

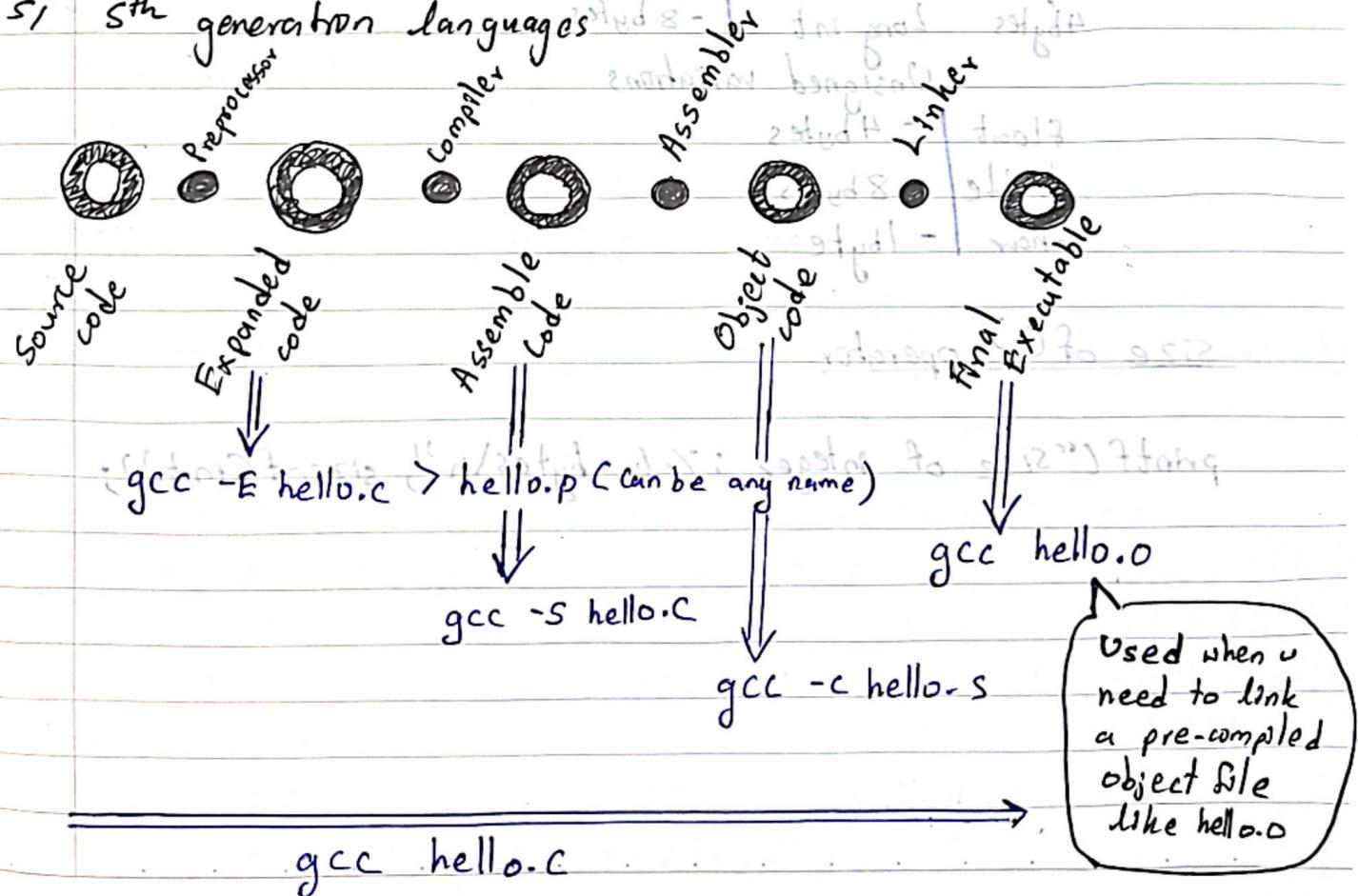
Data Structures & Program Design in C

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- 1/ 1st generation languages: Machine language (bit-level language)
- 2/ 2nd generation languages: Contains human readable notation.
 - Converted to ML using an assembler
- 3/ 3rd generation languages:
 - Contains series of English-like words that humans can understand easily.
 - Converted to ML using translator or interpreter
- 4/ 4th generation languages:
 - Higher level than 3rd gen
 - Used in more specific fields such as data and database manipulations.
 - More efficient than 3GL

