Project One Report Introduction to Operating Systems New Beginnings Spring 2018

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Description

For this assignment, I learned about the flow of control for system calls in xv6; how to add a new system call; how to access specific information for each active process; and how to use conditional compilation to enable and disable kernel features.

Deliverables

The following features were added to xv6:

• A system call tracing facility that, when enabled, prints the following information to the console:

```
<system call name> -> <system call return code>
```

- A new system call, date(), that returns the current UTC date and time.
- A new user command, date, that prints the current UTC date and time to standard output.
- Each process now records the value of the ticks global variable when that process is created. This value is used to calculate *elapsed time* for each process.
- A modification to the existing control—p mechanism, which displays debugging information, to include elapsed time for each process.

Implementation

System Call Tracing Facility

All the code for the system call tracing facility was conditionally compiled using the PRINT_SYSCALLS flag in the Makefile (line 07). The implementation modified syscall.c as follows:

- Lines 137 166 define an array of system call names, syscallnames[], indexed by system call number as defined in syscall.h.
- Lines 178 184 prints the name of the system call and the corresponding return value.

Date System Call

The following files were modified to add the date() system call.

- user.h. The user-side function prototype for the date() system call was added (line 27). The system call takes a pointer to a user-defined rtctime struct. The prototype is:
 - int date(struct rtcdate*);

The file date.h contains the rtcdate definition.

- syscall.h. The date() system call number was created by appending to the existing list (line 25).
- syscall.c. Modified to include the kernel-side function prototype (line 102); an entry in the function dispatch table syscalls[] (line 129); and an entry into the syscallnames[] array to print the system call name when the PRINT_SYSCALLS flag is defined. All prototypes here are defined as taking a *void* parameter as the function call arguments are passed into the kernel on the stack. Each implementation (e.g., sys_date()) retrieves the arguments from the stack according to the syntax of the system call.

- usys.S. The user-side stub for the new system call was added (line 33). This stub uses a macro that essentially just traps into kernel-mode.
- sysproc.c. Contains the kernel-side implementation of the system call in sys_date() (lines 96 108). This routine removes the pointer argument from the stack and passes it to the existing routine cmostime() in lapic.c (line 206). The pointer argument is expected to be a struct rtctime*. The routine cmostime() cannot fail so a success code is returned by sys_date().

Date User Command

The date user command is implemented in the file date.c. This command invokes the new date() system call to fill in the supplied rtcdate struct; passed by reference. The command then displays the date and time information on standard output. The return code from the system call is checked and handled as a user program does not know if a system call will succeed or fail.

control-p Modifications and Elapsed Time

The control—p console command prints debugging information to the console. The following modifications were made to capture and display elapsed time as part of the existing control—p debugging information.

- proc.h. A new field was added to struct proc named uint start_ticks for storing the time of creation (in *ticks*) for each process (line 69).
- proc.c. The routine allocproc() (line 42) was modified to correctly set start_ticks on process creation.
- procdump(). This routine in proc.c was modified to:
 - Print a header (line 512) to the console.
 - Calculate the *elapsed time* since process creation (lines 528 530). This section calculates elapsed time as seconds and hundredths of seconds as the granularity of the ticks variable is at thousandths of a second.
 - Include the elapsed time in the display of process information on the console (line 531).
 - Minor changes to the formatting of the existing information displayed (line 524)

Testing

Previous Output

For reference, the output when the project flag is set to 0 is included here:

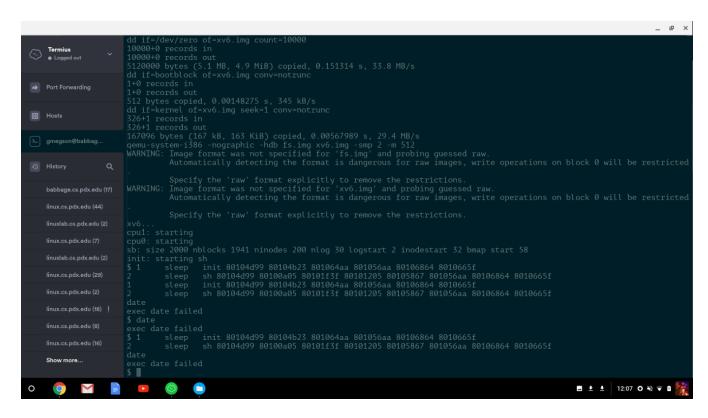


Figure 1: Original Output

System Call Tracing Facility

I tested this feature by modifying the Makefile to turn on PRINT_SYSCALLS flag, then booting the xv6 kernel, and observing the following output:

