

CS 3640: Introduction to Networks and Their Applications

Fall 2023, Lecture 2: The end-to-end principle

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You should have...

- Started working on Assignment 1 (due 8/31)

Today's class

1.

Recap

2.

Building a
picture of
the Internet

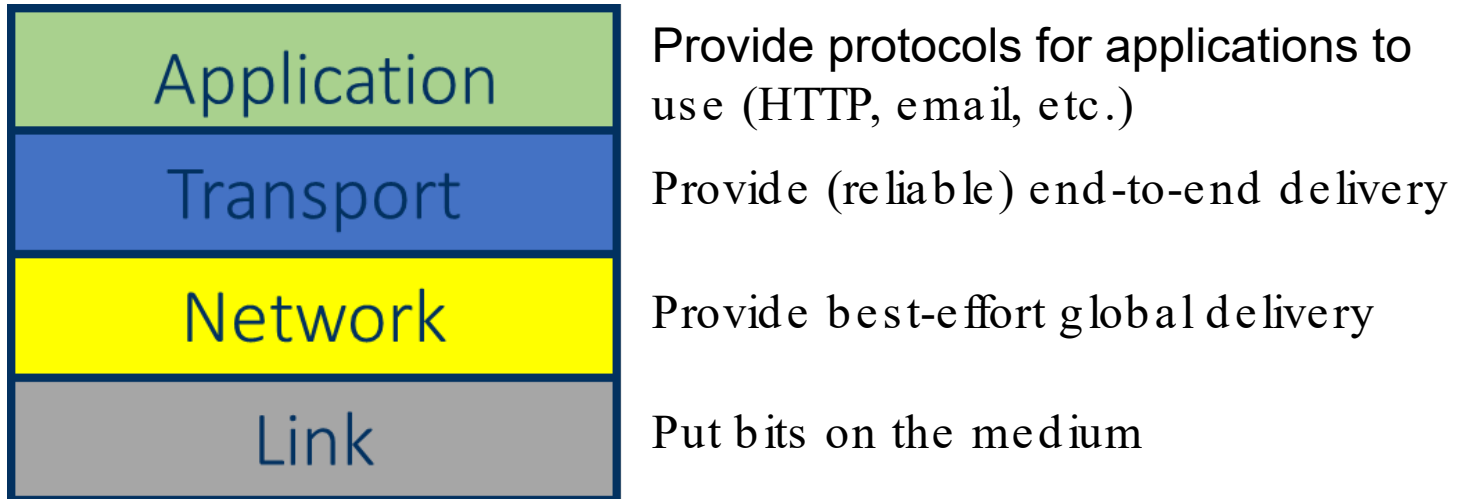
3.

The end-to-
end principle

Recap

- What is the Internet? How do the design goals of the Internet differ from a network?
- What are the four layers of the 4-layer model? What do they do?
- Why is reliability not built into the network layer?

Recap: Layering



Each layer only provides a service to the layer above it.

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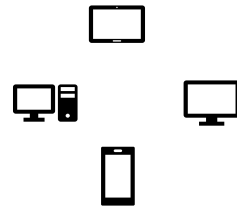
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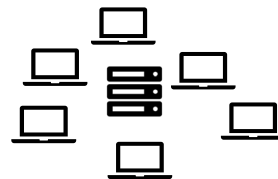
The end-to-
end principle

Building a picture of the Internet: End-hosts

- End-hosts are the devices using the Internet.
 - They exist on the “edge” of the Internet. They “host” Internet apps and services. 23 Billion end-hosts today!
- Two types of end-hosts: Servers and clients.
 - In the context of an app, the “server” provides the service and the “client” uses it.



Home network



Enterprise network

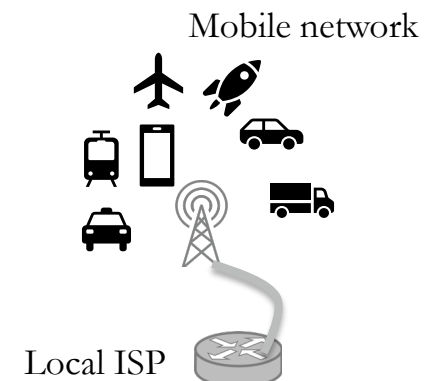
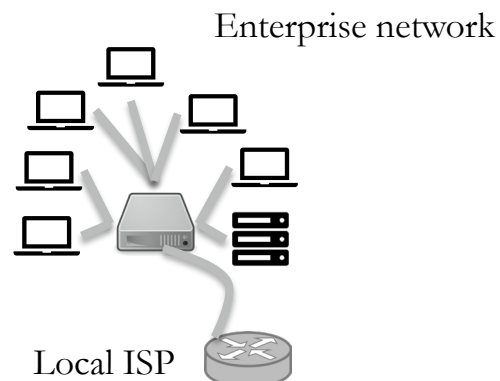
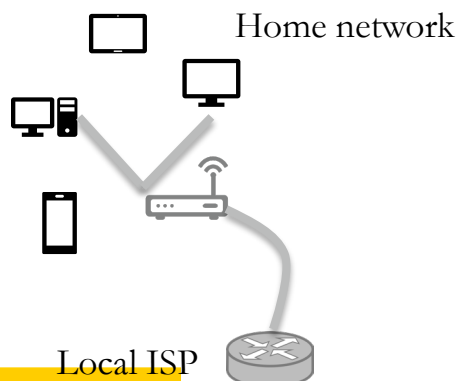


