, 4)
- 11.18. L = (3, 4, 4, 4, 4, 4)
- 11.19. L = (1, 2, 2, 3, 3, 4)
- 11.20. L = (3, 4, 5, 6, 6, 7)
- 11.21. L = (2, 3, 3, 3, 3, 3)
- 11.22. L = (2, 3, 4, 5, 5, 5)
- 11.23. L = (1, 2, 2, 2, 2, 2)
- 11.24. L = (3, 3, 4, 5, 5, 6)
- 11.25. L = (1, 2, 3, 4, 5, 5)
- 11.26. L = (2, 2, 2, 3, 3, 3)
- 11.27. L = (2, 3, 4, 4, 4, 4)
- 11.28. L = (1, 2, 3, 4, 4, 4)
- 11.29. L = (3, 3, 3, 4, 4, 4)
- 11.30. L = (3, 4, 4, 4, 5, 5)

- 12. Построить оптимальный двоичный код для заданного распределения вероятностей P.
 - 12.1. P = (0, 37; 0, 13; 0, 14; 0, 19; 0, 11; 0, 06)
 - 12.2. P = (0,06; 0,32; 0,17; 0,18; 0,03; 0,24)
 - 12.3. P = (0,62; 0,23; 0,09; 0,01; 0,03; 0,02)
 - 12.4. P = (0, 58; 0, 17; 0, 14; 0, 07; 0, 02; 0, 02)
 - 12.5. P = (0, 55; 0, 01; 0, 12; 0, 04; 0, 12; 0, 16)
 - 12.6. P = (0, 44; 0, 08; 0, 24; 0, 09; 0, 01; 0, 14)
 - 12.7. P = (0, 55; 0, 18; 0, 07; 0, 11; 0, 04; 0, 05)
 - 12.8. P = (0, 51; 0, 03; 0, 14; 0, 07; 0, 13; 0, 12)
 - 12.9. P = (0, 54; 0, 07; 0, 13; 0, 05; 0, 06; 0, 15)
 - 12.10. P = (0, 44; 0, 26; 0, 01; 0, 18; 0, 01; 0, 1)
- 12.11. P = (0.62; 0.08; 0.02; 0.01; 0.03; 0.24)
- 12.12. P = (0,04; 0,3; 0,17; 0,3; 0,03; 0,16)
- 12.13. P = (0, 45; 0, 01; 0, 26; 0, 17; 0, 03; 0, 08)
- 12.14. P = (0,06; 0,44; 0,31; 0,03; 0,04; 0,12)
- 12.15. P = (0, 28; 0, 32; 0, 13; 0, 18; 0, 06; 0, 03)
- 12.16. P = (0, 51; 0, 16; 0, 1; 0, 13; 0, 06; 0, 04)
- 12.17. P = (0, 6; 0, 09; 0, 12; 0, 02; 0, 08; 0, 09)
- 12.18. P = (0, 66; 0, 05; 0, 06; 0, 01; 0, 1; 0, 12)
- 12.19. P = (0,05; 0,63; 0,05; 0,13; 0,06; 0,08)
- 12.20. P = (0, 41; 0, 28; 0, 03; 0, 17; 0, 07; 0, 04)
- 12.21. P = (0, 27; 0, 04; 0, 05; 0, 17; 0, 12; 0, 35)
- 12.22. P = (0, 25; 0, 16; 0, 2; 0, 03; 0, 03; 0, 33)
- 12.23. P = (0, 18; 0, 47; 0, 02; 0, 21; 0, 01; 0, 11)
- 12.24. P = (0,3; 0,09; 0,23; 0,19; 0,05; 0,14)
- 12.25. P = (0, 54; 0, 2; 0, 09; 0, 07; 0, 06; 0, 04)
- 12.26. P = (0, 1; 0, 12; 0, 1; 0, 24; 0, 04; 0, 4)
- 12.27. P = (0, 45; 0, 33; 0, 01; 0, 13; 0, 01; 0, 07)
- 12.28. P = (0,61; 0,22; 0,05; 0,08; 0,01; 0,03)
- 12.29. P = (0, 16; 0, 05; 0, 46; 0, 07; 0, 14; 0, 12)
- 12.30. P = (0, 34; 0, 37; 0, 04; 0, 09; 0, 1; 0, 06)

