



# Normal Forms

# Normal Forms

## » Chomsky Normal Form (CNF)

» **Def.:** A context-free grammar  $G = (V, \Sigma, R, S)$  is said to be in Chomsky Normal Form (CNF), iff every rule in  $R$  is of one of the following forms:

»  $X \rightarrow a$  where  $a \in \Sigma$ , or

»  $X \rightarrow BC$  where  $B$  and  $C \in V$



# Example1

» Which one is Chomsky Normal Form ? Why?

$$S \rightarrow AS$$

$$S \rightarrow a$$

$$A \rightarrow SA$$

$$A \rightarrow b$$

$$S \rightarrow AS$$

$$S \rightarrow AAS$$

$$A \rightarrow SA$$

$$A \rightarrow aa$$



$$S \rightarrow AS$$

$$S \rightarrow a$$

$$A \rightarrow SA$$

$$A \rightarrow b$$

Chomsky  
Normal Form

$$S \rightarrow AS$$

$$S \rightarrow AAS$$

$$A \rightarrow SA$$

$$A \rightarrow aa$$

Not Chomsky  
Normal Form



# Converting to Chomsky Normal Form

» There exists 5-steps algorithm to convert a CFG  $G$  into a new grammar  $G_c$  such that:  $L(G) = L(G_c) - \{\epsilon\}$

» *convertChomsky*( $G$ :CFG) =

- » 1-  $G' = \text{removeEps}(G\text{:CFG})$   $S \rightarrow \epsilon$
- » 2-  $G'' = \text{removeUnits}(G'\text{:CFG})$   $A \rightarrow B$
- » 3-  $G''' = \text{removeMixed}(G''\text{:CFG})$   $A \rightarrow aB$
- » 4-  $G'^{\nu} = \text{removeLong}(G'''\text{:CFG})$   $S \rightarrow ABCD$
- » 5-  $G'^{\nu} : L(G) = L(G'^{\nu}) - \{\epsilon\}$

$G^{\nu} = \text{atmostoneEps}(G'^{\nu}\text{:CFG}) \quad S^* \rightarrow S, S \rightarrow \epsilon$



# Remove Epsilon

- » Find the set  $N$  of nullable variables in  $G$ .
- »  $X$  is nullable iff either  $X \rightarrow \varepsilon$  or
- »  $(X \rightarrow A, A \rightarrow \varepsilon) : X \rightarrow \varepsilon$
- » **Ex1:  $G$ :**  $S \rightarrow aACa$
- »  $A \rightarrow B \mid a$
- »  $B \rightarrow C \mid c$
- »  $C \rightarrow cC \mid \varepsilon$
- » **Now,** since  $C \rightarrow \varepsilon$ ,  $C$  is nullable
- » since  $B \rightarrow C$ ,  $B$  is nullable
- » since  $A \rightarrow B$ ,  $A$  is nullable
- » **Therefore**  $N = \{A, B, C\}$



# Remove Epsilon

» According to N, the rules will be:

»  $S \rightarrow aACa$

»  $A \rightarrow B \mid a \mid \epsilon$

»  $B \rightarrow C \mid c \mid \epsilon$

»  $C \rightarrow cC \mid \epsilon$

» Now, removeEps returns  $G'$ :

»  $S \rightarrow aACa \mid aAa \mid aCa \mid aa$

»  $A \rightarrow B \mid a$

»  $B \rightarrow C \mid c$

»  $C \rightarrow cC \mid c$



# Remove Unit Productions

- » **Def.:** unit production is a rule whose right hand side consists of a single nonterminal symbol. Ex:  $A \rightarrow B$
- » Remove any unit production from  $G'$ .
- » **Ex1(continue),  $G'$ :**
  - »  $S \rightarrow aACa \mid aAa \mid aCa \mid aa$
  - »  $A \rightarrow B \mid a$
  - »  $B \rightarrow C \mid c$
  - »  $C \rightarrow cC \mid c$





# Remove Unit Productions

- » **Now** by Applying *removeUnits*( $G':CFG$ ) :
- » Remove  $A \rightarrow B$  But  $B \rightarrow C \mid c$ , so Add  $A \rightarrow C \mid c$
- » Remove  $B \rightarrow C$  Add  $B \rightarrow cC$  ( $B \rightarrow c$ , already there)
- » Remove  $A \rightarrow C$  Add  $A \rightarrow cC$  ( $A \rightarrow c$ , already there)
- » **So** *removeUnits* returns  $G''$  :
  - »  $S \rightarrow aACa \mid aAa \mid aCa \mid aa$
  - »  $A \rightarrow a \mid c \mid cC$
  - »  $B \rightarrow c \mid cC$
  - »  $C \rightarrow cC \mid c$