Mobile network

- **Discuss: What layers should be implemented on end-hosts?**

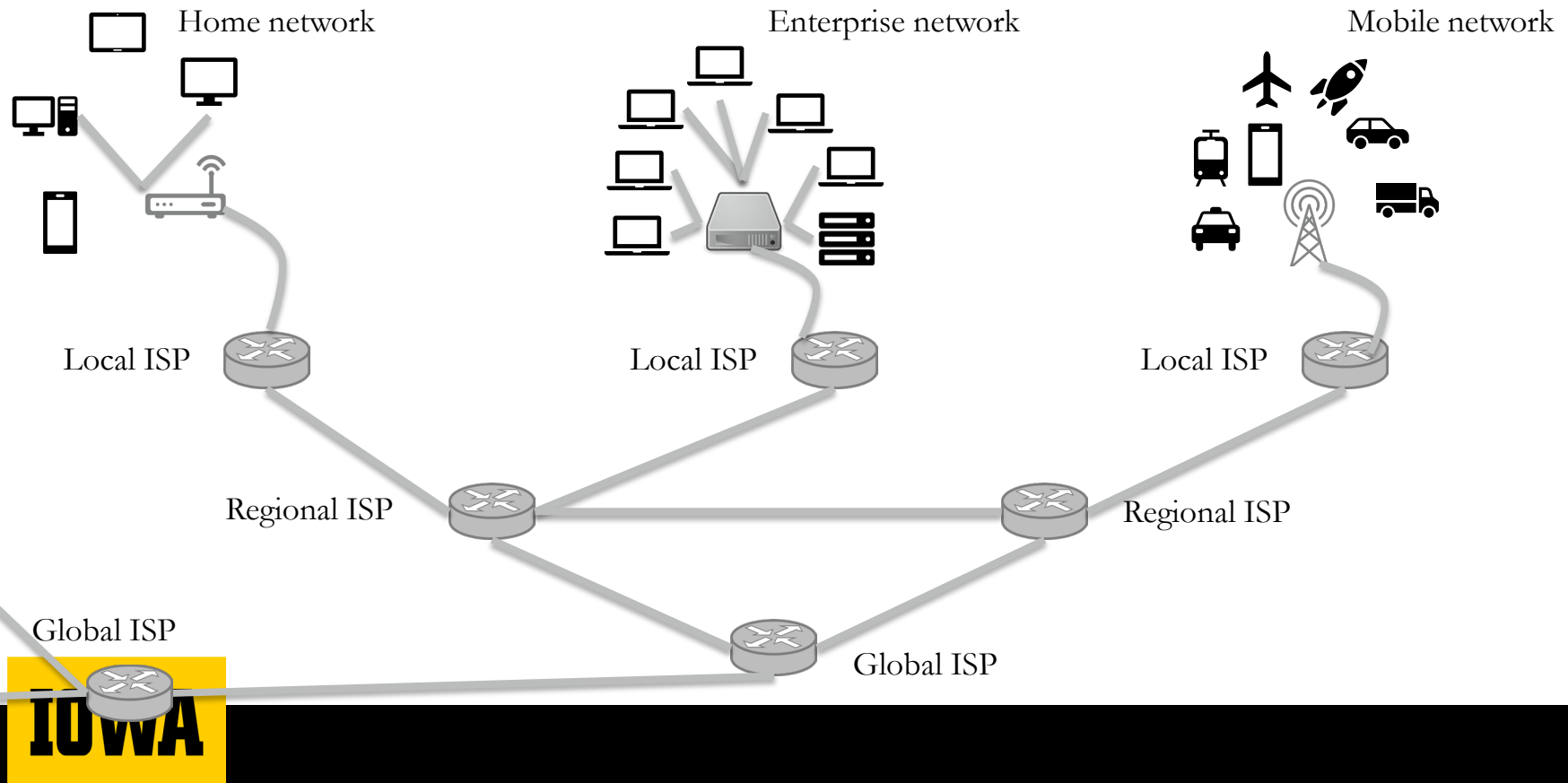
Building a picture of the Internet: Access Networks

- Access networks connect end -hosts to the Internet infrastructure.
 - This is what your Internet Service Provider (ISP) does. They may provide cable, DSL, dial-up, or satellite access networks.
 - Your modem facilitates access to this network via coax cables, radio, optic fiber, or other wide-area network transmission media.
 - You can connect to your modem using ethernet, Wi-Fi, or other local-area network transmission media.
- **Discuss: What layers should be implemented on access network infrastructure?**



Building a picture of the Internet: The network core

- This is the core infrastructure of the Internet
 - The core is where your local ISP connects to other “larger” ISPs. The larger ISPs do the same. Eventually, we end up with a large hierarchical network of networks – the Internet.
- **Discuss: What layers should be implemented on the network core?**



Building a picture of the Internet: Encapsulation

- Data is broken into units at each layer of the Internet model
 - These are called Protocol Data Units (PDUs)



Segments (TCP) Datagrams (UDP)



Packets



Frames

Which app (i.e., port) is the sender/receiver?
Information that enables check for corrupted data.
Multiple segments/datagrams (+headers) if over 64K
Information to help application layer correctly re-assemble data (if TCP).

