

# Negotiation Strategy for Self-Organizing Stable Matchings

Maxime Morge<sup>‡</sup>    **Gauthier Picard<sup>†</sup>**

<sup>†</sup>ISCOD team, Ecole des Mines de Saint-Etienne  
picard@emse.fr

<sup>‡</sup>SMAC team, Laboratoire d'Informatique Fondamentale de Lille  
Maxime.Morge@lifl.fr

December 15, 2010

# Introduction

## Our objectives

- ▶ Stable Matching Problems with **no central orchestrator**
- ▶ Multi-agent framework to distributively solve assignment problems
- ▶ Assignments are viewed as **emergent** phenomena resulting from local agent negotiations: *stability, social optimal welfare*

## Casanova

- ▶ Distributed method to solve the **stable marriage problem**
- ▶ Agent behaviors leading negotiation processes to *social optimal assignments*
- ▶ *Minimal concession strategy*
- ▶ Preserves the **privacy**
- ▶ Improves the **optimality** of the solution and the **equity** amongst partners

# Stable Marriage Problem



**Alain**

Barbara  $\succ$  Alice  $\succ$  Carla



**Bob**

Carla  $\succ$  Barbara  $\succ$  Alice



**Charles**

Alice  $\succ$  Carla  $\succ$  Barbara



**Alice**

Bob  $\succ$  Alain  $\succ$  Charles



**Barbara**

Charles  $\succ$  Bob  $\succ$  Alain



**Carla**

Alain  $\succ$  Charles  $\succ$  Bob

# Stable Marriage Problem: Definition

- ▶ Subclass of stable matching problems [GS62]
- ▶ Two finite sets of participants: men and women (or *communities*)

## Definition (Stable Marriage Problem)

A **stable marriage problem** of size  $n$  (with  $n \geq 1$ ) is a couple  $SM = \langle X, Y \rangle$  where:

- ▶  $X = \{x_1, \dots, x_n\}$  is a set  $n$  men ranking women in a strict and complete order forming his preference list.  
 $\forall 1 \leq i \leq n, x_i = (y_i^0, \dots, y_i^{n-1})$
- ▶  $Y = \{y_1, \dots, y_n\}$  is a set  $n$  women ranking men in a strict and complete order forming her preference list.  
 $\forall 1 \leq i \leq n, y_i = (x_i^0, \dots, x_i^{n-1})$

# Matchings



**Alain**

Barbara  $\succ$  Alice  $\succ$  Carla



**Bob**

Carla  $\succ$  Barbara  $\succ$  Alice



**Charles**

Alice  $\succ$  Carla  $\succ$  Barbara



**Alice**

Bob  $\succ$  Alain  $\succ$  Charles



**Barbara**

Charles  $\succ$  Bob  $\succ$  Alain



**Carla**

Alain  $\succ$  Charles  $\succ$  Bob

Alain–Barbara

Bob–Alice

Charles–Carla

# Unstable Matchings



**Alain**

Barbara  $\succ$  Alice  $\succ$  Carla



**Bob**

Carla  $\succ$  Barbara  $\succ$  Alice



**Charles**

Alice  $\succ$  Carla  $\succ$  Barbara



**Alice**

Bob  $\succ$  Alain  $\succ$  Charles



**Barbara**

Charles  $\succ$  Bob  $\succ$  Alain



**Carla**

Alain  $\succ$  Charles  $\succ$  Bob

Alain–Barbara

Bob–Alice

Charles–Carla

# Stable Matchings



**Alain**

Barbara  $\succ$  Alice  $\succ$  Carla



**Bob**

Carla  $\succ$  Barbara  $\succ$  Alice



**Charles**

Alice  $\succ$  Carla  $\succ$  Barbara



**Alice**

Bob  $\succ$  Alain  $\succ$  Charles



**Barbara**

Charles  $\succ$  Bob  $\succ$  Alain



**Carla**

Alain  $\succ$  Charles  $\succ$  Bob

Alain–Alice  
Bob–Barbara  
Charles–Carla

Charles–Alice  
Alain–Barbara  
Bob–Carla

Bob–Alice  
Charles–Barbara  
Alain–Carla

# Stable Matchings: Definition

## Definition (Matching)

Let  $SM = \langle X, Y \rangle$  be a stable marriage problem of size  $n$  (with  $n \geq 1$ ). A **matching** for  $SM$  is a  $n$ -uplet  $M = \langle m_1, \dots, m_n \rangle$  of  $n$  marriages where each  $m_i$  (with  $1 \leq i \leq n$ ) is a couple  $(x_i, y_i) \in X \times Y$  such that the matching is complete, i.e. each individual is married. Formally,  
 $\forall x \in X \exists! y \in Y (x, y) \in M$ .

## Definition (Stable Matching)

Let  $SM = \langle X, Y \rangle$  be a stable marriage problem of size  $n$  (with  $1 \leq n$ ). and  $M$  a matching for  $SM$ .  $M$  is **stable** iff:

$\forall (x_i, y_i) \in M, \nexists (x_j, y_j) \in M \ x_j \succ_{y_i} x_i \text{ and } y_j \succ_{x_i} y_i$

$\Rightarrow$  No pair of people like each other better than their current spouse

$\Rightarrow$  Find a stable assignment that is *optimal* with respect to a metric that depends on the preferences of the agents



# Social Welfare

## Social Welfare

- ▶ Aggregation of individual preferences can be modelled using the notion of social welfare as studied in Welfare Economics and Social Choice Theory [ASS02]
- ▶ individual agents evaluate their satisfaction using utility functions mapping assignments to numerical values
- ▶ The social welfare theory is used to evaluate the matching, considering the welfare of each person [ASS02]

## Definition (Utility function)

Let  $SM = \langle X, Y \rangle$  be a stable marriage problem of size  $n$  (with  $n \geq 1$ ),  $z = (t_i^0, \dots, t_i^k, \dots, t_i^{n-1})$  an individual agent and  $T$  be the potential partners of  $z$ . The **utility function** of the agent  $z$  is a function  $u_z : T \rightarrow \mathbb{R}$ . If the matching assigns  $z$  with  $t_i^k$ , then  $u_z(t_i^k) = \frac{(n-1)-k}{n-1}$ .

# Social Welfare in SM



**Alain**

Barbara  $\succ$  Alice  $\succ$  Carla



**Bob**

Carla  $\succ$  Barbara  $\succ$  Alice



**Charles**

Alice  $\succ$  Carla  $\succ$  Barbara



**Alice**

Bob  $\succ$  Alain  $\succ$  Charles



**Barbara**

Charles  $\succ$  Bob  $\succ$  Alain



**Carla**

Alain  $\succ$  Charles  $\succ$  Bob

Alain–Alice  
Bob–Barbara  
Charles–Carla

$$U_M = 1.5$$

$$U_W = 1.5$$

Charles–Alice  
Alain–Barbara  
Bob–Carla

$$U_M = 3$$

$$U_W = 0$$

Bob–Alice  
Charles–Barbara  
Alain–Carla

$$U_M = 0$$

$$U_W = 3$$

# Social Welfare in SM: Definition

## Definition (Social welfare)

Let  $SM = \langle X, Y \rangle$  be a stable marriage problem of size  $n$  (with  $n \geq 1$ ) and  $M$  a matching for  $SM$ .

