



UNIVERSITY INSTITUTE OF ENGINEERING

Advanced Database Management System Experiment 2.1

23CSP-333

Submitted To:

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Consider the following questions and answer accordingly.

1. Consider a relation R having attributes as R(ABCD), functional dependencies are given below:

AB->C, C->D, D->A

Identify the set of candidate keys possible in relation R. List all the set of prime and non prime attributes.

2. Relation R(ABCDE) having functional dependencies as:

A->D, B->A, BC->D, AC->BE

Identify the set of candidate keys possible in relation R. List all the set of prime and non prime attributes.

3. Consider a relation R having attributes as R(ABCDE), functional dependencies are given below:

B->A, A->C, BC->D, AC->BE

Identify the set of candidate keys possible in relation R. List all the set of prime and non prime attributes.

4. Consider a relation R having attributes as R(ABCDEF), functional dependencies are given below:

A->BCD, BC->DE, B->D, D->A

Identify the set of candidate keys possible in relation R. List all the set of prime and non prime attributes.

- **5.** Consider a relation schema R(W, X, Y, Z) with the following functional dependencies:
 - 1. $X \rightarrow Y$
 - 2. $WZ \rightarrow X$
 - 3. $WZ \rightarrow Y$
 - 4. $Y \rightarrow W$
 - 5. $Y \rightarrow X$
 - 6. $Y \rightarrow Z$

Tasks:

- 1. Identify all the candidate keys of R.
- 2. List the prime and non-prime attributes.
- 3. Determine the highest normal form of the relation R with proper justification.
- **6.** Consider a relation schema R(A, B, C, D, E, F) with the following functional dependencies:

$$A \to BC, D \to E, BC \to D, A \to D$$

Tasks:

- 1. Find all the candidate keys of R.
- 2. List the prime and non-prime attributes.
- 3. Determine the highest normal form of relation R with proper justification.



Answers of the above given questions are as follows:

1. Given: R(A B C D) with FDs: $AB \rightarrow C$, $C \rightarrow D$, $D \rightarrow A$.

Candidate keys: AB, BC, BD.

(Checks: $AB^+ = ABCD$; $BC^+ = B$, $C \rightarrow AD \rightarrow ABCD$; $BD^+ = B$, $D \rightarrow A \rightarrow ABCD$)

 $AB \rightarrow C \rightarrow ABCD$. All are minimal.)

Prime attributes: A, B, C, D (every attribute appears in some candidate

key).

Non-prime attributes: none.

Highest normal form: 3NF (all FDs either have a superkey on the left — e.g. AB→C — or have a prime attribute on the right; violates BCNF

because $C \rightarrow D$ (and $D \rightarrow A$) have non-superkey determinants).

2. Given: R(A B C D E) with FDs:

 $A \rightarrow D, B \rightarrow A, BC \rightarrow D, AC \rightarrow BE.$

Candidate keys: AC, BC.

(Checks: $AC^+ = \{A,C\} \rightarrow BE$ (by $AC \rightarrow BE$) and $A \rightarrow D \Rightarrow \{A,B,C,D,E\}$.

 $`BC^+ = \{B,C\} \rightarrow A \text{ (by } B \rightarrow A) \rightarrow D \text{ (by } A \rightarrow D) \text{ and } AC \rightarrow BE \text{ gives } E \Rightarrow all$

attributes.)

Prime attributes: A, B, C.

Non-prime attributes: D, E.

Highest normal form: 1NF.

Reason: A \rightarrow D is a partial dependency (A is a proper subset of the

candidate key AC and D is non-prime), so the relation violates 2NF (hence also not in 3NF/BCNF).

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3. Given: R(A B C D E) with FDs:

 $B \rightarrow A, A \rightarrow C, BC \rightarrow D, AC \rightarrow BE.$

Candidate keys: A, B.

(Checks: $A^+ = A \rightarrow C$; $AC \rightarrow BE \Rightarrow B,E$; $BC \rightarrow D \Rightarrow D$ so $A^+ = ABCDE$.

 $B^{\scriptscriptstyle +} = B \to A; \, A \to C \Rightarrow C; \, AC \to BE \Rightarrow E; \, BC \to D \Rightarrow D \text{ so } B^{\scriptscriptstyle +} =$

ABCDE.)

Prime attributes: A, B.

Non-prime attributes: C, D, E.

Highest normal form: BCNF (every FD has a superkey as determinant: A

and B are keys, and BC, AC are supersets of keys).

4. Given: R(A B C D E F) with FDs:

 $A \rightarrow B C D, BC \rightarrow D E, B \rightarrow D, D \rightarrow A.$

Candidate keys: AF, BF, DF.

(Reason: $A^+ = \{A,B,C,D,E\}$ so $AF^+ =$ all attributes; similarly B^+ and D^+

each give {A,B,C,D,E}, so adding F yields the whole relation. F must be

included because no FD produces F.)

Prime attributes: A, B, D, F.

Non-prime attributes: C, E.

Highest normal form: 1NF.

(Why: e.g. $A \rightarrow C$ is a partial dependency — A is a proper subset of the candidate key AF and determines non-prime C — so 2NF is violated; hence relation is not in 2NF/3NF/BCNF.)

5. Given: R(W X Y Z) with FDs:

 $X \rightarrow Y, WZ \rightarrow X, WZ \rightarrow Y, Y \rightarrow W, Y \rightarrow X, Y \rightarrow Z.$

Candidate kevs: X, Y, WZ.

(Checks: $X^+ = X \rightarrow Y \rightarrow \{W,Z\}$ so $X^+ = \{W,X,Y,Z\}$.

 $Y^{+} = Y \rightarrow W, X, Z \text{ so } Y^{+} = \{W, X, Y, Z\}.$

minimal.)

Prime attributes: W, X, Y, Z (every attribute appears in some candidate key).

Non-prime attributes: none.

Highest normal form: BCNF — every FD's left side is a superkey (X, Y,

and WZ are all keys).

6. Given: R(A B C D E F) with FDs:

 $A \rightarrow BC, D \rightarrow E, BC \rightarrow D, A \rightarrow D.$

Candidate key(s): AF.

(Reason: $A^+ = \{A \rightarrow BC \rightarrow D \rightarrow E\} = \{A,B,C,D,E\}$, so adding F gives $AF^+ = \{A,B,C,D,E,F\}$. F is not produced by any FD, so every key must include F; A is required to reach the other attributes, so AF is the minimal key.)

Prime attributes: A, F.

Non-prime attributes: B, C, D, E.

Highest normal form: 1NF.

(Why: A \rightarrow BC is a partial dependency because A is a proper subset of the candidate key AF and determines non-prime attributes B and C, so the

relation violates 2NF — hence it is only in 1NF.)