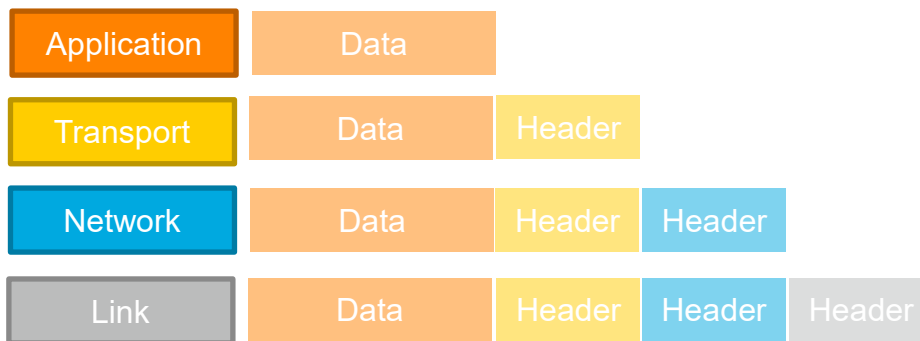
Which computer (i.e., IP address) is the sender/receiver?
How long before we destroy this packet (TTL)?
Information that enables check for corrupted data.
Multiple packets (+headers) if over 64K
Information to help transport layer correctly re-assemble data.

Which hardware interface (i.e., MAC address) is the sender/receiver?
Information that enables check for corrupted data.
Size limit decided by transmission medium.
Multiple frames (+headers) if over size limit.

Any guesses for why 64K is the size limit of a packet?

Building a picture of the Internet: Packets

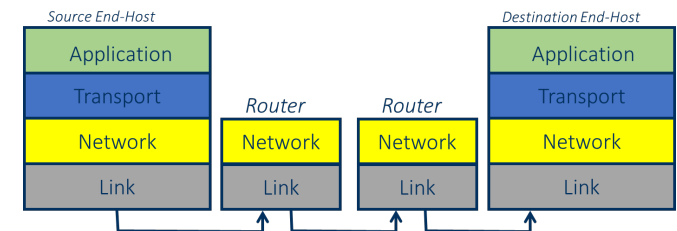
- Segments, datagrams, packets, and frames are units of data.
 - Segments and datagrams reside un-encapsulated only on the end hosts.
 - Frames only exist on a link.
 - What makes packets special?
 - They are seen un-encapsulated at every router between a source and destination because they exist at the network layer.



Segments (TCP)/Datagrams (UDP)

