

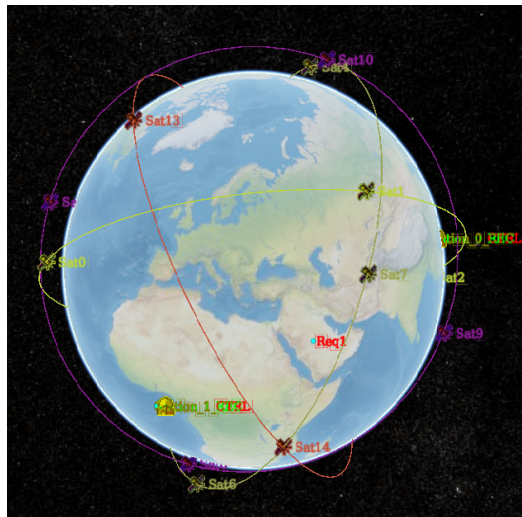
# Orbit Slot Allocation in Earth Observation Constellations

Sara Maqrot   Stéphanie Roussel   **Gauthier Picard**   Cédric Pralet  
ONERA/DTIS, Université de Toulouse, France

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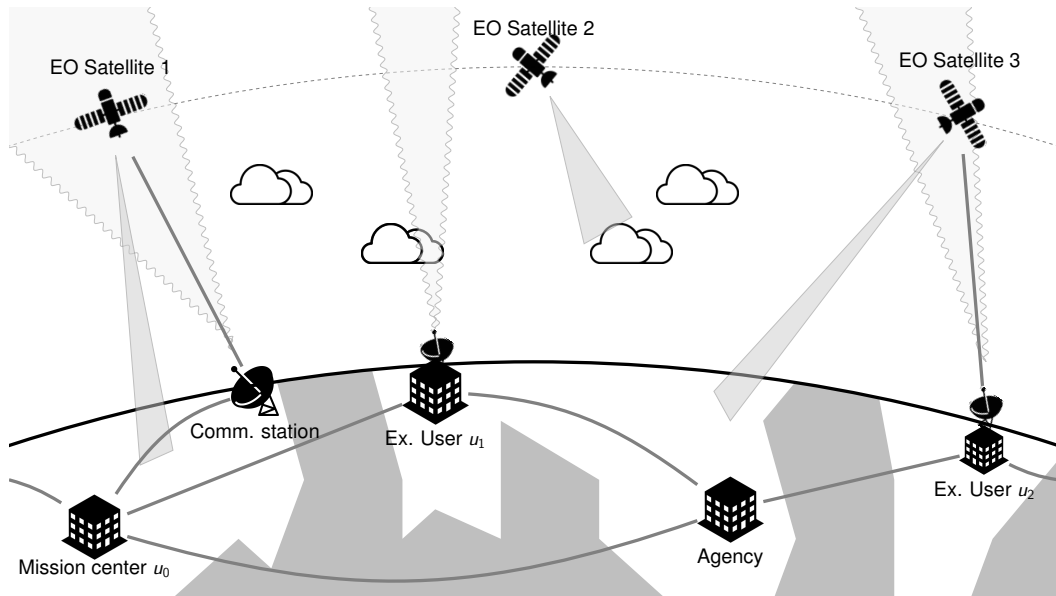
**LiChIE** projet for a new constellation for observing Earth (ADS, INRIA, ONERA, IXBLUE, EREMS)

- **Huge number of image requests** to (a single) mission center
  - More and more **complex requests** (periodic, systematic, etc.)
  - **Overloading** and **over-constraining**
  - **No guarantee** for end-users
- ⇒ New paradigm: **orbit slot ownership**



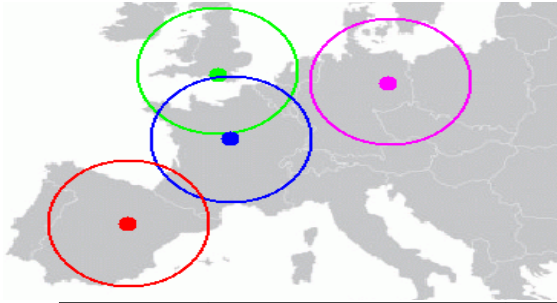
# Introduction

## Applicative Context and Motivation



# Introduction

## Illustrative Example

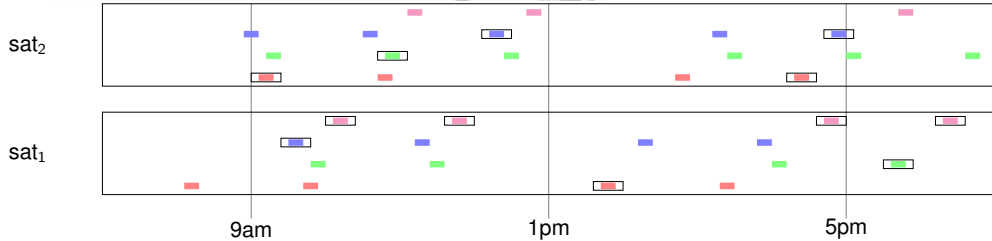


$r_4$  (global-30min)

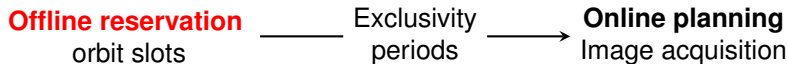
$r_3$  (periodic-4h)

$r_2$  (periodic-4h)

$r_1$  (periodic-4h)

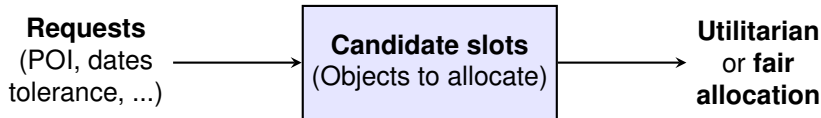


- **Problem** : exploitation of the same constellation by several stakeholders



- **Current allocation scheme**: first come, first served

- **Objective**



# Today's Menu

- 1 Introduction
- 2 Core Concepts and Problem Definition
- 3 Constraint Programming Model
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# General Definitions

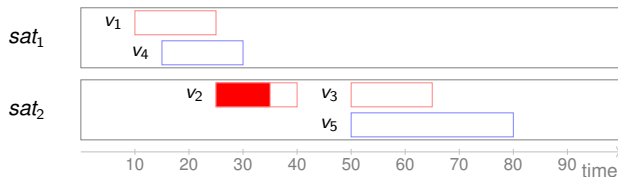
## Satellite reservation window

A *satellite reservation window*  $w$  is defined by

- a satellite  $sat_w$
- a time window  $[start_w, end_w]$
- an individual score  $\omega_w$

