

00. INTRODUCTION

data compression

- types of compression
 - lossless compression** - can recover the contents
 - lossy compression** - lose some quality - cannot convert back to the higher-quality version
- examples
 - sparse binary string - storing positions of 1s
 - equal number of 0/1s - $L \geq \log_2 \binom{64}{32} \approx 60.7$
 - english text - using relative frequency
 - morse code is NOT binary (contains spaces)
- info theory uses **probabilistic models** (letter frequency, sequence probabilities)
- 2 distinct approaches to compression:
 - variable length** - map more probable sequences to shorter binary strings
 - fixed length** - map most probable sequences to strings of a given length
 - insufficient strings for low-probability sequences
 - tradeoff between length/failure probability

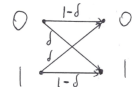
information theory concepts

- speed: **rate** $\rightarrow \frac{k}{n}$ (mapping k bits to n bits)
- reliability: $\mathbb{P}[\text{error}] = \mathbb{P}[\text{estimated msg} \neq \text{true msg}]$
- source coding theorem** \rightarrow the fundamental compression limit is given by a source-dependent quantity known as the **(Shannon) entropy** H . The (average) storage length can be arbitrarily close to H , but can never be any lower than H .
 - H is a property of the *probability distribution*
- channel coding theorem** \rightarrow there exists a channel-dependent quantity called the **(Shannon) capacity** C such that arbitrarily small error probability can be achieved only for rates $< C$
 - can achieve $\mathbb{P}[\text{error}] \leq \epsilon \iff \text{rate} < C$

data communication example

- a "transmitter" sends a sequence of 0s and 1s
- a "receiver" sends a sequence *with some corruptions*

channel transition diagram



- each bit is flipped independently with probability $\delta \in (0, \frac{1}{2})$

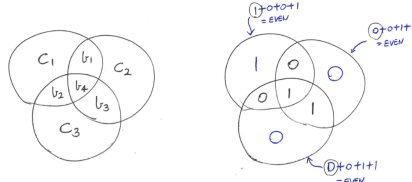
naive

- uncoded communication** - $\mathbb{P}[\text{correct}] = (1 - \delta)^N$
- repetition code** - transmit "000" for "0", "111" for "1"
 - $\mathbb{P}[\text{correct}] = [(1 - \delta)^3 + 3\delta(1 - \delta)^2]^N$
 - more reliable but 3x slower!

Hamming code

- able to correct one bit flip
- maps binary string of length 4 to binary string of length 7

- fill in $b_1 b_2 b_3 b_4$ and assign $c_1 c_2 c_3$ such that the sum of bits in each circle is even



- $\mathbb{P}[\text{correct}] \geq \mathbb{P}[\leq 1 \text{ bit flips}] = (1 - \delta)^7 + 7\delta(1 - \delta)^6$
- with $\delta = 1$: Shannon capacity $C \approx 0.531$

01. INFORMATION MEASURES

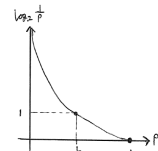
information of an event

- entropy** \rightarrow measure of "uncertainty" or "information" in a random variable
- given event A with some $\mathbb{P}[A] = p$, how much "information" learned by being told A occurred?
 - only $\mathbb{P}[A]$ matters
- if A occurs with probability p , then $\text{Information}(A) = \psi(p)$ for some function $\psi(\cdot)$

axioms for $\psi(\cdot)$

$$\psi(p) = \log_b \frac{1}{p} \text{ (for some base } b > 0)$$

we gain $\log_2 \frac{1}{p}$ "bits" of info if a probability- p event occurs.



- only $\psi(p) = \log_b \frac{1}{p}$ satisfies all axioms
- we focus on $b = 2$
 - information measured in bits
- all choices of b are equivalent up to scaling by a universal constant
 - e.g. # of nats = $\log_e 2 \times$ # of bits