- ▶ The **utilitarian welfare** considers the welfare of the whole society:  

$$sw_u(X \cup Y) = \sum_{z \in X \cup Y} u_z(p_M(z)).$$
- ▶ The **male welfare** considers the welfare of the men:  

$$sw_u(X) = \sum_{x \in X} u_x(p_M(x)).$$
- ▶ The **female welfare** considers the welfare of the women:  

$$sw_u(Y) = \sum_{y \in Y} u_y(p_M(y)).$$
- ▶ The **equity welfare** considers the fairness among partners' welfare in every marriage:  $sw_e(X \cup Y) = 1 - \frac{|sw_u(X) - sw_u(Y)|}{n}.$

# Social Welfare in SM: Definition (cont.)

## Social welfare $\approx$ measure of the global quality of a matching

- ▶ e.g. the utilitarian social welfare a suitable indicator when we are considering the mean of agent's utilities whatever they are male or woman
- ▶ e.g. the male welfare (resp. female welfare) measures the quality of a matching from the viewpoint of the men (resp. women)
- ▶ e.g. the equity welfare may be a suitable indicator when we have to satisfy both the men and the women

# Gale-Shapley Algorithm

## First Algorithm for SMP [GS62]

- ▶ Centralized algorithm
- ▶ Always finds a stable matching for any instance of the SMP
- ▶ Produces a matching in which each man has the best partner he can have

```

assign each person to be free
while a man  $x$  is free do
     $y \leftarrow$  first woman on  $x$ 's list
    //  $x$  proposes to  $y$ 
    if a man  $x_y$  is already married to  $y$  then
        | assign  $x_y$  to be free
    assign  $x$  and  $y$  to be married to each other
    for each successor  $x_y$  of  $x$  on  $y$ 's list do
        | delete  $x_y$  from  $y$ 's list
        | delete  $y$  from  $x_y$ 's list
  
```

**Algorithm 1:** The men-propose Gale-Shapley algorithm for SM

# GS Algorithm Sample Execution



Alice



Charles



Bob



Alain

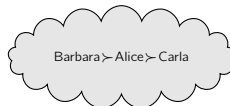
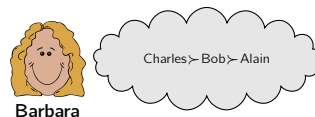
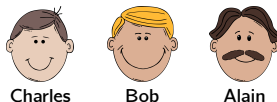


Barbara

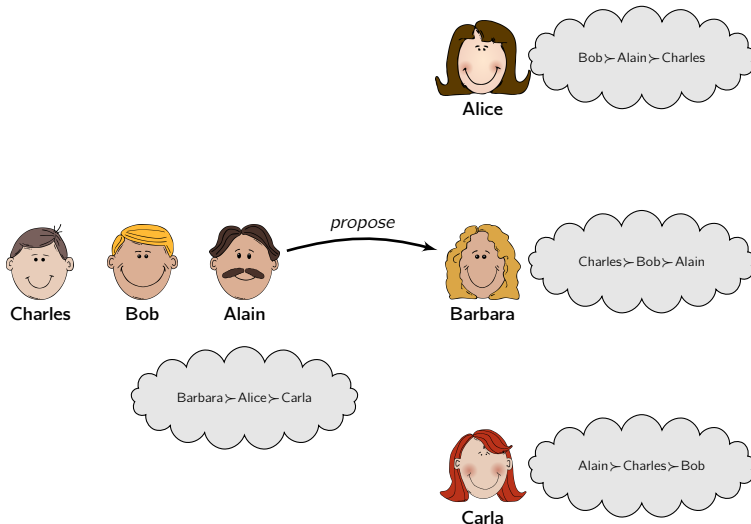


Carla

# GS Algorithm Sample Execution

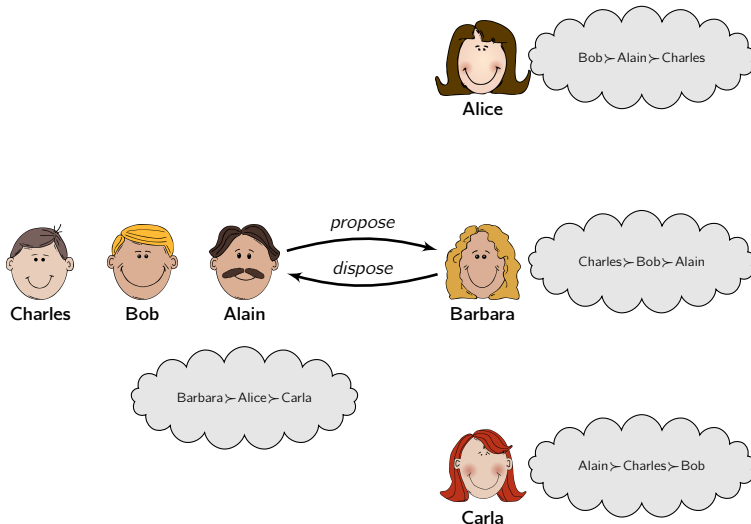


# GS Algorithm Sample Execution

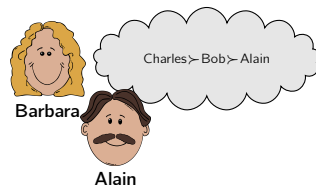
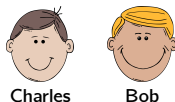