## Allocated orbit slot

An *allocated orbit slot*  $o$  within satellite reservation window  $w$  corresponds to a time window  $[start_o, end_o]$  included in  $[start_w, end_w]$



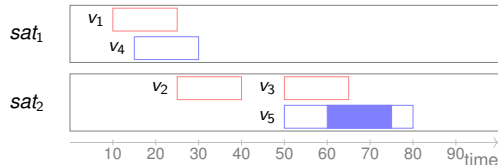
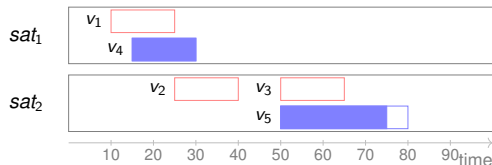


# Global Allocation Request

## Global allocation request

A (multi-mode) *global allocation request*  $r$  is defined by

- a set of satellite reservation windows  $\mathcal{V}_r$
- a minimum duration  $minSlotDur_r$  for each slot
- a list of allocation modes  $\mathcal{M}_r = [\mathcal{M}_{r,1}, \dots, \mathcal{M}_{r,K}]$ , where for all  $m \in \mathcal{M}_r$ 
  - $m$  is an alternative to fulfill  $r$
  - $globalDur_m$  is a global duration required over all orbit slots reserved in  $\mathcal{V}_r$

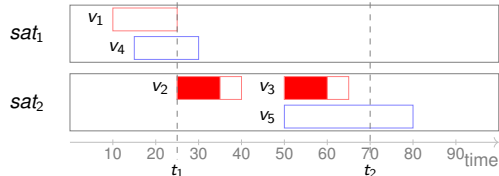
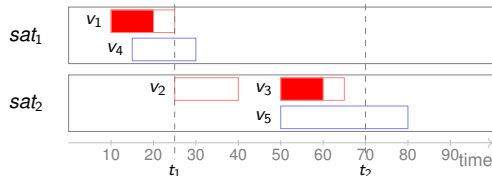


# Time-tagged Allocation Request

## Time-tagged allocation request

A *time-tagged allocation request*  $r$  is defined by

- a set of satellite reservation windows  $\mathcal{V}_r$
- a minimum duration  $minSlotDur_r$  for each slot
- a set of time references  $\mathcal{T}_r$ , with for each time reference  $t \in \mathcal{T}_r$  a subset  $\mathcal{V}_t^r \subseteq \mathcal{V}_r$  that defines the reservation windows associated with  $t$
- a list of allocation modes  $\mathcal{M}_r = [\mathcal{M}_{r,1}, \dots, \mathcal{M}_{r,K}]$ , where each allocation mode  $m \in \mathcal{M}_r$  is defined by a subset  $\mathcal{T}_m \subseteq \mathcal{T}_r$  of time references around which an orbit slot must actually be reserved



# Orbit Slot Allocation Problem

## Orbit Slot Allocation Problem

An *Orbit Slot Allocation Problem* (OSAP) is defined by

- a set of satellites  $\mathcal{S}$
- a set of requests  $\mathcal{R} = \mathcal{R}_G \cup \mathcal{R}_T$ 
  - $\mathcal{R}_G$  a set of global allocation requests
  - $\mathcal{R}_T$  a set of tagged-time allocation requests

## Solution for an OSAP

A *solution*  $\mathcal{A}$  for an OSAP is defined by one allocation  $\mathcal{A}_r$  for each  $r \in \mathcal{R}$

A solution is said to be feasible if and only if

- for every request  $r$ , allocation  $\mathcal{A}_r$  satisfies request  $r$
- for each satellite  $s \in \mathcal{S}$ , there is no overlapping between the orbit slots booked over  $s$  for all allocations in  $\{\mathcal{A}_r \mid r \in \mathcal{R}\}$

# Orbit Slot Allocation Problem

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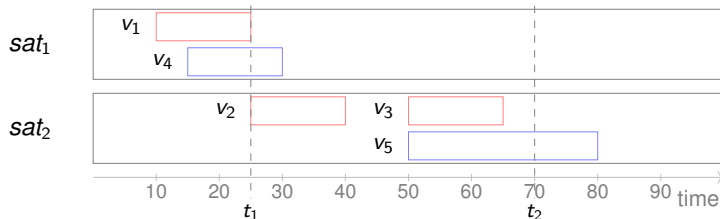
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# Example



- a time-tagged allocation request A (in red)
  - $minSlotDur_A = 10$
  - $\mathcal{V}_A = \{v_1, v_2, v_3\}$  with  $v_1 = [10, 25]$ ,  $v_2 = [25, 40]$ ,  $v_3 = [50, 65]$
  - $\mathcal{V}_{r_1}^{t_1} = \{v_1, v_2\}$  and  $\mathcal{V}_{r_1}^{t_2} = \{v_3\}$
  - 3 modes  $a_1$ ,  $a_2$  and  $a_3$ , with  $\mathcal{T}_{a_1} = \emptyset$ ,  $\mathcal{T}_{a_2} = \{t_1\}$  and  $\mathcal{T}_{a_3} = \{t_1, t_2\}$
- a global allocation request B (in blue)
  - $minSlotDur_B = 15$
  - $\mathcal{V}_B = \{v_4, v_5\}$  with  $v_4 = [15, 30]$ ,  $v_5 = [50, 80]$
  - 3 modes  $b_1$ ,  $b_2$  and  $b_3$ , with global duration of 0, 15 and 40

# How to Assess Allocation Quality?

## Mode reward for global allocation request

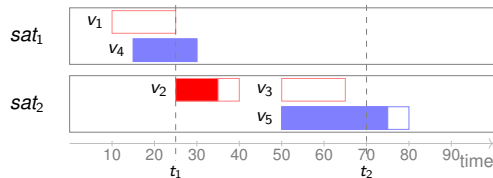
For a global allocation request  $r$  and a possible mode  $m \in \mathcal{M}_r$ , the reward  $\Omega_m$  associated with  $m$  corresponds to quantity  $globalDur_m$

## Mode reward for time-tagged allocation request

For a time-tagged allocation request  $r$  and a possible mode  $m \in \mathcal{M}_r$ , the reward  $\Omega_m$  associated with mode  $m$  corresponds to quantity  $|\mathcal{T}_m| \cdot minSlotDur_r$ , that is to the total satellite time required by  $m$  over all its relevant time references

