2. (7, I); (8, G)

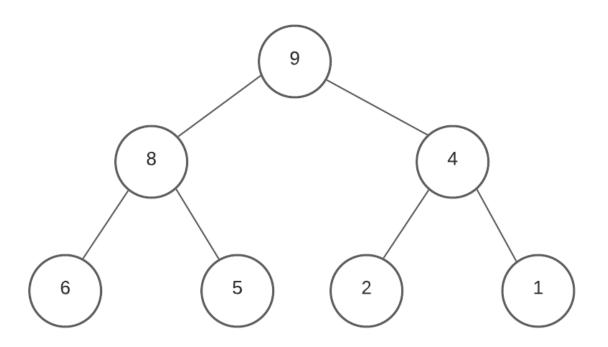
3. Deve-se usar a fila de prioridades por ser um simulador. Portanto o voo será retirado com o menor time-stamp e assim consecutivamente.

4. Pode estar armazenada na raiz da árvore.

5. Sim, pois para uma árvore ser um heap, ela deve ser uma árvore binária completa. Então a árvore T deve ser considerada um heap por ser completa.

6. Não se considera pois, sem criar o filho da esquerda antes, não será possível ter o filho da direita.

<del>7.</del>



Prefixado: 9, 8, 6, 5, 4, 2, 1

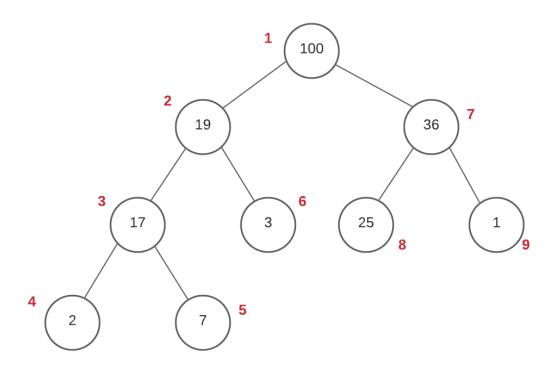
Interfixado: 6, 8, 5, 9, 2, 4, 1

Pós-fixado: 6, 5, 8, 2, 4, 1, 9

8. Prefixado: 1, 2, 4, 8, 9, 5, 10, 11, 3, 6, 12, 13, 7, 14, 15

Interfixado: 8, 4, 9, 2, 10, 5, 11, 1, 12, 6, 13, 3, 14, 7, 15

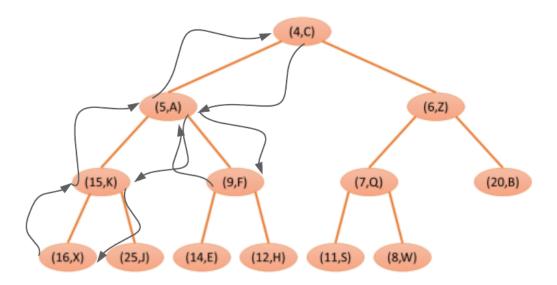
Pós fixado: 8, 9, 4, 10, 11, 5, 2, 12, 13, 6, 14, 15, 7, 3, 1



100 -> 19 -> 17-> 2-> 7 -> 3 -> 36 -> 25 -> 1

10. Um exemplo de heap que prova que ela está errada é o CompleteBinaryTree.

O heap que foi fornecido para provar como errado é a imagem do heap que é representado pela lista de [x, 1,5,2,8,9,1,6]. esta pilha não irá produzir chaves em ordem crescente quando um caminhamento pós fixado é usado.



12.

<del>13.</del>

- 14. O hashing duplo consegue tolerar um fator de carga superior a 1 e o que não consegue tolerar o fator de carga é o endereçamento aberto.
- 15. Se usarmos a soma dos componentes, a função hash seria x0 + x1 + x2 + x3 + x4 mod N.

<del>16.</del>

11-entry hash table: h(i) = (3i+5) mod 11

empty hash table:

<del>0, 0, 0, 0, 0, 0, 0, 0, 0, 0</del>

hash key 12:

<del>[], [], [], [], [], [], [], [12], [], []</del>

hash key 44

```
<del>[], [], [], [], [], [44], [], [], [12], [], []</del>
hash key 13
[13], [], [], [], [], [44], [], [], [12], [], []
hash key 88
[13], [], [], [], [[44, 88]], [], [], [12], [], []
hash key 23
[13], [], [], [], [[44, 88]], [], [], [12, 23], [], []
hash key 94
[13], [94], [], [], [], [[44, 88]], [], [], [[12, 23]], [], []
hash key 11
[13], [94], [], [], [], [[44, 88, 11]], [], [], [[12, 23]], [], []
hash key 39
[13], [[94, 39]], [], [], [[, [[44, 88, 11]], [], [], [[12, 23]], [], []
hash key 20
[13], [[94, 39]], [], [], [], [[44, 88, 11]], [], [], [[12, 23]], [], [20]
hash key 16
[13], [[94, 39]], [], [], [[44, 88, 11]], [], [[, [[12, 23]], [16], [20]
hash key 5
[13], [[94, 39]], [], [], [], [[44, 88, 11]], [], [], [[12, 23]], [[16, 5]], [20]
<del>17.</del>
11-entry hash table: h(i) = (3i+5) \mod 11
empty hash table:
```