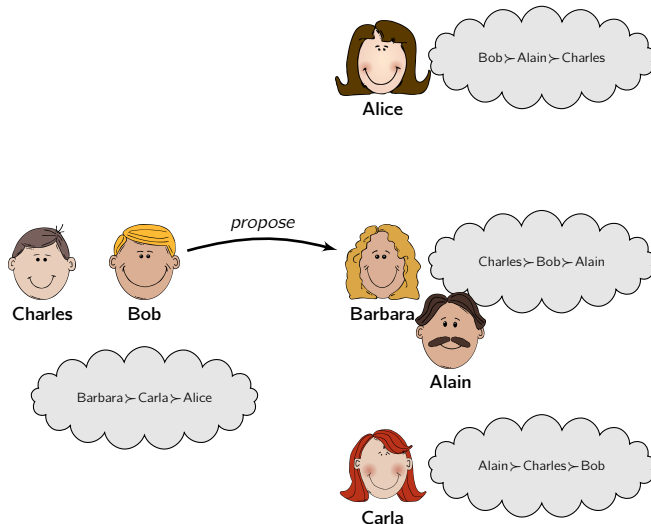
# GS Algorithm Sample Execution



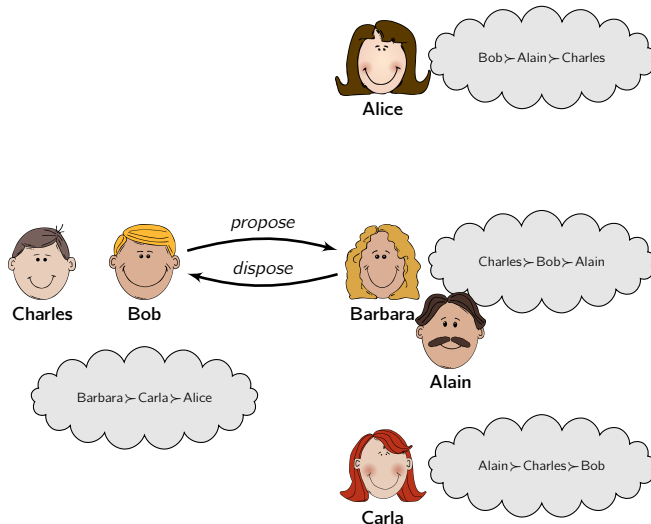
# GS Algorithm Sample Execution



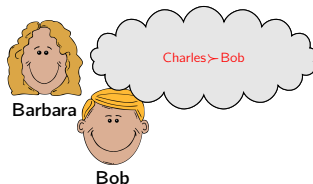
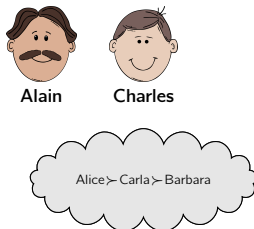
# GS Algorithm Sample Execution



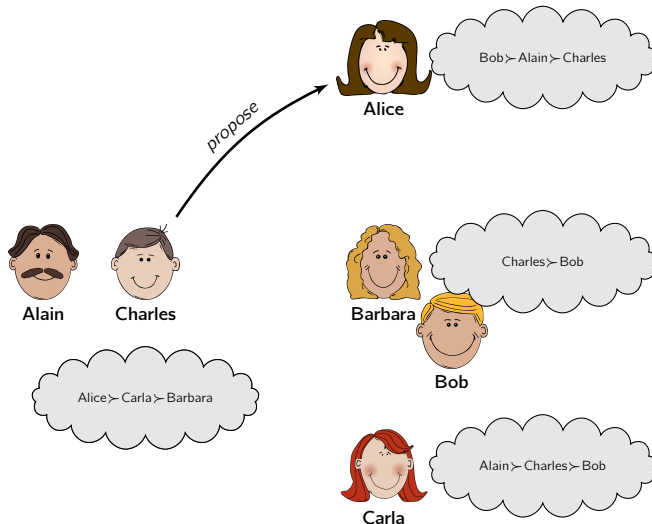
# GS Algorithm Sample Execution



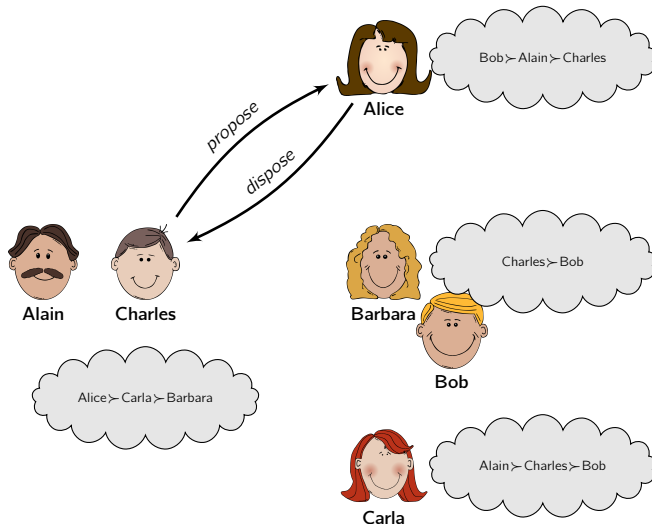
# GS Algorithm Sample Execution



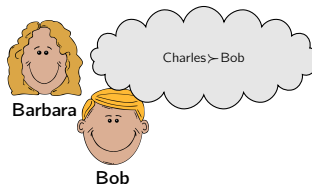
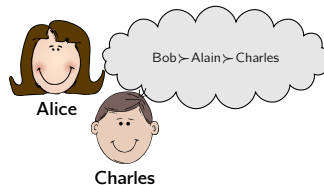
# GS Algorithm Sample Execution