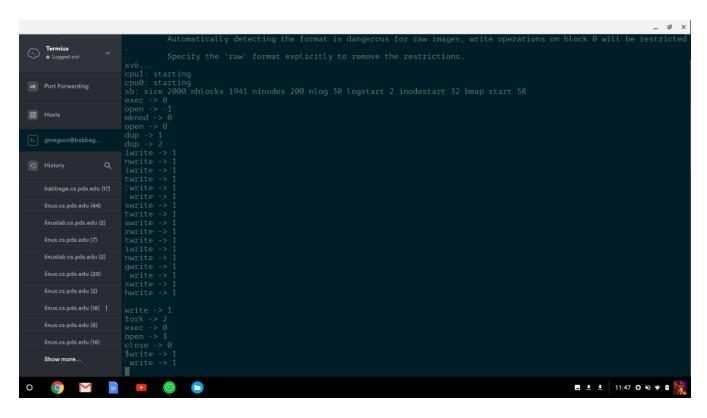


Figure 2: System Call Tracing Facility

The system call trace correctly displays invoked system calls. The standard output is interleaved with the trace output. The output "init: starting sh" is printed by the init() process (init.c) and the "\$" is printed by the shell process (sh.c).

This test PASSES.

Date System Call and User Command

I am going to use the date command to test both the date() system call and date command, as I can't directly invoke a system call from the shell. My testing will invoke my date command in xv6 and then invoke the corresponding Linux date command and see if the former closely matches the latter.

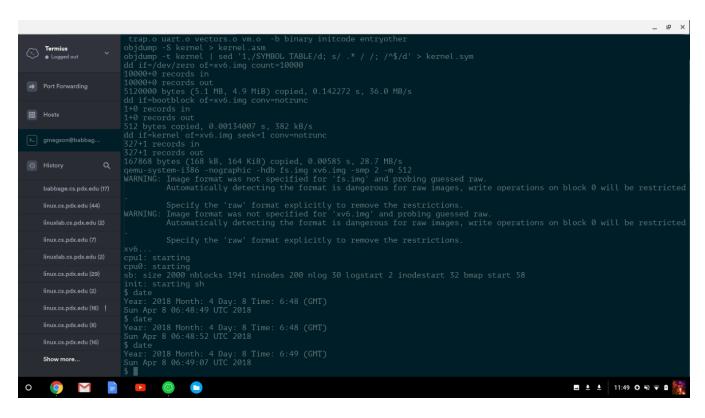


Figure 3: Date Test

The output from my date command closely matches the output of the Linux date command, except for a discrepancy in the precision. The struct rtcdate does not have fields more precise than seconds. This test shows that the date command works correctly, along with the date system call, since the command prints out all of the information extracted by the system call.

This test PASSES.

control-p and Elapsed Time

The test for these will be split into two phases. My first test will show that control—p is outputting the correct information, while my second test will use the first test to show that the elapsed time is correct.

Here is the output of the first test:

The control—p output indicates that there are two processes running in xv6. This is correct. The first process is the initial process, here named "init", with a PID of 1. The second process is the shell, named "sh", with a PID of 2 (as it is the second process created). Note that the PCs appear to be correct, as they correspond to valid addresses in the kernel asm file and the code for printing this information was not modified.

This sub-test PASSES.