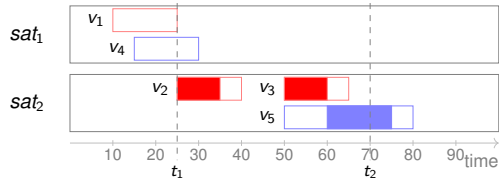
- Mode utility
  - *utilitarian allocation*: maximizing  $u(\mathcal{A}) = \sum_{r \in \mathcal{R}} \Omega_{m(\mathcal{A}_r)}$
  - *fair (leximin) allocation*: maximizing  $\vec{u}(\mathcal{A}) = [\Omega_{m(\mathcal{A}_{r_1})}, \dots, \Omega_{m(\mathcal{A}_{r_n})}]$
- Window utility :  $u^{slot}(\mathcal{A}) = \sum_{r \in \mathcal{R}} \sum_{v \in \mathcal{A}_r} \omega_v$

# Example



(a) Utilitarian-optimal allocation  $\mathcal{A}_{util}$

- modes  $a_2$  for  $A$  and  $b_3$  for  $B$
- $u(\mathcal{A}_{util}) = 10 + 40 = 50$
- $\vec{u}(\mathcal{A}_{util}) = [10, 40]$



(b) Leximin-optimal allocation  $\mathcal{A}_{fair}$

- modes  $a_3$  for  $A$  and  $b_2$  for  $B$
- $u(\mathcal{A}_{util}) = 20 + 15 = 35$
- $\vec{u}(\mathcal{A}_{util}) = [15, 20]$

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# Utilitarian CP Encoding

$$\begin{aligned} & \text{maximize} && \sum_{r \in \mathcal{R}} \sum_{m \in \mathcal{M}_r} \Omega_m \cdot x_m \\ & \text{s.t.} && (2), (3), (4), (5), (6) \end{aligned} \quad (1)$$

- $itv_v$ : interval variables with minimum size of  $minSlotDur_r$
- $x_m$ : boolean variable for choosing mode  $m$

$$\forall r \in \mathcal{R}, \quad \sum_{m \in \mathcal{M}_r} x_m = 1 \quad (2)$$

$$\forall s \in S, \quad noOverlap(\{itv_v | v \in \bigcup_{r \in \mathcal{R}} \mathcal{V}_r \wedge sat_v = s\}) \quad (3)$$

$$\forall r \in \mathcal{R}_T, \forall m \in \mathcal{M}_r, \forall t \in \mathcal{T}_m, \quad \sum_{v \in \mathcal{V}_r^t} presenceOf(itv_v) \geq x_m, \quad (4)$$

$$\forall r \in \mathcal{R}_T, \forall t \in \mathcal{T}_r, \quad \sum_{v \in \mathcal{V}_t} presenceOf(itv_v) \leq 1 \quad (5)$$

$$\forall r \in \mathcal{R}_G, \forall m \in \mathcal{M}_r, \quad \sum_{v \in \mathcal{V}_r} lengthOf(itv_v) \geq x_m \cdot globalDur_r \quad (6)$$

# Leximin CP Encoding

## Intuition

- Solve as many CP optimization problems as requests
  - The  $k$ -th CP problem allows to compute the  $k$ -th component of the sorted leximin vector  $\vec{u} = [u_1, \dots, u_n]$
  - The objective is to lexicographically maximize vector  $\Lambda = [\Lambda_1, \dots, \Lambda_n]$  obtained after ordering  $[u_1, \dots, u_n]$  following an increasing order
- 
- Variables
    - $\lambda \in [\Lambda_{K-1}, \max_{r \in \mathcal{R}} Z_r]$  is a real variable representing the utility obtained at level  $K$  in  $\Lambda$
    - $y_{rk}$  is a binary variable equal to 1 if request  $r \in \mathcal{R}$  plays the role of the request associated with level  $k \in [1..K-1]$  in  $[\Lambda_1, \dots, \Lambda_{K-1}]$ , 0 otherwise
    - $u_r$  is a real variable in  $[0, Z_r]$  representing the utility of request  $r$

# Leximin CP Encoding (cont.)

- Determining the  $K$ th level

$$\text{maximize } \lambda \quad (7)$$

$$\text{s.t. } (2), (3), (4), (5), (6)$$

$$\forall r \in \mathcal{R}, \quad u_r = \sum_{m \in \mathcal{M}_r} \Omega_m \cdot x_m \quad (8)$$

$$\forall k \in [1..K-1], \quad \sum_{r \in \mathcal{R}} y_{rk} = 1 \quad (9)$$

$$\forall r \in \mathcal{R}, \quad \sum_{k \in [1..K-1]} y_{rk} \leq 1 \quad (10)$$

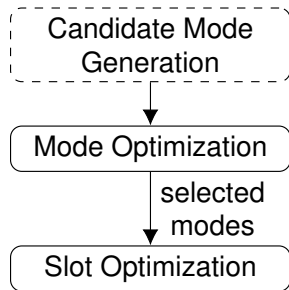
$$\forall r \in \mathcal{R}, \quad \lambda \leq u_r + M \sum_{k \in [1..K-1]} y_{rk} \quad (11)$$

$$\forall r \in \mathcal{R}, \quad u_r \geq \sum_{k \in [1..K-1]} \Lambda_k \cdot y_{rk} \quad (12)$$

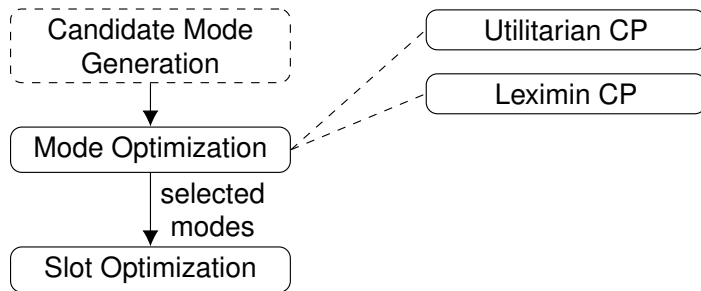
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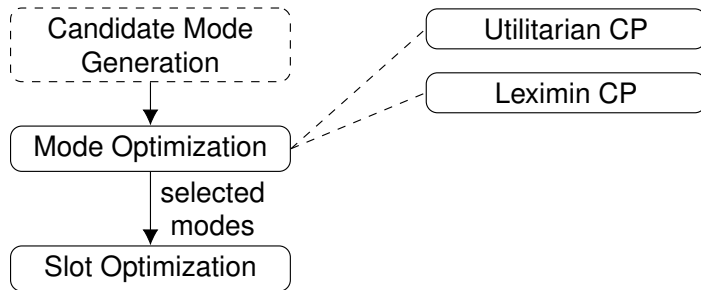
# Optimization Architecture



# Optimization Architecture

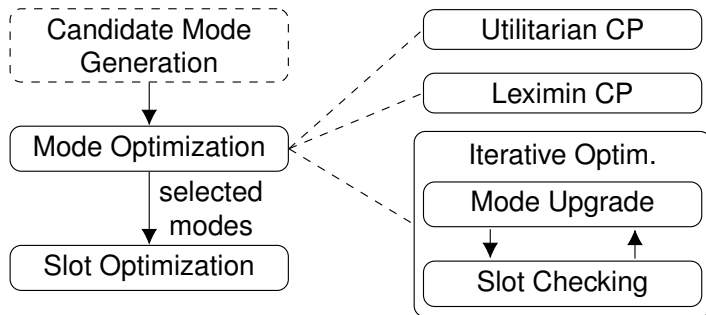


# Optimization Architecture



- Major issue: optimal CP approaches won't scale up!