# GS Algorithm Sample Execution

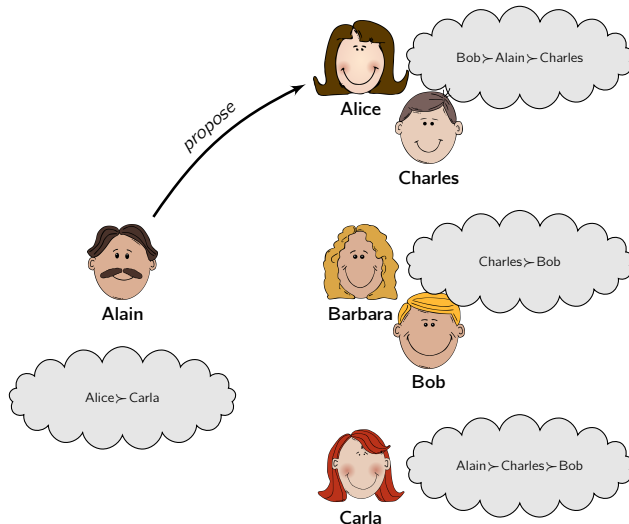


# GS Algorithm Sample Execution

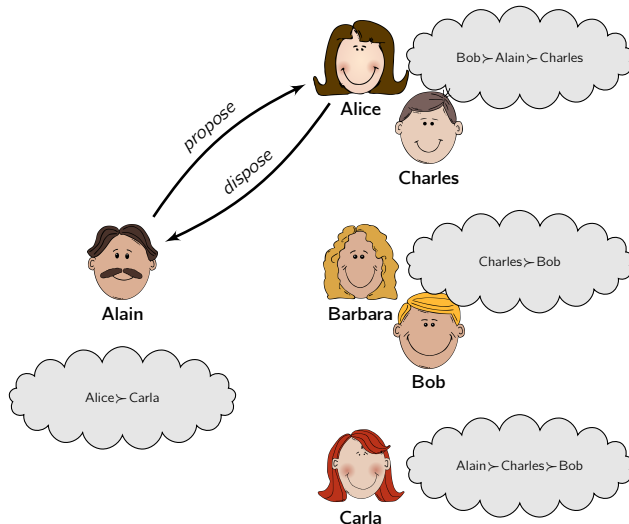




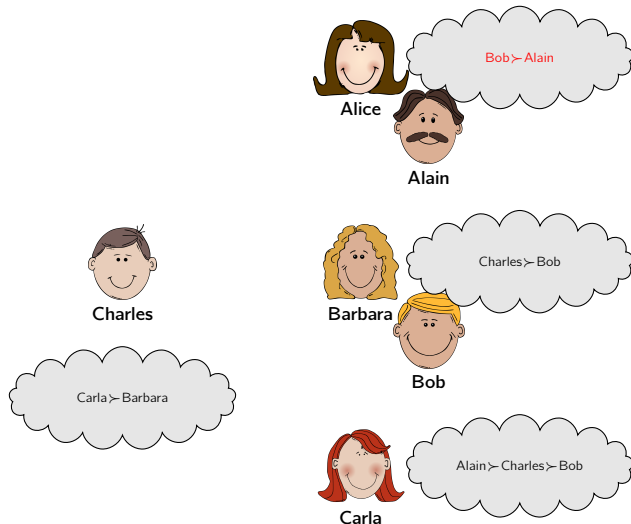
# GS Algorithm Sample Execution



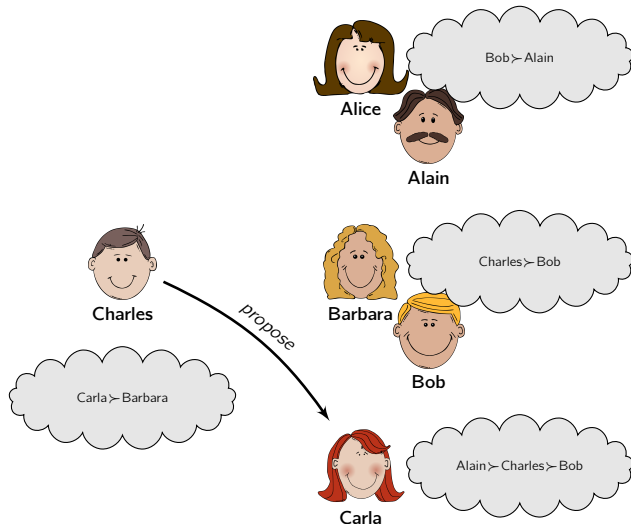
# GS Algorithm Sample Execution



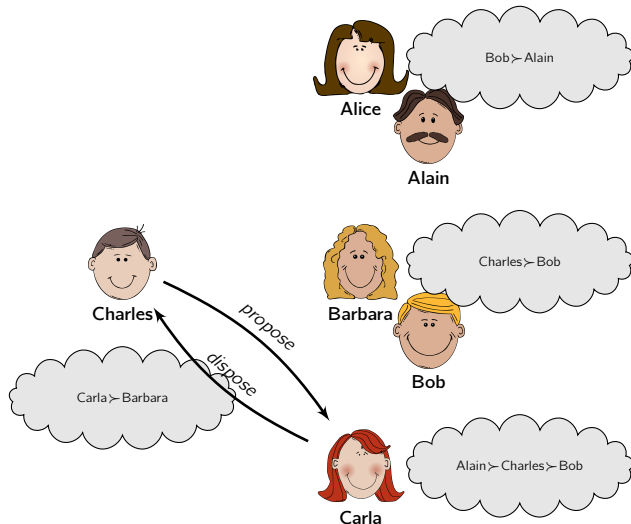
# GS Algorithm Sample Execution



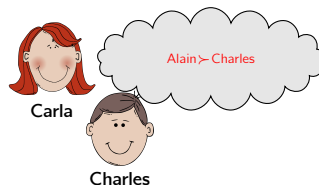
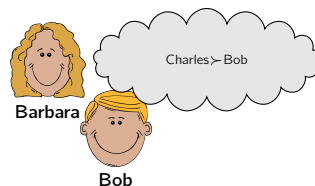
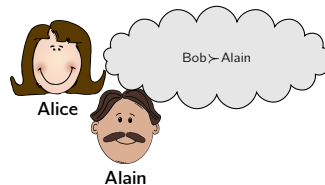
# GS Algorithm Sample Execution



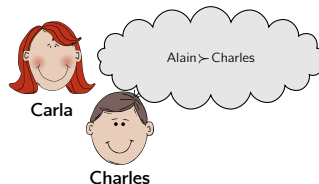
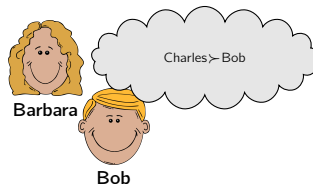
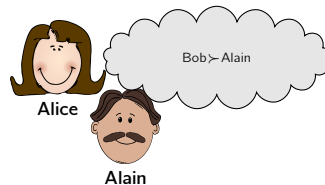
# GS Algorithm Sample Execution



# GS Algorithm Sample Execution



# GS Algorithm Sample Execution



# Gale-Shapley Algorithm (cont.)

## Distributed Extended Gale-Shapley Algorithm (DisEGS) [BM05]

- ▶ Distributed version of the GS algorithm
- ▶ Each man (and woman) is represented by an agent
- ▶ Agents exchange messages (*propose*, *accept* and *delete*) as to reproduce the GS algorithm and find a stable assignment
- ▶ Agents do not follow the same procedure:
  - ▶ men propose
  - ▶ women accept or decline proposals
- ▶ each agent keeps its own preferences → privacy



# Casanova Algorithm

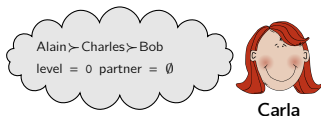
## Casanova Characteristics

- ▶ Multi-agent algorithm: men and women are agents
- ▶ Privacy preserving: preferences are not shared
- ▶ Equity preserving: men and women follow the same protocol
- ▶ Optimal wrt social utilitarian welfare

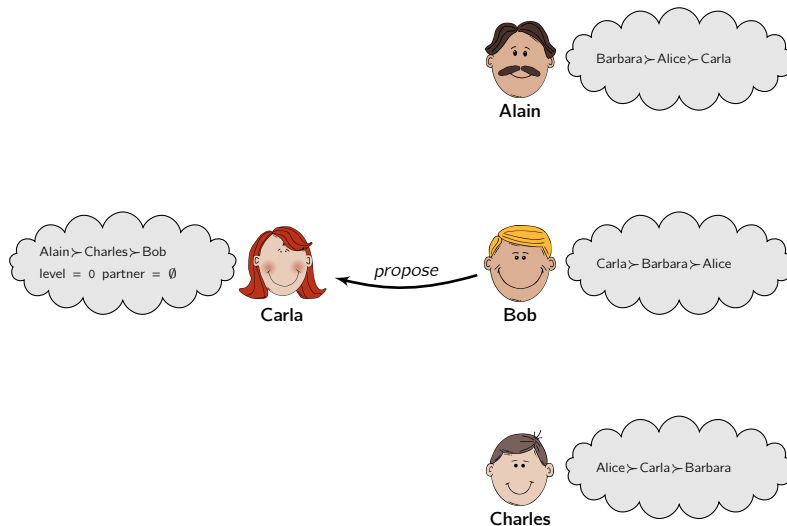
## Casanova Principles

- ▶ Agents send concurrently proposals and reply with acceptance or rejections
- ▶ Agents use the minimal concession strategy [RZ94, MM10]
  - ▶ from the most preferred partners
  - ▶ to the less preferred ones
  - ▶ but still trying to attract the preferred ones