Packets

Frames

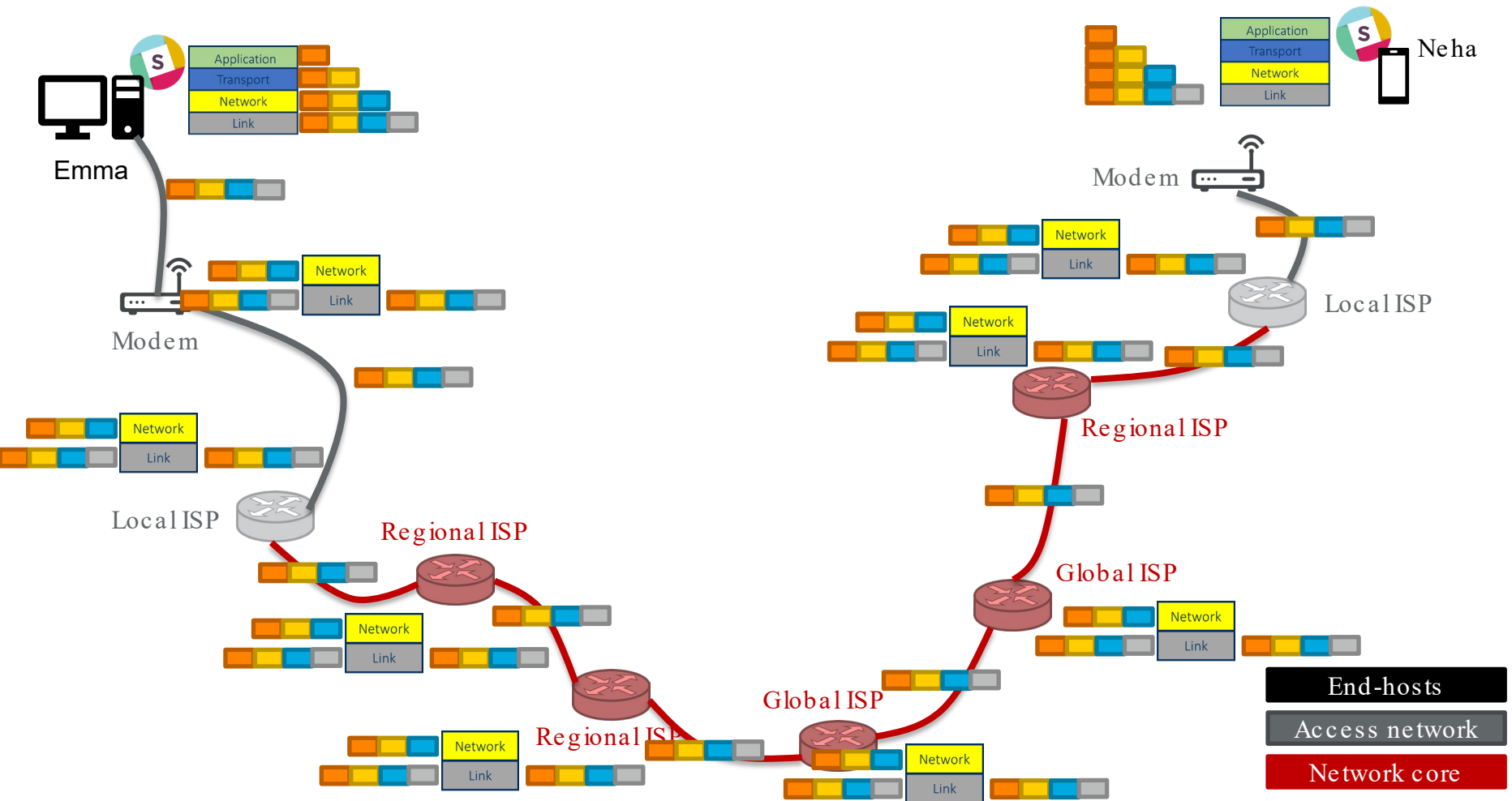


Building a picture of the Internet: Packets

- A packet is a unit of data on the network layer.
 - Think of each packet as a piece of mail.
 - It has control information such as: source IP address, destination IP address, and a “time to live” (TTL). This control information is stored in its 20 byte “header”.
 - Mail analogy: To/From addresses and postage stamps on the envelope.
- The payload of a packet has the segment/datagram to be read by the recipient.
 - At intermediate routers, the header is examined to identify the next router the link layer should send the packet to.
 - Mail analogy: Intermediate postal centers forward the envelope to the next postal center.
 - When a packet reaches its destination, the payload is delivered to the transport layer above it.
 - Mail analogy: Mail is processed by its receiver after delivery.

In 72ms a packet goes from NYC to Sydney. In its lifetime it travels under the oceans, through war-torn lands, dictatorships, and democracies.

The exciting life of a packet!



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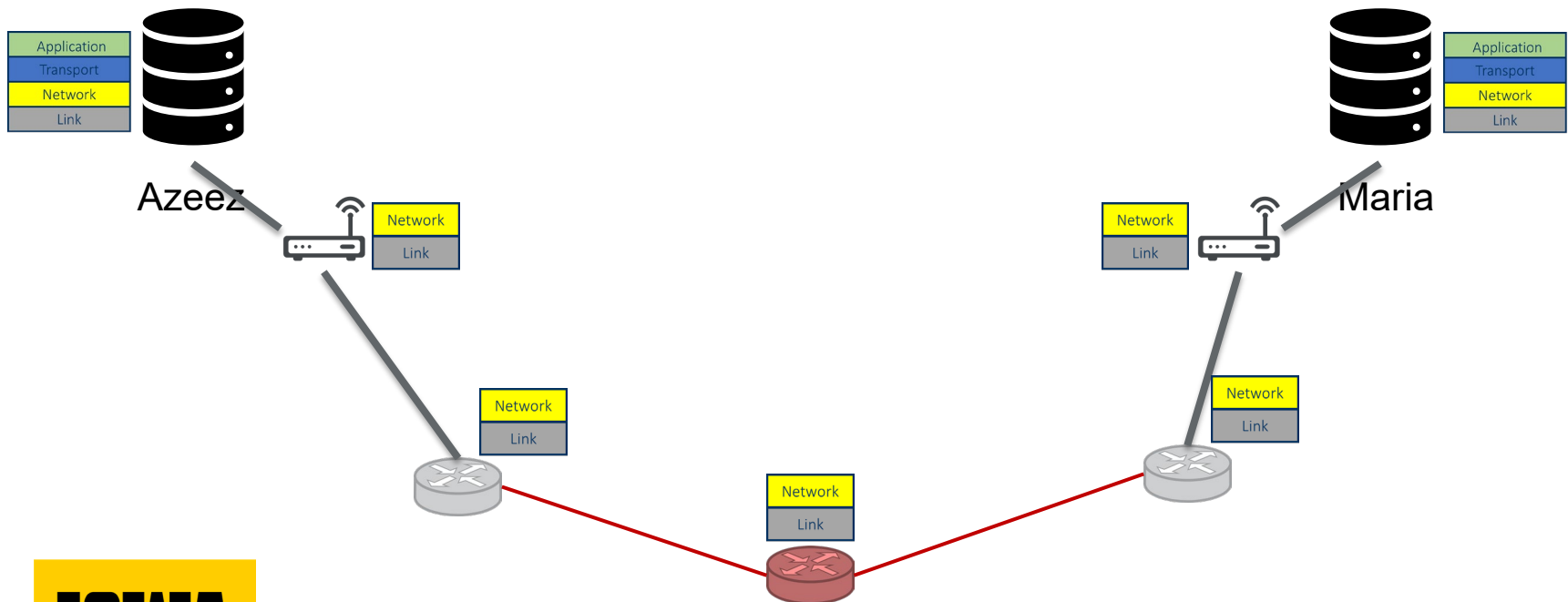
The end-to-
end principle

