Ans: 1.(1) Here, N=4, M=2

Given; Atteast one muddy child. The children don't know M's value.

Let d denote the set of children with mud [#d=2] c denote " " " " without mud.[#c=2].

ROUND 1 :-

All of d and c say "No" since they don't know if they are muddy. This is because, none of them knows the mo. of muddy foreheads.

- Inference after tround 1: # muddy forcheads = >1. [This is because, if there was exactly one muddy forchead, the child with muddy forchead would see no other muddy forchead and could answer "Yes". ROUND 2:could answer "yes".
  - · The children in a set answer "Yes" [they know if they are muddy]
  - · The children in a answer "No" [ they don't yet know if they were muddy

Reason: The children in a know that there are > 1 muddy foreheads and each of the & children in a can see only other child in d's muddy forehead.

The children in c can see two muddy forcheads (of those in a) and cannot conclude anything, about their own.

ROUND 3:-

- · Inference from Round 2: # muddy forcheads = 2. [since, both the children in of know they are muddy].
- each of them know whether each of them have mud on their own forehead · The children in d answer "Yes" · The children in c answer "Yes"

Thus, after Round 3, all Mesponses converge.

Ans:1(2) N=4, M>0, M= even\_

Notations and symbols: same as 1.(1)

ROUND 1 :-

All of d and c say No since mone of H

M=2 :-

Same as Ans:1(1).

M=4] :-

Round 1:-

All of d and c say "No" since all see 3 muddy forcheads.

Kound 2:

- · Inference from Round 1: # muddy foreheads > 1 [ same as 1.(1)].
- All of d and c say "No" since they are not yet able to figure out if they are muddy because all of them see 3 muddy foreheads and the total no.

  3 or 4.

Round 3:

- · Inference from Round 2: # muddy foresheads > 2 because o.w children in d
- · Still none of d and c can consider since they we not properly able to figure out if they themselves have mud on their foreheads. All say "No"

Round 4:

- · Inference from Round 3: # muddy foreheads > 3.
- · All of d and a answer "yes" since they are now swee that all of them have mud on their forcheads. .: Responses converge after Round 4.

An: 2.

Our Rostful web service will mainly do the following two tasks:—

create a new registreation for an user who had completed one dose of vaccination or has not taken any dose of vaccination

· return the vaccine information for the users who have administered a single dose of vaccine [information about the type/vial].

## Two resources:-

- · Home Person
- · Conta Vaccine

Design of the RESTful API:-

- Object Model:—
   Person, Vaccine
- · Creating Model URIS:

-/person/frame}

-/person/grames//peacine

- / person/{p-name}/vaccine/{vaccine-name}

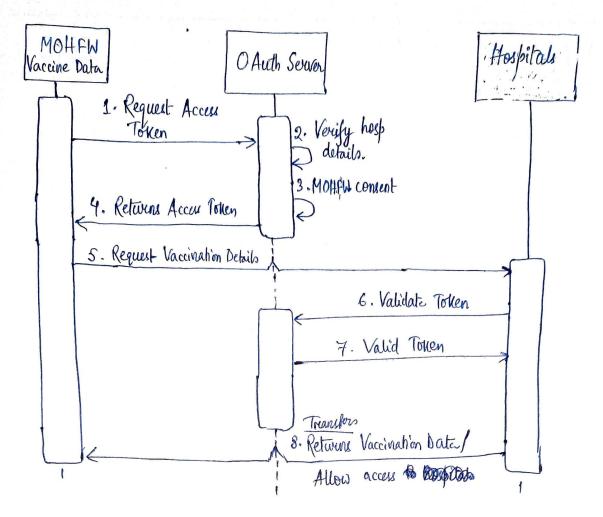
· Determining representations:

- Represent all person's information as an XML/JSON
- Represent all vaccine names as an XML/JSON.

· Assigning HTTP Methods:

- cheate registration for users who are new or completed one dose -> HTTP
- Melturn information about type of vaccine taken -> HTTP GET

Here, 'POST' is non-idempotent-'GET' is idempotent. Ans: 3.