- 13. Построить оптимальный q-ичный код (q=3) для заданного распределения вероятностей P.
 - 13.1. P = (0,09; 0,04; 0,23; 0,2; 0,28; 0,16)
 - 13.2. P = (0, 57; 0, 26; 0, 07; 0, 01; 0, 05; 0, 04)
 - 13.3. P = (0,62; 0,03; 0,22; 0,05; 0,02; 0,06)
 - 13.4. P = (0, 27; 0, 38; 0, 09; 0, 07; 0, 09; 0, 1)
 - 13.5. P = (0, 58; 0, 06; 0, 1; 0, 08; 0, 11; 0, 07)
 - 13.6. P = (0,08; 0,17; 0,18; 0,17; 0,2; 0,2)
 - 13.7. P = (0, 31; 0, 05; 0, 13; 0, 09; 0, 24; 0, 18)
 - 13.8. P = (0,01; 0,62; 0,06; 0,18; 0,02; 0,11)
 - 13.9. P = (0, 47; 0, 35; 0, 04; 0, 04; 0, 02; 0, 08)
 - 13.10. P = (0, 56; 0, 29; 0, 01; 0, 06; 0, 01; 0, 07)
- 13.11. P = (0, 55; 0, 05; 0, 05; 0, 22; 0, 09; 0, 04)
- 13.12. P = (0, 37; 0, 32; 0, 2; 0, 01; 0, 03; 0, 07)
- 13.13. P = (0,65; 0,15; 0,01; 0,01; 0,07; 0,11)
- 13.14. P = (0, 1; 0, 14; 0, 18; 0, 22; 0, 23; 0, 13)
- 13.15. P = (0, 15; 0, 25; 0, 27; 0, 17; 0, 06; 0, 1)
- 13.16. P = (0, 61; 0, 11; 0, 1; 0, 06; 0, 04; 0, 08)
- 13.17. P = (0, 36; 0, 29; 0, 22; 0, 04; 0, 04; 0, 05)
- 13.18. P = (0, 15; 0, 55; 0, 06; 0, 13; 0, 07; 0, 04)
- 13.19. P = (0,66; 0,09; 0,07; 0,01; 0,09; 0,08)
- 13.20. P = (0, 5; 0, 19; 0, 02; 0, 04; 0, 05; 0, 2)
- 13.21. P = (0, 2; 0, 47; 0, 13; 0, 09; 0, 04; 0, 07)
- 13.22. P = (0,41; 0,3; 0,16; 0,07; 0,01; 0,05)
- 13.23. P = (0, 33; 0, 24; 0, 04; 0, 09; 0, 14; 0, 16)
- 13.24. P = (0, 53; 0, 11; 0, 13; 0, 01; 0, 15; 0, 07)
- 13.25. P = (0,07; 0,47; 0,22; 0,12; 0,05; 0,07)
- 13.26. P = (0, 12; 0, 43; 0, 28; 0, 08; 0, 05; 0, 04)
- 13.27. P = (0,61; 0,1; 0,15; 0,05; 0,05; 0,04)
- 13.28. P = (0, 5; 0, 13; 0, 11; 0, 04; 0, 02; 0, 2)
- 13.29. P = (0,63; 0,17; 0,05; 0,04; 0,04; 0,07)
- 13.30. P = (0, 23; 0, 4; 0, 05; 0, 16; 0, 08; 0, 08)

14. Построить по методу Хэмминга кодовое слово для сообщения α .