5/ 5th generation languages



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) // is ignored by compiler & constitute part

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(Comments : // This is a comment // * Ignored by the compiler
 / * when the executable is
 . notation of // This is a comment // being created. //
 addressed as part of both lines)

Variables : name in C, as mentioned :-

- can contain letters, digits & underscores
- must begin with a letter, \$ or underscore.
- case sensitive
- no white spaces & special characters
- cannot be a keyword.

Variable types,

Primitive : int : short int	- 4 bytes	
4 bytes	Long int	- 8 bytes

Unsigned variations

Float	- 4 bytes
-------	-----------

double	- 8 bytes
--------	-----------

char	- 1 byte
------	----------

size of () operator

```
printf("size of integer : %lu bytes\n", sizeof(int));
```

0. offset 32p

1. offset 2 - 32p

2. offset 3 - 32p

3. offset 22p

Waste bytes
start of base
Elongated-size is
16 bits to 32 bits
offset width

%[flags][width].[precision][length]specifier

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COMMON FORMAT SPECIFIERS

1) %d int Signed decimal integer

2) %i int Same as %d, signed decimal integer

3) %u unsigned int Unsigned decimal integer

4) %f float/double Decimal floating-point (6 digit by default)

5) %F float/double Same as %f, but capitalized for INF/NAN

6) %e float/double Scientific notation (eg: 1.23e+03)

7) %E float/double Scientific notation with E (eg: 1.23E+03)

8) %g float/double Uses %f or %e depending on value precision

9) %G float/double Like %g, but uses %P/%E

10) %x unsigned int Hexadecimal (lowercase)

11) %X unsigned int Hexadecimal (uppercase)

12) %o unsigned int Octal

13) %c char Single character

14) %s char* String (null-terminated)

15) %p void* Pointer address

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LENGTH MODIFIERS

1) h repeating signed char/unsigned char %hd, %hu, %hhx

2) l short/unsigned short %ld, %lu, %lx

3) ll long/unsigned long %lld, %llu

4) L long double %Lf, %Le, %Lg

5) z size_t %zu, %zd

6) j intmax_t/uintmax_t %jd, %ju

7) t ptrdiff_t %td, %td

SPECIAL FLAGS (for printf())

1) - Left justify in field

2) + Force (+) for positive numbers

3) <space> space before positive numbers

4) 0 Pad with zeros

5) # Alternate formatting
(hexadecimal prefix, decimal point)

(Alternate form) to read a string in base 16

Applies to,

- %o → Prefix with 0 | printf("%o\n", 10); 8
- %n / %X → Prefix with 0x / 0X | printf("%#n\n", 255); 0xFF
- %f, %g, %e → Include decimal always | printf("%#.g\n", 3.0); 3.00000

WIDTH & PRECISION

1) %.5d at least 5 characters wide

2) %.2f 2 digits after decimal point

3) %.8.3f 8 characters wide, 3 after decimal

4) %.*.f width and precision specified at runtime (use with args)

SCANF - SPECIFIC FORMAT SPECIFIERS

1) %-d Read a signed decimal integer

2) %i (obsolete) Read an integer (detects base automatically: 0 for octal, 0x for hex)

3) %u Read an unsigned decimal integer

4) %.f Read a float

5) %.lf Read a double (matters in scanf, unlike in printf (lf))

6) %Lf Read a long double

7) %c Read a character (does not skip whitespace)

8) %s Read a string until whitespace

9) %[...]. Character set - read only characters in the bracket

10) %[^...]. Inverse character set - read until any character in brackets appears.

11) %p Read a pointer address

12) %%. Match a literal % in input

ESCAPE SEQUENCES (for outputs)

1) \n : next line

2) \t : adds tab space

3) \a : produce a sound or audible alert

4) \"": Print \ or double quotes "

Char: Display one character from ASCII code (128 characters)

(7 bits) (and not 8)

Nowadays there are more than 128 due to extended ASCII.