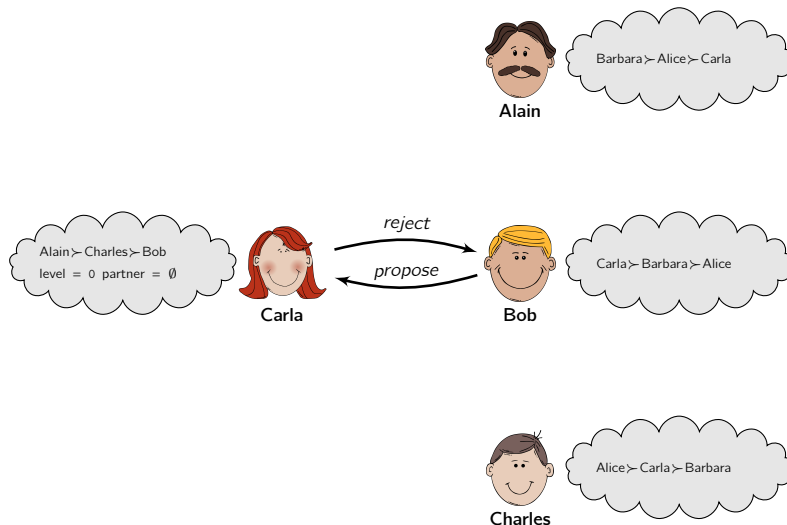
# Casanova Sample Execution



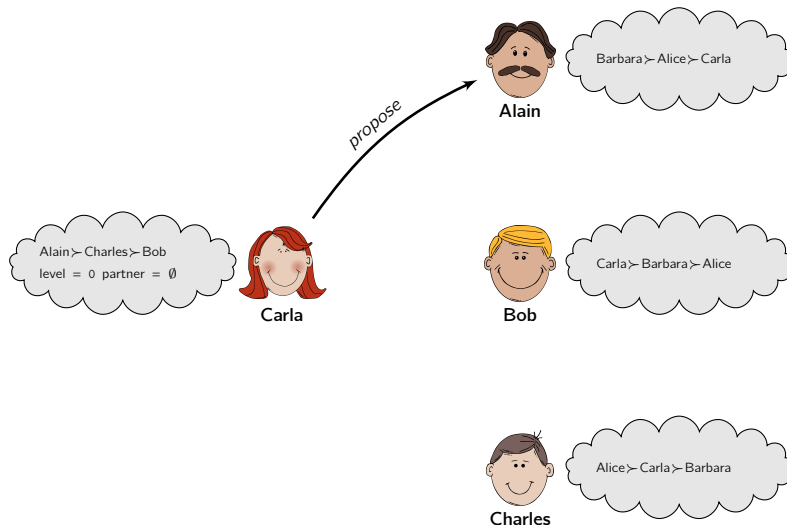
# Casanova Sample Execution



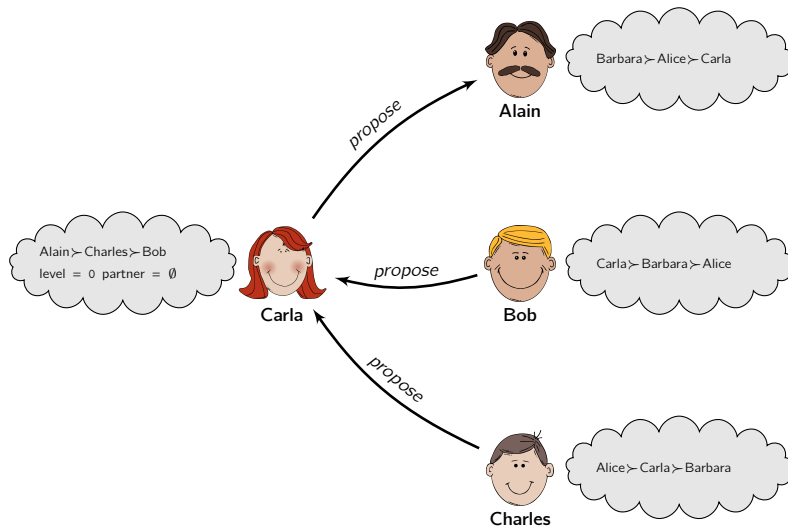
# Casanova Sample Execution



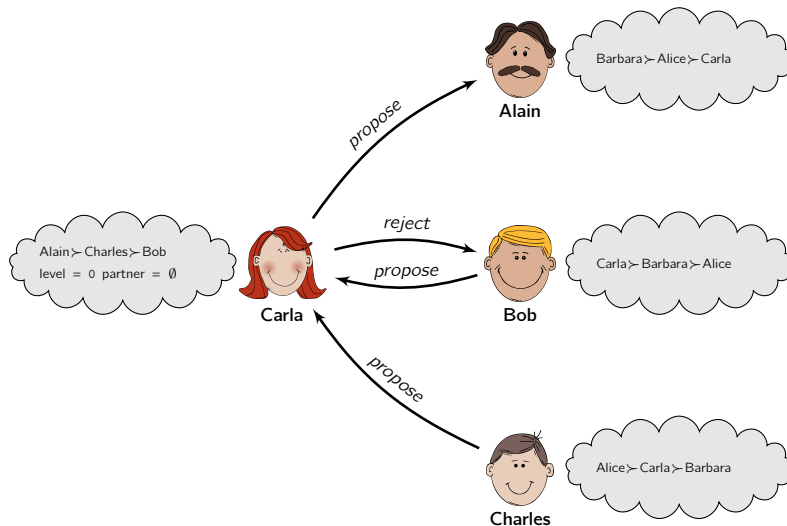
# Casanova Sample Execution



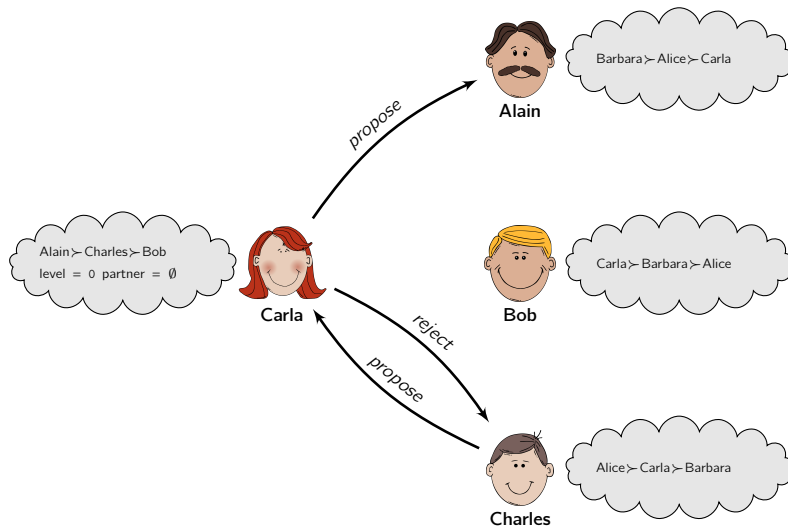
# Casanova Sample Execution



# Casanova Sample Execution

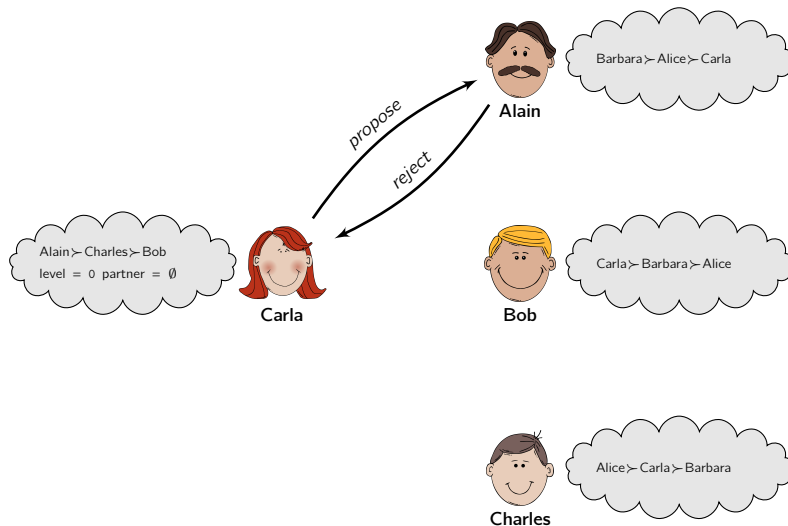


# Casanova Sample Execution

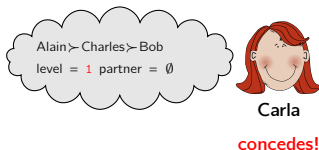




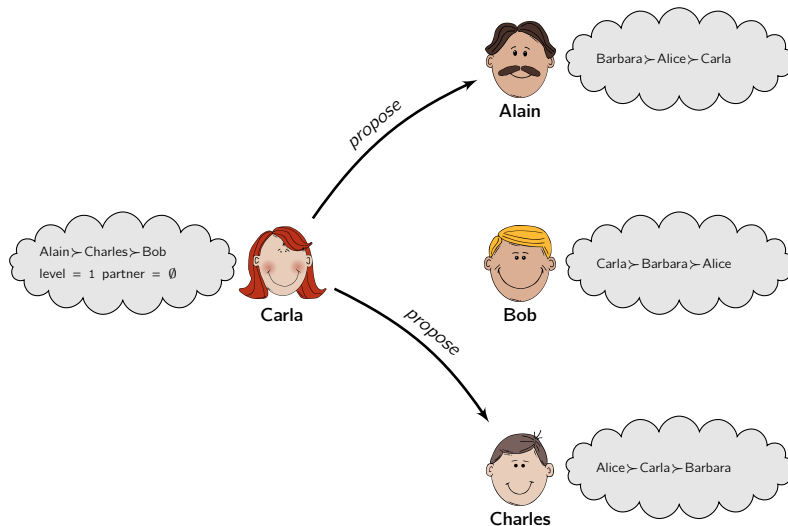
# Casanova Sample Execution



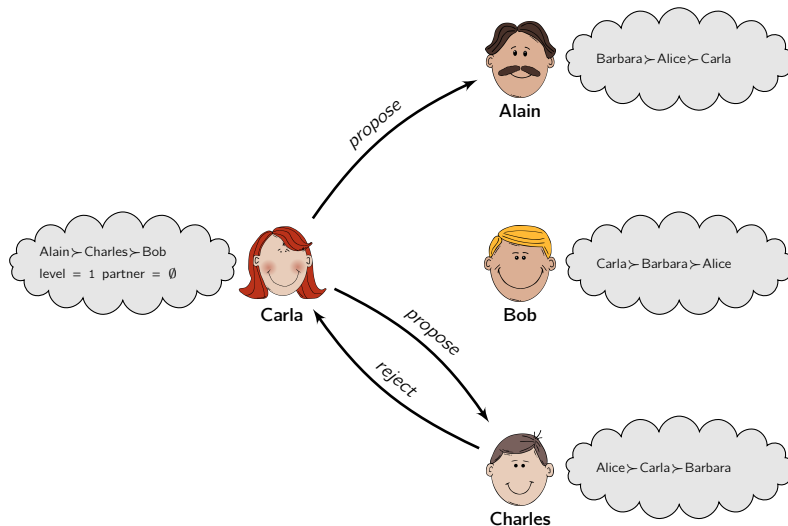
# Casanova Sample Execution



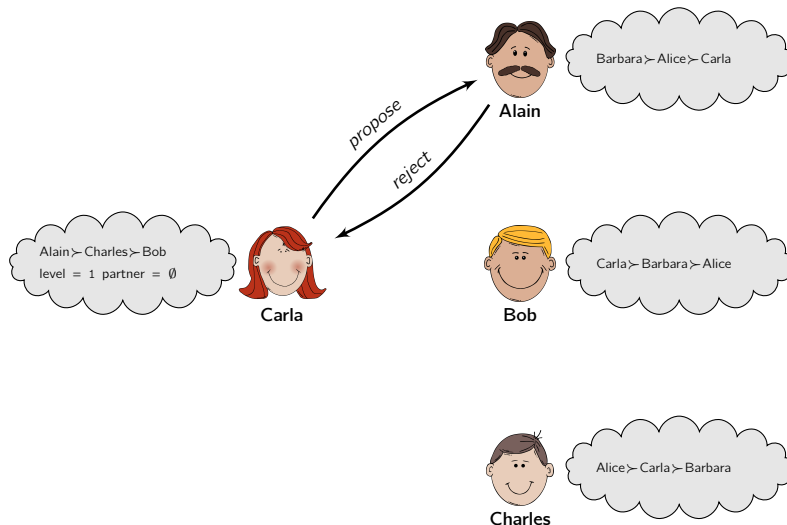
# Casanova Sample Execution



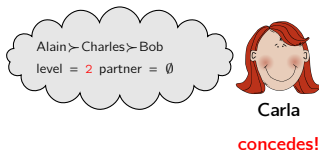
# Casanova Sample Execution



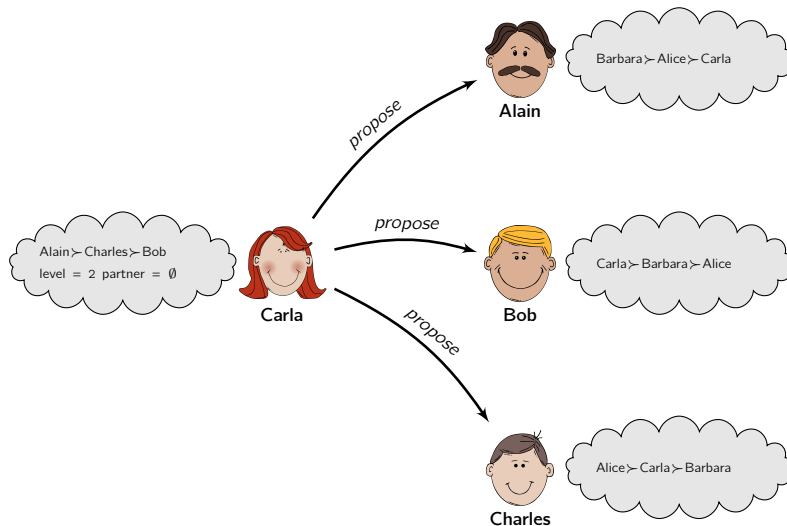