- 14.1. $\alpha = 10000$
- 14.2. $\alpha = 01000$
- 14.3. $\alpha = 11000$
- 14.4. $\alpha = 00100$
- 14.5. $\alpha = 10100$
- 14.6. $\alpha = 01100$
- 14.7. $\alpha = 11100$
- 14.8. $\alpha = 00010$
- 14.9. $\alpha = 10010$
- 14.10. $\alpha = 01010$
- 14.11. $\alpha = 11010$
- 14.12. $\alpha = 00110$
- 14.13. $\alpha = 10110$
- 14.14. $\alpha = 01110$
- 14.15. $\alpha = 11110$
- 14.16. $\alpha = 00001$
- 14.17. $\alpha = 10001$
- 14.18. $\alpha = 01001$
- 14.19. $\alpha = 11001$
- 14.20. $\alpha = 00101$
- 14.21. $\alpha = 10101$
- 14.22. $\alpha = 01101$
- 14.23. $\alpha = 11101$
- 14.24. $\alpha = 00011$
- 14.25. $\alpha = 10011$
- 14.26. $\alpha = 01011$
- 14.27. $\alpha = 11011$
- 14.28. $\alpha = 00111$
- 14.29. $\alpha = 10111$
- 14.30. $\alpha = 01111$

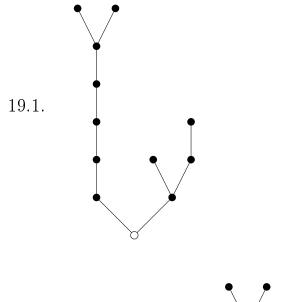
- 15. По кодовому слову β , построенному по методу Хэмминга, восстановить исходное сообщение, если известно, что произошло не более одной ошибки.
 - 15.1. $\beta = 100000000$
 - 15.2. $\beta = 001000000$
 - 15.3. $\beta = 111000000$
 - 15.4. $\beta = 010100000$
 - 15.5. $\beta = 101100000$
 - 15.6. $\beta = 000010000$
 - 15.7. $\beta = 110010000$
 - 15.8. $\beta = 011010000$
 - 15.9. $\beta = 100110000$
- 15.10. $\beta = 001110000$
- 15.11. $\beta = 1111110000$
- 15.12. $\beta = 010001000$
- 15.13. $\beta = 101001000$
- 15.14. $\beta = 000101000$
- 15.15. $\beta = 110101000$
- 15.16. $\beta = 011101000$
- 15.17. $\beta = 100011000$
- 15.18. $\beta = 001011000$
- 15.19. $\beta = 111011000$
- 15.20. $\beta = 010111000$
- 15.21. $\beta = 1011111000$
- 15.22. $\beta = 000000100$
- 15.23. $\beta = 110000100$
- 15.24. $\beta = 011000100$
- 15.25. $\beta = 100100100$
- 15.26. $\beta = 001100100$
- 15.27. $\beta = 111100100$
- 15.28. $\beta = 010010100$
- 15.29. $\beta = 101010100$
- 15.30. $\beta = 000110100$

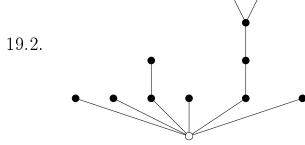
- 16. Для кода C определить, сколько ошибок он обнаруживает и сколько исправляет.
 - 16.1. $C = \{11010010, 01010100, 01101100, 10100000\}$
 - 16.2. $C = \{10110000, 11011010, 01000010, 00110110\}$
 - 16.3. $C = \{11011000, 00001101, 00001000, 11000111\}$
 - 16.4. $C = \{00000101, 10110000, 01100111, 01001110\}$
 - 16.5. $C = \{01101110, 10000100, 01010110, 11100011\}$
 - 16.6. $C = \{10111011, 11110110, 00110000, 00000010\}$
 - 16.7. $C = \{10000111, 01111010, 11100100, 00000101\}$
 - 16.8. $C = \{10101000, 01111100, 11011111, 10000000\}$
 - 16.9. $C = \{101111110, 01000010, 10100010, 00101000\}$
- 16.10. $C = \{01100110, 01001010, 11001001, 00010011\}$
- 16.11. $C = \{01100011, 00000101, 00011001, 01010011\}$
- 16.12. $C = \{00010000, 00101101, 01010110, 11100100\}$
- 16.13. $C = \{01111001, 01011111, 01101111, 11101011\}$
- 16.14. $C = \{111111110, 01110110, 10100110, 00111001\}$
- 16.15. $C = \{00011101, 10110100, 00101001, 11000100\}$
- 16.16. $C = \{10010101, 01010100, 00100001, 10100101\}$
- 16.17. $C = \{11101001, 10110110, 11000001, 11100010\}$
- 16.18. $C = \{00100111, 10100011, 11010101, 11001010\}$
- 16.19. $C = \{11110001, 11010011, 00010011, 01101110\}$
- 16.20. $C = \{10000101, 01000101, 10000000, 00110110\}$
- 16.21. $C = \{11101011, 11001110, 00001100, 01001011\}$
- $16.22.\ C = \{00000111, 01101101, 01011110, 00010011\}$
- $16.23.\ C = \{11010011, 10110101, 111111010, 10101100\}$
- $16.24. C = \{101111100, 10000000, 10011101, 11100101\}$
- 16.25. $C = \{00100000, 01010001, 11001101, 10010101\}$
- 16.26. $C = \{00101100, 01100111, 01001101, 10101000\}$
- $16.27. C = \{11000100, 00000100, 01001011, 01110110\}$
- 16.28. $C = \{00011110, 00100000, 01000110, 011111111\}$
- $16.29.\ C = \{111111110, 01000011, 11110010, 10010110\}$
- 16.30. $C = \{01110001, 10111111, 01100010, 10110000\}$