* 2 ways to assign characters,

1) char a = 'A'; (use single quotes) (double quotes for strings)

2) char a = 65; (using ASCII decimal value)

OPERATORS

operator (+, =, -, *, /)

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Operands (thing which the operator works on)

1) '=' represents assignment operator

$L.H.S <= R.H.S$

ex: int a = 5;

$\Sigma \leftarrow f_{n1} / f_{n2} = f_{n3} : \Sigma$

2) "==" used to check equivalence. (Relational operator)

True (1) or False (0)

takes 1 byte

ie: In C all numbers ~~are~~ except 0 is true.

Arithmetic Operators

$\Sigma \leftarrow f_{n1} / f_{n2} = f_{n3}$

BINARY OPERATORS (Needs 2 operands)

(+), (-), (*), (/), (%)

* % works only for int types. (Modulo)

ex: int a = 17, b = 5 $a \% b = 2$

UNARY OPERATORS (Needs only 1 operand)

Increment (++) , Decrement (--), Unary plus (+), Unary Minus (-)

$a++ \leftarrow$ post setting ex: int a = 5; **Output** a = 6 b = 5

~~start here~~ \leftarrow int b = a++; ie: b = a++; b = a
a = a + 1

$++a \leftarrow$ pre setting ex: int a = 5 **Output** a = 6 b = 6

int b = ++a ie: b = ++a : a = a + 1.

b = a

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* Same for --a & a--

~~++a++~~ ~~*~~ ERROR

Integer division: When 2 numbers are divided, only the whole number is displayed.

$$17/5$$

$$\text{ex: } \text{int} = \text{int}/\text{int} \Rightarrow 3$$

$$<2.1.9> <2.1.1>$$

$$(2 = 0 \text{ mod } 5)$$

$$\text{double} = \text{int}/\text{int} \Rightarrow 3.000000$$

(Integers have 10 digits, so we have 10 digits of float "1.5")

$$\text{double} = \text{double}/\text{int} = \text{int}/\text{double} \Rightarrow 3.400000$$

$$\text{double} = \text{double}/\text{double} \Rightarrow 3.400000$$

$$\text{int} = \text{double}/\text{double} \Rightarrow 3$$

$$\text{double}(17/5) \Rightarrow 3 \quad \text{BINARY} \quad 2907A9390$$

$$\text{double} c = 17.0/5.0 \Rightarrow 3.4, (1), (+), (-), (*)$$

RELATIONAL OPERATORS (True (1) or False (0))

(>), (>=), (<), (<=), (==), (!=)

LOGICAL OPERATORS

(-) AND (PP), OR (II), NOT (!) ← Unary logical operator

Relational Operators \longrightarrow Logical truth

↑
Works on
Logical Operators

ie: NAND & NOR universal operators. (Functionally Complete)
 (can do any logical operation using only NAND or NOR)

BIT MANIPULATION (works only for char & int types)

- * Used to manipulate large numbers by breaking them up.

1) & -bitwise AND: int a=6, b=4

$$\begin{array}{r} 00110 \leftarrow 6 \\ 00100 \leftarrow 4 \\ \hline 00100 \rightarrow 4 \end{array}$$

6 & 4 [OUTPUT] 4

2) | -bitwise OR: int a=6, b=4

$$\begin{array}{r} 00110 \leftarrow 6 \\ 00100 \leftarrow 4 \\ \hline 00110 \rightarrow 6 \end{array}$$

6 | 4 [OUTPUT] 6

3) << -leftshift: 6 << 2

$$\begin{array}{r} 0110 \xrightarrow{158421} 6 \\ 11000 \xrightarrow{8421} 24 \end{array}$$

[OUTPUT] 24

4) >> -rightshift: 6 >> 2

$$\begin{array}{r} 0110 \xrightarrow{8421} 6 \\ 0001 \xrightarrow{8421} 1 \end{array}$$

[OUTPUT] 1

5) ^ -bitwise XOR: int a=6, b=4

$$\begin{array}{r} 00110 \leftarrow 6 \\ 00100 \leftarrow 4 \\ \hline 00010 \rightarrow 2 \end{array}$$

6 ^ 4 [OUTPUT] 2

6) ~ -complement (unary operator): ~6

$$\begin{array}{r} 00110 \rightarrow 6 \\ 11001 \rightarrow -7 \end{array}$$

2's complement
due to signed

* If 6 was named unsigned then its
Is complement.