# Remove Mixed

- » **Def.:** mixed is a rule whose right hand side consists of combination of terminals or terminals with nonterminal symbol.
- » Create a new nonterminal  $T_a$  for each terminal  $a \in \Sigma$
- » For each  $T_a$ , add the rule  $T_a \rightarrow a$
- » **Ex1(continue),  $G''$  :**
  - »  $S \rightarrow aACa \mid aAa \mid aCa \mid aa$
  - »  $A \rightarrow a \mid c \mid cC$
  - »  $B \rightarrow c \mid cC$
  - »  $C \rightarrow cC \mid c$



# Remove Mixed

» Now, by apply *removeMixed*( $G''$ :CFG),  $G'''$  :

»  $S \rightarrow T_a A C T_a \mid T_a A T_a \mid T_a C T_a \mid T_a T_a$

»  $A \rightarrow a \mid c \mid T_c C$

»  $B \rightarrow c \mid T_c C$

»  $C \rightarrow T_c C \mid c$

»  $T_a \rightarrow a$

»  $T_c \rightarrow c$



# Remove Long

- » **Def.:** long is a rule whose right hand side consists of more than two nonterminal symbol.
- » **R:**  $A \rightarrow BCDE$
- » By remove long, it will be:
  - »  $A \rightarrow BM_2$
  - »  $M_2 \rightarrow CM_3$
  - »  $M_3 \rightarrow DE$
- » **Ex1(continue),  $G''$  :**
  - »  $S \rightarrow aACa \mid aAa \mid aCa \mid aa$
  - »  $A \rightarrow a \mid c \mid cC$
  - »  $B \rightarrow c \mid cC$
  - »  $C \rightarrow cC \mid c$



# Remove Long

» Ex1(continue),  $G'''$  :

»  $S \rightarrow T_a A C T_a \mid T_a A T_a \mid T_a C T_a \mid T_a T_a$

»  $A \rightarrow a \mid c \mid T_c C$

»  $B \rightarrow c \mid T_c C$

»  $C \rightarrow T_c C \mid c$

»  $T_a \rightarrow a$

»  $T_c \rightarrow c$



# Remove Long

» Now, by apply *removeLong*( $G''' : CFG$ ),  $G'^v$  :

»  $S \rightarrow T_a S_1 \mid T_a S_3 \mid T_a S_4 \mid T_a T_a$

»  $S_1 \rightarrow AS_2 \quad S_2 \rightarrow CT_a \quad S_3 \rightarrow AT_a \quad S_4 \rightarrow CT_a$

»  $A \rightarrow a \mid c \mid T_c C$

»  $B \rightarrow c \mid T_c C$

»  $C \rightarrow T_c C \mid c$

»  $T_a \rightarrow a$

»  $T_c \rightarrow c$



# Add Epsilon

» **Finally**, *atmostoneEps*( $G'^\nu$ :CFG) does not apply in Ex1,  $G^\nu = G'^\nu$  since  $L$  does not contain  $\varepsilon$  ( $S$  is not nullable).

»  $S \rightarrow T_a S_1 \mid T_a S_3 \mid T_a S_4 \mid T_a T_a$

»  $S_1 \rightarrow A S_2 \quad S_2 \rightarrow C T_a \quad S_3 \rightarrow A T_a \quad S_4 \rightarrow C T_a$

»  $A \rightarrow a \mid c \mid T_c C$

»  $B \rightarrow c \mid T_c C$

»  $C \rightarrow T_c C \mid c$

»  $T_a \rightarrow a$

»  $T_c \rightarrow c$



# Example2

» Convert the following CFG to CNF:

»  $S \rightarrow ABC$

»  $A \rightarrow aC \mid D$

»  $B \rightarrow bB \mid \varepsilon \mid A$

»  $C \rightarrow Ac \mid \varepsilon \mid Cc$

»  $D \rightarrow aa$





# 1-Remove epsilon

- »  $N = \{A, B, C\}$  So
- » add  $S \rightarrow AB \mid BC \mid AC, A \rightarrow a, B \rightarrow b, C \rightarrow c$
- » delete  $B \rightarrow \epsilon, C \rightarrow \epsilon$
- » The result is:
- »  $S \rightarrow ABC \mid AB \mid BC \mid AC$
- »  $A \rightarrow aC \mid D \mid a$
- »  $B \rightarrow bB \mid A \mid b$
- »  $C \rightarrow Ac \mid Cc \mid c$
- »  $D \rightarrow aa$