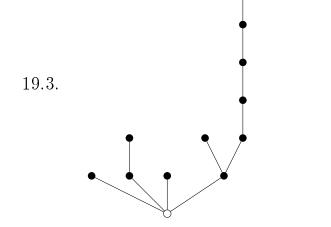
- 17. Определить, сколько ошибок обнаруживает и сколько исправляет код с характеристической функцией f.
 - 17.1. $x_2(x_3 \oplus x_1)(x_3 \oplus x_4)$
 - 17.2. $x_4(x_2 \sim x_1)(x_3 \sim x_2)$
 - 17.3. $x_4 \downarrow ((x_1 \sim x_3) \mid (x_2 \sim x_1))$
 - 17.4. $x_3 \downarrow ((x_1 \sim x_2) \lor (x_1 \sim x_4))$
 - 17.5. $\overline{(x_1 \oplus x_3)}(x_4 \downarrow (x_2 \sim x_1))$
 - 17.6. $x_3((x_4 \oplus x_1) \downarrow (x_2 \oplus x_1))$
 - 17.7. $((x_1 \oplus x_2) \downarrow (x_1 \sim x_3))x_4$
 - 17.8. $((x_3 \sim x_1) \lor (x_1 \sim x_4)) \downarrow x_2$
 - 17.9. $(x_4 \sim x_3)(\overline{x}_1 \downarrow (x_2 \oplus x_4))$
- 17.10. $(x_1 \lor (x_4 \oplus x_3)) \downarrow (x_4 \sim x_2)$
- 17.11. $((x_2 \oplus x_3) \lor (x_1 \sim x_3)) \downarrow x_4$
- 17.12. $(x_4 \mid (x_3 \sim x_2)) \perp (x_2 \oplus x_1)$
- 17.13. $((x_4 \oplus x_1) \to x_3) \downarrow (\overline{x}_1 \sim \overline{x}_2)$
- 17.14. $(x_2 \mid (x_3 \oplus x_1)) \downarrow (x_4 \oplus x_1)$
- 17.15. $(x_1 \to (x_3 \oplus x_4)) \downarrow (x_4 \sim x_2)$
- 17.16. $x_4 \downarrow ((x_1 \sim x_3) \mid (x_2 \oplus x_3))$
- 17.17. $(x_3 \oplus x_4) \downarrow (x_2 \lor (x_1 \oplus x_3))$
- 17.18. $x_1 \downarrow ((x_3 \oplus x_4) \mid (x_4 \oplus x_2))$
- 17.19. $x_3 \downarrow ((x_4 \sim x_1) \to (x_4 \sim x_2))$
- 17.20. $(x_1 \oplus x_2) \downarrow (x_3 \lor (x_2 \oplus x_4))$
- 17.21. $(x_2 \downarrow (x_1 \oplus x_4))(x_1 \oplus x_3)$
- 17.22. $x_2 \downarrow ((x_4 \sim x_1) \mid (x_3 \oplus x_1))$
- 17.23. $(x_2 \to (x_4 \oplus x_1)) \downarrow (x_3 \oplus x_1)$
- 17.24. $(x_3 \sim x_1) \downarrow ((x_1 \sim x_4) \rightarrow x_2)$
- 17.25. $(x_1 \lor (x_4 \sim x_3)) \downarrow (x_3 \oplus x_2)$
- 17.26. $((x_1 \sim x_4) \mid (x_4 \oplus x_3)) \downarrow x_2$
- 17.27. $x_4((x_3 \oplus x_2) \downarrow (x_2 \sim x_1))$
- 17.28. $x_1(x_3 \oplus x_2)(x_2 \sim x_1 \sim x_4)$
- 17.29. $(x_2 \oplus x_3) \downarrow ((x_1 \oplus x_2) \vee x_4)$
- 17.30. $((x_4 \oplus x_1) \mid x_3) \downarrow (x_1 \sim x_2)$