# Casanova Sample Execution



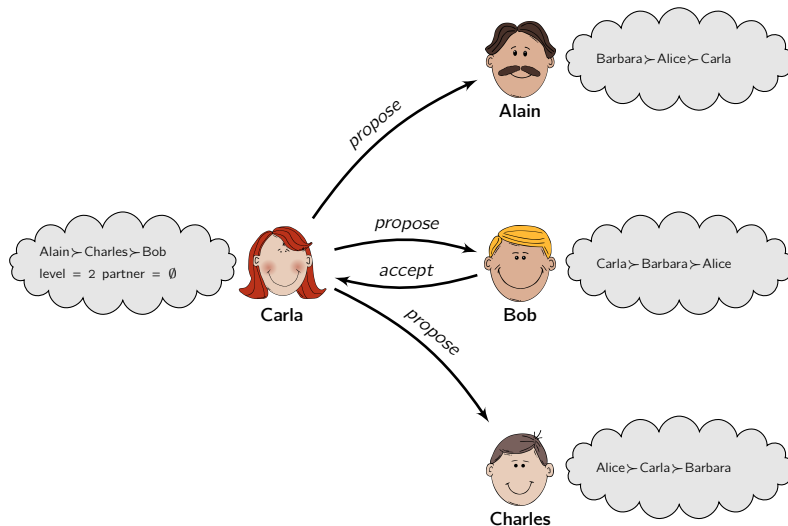
# Casanova Sample Execution



# Casanova Sample Execution

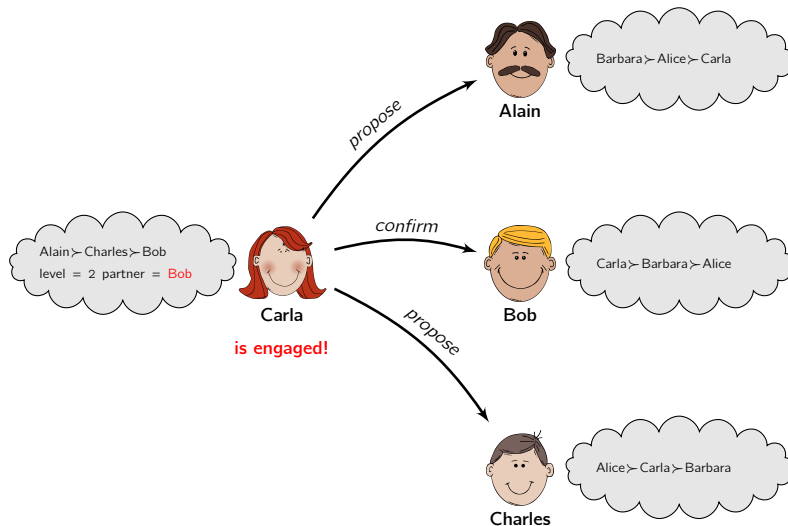


# Casanova Sample Execution

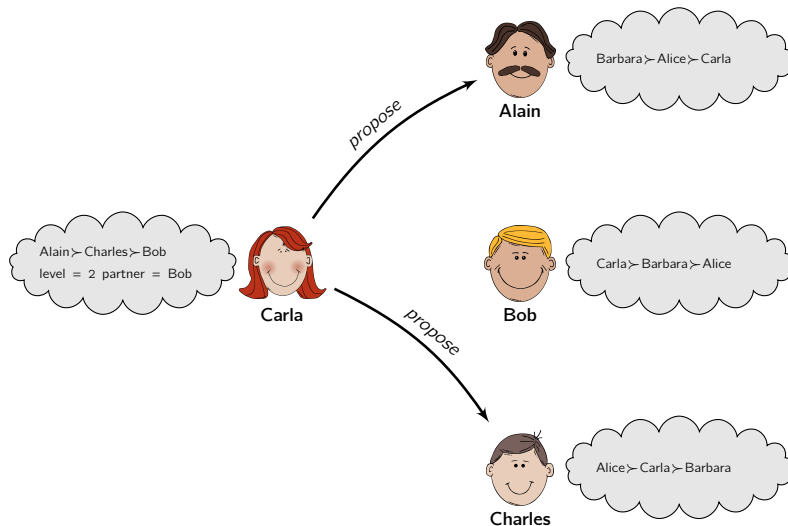




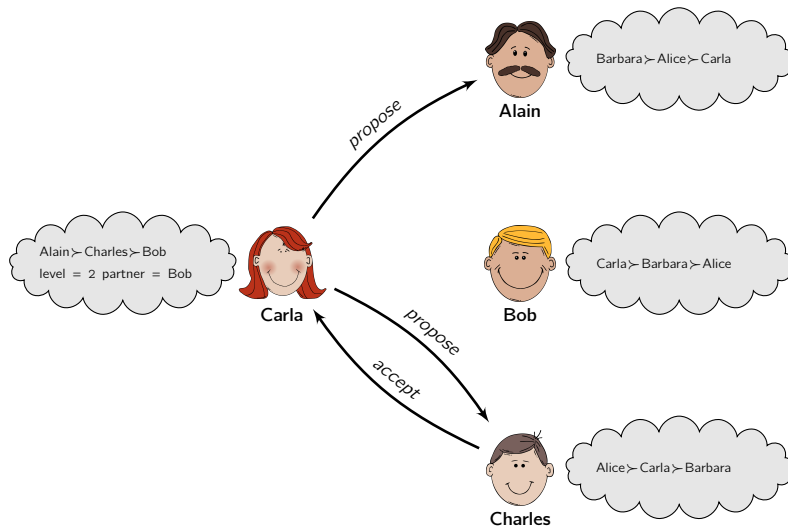
# Casanova Sample Execution



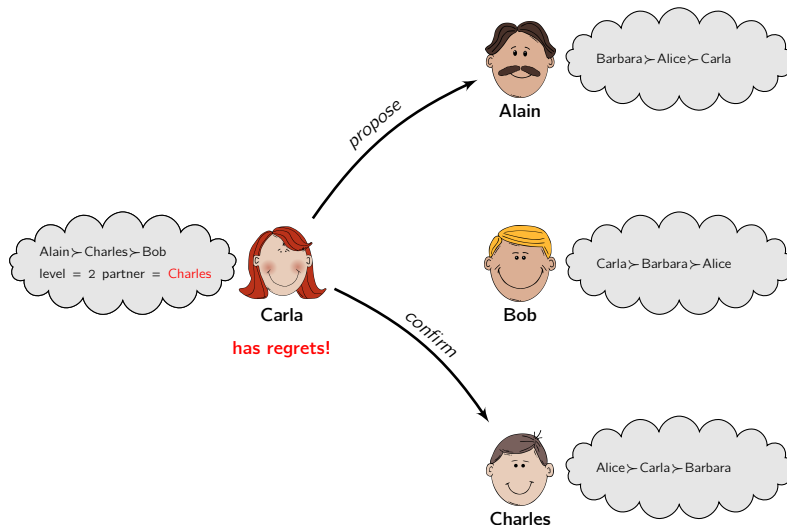
# Casanova Sample Execution



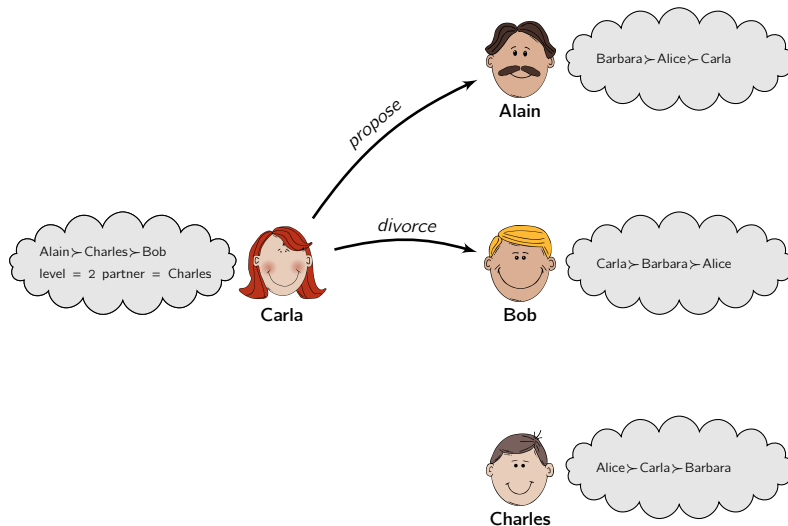
# Casanova Sample Execution



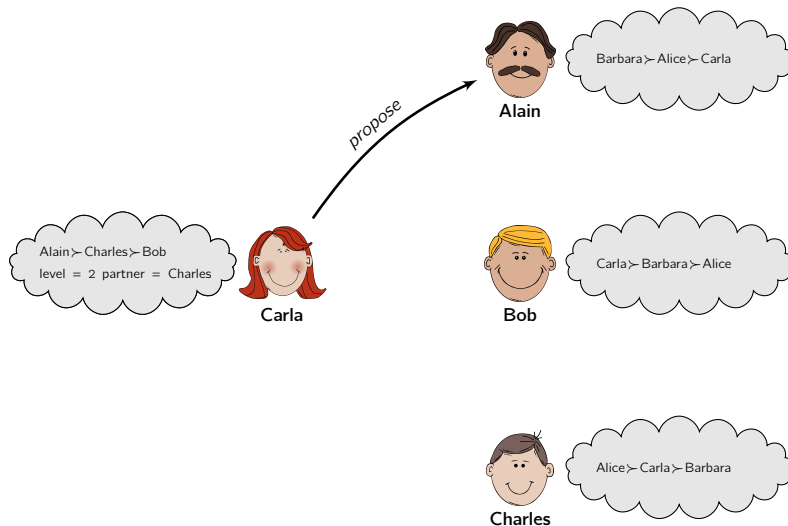
# Casanova Sample Execution



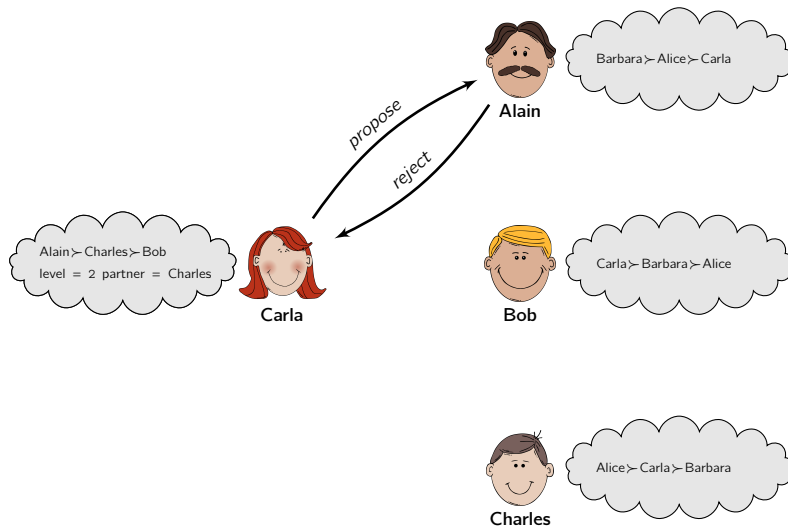
# Casanova Sample Execution



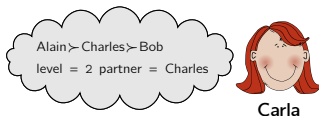
# Casanova Sample Execution



# Casanova Sample Execution

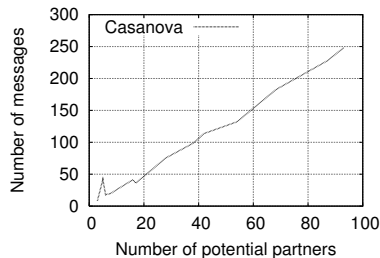