# Optimization Architecture

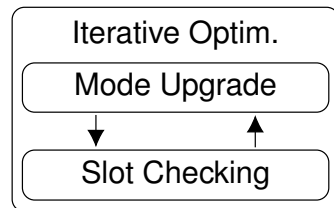


- Major issue: optimal CP approaches won't scale up!
- Proposal: iterative heuristic approach based on mode quality upgrade

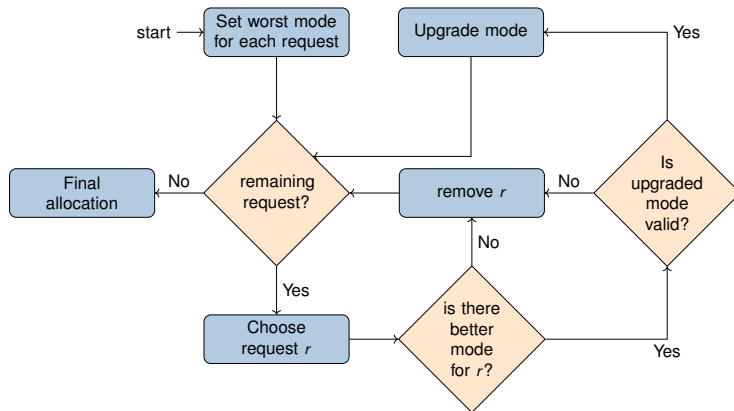


# Iterative Optimization

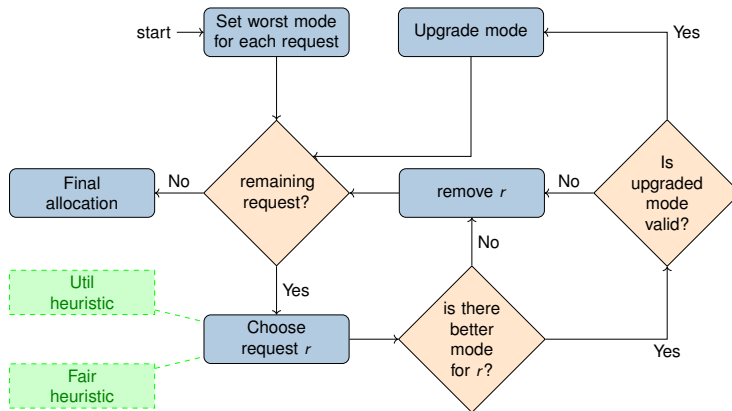
- **Purposes:** balance between utilitarianism and fairness, and scale up
- **Mode Upgrade** produces allocations so that
  - Constraint (2) is satisfied
  - Criteria (1) and (7) are optimized
- **Slot Checking** layer checks Constraints (2)– (6)



# Iterative Optimization (cont.)



# Iterative Optimization (cont.)



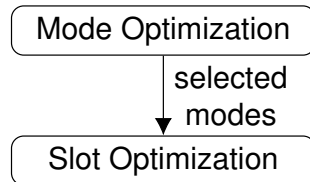
- $h^{util}$  selects the request whose next mode increases the most the global utility of the allocation
- $h^{fair}$  selects the request with the smallest utility

# Slot Optimization

- **Purpose:** Optimizing the **slots** for the modes that have been selected at the Mode Optimization step
- **Maximizing the window utility**

$$\begin{aligned} &\text{maximize} && \sum_{r \in \mathcal{R}} \sum_{v \in \mathcal{V}_r} \omega_v \cdot \text{presenceOf}(\text{itv}_v) \quad (13) \\ &\text{s.t.} && (2), (3), (4), (5), (6) \end{aligned}$$

with  $x_m = 1$  iff  $m$  is selected by the Mode Optimization module



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# Experimental Setup

## Constellation

- Low-Earth Orbit (500km altitude)
- 8 mono-satellite orbital planes with a 60 degrees inclination

## Requests

- ① randomly selecting a subset of national capitals  $\mathcal{C}$
- ② randomly parameterize requests
  - Global requests
    - most preferred global duration in  $[2h, 4h]$
    - less preferred modes remove 30min down to 0 ( $\mathcal{M}_{r,1}$ )
    - minimum slot duration in  $[2\text{min}, 4\text{min}]$
  - Time-tagged requests
    - 2 time references patterns:  $[8\text{am}, 12\text{pm}, 4\text{pm}, 8\text{pm}]$  and  $[9\text{am}, 1\text{pm}, 5\text{pm}]$
    - 1-hour tolerance:  $[t - 1\text{hour}, t + 1\text{hour}]$
    - less preferred mode  $\mathcal{M}_{r,1}$  has an empty set of time references
    - more preferred modes add one random time reference
- ③ randomly picking a national capital in  $\mathcal{C}$  as a ground station for each request

# Experimental Setup (cont.)

## Orderbooks

- 5 different order book instances per configuration (defined by  $|\mathcal{R}_G|$  and  $|\mathcal{R}_T|$ )
- For slot optimization, the reward is linear in  $[0,1]$
- The number of requests we consider is larger than current realistic data

## Computing Environment

- Solvers are coded in Java 1.8
- Execution on 20-core Intel(R) Xeon(R) CPU E5-2660 v3 @ 2.60GHz, 62GB RAM, Ubuntu 18.04.5 LTS
- CP Optimizer included in IBM ILOG CPLEX Studio 20.1 is used by the solvers through the Java API, with timeouts

