```
-- Translating EBNF to ParserLib
data Expr
    = Num Integer
    | Bln Bool
    | Var String
    | Prim2 Op2 Expr Expr
    -- Let [(name, rhs), ...] body
    | Let [(String, Expr)] Expr
    | Lambda String Expr -- Lambda var body
    | App Expr Expr -- App func arg
    -- Cond test then-branch else-branch
    | Cond Expr Expr Expr
    deriving (Eq, Show)
data Op2 = Eq | Lt | Plus | Minus | Mul |
   Div | Mod
    deriving (Eq, Show)
-- Z : : = X | Y
7. = X < I > Y
-- block ::= "[" X "]"
block = between (char '[' *> whitespaces)
                (char ']' *> whitespaces) X
-- cmp ::= "==" | "<"
-- stringToOp is MY helper that maps "=="
   -> Eq, "<" to Lt, etc.
cmp = do
    op <- operator "==" <|> operator "<"
    return . Prim2 $ stringToOp op
-- infix ::= arith [ cmp arith ]
-- Since [...] means "zero or one", rather
   than "zero or more", we actually rewrite
    this as 'arith | arith cmp arith'
infix = (do)
            1 <- arith
            c <- cmp
            r <- arith
            return $ c l r) <|> arith
-- atom ::= "(" block ")" | literal | var
atom = literal <|> (fmap Var var) <|>
   between
        (char '(' *> whitespaces)
        (char ')' *> whitespaces)
        block
```

```
-- cond ::= "if" block "then" block "else"
   block
cond = do
   keyword "if"
    a <- block
    keyword "then"
    b <- block
    kevword "else"
    c <- block
    return $ Cond a b c
-- equation ::= var "=" block ";"
equation = do
    v <- var
    char '=' *> whitespaces
    e <- block
    char ';' *> whitespaces
    return (v, e)
-- literal ::= integer | boolean
literal = int <|> bool
-- boolean ::= "True" | "False"
bool = do
   b <- keyword "True" <|> keyword "False"
    return . Bln $ if b == "True" then True
        else False
int = fmap Num integer
-- identifier: any alphabetical string
   EXCLUDING the ones apassed in the list
var = identifier ["if", "then", "else", "
   let", "in", "True", "False"]
-- Lab 11, parses:
        S ::= identifier | "(" S { S } ")"
data SExpr = Ident String | List [SExpr]
   deriving (Eq, Show)
sexpr :: Parser SExpr
sexpr = fmap Ident (identifier []) <|>
            between
                (char '(' *> whitespaces)
                (char ')' *> whitespaces)
                (do
                    s <- some sexpr
                    return (List s))
```

```
mainParser :: Parser SExpr
mainParser = whitespaces *> sexpr <* eof</pre>
-- Miscellaneous interpreter things
data Value = VN Integer
          I VB Bool
           | VClosure (Map String Value)
              String Expr
   deriving (Eq, Show)
data Error = VarNotFound | TypeError |
   DivByZero deriving (Eq, Show)
-- Core interpreter logic.
interp :: Map String Value -> Expr ->
   Either Error Value
-- Base cases, unwrap value.
interp _ (Num i) = pure $ VN i
interp _ (Bln b) = pure $ VB b
-- Evaluate an if condition, then evaluate
   the branch taken.
interp env (Cond a b c) = do
   cond <- interp env a
   assertBool cond
   interp env $ case cond of
        VB True -> b
        VB False -> c
-- Recursively evaluate all bindings in a
   let block.
interp env (Let [] blk) = interp env blk
interp env (Let ((name, expr):xs) blk) = do
   v <- interp env expr
   interp (Map.insert name v env) (Let xs
       blk)
-- Apply an argument to lambda.
interp env (App f e) = do
   lambda <- interp env f
    case lambda of
        VClosure closure v body -> do
            arg <- interp env e
           interp (Map.insert v arg
               closure) body
        _ -> Left TypeError
```

```
-- Lab 12, mutual recursion
     let f x = expr1
           g x = expr2
          h x = expr3
       in body
  is represented by
       LetRecFuns [ ("f", "x", expr1),
                    ("g", "x", expr2),
                    ("h", "x", expr3) ]
-}
interp env (LetRecFuns defs body) = interp
   (buildEnv env defs) body
  where
    buildEnv env [] = env
    buildEnv env ((fName, var, expr):defs)
       = f E n v
      where
        fEnv = buildEnv (Map.insert fName f
            env) defs
-- The trick is here, for each function
   define two closures that takes the other
    one in; this abuses the fact that
   Haskell *does* have forward references.
   We need to do this recursively, though.
        f = VClosure (buildEnv (Map.insert
           fName f' env) defs) var expr
        f' = VClosure (buildEnv (Map.insert
            fName f env) defs) var expr
-- EBNF refresher
-- either X or Y
XIY
-- val, followed by ZERO or MORE 'op val'
   pairs. Use chain11 in ParserLib for left
   -associative, and chainr1 for right-
   associative.
val { op val }
-- same as above, except ONE or NONE
val [ op val ]
-- terminal string
"if", e.g. "if" block "else" block "then"
```