The three design principles of the Internet

- Layering
 - What modules does the Internet need?
- End-to-end
- Fate-sharing

Let's design a file transfer app

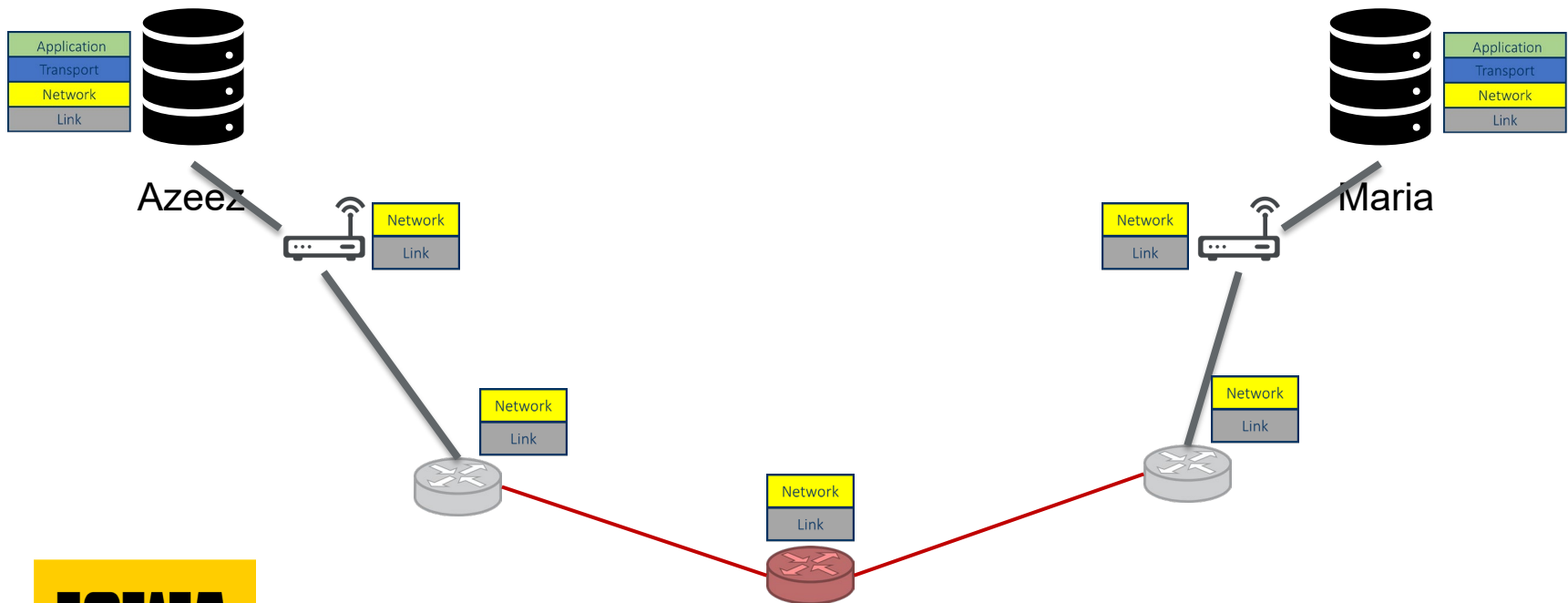
- Scenario: Azeez has to transfer large files to Maria. They need reliability guarantees. How should they implement it?
 - Errors may come from hosts or the network (packets can get corrupted or lost at each node).
 - At the network layer (per hop) or transport layer (per file)? Why?



Let's design a file transfer app

- Solution 1: Make each hop reliable.

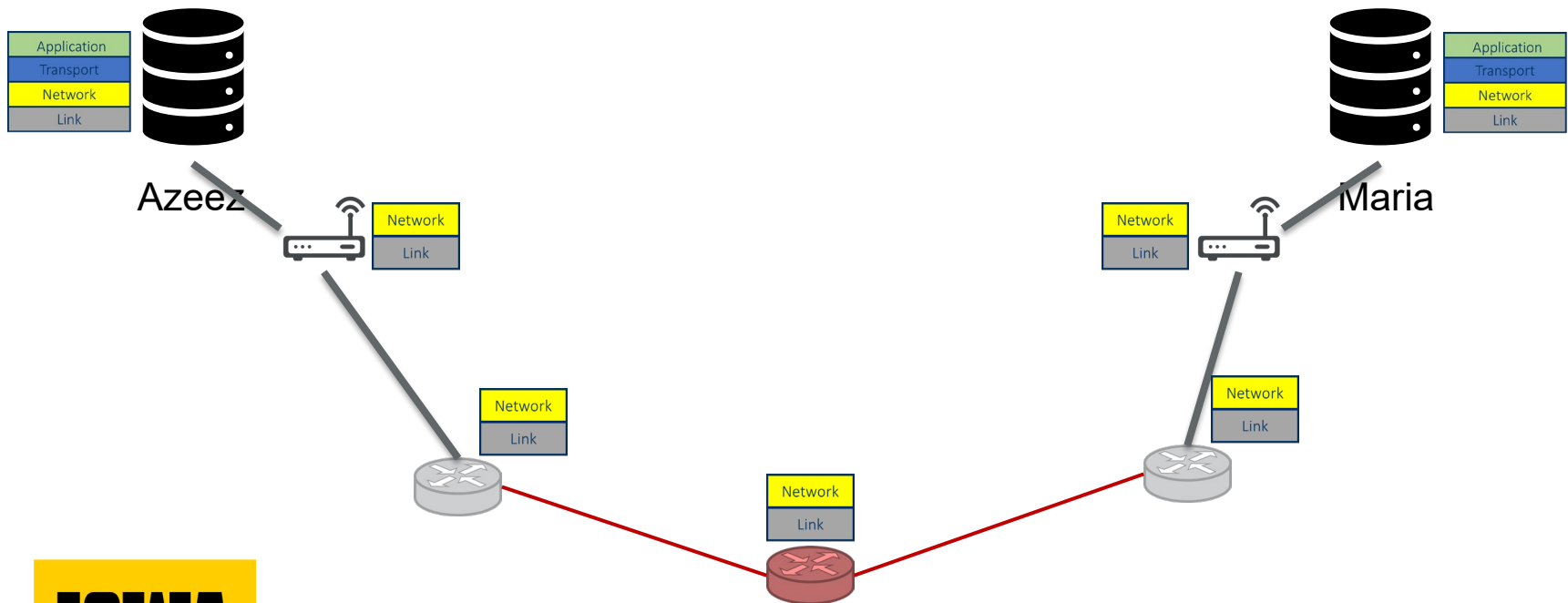
Discuss: Good idea?



