

Pre-tutorial questions

Do you know the basic concepts of this week's lecture content? These questions are only to test yourself. They will not be explicitly discussed in the tutorial, and no solutions will be given to them.

1. Running times

- (a) What is an algorithm's worst case running time?
- (b) What does it mean when we say that the algorithm runs in polynomial time?
- (c) An algorithm is efficient?
- (d) Look at the table with the running times in the slides. Can you explain the reason for these numbers?

2. Asymptotic analysis

- (a) Can you explain the ideas of the $O(\cdot)$, $\Theta(\cdot)$ and $\Omega(\cdot)$ functions? Why do we use these functions?
- (b) What does it mean that these functions are transitive and additive?

3. Basic data structures

- (a) What are linked lists, queues, stacks and balanced binary trees? What sort of operations are usually supported by these structures?
- (b) What is the height of a balanced binary tree containing n elements?

Tutorial

Problem 1

Sort the following functions in increasing order of asymptotic growth

$$n, n^3, n \log n, \frac{n}{\log n}, \frac{n}{\log^2 n}, \sqrt{n}, \sqrt{n^3}$$

Solution: Answer should be:

$$\sqrt{n}, \frac{n}{\log^2 n}, \frac{n}{\log n}, n, n \log n, \sqrt{n^3}, n^3$$

Problem 2

Sort the following function in decreasing order of asymptotic growth

$$n^{1.5}, 2^n, \frac{n}{2^n}, \frac{2^n}{n^{10}}, n!, 1.5^n, 2^{\log n}$$

Solution: Answer should be:

$$n!, 2^n, \frac{2^n}{n^{10}}, 1.5^n, n^{1.5}, 2^{\log n}, \frac{n}{2^n}$$

Problem 3

Which of the following is largest asymptotically

$$\log_3 n, \log_2 n, \log_{10} n$$

Solution: They are all equal. This is because:

$$\log_b a = \frac{\log_d a}{\log_d b}$$

Problem 4

Prove that the sum $S(n)$ of the first n natural numbers is $n(n+1)/2$.

Solution: Proof is by induction on n . Base case: If $n = 1$ then the claim is trivially true. Induction hypothesis: Assume that the sum of the first n natural is $n(n+1)/2$. Induction step: Next we want to prove the statement for $n+1$. From the definition of $S(n)$ we know $S(n+1) = S(n) + n + 1$, and from the induction hypothesis we know $S(n) = n(n+1)/2$, and therefore $S(n+1) = S(n) + n + 1 = n(n+1)/2 + n + 1 = (n+2)(n+1)/2$.

Problem 5

Imagine a program A running with time complexity $\Theta(f(n))$, taking t seconds for an input of size m . What would your estimation be for the execution time for an input of size $2m$ for the following functions: n , $n \log n$, n^2 or n^3 .

Solution: $n \rightarrow 2m = 2t$ seconds
 $n \log n \rightarrow 2m \log 2m = 2t(1 + \log 2 / \log m)$ seconds
 $n^2 \rightarrow 4m^2 = 4t$ seconds
 $n^3 \rightarrow 8m^3 = 8t$ seconds

Problem 6

1. For each of the following pseudo-code fragments. give an upper bound for their running time using the big-Oh notation.
2. The upper bound captures the worst case running time of an algorithm. Some instances might require the algorithm to perform more steps than others. Some instances might allow the algorithm to terminate faster. A lower bound for the running time is a “best-case” running time in the worst-case. If an algorithm is $\Omega(f(n))$ for example, then there exists an instance that will take *at least* $f(n)$ steps. For the second algorithm shown below, give a lower bound for the running time.

Algorithm 1 Stars

```
1: for  $i = 1, \dots, n$  do  
2:   print “*”  $i$  times  
3: end for
```

Solution: first iteration prints 1 star, second prints two, third prints 3 and so on. Total number of stars is

$$1 + 2 + \cdots + n = \sum_{j=1}^n j = \frac{n(n+1)}{2} = O(n^2)$$

Algorithm 2 CHECK NUMBERS

```

1: procedure CHECKNUMBERS( $A, B$ )
     $\triangleright A$  and  $B$  are two lists of integers
2:   count = 0
3:   for  $i = 1, \dots, n$  do
4:     for  $j = i \dots m$  do
5:       if  $A[i] \geq B[j]$  then
6:         count = count + 1
7:         break
8:       end if
9:     end for
10:  end for
11: end procedure

```

Solution: An upper bound is $O(n \cdot m)$. The lower bound is $\Omega(n \cdot m)$. Ask the students to come up with a specific example input for this (which is easy).

Problem 7

Given an array A consisting of n integers $A[0], A[1], \dots, A[n-1]$, we want to compute the upper triangle matrix

$$C[i][j] = \frac{A[i] + A[i+1] + \cdots + A[j]}{j - i + 1}$$

for $0 \leq i \leq j < n$. Consider the following algorithm for computing C :

Algorithm 3 summing-up

```

1: function SUMMING-UP( $(A)$ )
2:   for  $i = 0, \dots, n-1$  do
3:     for  $j = i, \dots, n-1$  do
4:       add up entries  $A[i]$  through  $A[j]$  and divide by  $j - i + 1$ 
5:       store result in  $C[i][j]$ 
6:     end for
7:   end for
8:   return  $C$ 
9: end function

```

1. Using the O -notation, upperbound the running time of SUMMING-UP.
2. Using the Ω -notation, lowerbound the running time of SUMMING-UP.

Solution:

1. The number of iterations is $n + n - 1 + \cdots + 1 = \binom{n}{2}$, which is bounded by n^2 . In the iteration corresponding to indices (i, j) we need to scan $j - i + 1$ entries from A , so it takes $O(j - i + 1) = O(n)$. Thus, the overall time is $O(n^3)$.
2. In an implementation of this algorithm, Line 3 would be computed with a for loop; when $i < \frac{1}{4}n$ and $j > \frac{3}{4}n$, this loop would iterate least $n/2$ times, which takes $\Omega(n)$ time. There are $n^2/16$ pairs (i, j) of this kind, which is $\Omega(n^2)$. Thus, the overall time is $\Omega(n^3)$.