[13], [94], [39], [16], [], [44], [88], [11], [12], [23], [20]

hash key 16

hash key 5

```
[13], [94], [39], [16], [5], [44], [88], [11], [12], [23], [20]
```

<del>18.</del>

11-entry hash table:  $h(i) = (3i+5) \mod 11$ 

Quadratic Probing:  $h(i) = (3i+5) + j^2 \mod 11$  for j = 0,1,2...

empty hash table:

0, 0, 0, 0, 0, 0, 0, 0, 0

hash key 12 (j = 0)

<del>[], [], [], [], [], [], [], [], [], []</del>

hash key 44 (j = 0)

<del>[], [], [], [], [], [44], [], [], [12], [], []</del>

hash key 13 (j = 0)

[13], [], [], [], [], [44], [], [], [12], [], []

hash key 88 (j = 1)

[13], [], [], [], [44], [88], [], [12], [], []

hash key 23 (j = 1)

[13], [], [], [], [44], [88], [], [12], [23], []

hash key 94 (j = 0)

[13], [94], [], [], [], [44], [88], [], [12], [23], []

hash key 11 (j = 3)

[13], [94], [], [11], [], [44], [88], [], [12], [23], []

hash key 39 (j = 1)

[13], [94], [39], [11], [], [44], [88], [], [12], [23], []

```
hash key 20 (j = 0)
[13], [94], [39], [11], [], [44], [88], [], [12], [23], [20]
hash key 16 (j = 3)
[13], [94], [39], [11], [], [44], [88], [16], [12], [23], [20]
hash key 5 (i = ?)
all attempts for j = 1 to 10, mod 11 does not equal the empty bucket number 4.
<del>19.</del>
11 entry hash table: h(i) = [(3i+5) + i \times [7 + (k \mod 7)]] \mod 11, for i = 0,1,2...
empty hash table:
<del>0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0</del>
hash key 12:
<del>[], [], [], [], [], [], [], [12], [], []</del>
hash key 44
<del>[], [], [], [], [], [44], [], [], [12], [], []</del>
hash key 13
[13], [], [], [], [], [44], [], [], [12], [], []
hash key 88
h(i) = [(3i+5) + (i \times [7 - (k \mod 7)])] \mod 11, for i = 3
h(88) = [(3(88)+5) + (3 \times [7 + (88 \mod 7)])] \mod 11
h(88) = [269 + (3 \times [7 - (88 \mod 7)])] \mod 11
h(88) = [269 + (3 \times [7 - 4])] \mod 11
h(88) = [269 + 9] \mod 11
h(88) = 278 \mod 11 = 3
[13], [], [], [88], [], [44], [], [], [12], [], []
hash key 23
```

```
h(i) = [(3i+5) + (j \times [7 - (k \mod 7)])] \mod 11, for j = 1
h(23) = [(3(23) + 5) + (1 \times [7 - (23 \mod 7)])] \mod 11
h(23) = [74 + (1 \times [7 - (23 \mod 7)])] \mod 11
h(23) = [74 + (1 \times [7 - 2])] \mod 11
h(23) = [74 + 5)] \mod 11
h(23) = 79 \mod 11 = 2
[13], [], [23], [88], [], [44], [], [], [12], [], []
hash key 94
[13], [94], [23], [88], [], [44], [], [], [12], [], []
hash kev 11
h(i) = [(3i+5) + (i \times [7 + (k \mod 7)])] \mod 11, for i = 4
h(11) = [(3(11) + 5) + (4 \times [7 - (11 \mod 7)])] \mod 11
h(11) = [38 + (4 \times [7 - (11 \mod 7)])] \mod 11
h(11) = [38 + (4 \times [7 - 4])] \mod 11
h(11) = [38 + (4 \times 3)] \mod 11
h(11) = 50 \mod 11 = 6
[13], [94], [23], [88], [], [44], [11], [], [12], [], []
hash key 39
h(i) = [(3i+5) + (i \times [7 - (k \mod 7)])] \mod 11, for i = 1
h(39) = [(3(39) + 5) + (1 \times [7 - (39 \mod 7)])] \mod 11
h(39) = [122 + (1 \times [7 - (39 \mod 7)])] \mod 11
h(39) = [122 + (1 \times [7 - 4])] \mod 11
h(39) = [122 + (1 \times 3)] \mod 11
h(39) = 125 \mod 11 = 4
[13], [94], [23], [88], [39], [44], [11], [], [12], [], []
hash key 20
[13], [94], [23], [88], [39], [44], [11], [], [12], [], [20]
hash key 16
[13], [94], [23], [88], [39], [44], [11], [], [12], [16], [20]
hash key 5
h(i) = [(3i+5) + (j \times [7 - (k \mod 7)])] \mod 11, for j = 10
h(i) = [(3(5) + 5) + (10 \times [7 - (5 \mod 7)])] \mod 11
```

```
\begin{aligned} &h(i) = [20 + (10 \times [7 - (5 \mod 7)])] \mod 11 \\ &h(i) = [20 + (10 \times [7 - 5])] \mod 11 \\ &h(i) = [20 + (10 \times 2)] \mod 11 \\ &h(i) = 40 \mod 11 = 7 \\ &[13], [94], [23], [88], [39], [44], [11], [5], [12], [16], [20] \end{aligned}
```

<del>20.</del>

21. A principal diferença entre os dois é que o TAD mapa deve ter chaves únicas, já o dicionário pode ter múltiplas entradas com a mesma chave.

22.

<del>23.</del>