OPERATOR PRECEDENCE

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Associativity

- 1/ Postfix ($() [] \rightarrow . \quad ++ \quad --$)
- 2/ Unary ($+ - ! \sim \quad ++ \quad -- \quad (\text{type})^* \quad \& \quad \text{size of}$) Right to left
/Prefix
- 3/ Multiplicative ($* \quad / \quad \% \quad \&$)
- 4/ Additive ($+ \quad - \quad += \quad -=$)
- 5/ Shift ($<< \quad >>$)
- 6/ Relational ($< \quad <= \quad > \quad >=$)
- 7/ Equality ($= \quad != \quad == \quad !=$)
- 8/ Bitwise AND $\&$
- 9/ Bitwise XOR \oplus
- 10/ Bitwise OR $\|$
- 11/ Logical AND $\& \quad \&=$
- 12/ Logical OR $\| \quad \| =$
- 13/ Conditional $? \quad ?: \quad ?::$ Right to left
- 14/ Assignment $= \quad /= \quad *= \quad += \quad -= \quad *== \quad /== \quad +== \quad -==$ Right to left
- 15/ Comma

Example 1: int $x=4, y=5, z=2;$

int result = $x++ * --y + z++ / x;$

$$= \underbrace{(x++)}_4 * \underbrace{(--y)}_4 + \underbrace{(z++)}_2 / \underbrace{x}_5$$

$n = n+1$ done after the operation

$$= 4 * 4 + 2 / 5$$

$$= 16 + 2 / 5$$

$$= \underline{\underline{16}}$$

Ans (using go program) transformation = 16

Example 2: int $a = 3, b = 4, c = 2; \text{int result} = ++a * b -- + c++ - a;$

$$= (\cancel{+} a) * (\cancel{b--}) + (\cancel{c++}) - \cancel{a}$$

$$= 4 * 4 + 2 - 4$$

$$= 16 - 2$$

$$= \underline{\underline{14}}$$

Example 3: int $x = 5, y = 10, z = 4;$

$$\text{int result} = x \& y | z \& \neg z$$

$$= (\cancel{x++}) \& y | (\cancel{z++}) \wedge (\cancel{-z})$$

$$= 5 \& 10 | 4 \wedge 5$$

$$= \frac{\begin{array}{r} 0101 \\ 1010 \\ \hline 0000 \end{array}}{\frac{\begin{array}{r} 100 \\ 101 \\ \hline 001 \end{array}}{}}$$

CONSTANTS

- Read only variables. Value cannot be changed once declared.

2 ways to define constants

i.e. Variable are by convention named

1/ `const <data-type><var_name> = value;` in full capitals.

2/ `#define <const_name> <value>`

- Defined outside the main method

- Replaces the constant name with the data in the preprocessor step.

- Doesn't have to consider data types.

- Difference: 2) defines faster
2) doesn't consider variable type, it just replaces.
∴ it sacrifices data type information.

Enumerations (Different from array index) 4 bytes.

- lists that contain constant integer values.

enum <enum-name> {<options>};

- list of values has to be unique within the scope

Example: enum Grade {A, B, C, D};

```
int main()
{
```

OUTPUT

enum Grade T=A;
printf("%d\n", T);

0 (position of A in
the enumerator).

* To print actually "A" you
need pointers.

2) if {A=50, B, C, D} from which address will go to
Then B=51, C=52, D=53

3) if {A, B=50, C, D}

Then A=0, C=51, D=52

- You can also set different values to all the elements in the list.

4) if {A=50, B=50, C=50, D}

Then A=50, B=50, C=50, D=50

TYPE DEF

- Typedef keyword is used to define an alias for existing data types and enumerations.

typedef <data-type/enum> <new-name>;

Example : `typedef enum Grade{A, B=50, C=50, D=50}mark;`
`typedef int VALUE;`

int main()

{

 VALUE a = 45;

OUTPUT

50

mark T = A,

mark X = A,

printf("%d %d/n", X, T);

TYPE CONVERSIONS (type casting)

- When the operator has operands of different types, they are converted to common types.