When a particular hospital needs to transfer the resources on vaccination data, through a secure method, the DAuth Server ares the hospitals to login for a secure transaction/transfer of data and after verifying the hospital details, when the thoughther, requests access to the vaccine data, ithe hospitals validates the tonen and returns the data and allows access to the data for the data. This has been represented using a sequence diagram as above.

## Ans: 4. Class Diagnam for MOHFW scenario:

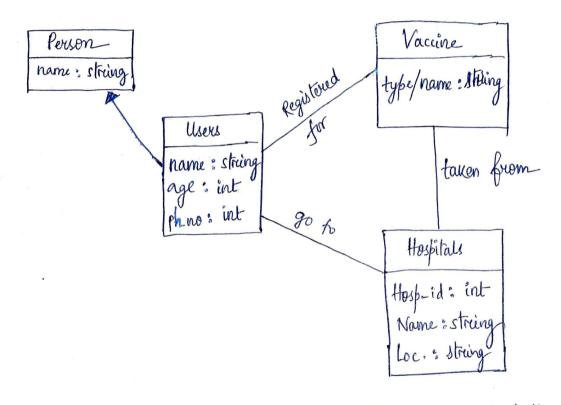


Fig: Class Diagnam to demonstrate vaccination situation

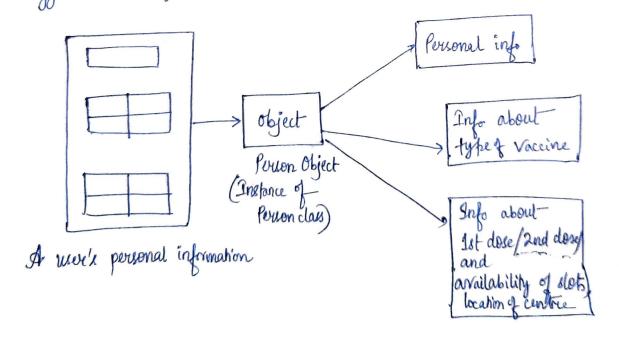
Here, the classes are:

- · Useus
- · Vaccine
- · Hospitals
- · Users Person is the "IS-A" Relationship.

acu los

ORM can be used in the following ways in this situation:

- From transfer of vaccine data from hospitals to MOFHW, the hospitals could simply provide its details and the data to MOFHW and the ORM abstracts details like verification through access towers, generating requests and realidating them etc.
- From weres' registration at the MOFHW portal for vaccination, the user's just enter credentials or generate request for & slot details, vaccine availability and the system provides the details to the user and abstracts all the intermediate steps like separating the user credentials to separate tables, fetching data related to the logged in user from the server etc. from the user.



Ans: 7.

(1) Here, we want to choose two processes as co-ordinators from a group of n independent processes.

We assume that every active process in the system has a priority no associated and the processes can send messages to all other processes in the system.

We want to elect the top two processes with the highest priority nos as the co-ordinators.

If any host thinks that the co-ordinator has failed, it tries to elect itself by sending a message to the highest numbered processors. If any of them answers, the host loses and each of the processors will can election and try to win thomselves.

If none of the highest priority no: Co-ordinators answer, then the host becomes the co-ordinator.

If a new processor arrives, it again calls for an election.

This continues two times

After getting the 1st co-ordinator, the above process repeats itself excluding the new co-ordinator to elect the next co-ordinator.

(2) Similarly, for me electing in priocess, continue the procedure in times keeping in mind whether new processes arrive and in that case, recall the election.

Ans: 6. Let us consider a vector clock consisting n elements  $V = (v_1, ..., v_n)$  which can be encoded by distinct prime nos.  $p_1, p_2, ..., p_n$ . A method could be to encode Here, actually, we have encoded the entire vector V to a unique number using, p1, p2 ..., pn. The encoded teptesentation of  $V = p_1^{v_1} \cdot p_2^{v_2} \cdots p_n^{v_n}$ This works well because the above product is distinct for any permutation of {v1, v2, ..., vn/ a single no. neduces the space needed to This encoding of n-entries to

represent vector clocks.