Method	Mode opt.	Slot opt.	Slot check
upgrade – util	n/a	300s	120s
upgrade – fair	n/a	300s	300s
cp – util / cp – fair	$300s \times  \mathcal{R} $	300s	300s

configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
$ \mathcal{R}_G $	$ \mathcal{R}_T $	$ \mathcal{M} $	$ \mathcal{V} $	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$
0	5	22.0	107.0	<b>1980.20</b> †	<b>4.43</b>	<b>1980.20</b> *	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>
0	10	44.6	218.2	3925.00†	<b>8.85</b>	<b>3953.40</b> *	8.66	3925.00	<b>8.85</b>	<b>3953.40</b>	8.66
0	15	67.2	326.2	6260.40†	13.15	<b>6288.80</b> *	12.96	6260.40	<b>13.16</b>	<b>6288.80</b>	12.96
0	20	90.0	439.6	8294.00†	<b>17.27</b>	<b>8322.40</b>	17.06	8294.00	17.25	<b>8322.40</b>	17.03
0	25	112.0	549.8	10313.20	21.09	<b>10341.60</b>	20.94	10313.20	<b>21.16</b>	10276.60	20.78
5	0	31.4	198.6	39874.00	<b>4.64</b>	39911.20	4.50	<b>42394.00</b>	4.31	42034.00	4.31
10	0	63.8	405.0	<b>44646.60</b>	9.20	42953.60	8.20	44286.60	<b>9.32</b>	44286.60	9.27
15	0	96.4	606.2	42109.20	13.29	42730.20	10.76	44291.60	13.51	<b>44420.00</b>	<b>13.90</b>
20	0	129.6	814.6	27927.20	9.80	40992.60	9.32	43131.20	<b>14.14</b>	<b>43409.00</b>	13.86
25	0	161.4	1018.2	28864.80	9.80	40489.20	9.30	39645.40	<b>13.87</b>	<b>43117.40</b>	13.23
5	5	53.8	311.0	39515.60	<b>8.97</b>	40998.60	8.43	42395.60	8.30	<b>44388.00</b>	7.23
10	10	109.6	627.0	43594.40	15.93	42664.40	15.52	44674.40	<b>15.98</b>	<b>47071.60</b>	13.98
15	15	165.2	944.0	34171.40	<b>23.29</b>	39015.00	20.00	46368.80	21.74	<b>47244.60</b>	19.10
20	20	219.4	1258.2	31823.00	<b>26.31</b>	41759.80	24.18	45223.60	25.16	<b>47728.40</b>	19.45
25	25	274.6	1572.8	29788.40	<b>28.87</b>	41641.00	26.05	46824.60	27.16	<b>47474.80</b>	21.24



configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
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25	0	161.4	1018.2	28864.80	9.80	40489.20	9.30	39645.40	<b>13.87</b>	<b>43117.40</b>	13.23
5	5	53.8	311.0	39515.60	<b>8.97</b>	40998.60	8.43	42395.60	8.30	<b>44388.00</b>	7.23
10	10	109.6	627.0	43594.40	15.93	42664.40	15.52	44674.40	<b>15.98</b>	<b>47071.60</b>	13.98
15	15	165.2	944.0	34171.40	<b>23.29</b>	39015.00	20.00	46368.80	21.74	<b>47244.60</b>	19.10
20	20	219.4	1258.2	31823.00	<b>26.31</b>	41759.80	24.18	45223.60	25.16	<b>47728.40</b>	19.45
25	25	274.6	1572.8	29788.40	<b>28.87</b>	41641.00	26.05	46824.60	27.16	<b>47474.80</b>	21.24

Time-tagged-only allocation requests:  
cp – util and upgrade – util provide the best mode utilitarian allocation

configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
$ \mathcal{R}_G $	$ \mathcal{R}_T $	$ \mathcal{M} $	$ \mathcal{V} $	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$
0	5	22.0	107.0	<b>1980.20†</b>	<b>4.43</b>	<b>1980.20*</b>	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>
0	10	44.6	218.2	3925.00†	<b>8.85</b>	<b>3953.40*</b>	8.66	3925.00	<b>8.85</b>	<b>3953.40</b>	8.66
0	15	67.2	326.2	6260.40†	13.15	<b>6288.80*</b>	12.96	6260.40	<b>13.16</b>	<b>6288.80</b>	12.96
0	20	90.0	439.6	8294.00†	<b>17.27</b>	<b>8322.40*</b>	17.06	8294.00	17.25	<b>8322.40</b>	17.03
0	25	112.0	549.8	10313.20	21.09	<b>10341.60</b>	20.94	10313.20	<b>21.16</b>	10276.60	20.78
5	0	31.4	198.6	39874.00	<b>4.64</b>	39911.20	4.50	<b>42394.00</b>	4.31	42034.00	4.31
10	0	63.8	405.0	<b>44646.60</b>	9.20	42953.60	8.20	44286.60	<b>9.32</b>	44286.60	9.27
15	0	96.4	606.2	42109.20	13.29	42730.20	10.76	44291.60	13.51	<b>44420.00</b>	<b>13.90</b>
20	0	129.6	814.6	27927.20	9.80	40992.60	9.32	43131.20	<b>14.14</b>	<b>43409.00</b>	13.86
25	0	161.4	1018.2	28864.80	9.80	40489.20	9.30	39645.40	<b>13.87</b>	<b>43117.40</b>	13.23
5	5	53.8	311.0	39515.60	<b>8.97</b>	40998.60	8.43	42395.60	8.30	<b>44388.00</b>	7.23
10	10	109.6	627.0	43594.40	15.93	42664.40	15.52	44674.40	<b>15.98</b>	<b>47071.60</b>	13.98
15	15	165.2	944.0	34171.40	<b>23.29</b>	39015.00	20.00	46368.80	21.74	<b>47244.60</b>	19.10
20	20	219.4	1258.2	31823.00	<b>26.31</b>	41759.80	24.18	45223.60	25.16	<b>47728.40</b>	19.45
25	25	274.6	1572.8	29788.40	<b>28.87</b>	41641.00	26.05	46824.60	27.16	<b>47474.80</b>	21.24

For small instances:

Utility-optimal and fairness-optimal allocations returned by cp – util and cp – fair

configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
$ \mathcal{R}_G $	$ \mathcal{R}_T $	$ \mathcal{M} $	$ \mathcal{V} $	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$
0	5	22.0	107.0	<b>1980.20†</b>	<b>4.13</b>	<b>1980.20*</b>	<b>4.13</b>	<b>1980.20</b>	<b>4.13</b>	<b>1980.20</b>	<b>4.13</b>
0	10	44.6	218.2	3925.00†	<b>8.85</b>	<b>3953.40*</b>	8.66	3925.00	<b>8.85</b>	<b>3953.40</b>	8.66
0	15	67.2	326.2	6260.40†	13.15	<b>6288.80*</b>	12.96	6260.40	<b>13.16</b>	<b>6288.80</b>	12.96
0	20	90.0	439.6	8294.00†	<b>17.27</b>	<b>8322.40</b>	17.06	8294.00	17.25	<b>8322.40</b>	17.03
0	25	112.0	549.8	10313.20	21.09	<b>10341.60</b>	20.94	10313.20	<b>21.16</b>	10276.60	20.78
5	0	31.4	198.6	39874.00	<b>4.64</b>	39911.20	4.50	<b>42394.00</b>	4.31	42034.00	4.31
10	0	63.8	405.0	<b>44646.60</b>	9.20	42953.60	8.20	44286.60	<b>9.32</b>	44286.60	9.27
15	0	96.4	606.2	42109.20	13.29	42730.20	10.76	44291.60	13.51	<b>44420.00</b>	<b>13.90</b>
20	0	129.6	814.6	27927.20	9.80	40992.60	9.32	43131.20	<b>14.14</b>	<b>43409.00</b>	13.86
25	0	161.4	1018.2	28864.80	9.80	40489.20	9.30	39645.40	<b>13.87</b>	<b>43117.40</b>	13.23
5	5	53.8	311.0	39515.60	<b>8.97</b>	40998.60	8.43	42395.60	8.30	<b>44388.00</b>	7.23
10	10	109.6	627.0	43594.40	15.93	42664.40	15.52	44674.40	<b>15.98</b>	<b>47071.60</b>	13.98
15	15	165.2	944.0	34171.40	<b>23.29</b>	39015.00	20.00	46368.80	21.74	<b>47244.60</b>	19.10
20	20	219.4	1258.2	31823.00	<b>26.31</b>	41759.80	24.18	45223.60	25.16	<b>47728.40</b>	19.45
25	25	274.6	1572.8	29788.40	<b>28.87</b>	41641.00	26.05	46824.60	27.16	<b>47474.80</b>	21.24

With more than 5 time-tagged requests:  
utilitarian and fair approaches converge to different optima

configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
$ \mathcal{R}_G $	$ \mathcal{R}_T $	$ \mathcal{M} $	$ \mathcal{V} $	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$
0	5	22.0	107.0	<b>1980.20†</b>	<b>4.43</b>	<b>1980.20*</b>	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>
0	10	44.6	218.2	3925.00†	<b>8.85</b>	<b>3953.40*</b>	8.66	3925.00	<b>8.85</b>	<b>3953.40</b>	8.66
0	15	67.2	326.2	6260.40†	13.15	<b>6288.80*</b>	12.96	6260.40	<b>13.16</b>	<b>6288.80</b>	12.96
0	20	90.0	439.6	8294.00†	<b>17.27</b>	<b>8322.40</b>	17.06	8294.00	17.25	<b>8322.40</b>	17.03
0	25	112.0	549.8	10313.20	21.09	<b>10341.60</b>	20.94	10313.20	<b>21.16</b>	10276.60	20.78
5	0	31.4	198.6	39874.00	<b>4.64</b>	39911.20	4.50	<b>42394.00</b>	4.31	42034.00	4.31
10	0	63.8	405.0	<b>44646.60</b>	9.20	42953.60	8.20	44286.60	<b>9.32</b>	44286.60	9.27
15	0	96.4	606.2	42109.20	13.29	42730.20	10.76	44291.60	13.51	<b>44420.00</b>	<b>13.90</b>
20	0	129.6	814.6	27927.20	9.80	40992.60	9.32	43131.20	<b>14.14</b>	<b>43409.00</b>	13.86
25	0	161.4	1018.2	28864.80	9.80	40489.20	9.30	39645.40	<b>13.87</b>	<b>43117.40</b>	13.23
5	5	53.8	311.0	39515.60	<b>8.97</b>	40998.60	8.43	42395.60	8.30	<b>44388.00</b>	7.23
10	10	109.6	627.0	43594.40	15.93	42664.40	15.52	44674.40	<b>15.98</b>	<b>47071.60</b>	13.98
15	15	165.2	944.0	34171.40	<b>23.29</b>	39015.00	20.00	46368.80	21.74	<b>47244.60</b>	19.10
20	20	219.4	1258.2	31823.00	<b>26.31</b>	41759.80	24.18	45223.60	25.16	<b>47728.40</b>	19.45
25	25	274.6	1572.8	29788.40	<b>28.87</b>	41641.00	26.05	46824.60	27.16	<b>47474.80</b>	21.24

With global allocation requests:  
Higher reward but difficult to solve by cp – fair and cp – util

configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
$ \mathcal{R}_G $	$ \mathcal{R}_T $	$ \mathcal{M} $	$ \mathcal{V} $	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$
0	5	22.0	107.0	<b>1980.20†</b>	<b>4.43</b>	<b>1980.20*</b>	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>
0	10	44.6	218.2	3925.00†	<b>8.85</b>	<b>3953.40*</b>	8.66	3925.00	<b>8.85</b>	<b>3953.40</b>	8.66
0	15	67.2	326.2	6260.40†	13.15	<b>6288.80*</b>	12.96	6260.40	<b>13.16</b>	<b>6288.80</b>	12.96
0	20	90.0	439.6	8294.00†	<b>17.27</b>	<b>8322.40</b>	17.06	8294.00	17.25	<b>8322.40</b>	17.03
0	25	112.0	549.8	10313.20	21.09	<b>10341.60</b>	20.94	10313.20	<b>21.16</b>	10276.60	20.78
5	0	31.4	198.6	39874.00	<b>4.64</b>	39911.20	4.50	<b>42394.00</b>	4.31	42034.00	4.31
10	0	63.8	405.0	<b>44646.60</b>	9.20	42953.60	8.20	44286.60	<b>9.32</b>	44286.60	9.27
15	0	96.4	606.2	42109.20	13.29	42730.20	10.76	44291.60	13.51	<b>44420.00</b>	<b>13.90</b>
20	0	129.6	814.6	27927.20	9.80	40992.60	9.32	43131.20	<b>14.14</b>	<b>43409.00</b>	13.86
25	0	161.4	1018.2	28864.80	9.80	40489.20	9.30	39645.40	<b>13.87</b>	<b>43117.40</b>	13.23
5	5	53.8	311.0	39515.60	<b>8.97</b>	40998.60	8.43	42395.60	8.30	<b>44388.00</b>	7.23
10	10	109.6	627.0	43594.40	15.93	42664.40	15.52	44674.40	<b>15.98</b>	<b>47071.60</b>	13.98
15	15	165.2	944.0	34171.40	<b>23.29</b>	39015.00	20.00	46368.80	21.74	<b>47244.60</b>	19.10
20	20	219.4	1258.2	31823.00	<b>26.31</b>	41759.80	24.18	45223.60	25.16	<b>47728.40</b>	19.45
25	25	274.6	1572.8	29788.40	<b>28.87</b>	41641.00	26.05	46824.60	27.16	<b>47474.80</b>	21.24