2 ways,

1/ Automatic :- Narrow operands converted to wide ones without losing information.

ex:- int → float, float → double

not prohibited
or illegal

2/ Explicit :- Wider operands converted to narrow operands.

<type-name> <expression>

• Generates a compiler warning
• Requires explicit type conversion

ex: double a = 68.7; | ex: double c = (double) a/b $b=2$
 $\ln b \quad b = (\ln b) \quad a;$ *this happens first.
*So a turns to a double.

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Importance:- Type safety

- Enhancing code readability

- Improved data manipulation

Problems :- Loss of precision

- Overflow or Underflow

- Unexpected behavior

CONTROL FLOW

- Statement : line of code commanding to do a specific task. (terminated by ;)

- Block : Set of statements when executed is syntactically equal to a single statement. (e.g. { ... }) (No terminating character)

IF - ELSE... (make decisions in the program flow)

if (<expression>)

{

 // true
 {

 else

* else is optional

{

 // false

} // this is called bracket nesting - i.e. statements!

with nested

else if (true) {
 // code
}

else if (false) {
 // code
}

SWITCH

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- Multiway decision that tests whether an expression matches one of the defined constant integer variables.

switch (<expression>)

Full through behaviour

- * Once one case matches, every other case below it gets executed.

- * Default case is executed when no cases match. (Not compulsory to have a default case).

- * Default can also be on the top too no strict order.
(Fall through happens through default as well.)

JUMP STATEMENTS

1 / break

Implicit jump

Explicit jump

2 / continue

* Based on a condition

* Explicitly give the command

3 / goto

you perform a jump

to jump.

4 / return

jump conditions

BREAK STATEMENT

- Terminates the immediately enclosing loop block or switch block. (Control passes to the statement immediately after the enclosing block)

- * If there is no enclosing block, compiler will give an error message.

Basically goes out of the curly brackets.

{
 break;
}

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TERNARY OPERATOR

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<condition> ? <if-true> : <if-false>

- * Can be easily used to assign values to variables, but the main issue is readability.

```
ex: int a = 26;  
    int b = 0;           b = 5  
    if ((b = (a > 25) ? 10 : 5) == 5)  
    {  
        printf("T\n");  
    }  
    else  
    {  
        printf("F\n");  
    }
```

Iterative STATEMENTS (loops)

- * All loops have 3 components,
 - 1/ Loop condition
 - 2/ Loop counter *defined outside of the loop (exception for FOR LOOPS)
 - 3/ Loop Body

1/ WHILE LOOP

```
while (<condition>){...}
```

```
for (<initialization>; <condition>; <counter>){...}
```

3/ DO-WHILE LOOP

```
do  
{  
    while (<condition>);
```

* even if the condition is not met, do-while executes at least once regardless.

Terminates through a semi colon.

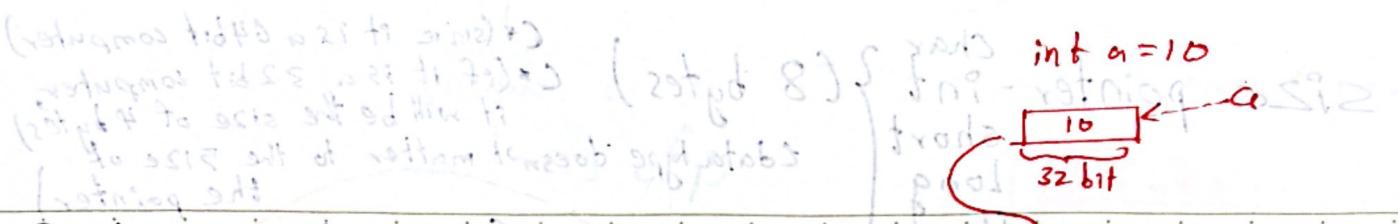
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GOTO STATEMENT (explicit jump)

```
goto <label-name>;  
<label-name>;
```

ex: for (int i=0; i<5; i++)

```
    {  
        A:  
        printf("%d\n", i);  
        }  
        B:  
        printf("%B\n");  
        goto A;  
    }
```