- $\psi(p) \geq 0$ (**non-negativity**)
- $\psi(1) = 0$ (**zero for definite events**)
- if $p \leq p'$, then $\psi(p) \geq \psi(p')$ (**monotonicity**)
 - the less likely an event is, the more information was learnt by the fact that it occurred
- $\psi(p)$ in continuous in p (**continuity**)
 - small change in probability: no drastic change in info
- $\psi(p_1 p_2) = \psi(p_1) + \psi(p_2)$ (**additivity under independence**) if A and B are independent events with probabilities p_1 and p_2 , then $\mathbb{P}[A \cap B] = p_1 p_2$, and the information learnt from both A and B occurring is the sum of the two individual amounts of information (because they are independent)
 - $\psi(\mathbb{P}[A_1 \cap A_2]) = \psi(\mathbb{P}[A_1]) + \psi(\mathbb{P}[A_2])$

information of a random variable - entropy

- let X be a discrete r.v. with pmf P_X
- if we observe $X = x$ then we have learnt $\log_2 \frac{1}{P_X(x)}$ bits of information

(Shannon) entropy

is the average *information/uncertainty* in X wrt P_X :

$$H(X) = \mathbb{E}_{X \sim P_X} \left[\log_2 \frac{1}{P_X(X)} \right] = \sum_x P_X(x) \log_2 \frac{1}{P_X(x)}$$

- binary entropy function** \rightarrow

$$H_2(p) = p \log_2 \frac{1}{p} + (1 - p) \log_2 \frac{1}{1 - p}$$

- e.g.

- binary source: $X \sim \text{Bernoulli}(p)$, $p \in (0, 1)$

$$\Rightarrow H(X) = H_2(p)$$
- uniform source: X is uniform on a finite set \mathcal{X}

$$P_X(x) = \frac{1}{|\mathcal{X}|}$$

$$\Rightarrow H(X) = \mathbb{E} \left[\log_2 \frac{1}{1/|\mathcal{X}|} \right] = \log_2 |\mathcal{X}|$$

- entropy \neq variance

- entropy depends *only* on the probability values

axiomatic view (Shannon)

X is a d.r.v. taking N values with $\mathbf{p} = (p_1, \dots, p_N)$. We consider a general information measure of the form

$$\Phi(\mathbf{p}) = \Phi(p_1, \dots, p_N)$$

only $\Phi(X) = \text{constant} \times H(X)$ satisfies all axioms.

- $\Psi(\mathbf{p})$ is continuous on \mathbf{p} (**continuity**)
- if $p_i = \frac{1}{N}$, then $\Psi(\mathbf{p})$ is increasing in N (**uniform case**)
 - uniformity over a larger set of outcomes always means more uncertainty
- (successive decisions)** $\Psi(p_1, \dots, p_N) = \Psi(p_1 + p_2, p_3, \dots, p_N) + (p_1 + p_2) \Psi(\frac{p_1}{p_1 + p_2}, \frac{p_2}{p_1 + p_2})$

variations

- joint entropy** of two random variables $(X, Y) \rightarrow$

$$H(X, Y) = \mathbb{E}_{(X, Y) \sim P_{XY}} \left[\log_2 \frac{1}{P_{XY}(X, Y)} \right] = \sum_{x, y} P_{XY}(x, y) \log_2 \frac{1}{P_{XY}(x, y)}$$

- conditional entropy** of Y given $X \rightarrow$

$$H(Y|X) = \mathbb{E}_{(X, Y) \sim P_{XY}} \left[\log_2 \frac{1}{P_{Y|X}(Y|X)} \right] = \sum_{x, y} P_{XY}(x, y) \log_2 \frac{1}{P_{Y|X}(y|x)} = \sum_x P_X(x) H(Y|X = x)$$

- on average, knowing X reduces uncertainty about Y ($H(Y|X) \leq H(Y)$), but seeing a *specific* outcome of X may increase uncertainty about Y ($H(Y|X = i) > H(Y)$ for some values of i)

properties of entropy

- $H(X) \geq 0$ (**non-negativity**)
 - $H(X) = 0 \iff X$ if deterministic
 - proof: information $\log_2 \frac{1}{p} \geq 0$ for $p \in [0, 1]$, so entropy is the average of a non-negative quantity, and itself is non-negative
- $H(X) \leq \log_2 |\mathcal{X}|$ (**upper bound**)
 - if X takes values on a finite alphabet \mathcal{X}
 - $H(X) = \log_2 |\mathcal{X}| \iff X \sim \text{Uniform}(\mathcal{X})$
 - implies $H(X|Y) \leq \log_2 |\mathcal{X}|$
- $H(X, Y) = H(X) + H(Y|X)$ (**chain rule**)
 - or $H(X, Y) = H(Y) + H(X|Y)$