

01. Use the recursion tree method to find an asymptotic upper bound for the given recurrence forms.

A.	$T(n) = \begin{cases} T(n/2) + cn^2 & n \geq 2 \\ c & n = 1 \end{cases}$
B.	$T(n) = \begin{cases} 3T(n/4) + cn^2 & n \geq 4 \\ 1 & n = 1 \end{cases}$
C.	$T(n) = \begin{cases} 2T(n/2) + 1 & n \geq 2 \\ 1 & n = 1 \end{cases}$
D.	$T(n) = \begin{cases} 3T(n/2) + cn & n \geq 2 \\ 1 & n = 1 \end{cases}$
E.	$T(n) = \begin{cases} 7T(n/2) + cn^2 & n \geq 2 \\ 1 & n = 1 \end{cases}$

I. For the above each recurrence

- A. Draw the tree
- B. Compute depth
- C. Compute cost per level
- D. Sum and produce the final upper bound

II. After drawing the recursion tree for the given recurrence:

- A. Explain how the cost at each level changes as the tree goes deeper.
- B. Identify which part of the tree (top, middle, or bottom) dominates the total cost, and justify your reasoning.

O2. Draw a recursion tree for the following recurrences and use it to obtain Big O.

- A. $T(n) = T(n/5) + T(4n/5) + n$
- B. $T(n) = T(n/3) + T(2n/3) + n$
- C. $T(n) = T(n-1) + \log n$
- D. $T(n) = 2T(n-1) + 1$