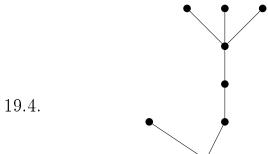
- 18. Построить плоское корневое дерево по его коду $\tilde{\alpha}$.
 - 18.1. $\tilde{\alpha} = 00000101111110100101101$
 - 18.2. $\tilde{\alpha} = 01001101010000111110011$
 - 18.3. $\tilde{\alpha} = 01000000110001101111111$
 - 18.4. $\tilde{\alpha} = 00011000010100011111111$
 - 18.5. $\tilde{\alpha} = 01000001111110100001111$
 - 18.6. $\tilde{\alpha} = 0011010101011000011111$
 - 18.7. $\tilde{\alpha} = 00010111100011101001101$
 - 18.8. $\tilde{\alpha} = 01010100001000110111111$
 - 18.9. $\tilde{\alpha} = 01010001001010001111111$
- 18.10. $\tilde{\alpha} = 0000000110100111101111$
- 18.11. $\tilde{\alpha} = 000101111001001001111101$
- 18.12. $\tilde{\alpha} = 0000110110110101010101$
- 18.13. $\tilde{\alpha} = 00100010101101010101111$
- 18.14. $\tilde{\alpha} = 01000001000001111111111$
- 18.15. $\tilde{\alpha} = 0101000011011100001111$
- 18.16. $\tilde{\alpha} = 0101001010110101001011$
- 18.17. $\tilde{\alpha} = 0101010101000110001111$
- 18.18. $\tilde{\alpha} = 0000010100111011110101$
- 18.19. $\tilde{\alpha} = 01010001010011111010011$
- 18.20. $\tilde{\alpha} = 0001110011001001100111$
- 18.21. $\tilde{\alpha} = 0100001010011011101011$
- 18.22. $\tilde{\alpha} = 0100110010101010110011$
- 18.23. $\tilde{\alpha} = 0101000111010100001111$
- 18.24. $\tilde{\alpha} = 00110101011001101101011$
- 18.25. $\tilde{\alpha} = 00000001111111011010011$
- 18.26. $\tilde{\alpha} = 0010110101010001011101$
- 18.27. $\tilde{\alpha} = 0101010000101101001111$
- 18.28. $\tilde{\alpha} = 0101001011001100001111$
- 18.29. $\tilde{\alpha} = 00110101010101010101011$
- 18.30. $\tilde{\alpha} = 0010001110110100110011$

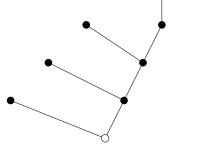
19. Построить код плоского корневого дерева.