upgrade – util **performs better in terms of utility** in most of the settings

configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
$ \mathcal{R}_G $	$ \mathcal{R}_T $	$ \mathcal{M} $	$ \mathcal{V} $	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$	$u$	$u^{slot}$
0	5	22.0	107.0	<b>1980.20†</b>	<b>4.43</b>	<b>1980.20*</b>	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>	<b>1980.20</b>	<b>4.43</b>
0	10	44.6	218.2	3925.00†	<b>8.85</b>	<b>3953.40*</b>	8.66	3925.00	<b>8.85</b>	<b>3953.40</b>	8.66
0	15	67.2	326.2	6260.40†	13.15	<b>6288.80*</b>	12.96	6260.40	<b>13.16</b>	<b>6288.80</b>	12.96
0	20	90.0	439.6	8294.00†	<b>17.27</b>	<b>8322.40</b>	17.06	8294.00	17.25	<b>8322.40</b>	17.03
0	25	112.0	549.8	10313.20	21.09	<b>10341.60</b>	20.94	10313.20	<b>21.16</b>	10276.60	20.78
5	0	31.4	198.6	39874.00	<b>4.64</b>	39911.20	4.50	<b>42394.00</b>	4.31	42034.00	4.31
10	0	63.8	405.0	<b>44646.60</b>	9.20	42953.60	8.20	44286.60	<b>9.32</b>	44286.60	9.27
15	0	96.4	606.2	42109.20	13.29	42730.20	10.76	44291.60	13.51	<b>44420.00</b>	<b>13.90</b>
20	0	129.6	814.6	27927.20	9.80	40992.60	9.32	43131.20	<b>14.14</b>	<b>43409.00</b>	13.86
25	0	161.4	1018.2	28864.80	9.80	40489.20	9.30	39645.40	<b>13.87</b>	<b>43117.40</b>	13.23
5	5	53.8	311.0	39515.60	<b>8.97</b>	40998.60	8.43	42395.60	8.30	<b>44388.00</b>	7.23
10	10	109.6	627.0	43594.40	15.93	42664.40	15.52	44674.40	<b>15.98</b>	<b>47071.60</b>	13.98
15	15	165.2	944.0	34171.40	<b>23.29</b>	39015.00	20.00	46368.80	21.74	<b>47244.60</b>	19.10
20	20	219.4	1258.2	31823.00	<b>26.31</b>	41759.80	24.18	45223.60	25.16	<b>47728.40</b>	19.45
25	25	274.6	1572.8	29788.40	<b>28.87</b>	41641.00	26.05	46824.60	27.16	<b>47474.80</b>	21.24

upgrade – fair **outputs quite good utilitarian allocations**

# Computation Time (ms)

configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
$ \mathcal{R}_G $	$ \mathcal{R}_T $	$ \mathcal{M} $	$ \mathcal{V} $	mode	slot	mode	slot	mode	slot	mode	slot
0	5	22.0	107.0	6.34	2.09	12.96	<b>1.63</b>	<b>4.24</b>	1.96	4.89	2.32
0	10	44.6	218.2	77.79	3.93	45.96	3.94	<b>9.73</b>	4.00	10.45	<b>3.79</b>
0	15	67.2	326.2	164.48	242.04	439.87	242.63	<b>9.41</b>	243.74	9.63	<b>241.06</b>
0	20	90.0	439.6	195.42	300.38	6000.21	300.09	<b>9.50</b>	300.13	10.19	300.36
0	25	112.0	549.8	759.76	300.16	7500.19	300.09	<b>237.50</b>	300.25	294.00	300.25
5	0	31.4	198.6	1500.57	300.04	1500.14	300.06	390.88	300.03	<b>339.86</b>	300.04
10	0	63.8	405.0	3002.12	300.08	3700.13	300.13	<b>1249.88</b>	300.06	1279.68	300.06
15	0	96.4	606.2	4502.34	300.06	4500.43	300.12	1849.59	300.11	<b>1834.56</b>	300.05
20	0	129.6	814.6	6003.48	300.07	6000.23	300.07	<b>2496.88</b>	300.08	2526.24	300.08
25	0	161.4	1018.2	7504.20	300.07	7500.22	300.07	3137.79	300.10	<b>3074.31</b>	300.09
5	5	53.8	311.0	1504.79	300.05	3000.13	300.04	<b>420.16</b>	300.09	433.25	300.06
10	10	109.6	627.0	3029.96	300.08	6000.21	300.09	<b>1206.61</b>	300.08	2026.38	300.06
15	15	165.2	944.0	8926.65	300.09	9000.24	300.14	<b>2410.81</b>	300.10	3530.65	300.08
20	20	219.4	1258.2	12011.00	300.15	12000.43	300.12	<b>3438.31</b>	300.16	4494.19	300.10
25	25	274.6	1572.8	15014.64	300.19	15000.44	300.14	<b>5187.83</b>	300.17	6019.73	300.11

# Computation Time (ms)

configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
$ \mathcal{R}_G $	$ \mathcal{R}_T $	$ \mathcal{M} $	$ \mathcal{V} $	mode	slot	mode	slot	mode	slot	mode	slot
0	5	22.0	107.0	6.34	2.09	12.96	<b>1.63</b>	<b>4.24</b>	1.96	4.89	2.32
0	10	44.6	218.2	77.79	3.93	45.96	3.94	<b>9.73</b>	4.00	10.45	<b>3.79</b>
0	15	67.2	326.2	164.48	242.04	439.87	242.63	<b>9.41</b>	243.74	9.63	<b>241.06</b>
0	20	90.0	439.6	195.42	300.38	6000.21	300.09	<b>9.50</b>	300.13	10.19	300.36
0	25	112.0	549.8	759.76	300.16	7500.19	300.09	<b>237.50</b>	300.25	294.00	300.25
5	0	31.4	198.6	1500.57	300.04	1500.14	300.06	390.88	300.03	<b>339.86</b>	300.04
10	0	63.8	405.0	3002.12	300.08	3700.13	300.13	<b>1249.88</b>	300.06	1279.68	300.06
15	0	96.4	606.2	4502.34	300.06	4500.43	300.12	1849.59	300.11	<b>1834.56</b>	300.05
20	0	129.6	814.6	6003.48	300.07	6000.23	300.07	<b>2496.88</b>	300.08	2526.24	300.08
25	0	161.4	1018.2	7504.20	300.07	7500.22	300.07	3137.79	300.10	<b>3074.31</b>	300.09
5	5	53.8	311.0	1504.79	300.05	3000.13	300.04	<b>420.16</b>	300.09	433.25	300.06
10	10	109.6	627.0	3029.96	300.08	6000.21	300.09	<b>1206.61</b>	300.08	2026.38	300.06
15	15	165.2	944.0	8926.65	300.09	9000.24	300.14	<b>2410.81</b>	300.10	3530.65	300.08
20	20	219.4	1258.2	12011.00	300.15	12000.43	300.12	<b>3438.31</b>	300.16	4494.19	300.10
25	25	274.6	1572.8	15014.64	300.19	15000.44	300.14	<b>5187.83</b>	300.17	6019.73	300.11