Let's design a file transfer app

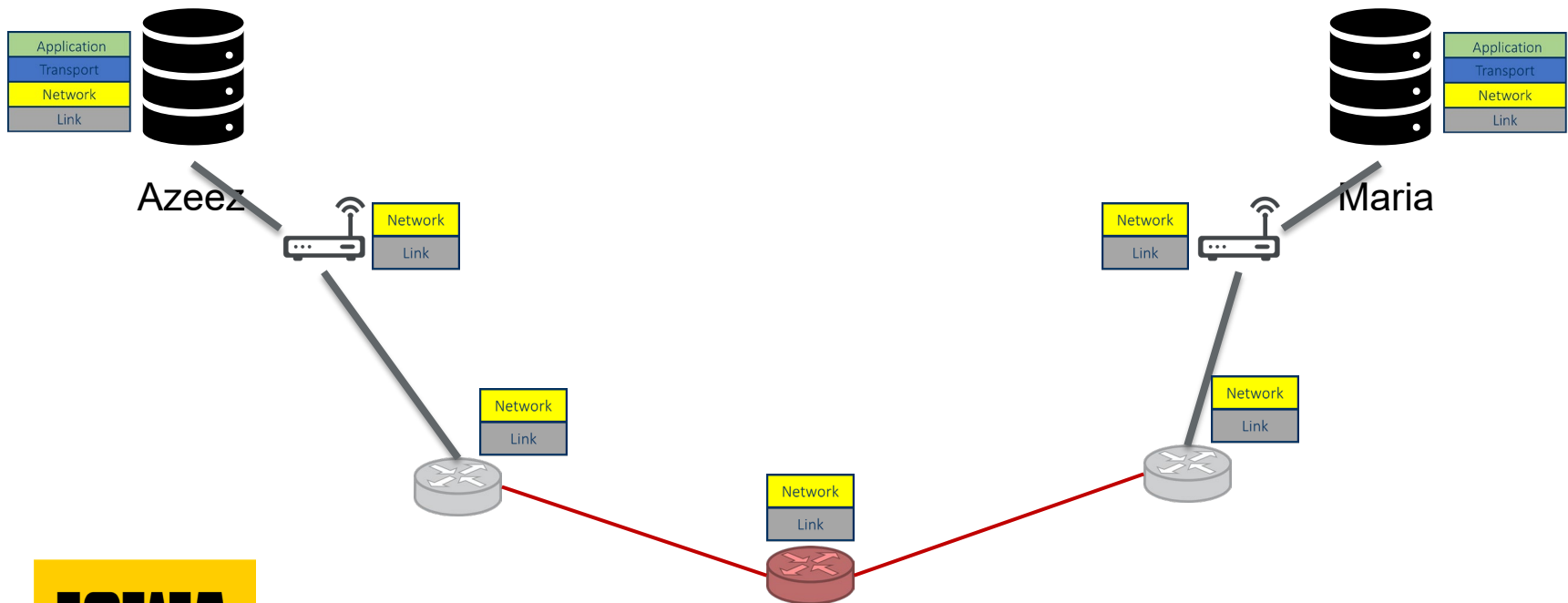
- Solution 2: Let hops be unreliable. Do a final check at the end-host and try the transfer again if required.

Discuss: Good idea?