block

```
-- Ambiguous grammars. Consider:
<stmt> ::= <cond> | "A" | "B" | "C"
<cond> ::= "if" <test> "then" <stmt>
                 | "if" <test> "then" <stmt>
                     "else" <stmt>
<test> ::= "T1" | "T2"
For the input "if T1 then if T2 then A else
    B", two interpretations possible:
1. if T1 then {if T2 then {A} else {B}}
2. if T1 then {if T2 then {A}} else {B}
In general case, CFG ambiguity is
   undecidable!
Parse tree (NOT ABSTRACT SYNTAX TREE) for
   #1 shown below.
                  stmt
                  cond
                  then
                       stmt
            test
            T1
                        cond
            test
                  then
                       stmt
                              else
                                   stmt
           T2
                                  В
For AST, get rid of all the useless
   terminals "if" "then" from the parse
   tree. AST is the form that ParserLib
   generates.
-- Monad stuff, Lab 9
postInc :: State Int Int
postInc = StateOf (i \rightarrow (i+1, i))
numberHelper :: BT v -> State Int (BT Int)
numberHelper Null = return Null
numberHelper (Node left _ right) = do
 left <- numberHelper left</pre>
```

curr <- postInc</pre>

right <- numberHelper right

return (Node left curr right)

```
-- Monad/Applicative/Functor shit from A2
instance Functor FreeGameMaster where
-- fmap :: (a -> b) -> FreeGameMaster a ->
   FreeGameMaster b
-- or fmap = liftM
   fmap f (Pure a) = Pure (f a)
   fmap f (GMAction lo hi fx) = GMAction
       lo hi (\msg -> fmap f (fx msg))
instance Applicative FreeGameMaster where
-- pure :: a -> FreeGameMaster a
   pure = Pure
-- (<*>) :: FreeGameMaster (a -> b) ->
   FreeGameMaster a -> FreeGameMaster b
-- or (<*>) = ap
   (Pure f) <*> rhs
                             = fmap f rhs
    (GMAction lo hi f) <*> rhs = GMAction
       lo hi (\msg -> (f msg) <*> rhs)
instance Monad FreeGameMaster where
-- return :: a -> FreeGameMaster a
   return = Pure
-- (>>=) :: FreeGameMaster a -> (a ->
   FreeGameMaster b) -> (FreeGameMaster b)
   (Pure a) >>= f
    (GMAction lo hi fx) >>= f = GMAction lo
        hi (\mbox{m} \rightarrow (fx \mbox{m}) >>= f)
-- Mutable memory monad
newtype MutM a = MkMutM (IntMap Value ->
   Either String (IntMap Value, a))
runMutM (MkMutM f) = f
instance Monad MutM where
   return a = MkMutM (\s -> Right (s, a))
   MkMutM f >>= k = MkMutM (\s -> case f s
        of Left msg -> Left msg
        Right (s', a) -> runMutM (k a) s')
getMem = MkMutM (\s -> Right (s, s))
modMem f = MkMutM (\s -> Right (f s, ()))
new val = do
   n <- fmap IntMap.size getMem
   store n val
   pure n
store a val = modMem (IntMap.insert a val)
```