## 2- Remove units

»  $S \rightarrow ABC \mid AB \mid BC \mid AC$

»  $A \rightarrow aC \mid D \mid a$  ( $A \rightarrow D$ ,  $D \rightarrow aa$ ) remove  $A \rightarrow D$

add  $A \rightarrow aa$

»  $B \rightarrow bB \mid A \mid b$  ( $B \rightarrow A$ ,  $A \rightarrow aC \mid aa \mid a$ ) remove  $B \rightarrow A$

add  $B \rightarrow aC \mid aa \mid a$

»  $C \rightarrow Ac \mid Cc \mid c$

»  $D \rightarrow aa$



## 2- Remove units

»  $S \rightarrow ABC \mid AB \mid BC \mid AC$

»  $A \rightarrow aC \mid aa \mid a$

»  $B \rightarrow bB \mid aC \mid aa \mid a \mid b$

»  $C \rightarrow Ac \mid Cc \mid c$

»  $D \rightarrow aa$



### 3- Remove Mixed

$$\gg S \rightarrow ABC / AB / BC / AC$$

$$\gg A \rightarrow T_a C \mid T_a T_a \mid a$$

$$\gg B \rightarrow T_b B \mid T_a C / T_a T_a \mid a \mid b$$

$$\gg C \rightarrow AT_c \mid CT_c \mid c$$

$$\gg D \rightarrow T_a T_a$$



## 4- Remove Long

$$\gg S \rightarrow AM_2 / \textcolor{red}{M}_2 \mid AC / AB$$

$$\gg M_2 \rightarrow BC$$

$$\gg A \rightarrow T_a C \mid T_a T_a \mid a$$

$$\gg B \rightarrow T_b B \mid T_a C \mid T_a T_a \mid a \mid b$$

$$\gg C \rightarrow AT_c \mid CT_c \mid c$$

$$\gg D \rightarrow T_a T_a$$

$$\gg T_a \rightarrow a$$

$$\gg T_b \rightarrow b$$

$$\gg T_c \rightarrow c$$



## 5- Add Epsilon: N/A

»  $S \rightarrow AM_2 | \textcolor{red}{M_2} | \textcolor{red}{BC} | AC | AB$

»  $M_2 \rightarrow BC$

»  $A \rightarrow T_a C | T_a T_a | a$

»  $B \rightarrow T_b B | T_a C | T_a T_a | a | b$

»  $C \rightarrow AT_c | CT_c | c$

»  $D \rightarrow T_a T_a$

»  $T_a \rightarrow a$

»  $T_b \rightarrow b$

»  $T_c \rightarrow c$



# Example3

» Convert the following CFG to CNF:

»  $A \rightarrow BAB \mid B \mid \varepsilon$

»  $B \rightarrow 00 \mid \varepsilon$

» 1- remove epsilon:

»  $A \rightarrow BAB \mid B \mid BB \mid AB \mid BA$

»  $B \rightarrow 00$

» 2- remove units:

»  $A \rightarrow BAB \mid 00 \mid BB \mid AB \mid BA$

»  $B \rightarrow 00$



» 3- remove mixed

»  $A \rightarrow BC \mid T_0T_0 \mid BB \mid AB \mid BA$

»  $C \rightarrow AB, B \rightarrow T_0T_0, T_0 \rightarrow 0$

» 4- Add epsilon (Optional, Generally this step is applicable for root variable i.e. S)

»  $A \rightarrow BC \mid T_0T_0 \mid BB \mid AB \mid BA \mid \varepsilon$

»  $C \rightarrow AB$

»  $B \rightarrow T_0T_0$

»  $T_0 \rightarrow 0$





## Example4

» Convert the following CFG to CNF:

$$S \rightarrow ABa$$

$$A \rightarrow aab$$

$$B \rightarrow Ac$$



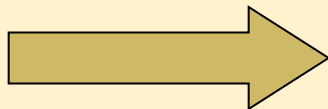
Introduce new variables for the terminals:

$$T_a, T_b, T_c$$

$$S \rightarrow ABa$$

$$A \rightarrow aab$$

$$B \rightarrow Ac$$



$$S \rightarrow ABT_a$$

$$A \rightarrow T_aT_aT_b$$

$$B \rightarrow AT_c$$

$$T_a \rightarrow a$$

$$T_b \rightarrow b$$

$$T_c \rightarrow c$$



Introduce new intermediate variable  $V_1$   
to break first production:

$$S \rightarrow ABT_a$$

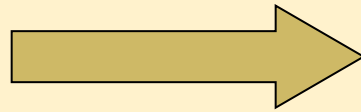
$$A \rightarrow T_a T_a T_b$$

$$B \rightarrow AT_c$$

$$T_a \rightarrow a$$

$$T_b \rightarrow b$$

$$T_c \rightarrow c$$



$$S \rightarrow AV_1$$

$$V_1 \rightarrow BT_a$$

$$A \rightarrow T_a T_a T_b$$

$$B \rightarrow AT_c$$

$$T_a \rightarrow a$$

$$T_b \rightarrow b$$

$$T_c \rightarrow c$$



Introduce intermediate variable:  $V_2$

$$S \rightarrow AV_1$$

$$V_1 \rightarrow BT_a$$

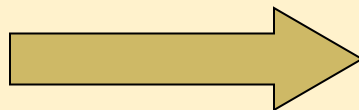
$$A \rightarrow T_a T_a T_b$$

$$B \rightarrow AT_c$$

$$T_a \rightarrow a$$

$$T_b \rightarrow b$$

$$T_c \rightarrow c$$



$$S \rightarrow AV_1$$

$$V_1 \rightarrow BT_a$$

$$A \rightarrow T_a V_2$$

$$V_2 \rightarrow T_a T_b$$

$$B \rightarrow AT_c$$

$$T_a \rightarrow a$$

$$T_b \rightarrow b$$

$$T_c \rightarrow c$$



# Final grammar in Chomsky Normal Form:

Initial grammar

$$S \rightarrow ABa$$

$$A \rightarrow aab$$

$$B \rightarrow Ac$$

$$S \rightarrow AV_1$$

$$V_1 \rightarrow BT_a$$

$$A \rightarrow T_a V_2$$

$$V_2 \rightarrow T_a T_b$$

$$B \rightarrow AT_c$$

$$T_a \rightarrow a$$

$$T_b \rightarrow b$$

$$T_c \rightarrow c$$



# Example5

» Convert the following CFG to CNF:

»  $E \rightarrow E + T$

»  $E \rightarrow T$

»  $T \rightarrow T * F$

»  $T \rightarrow F$

»  $F \rightarrow (E)$

»  $F \rightarrow \text{id}$



# 1- Remove units

» Remove  $E \rightarrow T$  , add  $E \rightarrow T * F \mid F$

» Remove  $E \rightarrow F$  , add  $E \rightarrow (E) \mid id$

» Remove  $T \rightarrow F$  , add  $T \rightarrow (E) \mid id$

» The result:

»  $E \rightarrow E + T \mid T * F \mid (E) \mid id$

»  $T \rightarrow T * F \mid (E) \mid id$

»  $F \rightarrow (E) \mid id$



## 2- Remove Mixed

»  $E \rightarrow E T_+ T \mid T T_* F \mid T_-(E T_+) \mid \text{id}$

»  $T \rightarrow T T_* F \mid T_-(E T_+) \mid \text{id}$

»  $F \rightarrow T_-(E T_+) \mid \text{id}$

»  $T_- \rightarrow ($

»  $T_+ \rightarrow )$

»  $T_+ \rightarrow +$

»  $T_* \rightarrow *$





### 3- Remove Long

»  $E \rightarrow E M_2 \mid T M_3 \mid T_(_ M_4 \mid \text{id}$

»  $T \rightarrow T M_3 \mid T_(_ M_4 \mid \text{id}$

»  $F \rightarrow T_(_ M_4 \mid \text{id}$

»  $M_2 \rightarrow T_+ T$

»  $M_3 \rightarrow T_* F$

»  $M_4 \rightarrow E T_)$

»  $T_(_ \rightarrow ($

»  $T_)_ \rightarrow )$

»  $T_+ \rightarrow +$

»  $T_* \rightarrow *$



## 5- Add Epsilon: N/A

- »  $E \rightarrow E M_2 \mid T M_3 \mid T_{(} M_4 \mid \text{id}$
- »  $T \rightarrow T M_3 \mid T_{(} M_4 \mid \text{id}$
- »  $F \rightarrow T_{(} M_4 \mid \text{id}$
- »  $M_2 \rightarrow T_{+} T$
- »  $M_3 \rightarrow T_{*} F$
- »  $M_4 \rightarrow E T_{)}$
- »  $T_{(} \rightarrow ($
- »  $T_{)} \rightarrow )$
- »  $T_{+} \rightarrow +$
- »  $T_{*} \rightarrow *$



» Thank you for various WEB Resources!!

