CSCI 430: Homework 1

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2.1-1

31	41	59	26	41	58
31	41	59	26	41	58
31	41	59	26	41	58
26	31	41	59	41	58
26	31	41	41	59	58
26	31	41	41	58	59

Blue = If moved

Red = Next to be sorted

2.1-2

Algorithm 1 Reverse Sort Sort

```
for j=2 to A.length do
key = A[j]
i = j-1
while i > 0 and A[i] < key do
A[i+1] = A[i]
i = i-1
end while
A[i+1] = key
end for
```

Algorithm 2 Search for v

```
locate = 0

for i = 1 to A.length do

if A[i] == v then

locate = i

end if

end for

if locate == 0 then

print"NIL"

else

print(locate)

end if
```

Initialization:

Show that 'locate', the variable that will hold the location of 'v' in the array, is 0. A number nonexistent in A[1...n].

Maintenance:

Show that the loop maintains. The body of the loop checks if A[1], A[2], and so on by 1 position to see if the current element is the same as the value of 'v' until A[n]. When an element satisfies the condition if changes 'locate' to the location in the array.

Termination:

Condition of for loop termination is that i > A.length or n. Because each iterator increase by 1, we must have i = (n+1) at that time. Substituting i for (n+1) in the loop we have either 'locate' is 0 or the i were A[i] is equal to v in A[1-n]. Then the if-else statement decides if locate is 0 to print "NIL" or to print the location of v in A[]. Hence the algorithm is correct.

1 2.1-4

Formal:

Input: 2 binary numbers, A and B, represented by A[0...n] and B[0...n] in binary form **Output**: Binary number C[0...(N+1)] where C = A + B.

Pseudo:

Algorithm 3 Binary Addition

```
egin{aligned} &carry = 0 \ &i = A.length \ & \mathbf{while} \ i < 0 \ \mathbf{do} \ & added = A[i] + B[i] + carry \ & C[i+1] = added\%2 \ & carry = added \div 2 \ & i = i-1 \ & \mathbf{end} \ \mathbf{while} \ & C[0] = carry \end{aligned}
```