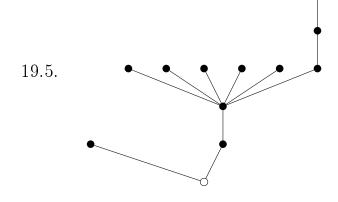


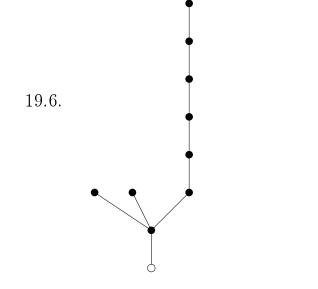


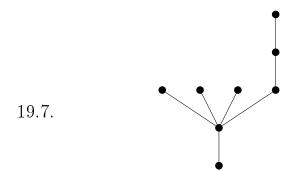


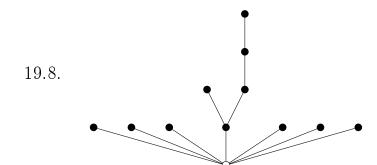


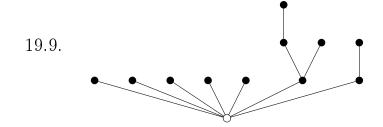


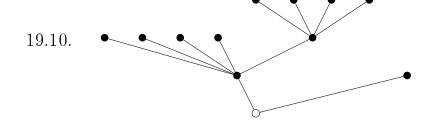


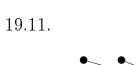


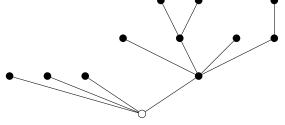




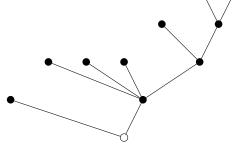


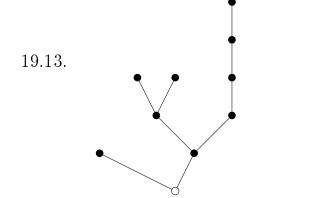


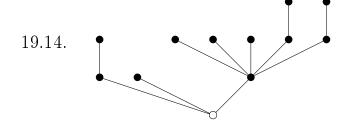


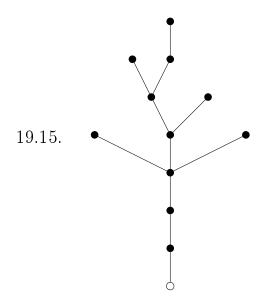


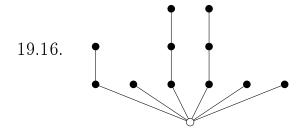


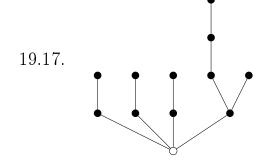


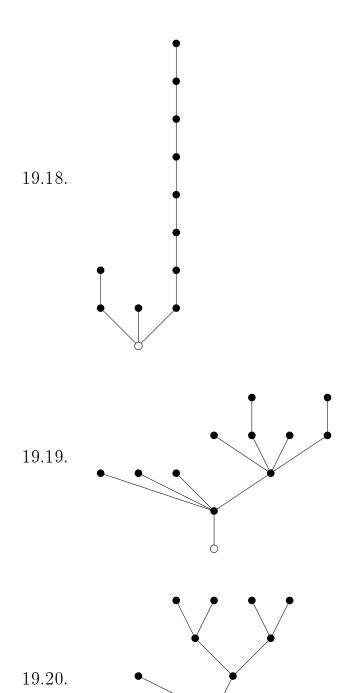




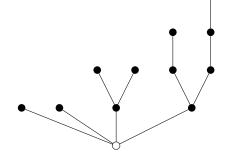




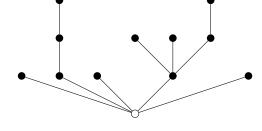




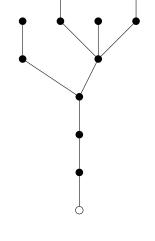




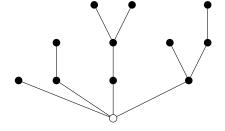
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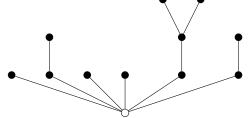
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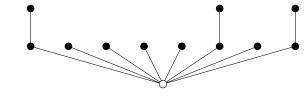
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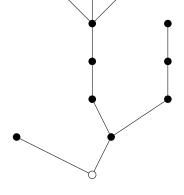




