

# Quiz 1

The following shows the implementation of the Euclidean Algorithm  
for (36, 21)

$$1) \quad 36 = 1 \cdot 21 + 15$$

$$2) \quad 21 = 1 \cdot 15 + 6$$

$$3) \quad 15 = 2 \cdot 6 + 3$$

$$4) \quad 6 = 2 \cdot 3 + \underline{0} \text{ Halt!}$$

Question: Using above to find an integer solution of

$$36x + 21y = 3$$

## Solution to Quiz 1 :

$$1) \quad 36 = 1 \cdot 21 + 15$$

$$2) \quad 21 = 1 \cdot 15 + 6$$

$$3) \quad 15 = 2 \cdot 6 + 3$$

$$4) \quad 6 = 2 \cdot 3 + \underline{0} \text{ Halt!}$$

$$3 \stackrel{(3)}{=} 15 - 2 \cdot 6$$

$$\stackrel{(2)}{=} 15 - 2 \cdot (21 - 1 \cdot 15)$$

$$= -2 \cdot 21 + 3 \cdot 15$$

$$\stackrel{(1)}{=} -2 \cdot 21 + 3 \cdot (36 - 1 \cdot 21)$$

$$= 3 \cdot 36 - 5 \cdot 21$$

This gives an integer solution

$$\begin{cases} x = 3 \\ y = -5 \end{cases}$$

## Situation

$$ax + by = R,$$

where  $a, b$  are positive integers, and

$R$  is the Last non-zero remainder  
in the Euclidean Algorithm for  $(a, b)$ .

We are able to find an integer solution under this situation.

Terminology: An  $S$ -linear combination of  $a$  and  $b$  is an expression

$$s \cdot a + t \cdot b \quad (s, t \in S)$$

We say  $R$  can be written as an  $S$ -linear combination of  $a$  and  $b$   
if there are  $s, t \in S$  such that  $s \cdot a + t \cdot b = R$ .

Proof. Carry out the Euclidean Algorithm as follows:

$$\begin{array}{ll} 0) & a = q_0 b + r_1 \quad \Rightarrow \quad r_1 \text{ is a } \mathbb{Z}\text{-linear combination of } a \text{ \& } b \\ 1) & b = q_1 r_1 + r_2 \quad \Rightarrow \quad r_2 \text{ is a } \mathbb{Z}\text{-linear combination of } b \text{ \& } r_1 \\ & \dots \quad \dots \quad \dots \\ t) & r_{t-1} = q_t r_t + R \quad \Rightarrow \quad R \text{ is a } \mathbb{Z}\text{-linear combination of } r_{t-1} \text{ \& } r_t \\ & r_t = q_{t+1} R + \underline{0}. \end{array}$$

Applying the following lemma to  $\left. \begin{array}{l} r_1 \\ r_2 \\ \vdots \\ R \end{array} \right\}$ , then we prove that

$R$  is a  $\mathbb{Z}$ -linear combination of  $a$  &  $b$  by induction.  $\square$

Lemma: If  $\alpha$  is a  $\mathbb{Z}$ -linear combination of  $a$  &  $b$ ,  
 $\beta$  is a  $\mathbb{Z}$ -linear combination of  $b$  &  $\alpha$ ,  
then  $\beta$  is a  $\mathbb{Z}$ -linear combination of  $a$  &  $b$ .

Summarize so far:

We can use Euclidean Algorithm to find an integer solution of

$$ax + by = R,$$

where  $a, b$  are positive integers, and

$R$  is the last non-zero remainder  
in the Euclidean Algorithm for  $(a, b)$ .

Extensions:

1) Negative  $a$  or  $b$

E.g.  $a > 0 \Rightarrow b < 0$

If  $x_0$  &  $y_0$  form an integer solution of  $|a|x + |b|y = R$ ,

then  $x_0$  &  $-y_0$  form an integer solution of  $ax + by = R$ .

2) Replace  $R$  by a multiple  $C$  of it.

E.g.  $C = m \cdot R$

If  $x_0$  &  $y_0$  form an integer solution of  $ax + by = R$ ,

then  $m \cdot x_0$  &  $m \cdot y_0$  form an integer solution of  $ax + by = C$ .