Upgrade-based heuristic methods clearly outperform optimal ones



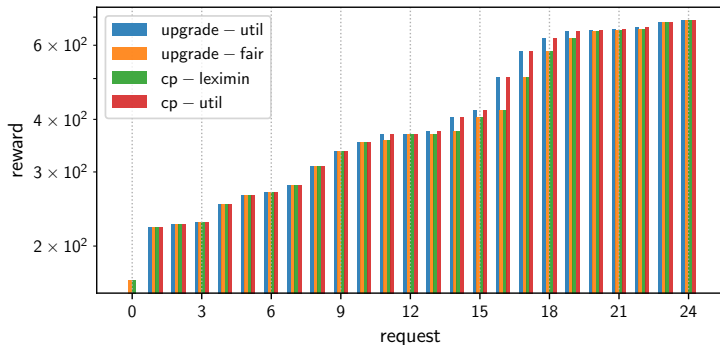
# Computation Time (ms)

configurations				cp – fair		cp – util		upgrade – fair		upgrade – util	
$ \mathcal{R}_G $	$ \mathcal{R}_T $	$ \mathcal{M} $	$ \mathcal{V} $	mode	slot	mode	slot	mode	slot	mode	slot
0	5	22.0	107.0	6.34	2.09	12.96	<b>1.63</b>	<b>4.24</b>	1.96	4.89	2.32
0	10	44.6	218.2	77.79	3.93	45.96	3.94	<b>9.73</b>	4.00	10.45	<b>3.79</b>
0	15	67.2	326.2	164.48	242.04	439.87	242.63	<b>9.41</b>	243.74	9.63	<b>241.06</b>
0	20	90.0	439.6	195.42	300.38	6000.21	300.09	<b>9.50</b>	300.13	10.19	300.36
0	25	112.0	549.8	759.76	300.16	7500.19	300.09	<b>237.50</b>	300.25	294.00	300.25
5	0	31.4	198.6	1500.57	300.04	1500.14	300.06	390.88	300.03	<b>339.86</b>	300.04
10	0	63.8	405.0	3002.12	300.08	3700.13	300.13	<b>1249.88</b>	300.06	1279.68	300.06
15	0	96.4	606.2	4502.34	300.06	4500.43	300.12	1849.59	300.11	<b>1834.56</b>	300.05
20	0	129.6	814.6	6003.48	300.07	6000.23	300.07	<b>2496.88</b>	300.08	2526.24	300.08
25	0	161.4	1018.2	7504.20	300.07	7500.22	300.07	3137.79	300.10	<b>3074.31</b>	300.09
5	5	53.8	311.0	1504.79	300.05	3000.13	300.04	<b>420.16</b>	300.09	433.25	300.06
10	10	109.6	627.0	3029.96	300.08	6000.21	300.09	<b>1206.61</b>	300.08	2026.38	300.06
15	15	165.2	944.0	8926.65	300.09	9000.24	300.14	<b>2410.81</b>	300.10	3530.65	300.08
20	20	219.4	1258.2	12011.00	300.15	12000.43	300.12	<b>3438.31</b>	300.16	4494.19	300.10
25	25	274.6	1572.8	15014.64	300.19	15000.44	300.14	<b>5187.83</b>	300.17	6019.73	300.11

All methods quickly achieve the slot optimization timeout

# Fairness

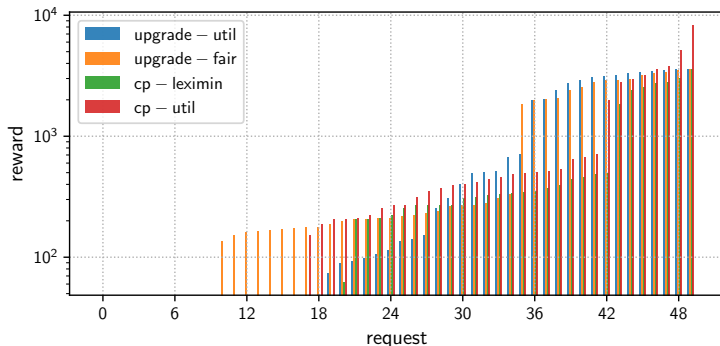
Utility profiles an instance with 25 time-tagged allocation requests



- All the methods behave quite **similarly**
- Such problems are **not too constrained**, given the constellation configuration
- The **fair approaches only serve one more request**

# Fairness (cont.)

Utility profiles an instance with 25 time-tagged allocation requests and 25 global allocation requests



- The **most rewarding** requests (25 to 49) are the **global allocation** ones
- Some requests **cannot be fulfilled** (even by leximin)
- upgrade – fair **serves more requests** (more than cp – fair, due to time budget)
- upgrade – util **is better than** cp – util on most of the high reward requests
- upgrade – fair and upgrade – util behave very well on the fairness side

# Today's Menu

- 1 Introduction
- 2 Core Concepts and Problem Definition
- 3 Constraint Programming Model
- 4 Optimization Framework and Algorithms
- 5 Experimental Evaluation
- 6 Conclusions and Perspectives**

# Conclusions

## To Sum Up

- We modeled a **novel problem (OSAP)** for allocating orbit slots
- We considered both **utilitarian and fairness** objectives
- We considered **two types of requests**: time-tagged requests and global requests
- We proposed an iterative **two-level optimization framework**
- We evaluated **four solution methods**
  - **Global allocation requests are the hardest** ones to fulfill
  - cp – util **and** cp – fair **do not scale** on larger instances
  - **Iterative upgrading methods result in good quality solutions** and are **3 times faster** on larger instances

## Future Research

- Investigating other types of requests (e.g. areas of interest)
- Investigating other types of mode selection (e.g. searching in the mode space)
- Exploring other iterative schemes (e.g. degrading instead of upgrading)

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**Thank you for your attention!**

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