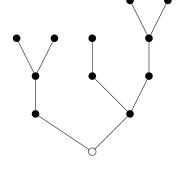
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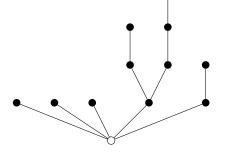
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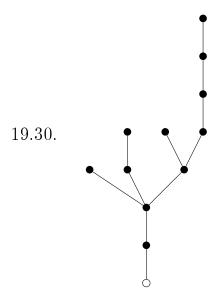


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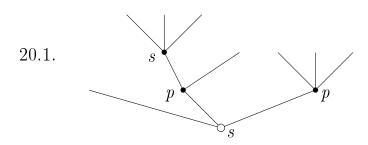


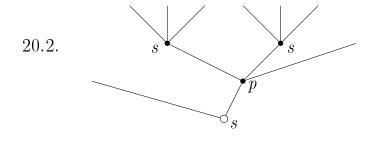
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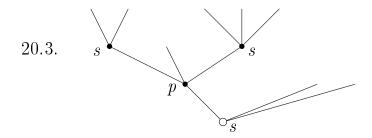


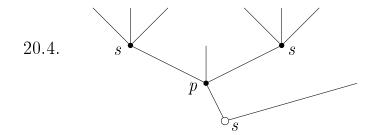


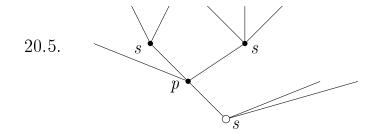
20. По диаграмме расщепления восстановить π -сеть.

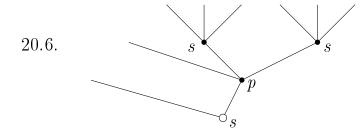


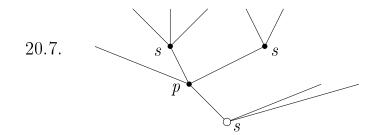


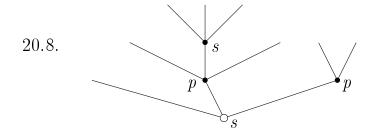


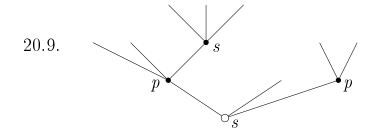


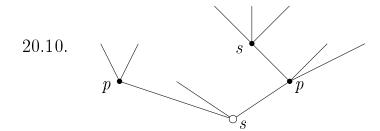




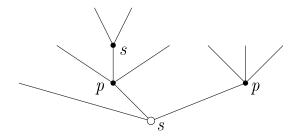




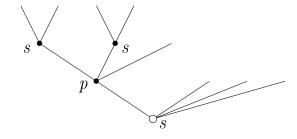




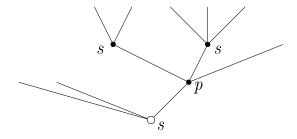




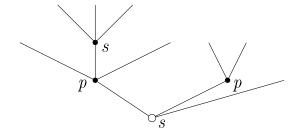
20.12.



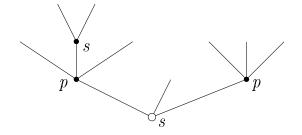
20.13.

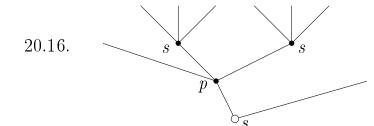


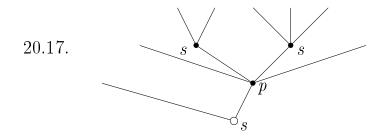
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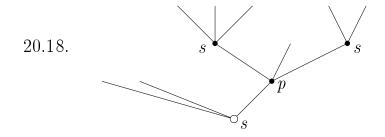


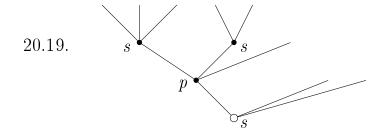
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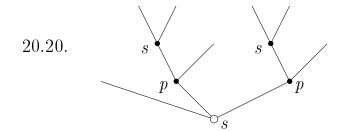




