USING RANDOMNESS IN COMPUTER SCIENCE

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[Lec.1] (1) Introduction to randomized algorithms (definition, cost,...) (2) Simpler (and faster) than deterministic: A randomized algorithm for Min-CUT (a) A simple $O(n \log n)$ -time randomized algorithm that succeeds with probability $\Theta(1/n^2)$ (b) A recursive $O(n^2)$ -time randomized algorithm that succeeds with probability $\Theta(1/\log n)$. (c) Increasing the probability of success to $1 - 1/n^c$ for any constant c > 0. (3) Less ressources than deterministic: A $\Theta(\log \log n)$ -bits randomized counter that counts up to n[Lec.2](4) Proving optimality of randomized algorithms: Yao's principle (a) Faster than deterministic: Randomized evaluation of OR-AND complete boolean (b) Yao's principle (c) A lower bound on randomized algorithm for AND-OR trees based on Yao's principle (d) Hard drive energy consumption optimization (i) Optimal deterministic algorithm [Lec.3](ii) Optimal randomized algorithm via Yao's principle (5) Randomized rounding (a) Approximation algorithms for Max-SAT (i) The "random solution" randomized algorithm (ii) A randomized algorithm based on linear programming (iii) Combining both (iv) Derandomizing these algorithms by the conditional expectation method (b) Approximation algorithms for Max-CUT (i) The "random solution" algorithm (ii) Derandomizing this algorithm by the conditional expectation method yields an analysis of the simple greedy algorithm (6) Sample and Guess [Lec.4](a) A polynomial-time randomized algorithm for testing polynomial identity (Zero-P) (b) Self-correction of integer multiplication (c) The class PCP(r,q); $NP \subseteq PCP(\log n, 1) \Leftrightarrow (\exists c)$ GAP-SAT_{1,c} is NP-hard (i.e. hardness of approximation result for Max-SAT) (d) Linearity testing and self-correction (e) (Application) One of the main step of the PCP theorem: $NP \subseteq PCP(poly(n), 1)$ [Lec.5] (i) QUADEQ is NP-complete (ii) A PCP $(n^2, 1)$ -verifier for QUADEQ (f) A constant-time approximation scheme for the size of a maximal matching in a [Lec.6] constant degree graph (7) Exhaustive guessing 1

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(8) Random Walks 1 (a) Walk-SAT

(a) A polynomial-time randomized $(1-\varepsilon)$ -approximation for Max-CUT in dense graph