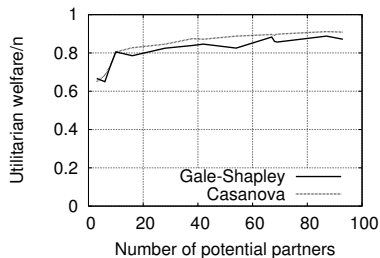
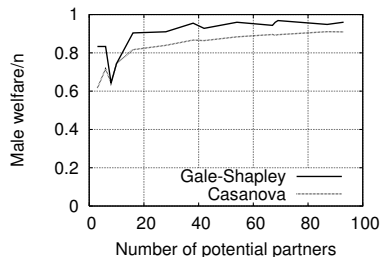
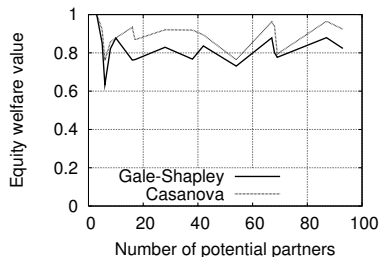


# Casanova Sample Execution





# Evaluation



# Conclusion

## Results

- ▶ Optimal matching = **emergent result** from the local interactions
- ▶ Realisation of the minimal concession strategy
- ▶ Preserves privacy
- ▶ Improves the optimality of the matching and the equity amongst the partners

## Future Works

- ▶ Realize more experiments for evaluating other metrics of social welfare
- ▶ The assignment problem we consider here should be generalized in two directions:
  - ▶ SM with incomplete list and total order (e.g. home swapping)
  - ▶ SM with incomplete list and indifference (e.g. carpool)

# Questions?

# References



K. Arrow, A. Sen, and K. Suzumura. *Handbook of Social Choice and Welfare*. Elsevier, 2002.



I. Brito and P. Meseguer. Distributed stable marriage problem. In *6th Workshop on Distributed Constraint Reasoning at IJCAI 05*, pages 135–147, 2005.



D. Gale and L. S. Shapley. College admissions and the stability of marriage. *American Mathematical Monthly*, 69:9–14, 1962.



Maxime Morge and Paolo Mancarella. Assumption-based argumentation for the minimal concession strategy. In Peter McBurney, Iyad Rahwan, Simon Parsons, and Nicolas Maudet, editors, *Proc. of ArgMAS 2009*, volume 6057 of *LNCS*, pages 114–133. Springer Verlag, Berlin, Germany, 2010.



Jeffrey S. Rosenschein and Gilad Zlotkin. *Rules of encounter: designing conventions for automated negotiation among Computers*. The MIT press series of artificial intelligence. MIT press, 1994.