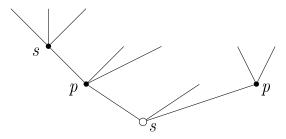




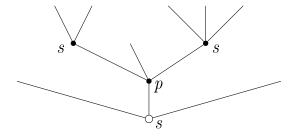




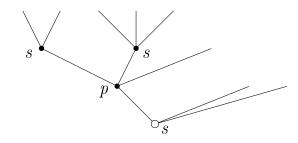




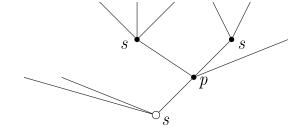
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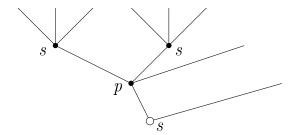
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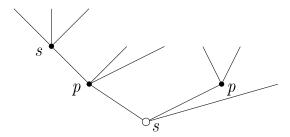
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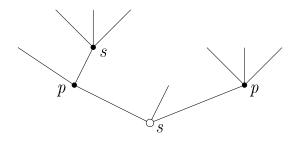
20.25.



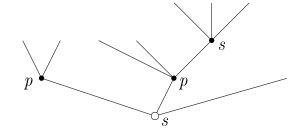




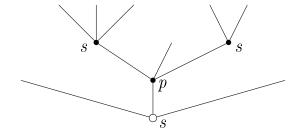
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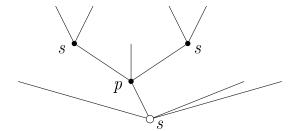
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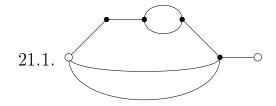
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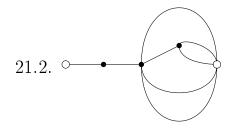


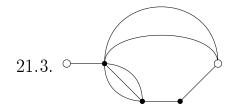
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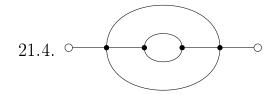


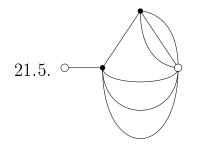
21. Для π -сети построить диаграмму расщепления.

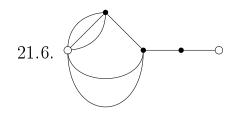




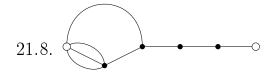


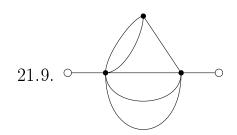






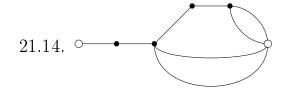


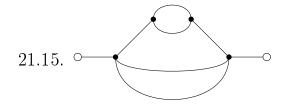


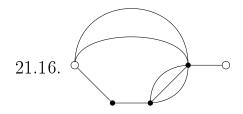


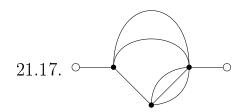


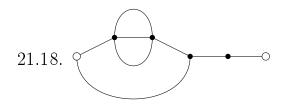


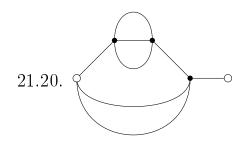


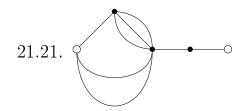


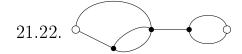


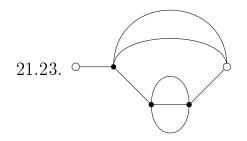


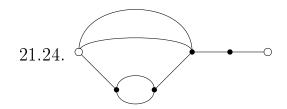


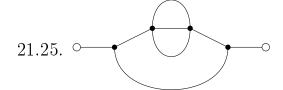


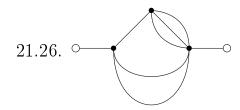


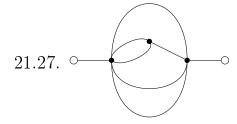


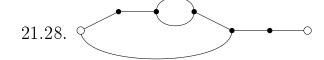


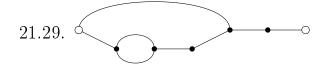


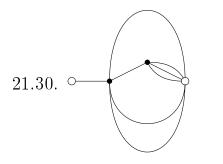












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