Ex: What if one of  $a$  &  $b$  is zero?

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### General Questions

$$ax + by = c$$

Q1: Is there any INTEGER solution?  $\checkmark_p$

Q2: If there is, find ONE such a solution.  $\checkmark_p$

Q3: Find ALL integer solutions.

Def. Let  $a$  and  $b$  be two integers.

The **greatest common divisor** of  $a$  and  $b$  is a natural number  $g \in \mathbb{N}$  satisfying the following properties:

- i)  $g$  is a common divisor of  $a$  and  $b$ , i.e.  $g|a$  &  $g|b$
- ii) If  $d$  is a common divisor of  $a$  and  $b$ , then  $d|g$

Notation:  $\text{GCD}(a, b)$ .

Rmk The properties i) & ii) together are called the **defining property** or the **universal property** of the notion "the greatest common divisor of  $a$  and  $b$ ".

## Prop (uniqueness of GCD)

There is at most **ONE** natural number  $g \in \mathbb{N}$  satisfying i) & ii).

Proof: Suppose  $g_1$  &  $g_2$  are GCD of  $a$  and  $b$ .

By i), we have  $g_1 \mid a$ ,  $g_1 \mid b$ ,  $g_2 \mid a$ ,  $g_2 \mid b$ .

By ii), we have  $g_1 \mid g_2$  and  $g_2 \mid g_1$ .

By ~~reflexive~~ property of  $\mid$ ,  $g_1 = g_2$ .  
*antisymmetric*





Prop : Let  $a$  &  $b$  be positive integers, say  $a \geq b$ . Carry out the Euclidean Algorithm for  $(a, b)$ . Then the last non-zero remainder is  $\text{GCD}(a, b)$ .

Pf: 0)  $a = q_0 \cdot b + r_1$   
 1)  $b = q_1 \cdot r_1 + r_2$   
 ...  
 t)  $r_{t-1} = q_t \cdot r_t + R$   
 $r_t = q_{t+1} \cdot R + \underline{0}$ .

To show  $R$  satisfies i) & ii)  
 i)  $R \mid a, R \mid b$   
 $R \mid r_t$   
 By t) & 2-out-of-3,  
 $R \mid r_{t-1}$

In general,

if  $R \mid r_i, r_{i+1}$ , then by  $r_{i-1} = q_i \cdot r_i + r_{i+1}$  & 2-out-of-3  
 we have  $R \mid r_{i-1}$

ii) If  $d|a$  &  $d|b$ , then  $d|R$ .

By  $a = q_0 \cdot b + r_1$  & 2-out-of-3,  $d|r_1$ .

In general,

if  $R|r_{i-1}, r_i$ , then by  $r_{i-1} = q_i \cdot r_i + r_{i+1}$  & 2-out-of-3  
we have  $d|r_{i+1}$

In particular,  $d|R$ .



By the uniqueness.  $\text{GCD}(a, b) = R$ .

EX: How to compute  $\text{GCD}(a, b)$   
when  $a, b \in \mathbb{Z}$ .

## Theorem.

Let  $a$ ,  $b$  &  $c$  be integers. The equation  $ax + by = c$  has an integer solution iff  $c$  is a multiple of  $\text{GCD}(a, b)$

Pf: If case follows from Euclidean Algorithm.

Only If case:

Assume  ~~$\text{GCD}(a, b) \nmid c$~~  and  $(x_0, y_0)$  is an integer solution of  $ax + by = c$ .

Then  $ax_0 + by_0 = c$ .

Since  $\text{GCD}(a, b) \mid a, b$ , by 2-out-of-3,

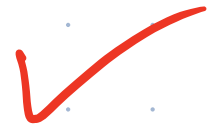
$$\text{GCD}(a, b) \mid c. \quad \text{---}$$



## General Questions

$$ax + by = c$$

Q1: Is there any INTEGER solution?



Q2: If there is, find ONE such a solution.



Q3: Find ALL integer solutions.

Some after-class reading suggestions:

- Today's topic are the GCD and the solvability of the linear Diophantine equation  $ax + by = c$ . Refer pp. 30–33 in the textbook.
- I encourage you to read the rest of Chapter 1 preparing for our next meeting.