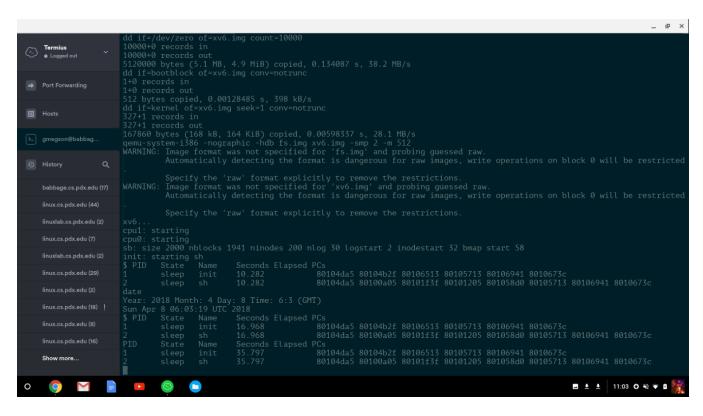


Figure 4: Control-p Test

For the second test, I will restart the kernel, and then press control—p several times, each press being within one second of the other. The results are shown below:

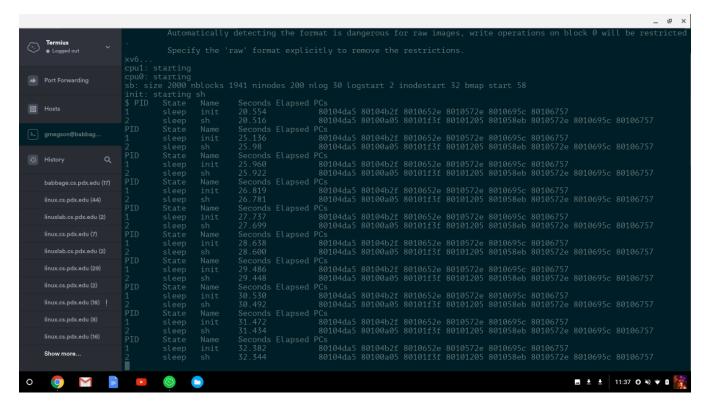


Figure 5: Elapsed Time Test

The elapsed time for the init process is 0.38 seconds higher than that of the sh process in all outputs. This makes sense as init starts before sh. Also, note that the elapsed times are steadily increasing by about one second with the same 0.38 s difference between init and sh.

This sub-test PASSES.

Because all sub-tests passed, this test PASSES.

Reference: Usertests and Forktest

Inluded here are screenshots of the usertests and forktest outputs.

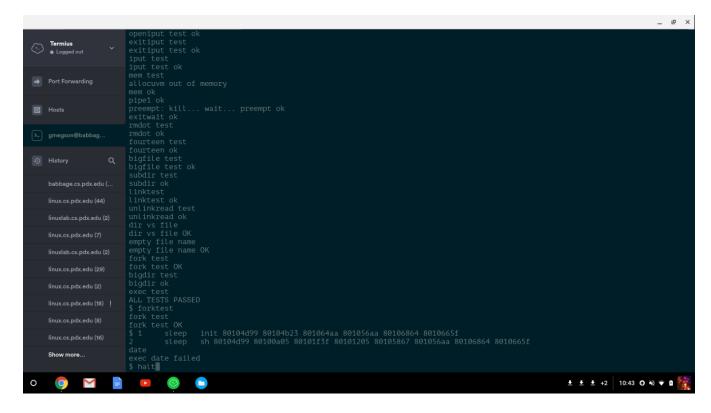


Figure 6: Project Flag to 0

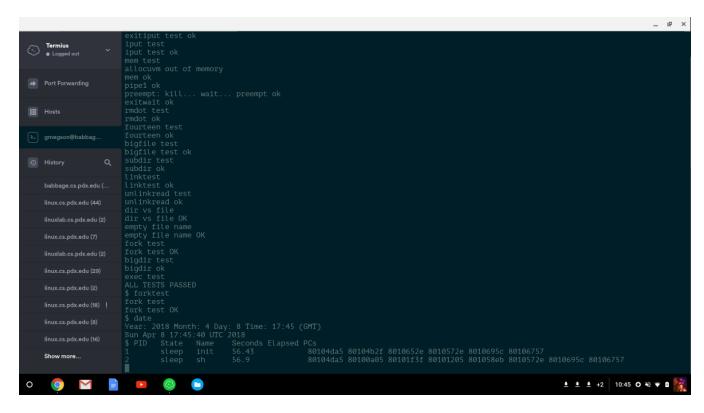


Figure 7: Project Flag to 1