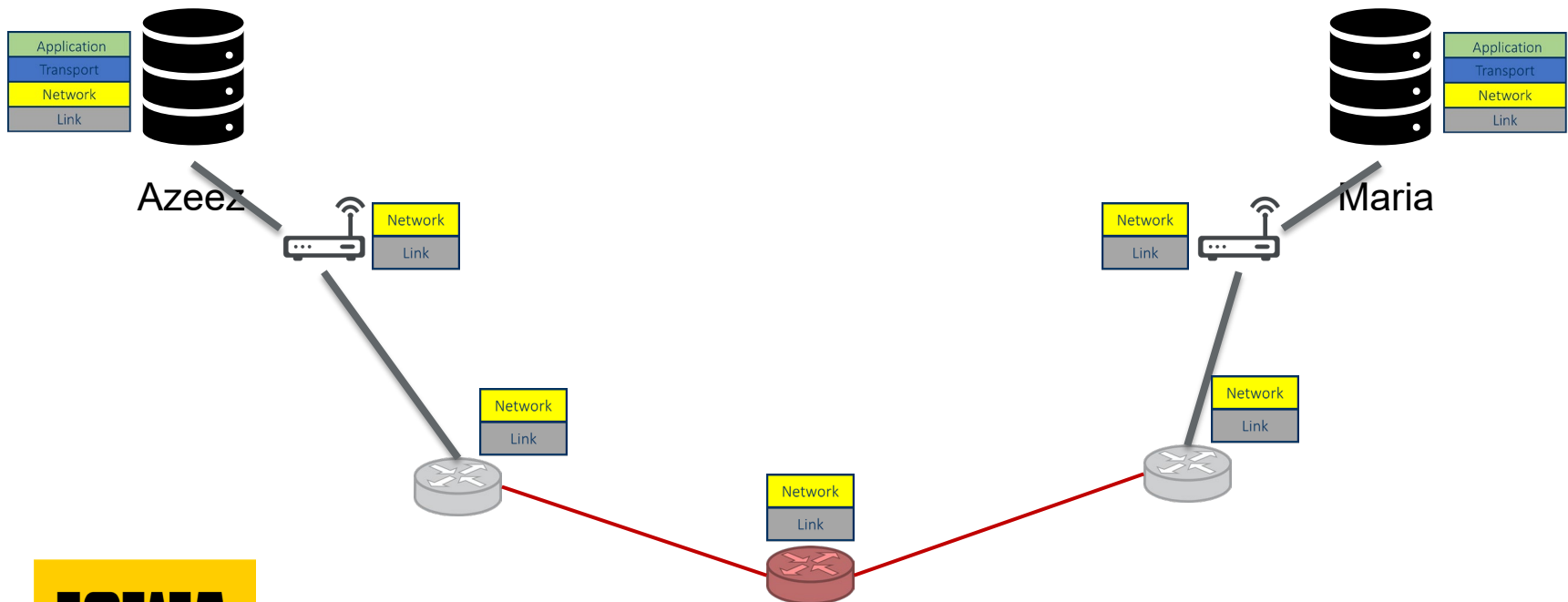
Let's design a file transfer app

- Solution 1: Make each hop reliable.
 - You cannot make anything on the network perfectly reliable. Requires that Maria trusts the functioning of network elements not in her control.
 - What happens if a router misbehaves? Maria will still always have to check files at her end. We introduced a redundant functionality to the network layer.



Let's design a file transfer app

- Solution 2: Let hops be unreliable. Do a final check at the end-host and try the transfer again if required.
 - Can fail only if Maria (the end-host) herself fails. She only has to trust her own functionality.



Layering vs. end-to-end principle

- Layering tells us how to compose functionality. But it never tells us what functionality should be implemented at which layer.
 - Example: Use HTTP over TCP over IP over Ethernet to download web pages. But where should encryption be implemented?
- The end-to-end principle gives us guidance on where functionality should be implemented.

The end-to-end principle

- Functionality should be implemented at the lower layer if and only if:
 - It can be implemented completely
 - It can be implemented correctly
 - It is not made redundant by a higher layer
- We can make exceptions for performance optimizations.
 - As long as it does not add a burden to applications that do not need the functionality.
- This is a principle, not a rule.
 - Violations exist on the Internet today.
 - Examples: Firewalls, proxies, and NATs. (We'll get to them later in the term).

A simplified version of the end-to-end principle

If you can implement a functionality on the end -hosts without requiring the network's help, do it.

Why?

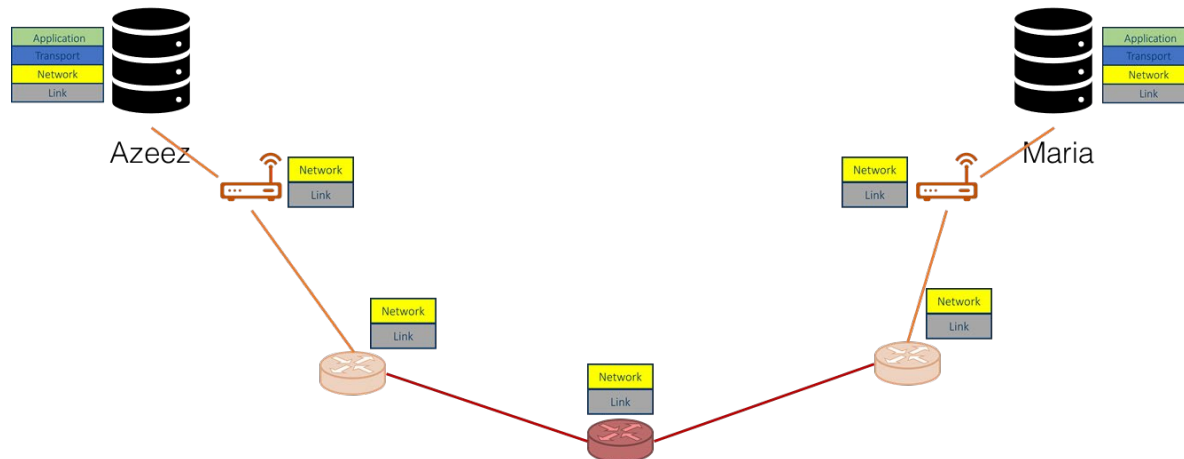
Because you can never rely on the network.

Why does the end-to-end principle make sense?

