

# # Sorting #

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- It is a process of arranging items systematically.

## \* Bubble Sort \*

- It is the simplest Algorithm that works by repeatedly swapping the adjacent element if they are in wrong order.

Example:

First Pass: 3, 1, 5, 4, 2

1, 3, 5, 4, 2

1, 3, 4, 5, 2

1, 3, 4, 2, 5

1, 3, 4, 2, 5

with first pass through the entire array, the largest element (here 5) came to the end

Second Pass: 1, 3, 4, 2, 5

1, 3, 2, 4, 5

1, 3, 2, 4, 5

with 2<sup>nd</sup> Pass 2<sup>nd</sup> largest element comes at second from the last index.

Third Pass: 1, 3, 2, 4, 5

1, 2, 3, 4, 5

1, 2, 3, 4, 5

sorted!!

don't need to compare again & again since it is already sorted.

Bubble sort is also known as sinking sort or Exchange sort.

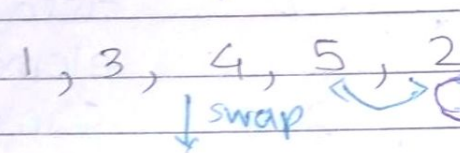
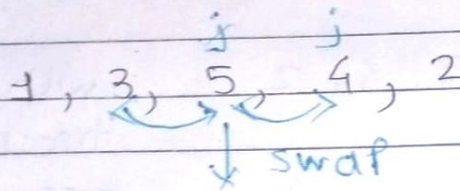
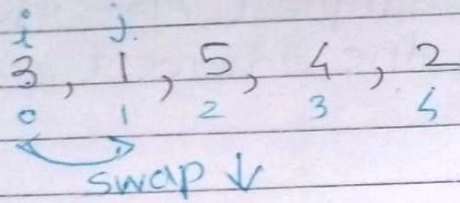


let's understand more deeply how Actually bubble sort works!!

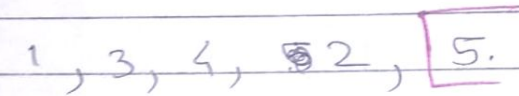
Counter

$i = 0$

1<sup>st</sup> pass

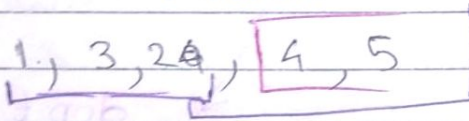
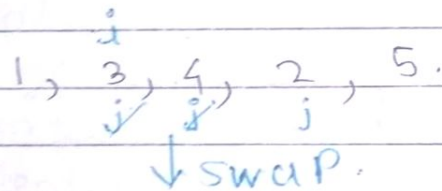


internal loop  
it runs  $(n-1)$  times.



$i = 1$

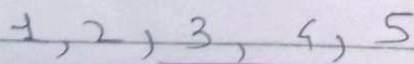
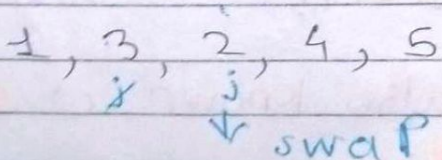
2<sup>nd</sup> pass



$j$  will only check this elements because after this already sorted

$i = 2$

3<sup>rd</sup> pass





## • Time Complexity

• Space Complexity :  $O(1)$  // constant.

⇒ since here no extra space is required.  
i.e. like copying the array etc. is not required.

Also known as inplace sorting algorithm

• Time Complexity :

1) Best case → Array is sorted.

↳  $O(N)$

$j = 0$   
1<sup>st</sup> pass

1, 2, 3, 4, 5  
i j j j

→ Only once it ran we don't need to check it again.

Note :- when  $j$  never swaps for value of  $i$ , it means array is sorted. Hence you can end the program.

★ Best Case Comparisons =  $N - 1 \Rightarrow \textcircled{N}$

Since in time complexity constants are ignored, we don't want exact time, we just want relationship i.e. mathematical function

2) worst case → sorting descending order to ascending

$j = 0$

1<sup>st</sup> pass

5, 4, 3, 2, 1  
j ↓

4, 5, 3, 2, 1  
j ↓

4, 3, 5, 2, 1  
j ↓

4, 3, 2, 5, 1  
j ↓

4, 3, 2, 1, 5

⇒  $(N - 1)$  swaps.



$j = 1$   
2<sup>nd</sup> pass

4, 3, 2, 1, 5  
↓  
j

3, 4, 2, 1, 5  
↓  
j

3, 2, 4, 1, 5  
↓  
j

3, 2, 1, 4, 5

$\Rightarrow (N-2)$  swaps

$j = 2$   
3<sup>rd</sup> pass

3, 2, 1, 4, 5  
↓  
j

2, 3, 1, 4, 5  
↓  
j

2, 1, 3, 4, 5

$\Rightarrow (N-3)$  swaps

$j = 3$   
4<sup>th</sup> pass

2, 1, 3, 4, 5  
↓  
j

1, 2, 3, 4, 5

$\Rightarrow (N-4)$  swaps

$$\text{total comparisons} = N-1 + N-2 + N-3 + N-4$$

$$= 4N - (1+2+3+4)$$

$$= 4N - \left[ \frac{N \times (N+1)}{2} \right]$$

$$= 4N - \frac{N^2 + N}{2}$$

$$= O\left(\frac{7N - N^2}{2}\right)$$

$$\text{Total comparisons} = O(N^2)$$

In time complexity constant & less dominating terms are ignored.



stable :  
Sorting  
Algo

(10) (20) (20) (30) (10)

↓ after sorting

(10) (10) (20) (20) (30)

⇒ Here after sorting the order is maintained  
⇒ order is same when value is same:

unstable :  
Sorting  
Algo

(10) (20) (20) (30) (10)

↓ after sorting

(10) (10) (20) (20) (30)

Ex:-

```
static void bubbleSort(int[] arr) {
    boolean swapped;
    for (i = 0; i < arr.length; i++) {
        swapped = false;
        for (j = 1; j < arr.length - i; j++) {
            if (arr[j] < arr[j-1]) {
                int temp = arr[j];
                arr[j] = arr[j-1];
                arr[j-1] = temp;
                swapped = true;
            }
        }
        if (!swapped) {
            break;
        }
    }
}
```