- In general, implementing functionality in the network:
 - Does not reduce end-host/application complexity.
 - Increases network complexity.
 - Imposes additional costs on all applications (even those that do not require the functionality).

When does violating the e2e principle make sense?

- Scenario: Highly unreliable networks.
 - There are six links between Azeez and Maria. Each has a fail rate of 10%
 - Probability of a packet not reaching Maria: 47%
 - Should we retransmit the entire file each time the network drops a packet?



Impact of the e2e principle

- The end-to-end principle is why reliability is not implemented in the network.
 - It is implemented at the end-host (in the transport layer via TCP).
- Reduces the functions of the stressed network core. Allows it focus on doing a few jobs well.
 - (Funnily enough) This results in better network reliability.
- More flexibility for applications.
 - They get to pick and choose what kind of functionality they want.
 - They don't pay a price for things they don't need.

Discussion

- Who benefits from the e2e principle? Which entities have an interest in violating it?

The three design principles of the Internet

- Layering
 - What modules does the Internet need?
- End-to-end
 - What functionality should go in which modules?
- Fate-sharing

Fate sharing

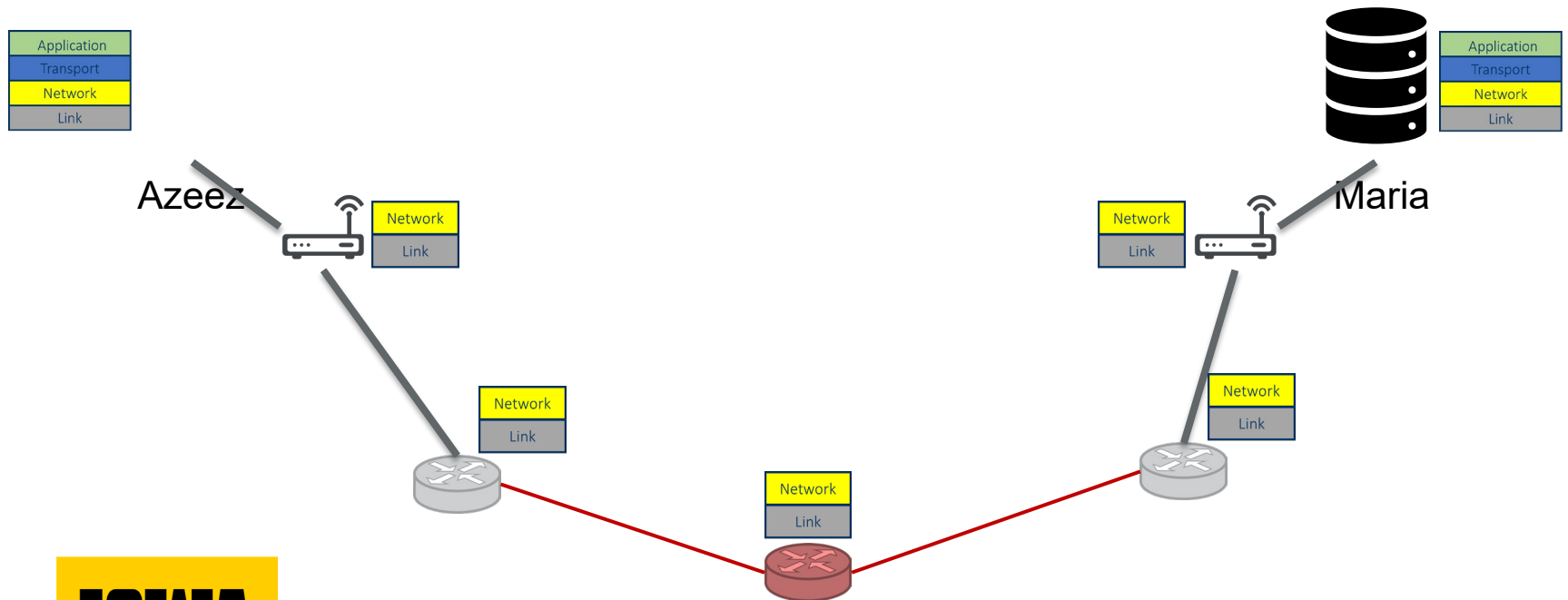
“A distributed system is one in which the failure of a computer you didn't even know existed can render your own computer unusable.” – Leslie Lamport

- Store network state only in the entities that rely on this state.
 - Critical state is lost only when the entities that depend on it fail.

We'll revisit this principle a lot when we explore how packets are routed on the Internet (the network layer).

Fate sharing

- Solution 2: Allow hops to be unreliable. Do a check on completion and try the transfer again if required.
- Why was this the better solution?
 - The success of the transfer relied only on Maria and Azeez – the two parties who benefitted from its completion.



The three design principles of the Internet

- Layering
 - What modules does the Internet need?
- End-to-end
 - What functionality should go in which modules?
- Fate-sharing
 - Who should store network state?

Highly recommended reading

- David Clark. ["The Design Philosophy of the DARPA Internet Protocols"](#).

What you should remember

- What are the building blocks of the Internet?
 - Which Internet layers do each of them use?
- What are the units of data at each layer?
 - What role do headers play at each level?
- What is the end-to-end principle?
 - What does it do that layering does not?
 - Why does it apply to most conditions?
 - When might we consider violating the principle?
 - How has it shaped the Internet?
- What is the fate-sharing principle?