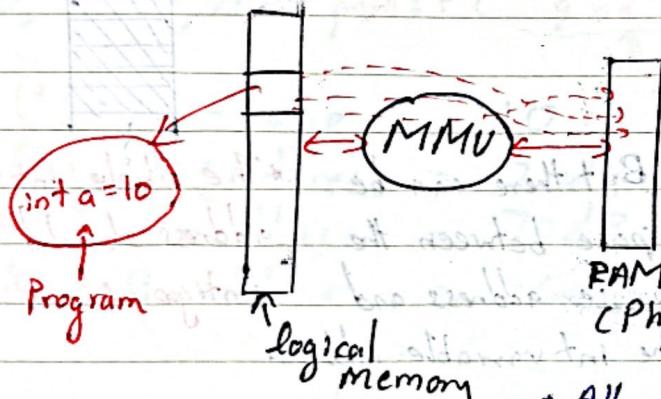
Random Variables Pointers

- * Holds memory address of a variable.
- * Declared using * (indirection or differencing operator)
- * Represents a logical memory address.

Memory has 2 forms: / Logical Memory \Rightarrow memory where information is stored as seen by the computer.

- * \leftarrow declare pointer
- $\&$ \leftarrow get the address of the variable

2/ Physical Memory \Rightarrow the actual RAM of the computer.



MMU - Memory Management Unit

* **Pageable for referencing physical memory**

* Allocation of logical memory to physical memory is done by mmu.

- * Program sees memory as consecutively numbered cells. *groups one after the other without gap in memory*
- * Cells can be manipulated individually or as contiguous groups.
- * Address operator ($\&$) gives the address of an object. (unary operator)

```
int a;
int *p;      <data type> *pointer-name // to declare
p = &a;       pointer-name = & <object> // to assign
```

To retrieve and assign values to a memory location
 $<\text{variable}> = ^\star \text{ptr.}$ // variable is assigned the value
 $^\star \text{ptr} = <\text{value}>$ // store the value

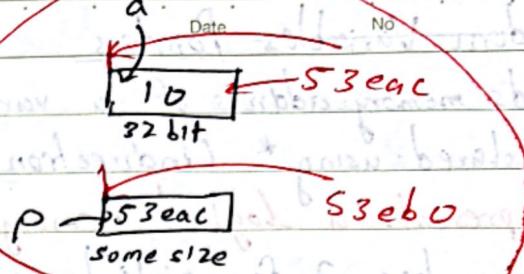
size of pointer - int } (8 bytes)
 short
 long
 float
 double

(since it is a 64 bit computer)
 (if it is a 32 bit computer,
 it will be the size of 4 bytes)
 (data type doesn't matter to the size of
 the pointer)

ex: int main ()

P holds the address to a integer variable.
 int a = 10;
 int *p = & a;

create pointer assign address of a to p.
 printf("%p\n", &a);
 printf("%p\n", &p);



- * You will get a compiler warning if you assign

a incompatible variable type to a pointer which is of a different variable type.

(Can manipulate the compiler to avoid warnings).

- * But there can be space between the pointer address and the int variable address.

- * The whole int variable address should be contiguous.

- * You need data types for pointers to

- * Advantage of using pointers is that you can do your operations in the same location where you initially created the variable without copying it several times. Which saves memory.

- * Every time you run the program, the OS assigns address to your variables but once it is done executing it is destroyed.

- * If you run the program again it will have different address.

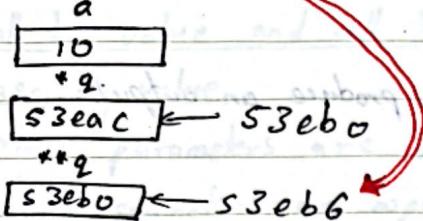
Dereference : printf("Pointed Value = %d\n", *p);

(go to address which p holds and print the value in that address). (Can be used to perform arithmetic operations)

$\&P$ - address of P
 P - contents of P

* P - value of the variable

int $**q$



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int $a = 10;$
int $*p = \&a;$
int $**q = \&p;$

Alias of a pointer

int $*p = \&a;$

int $*q = p;$

$p \rightarrow 10$
 $q \rightarrow 10$

* If you change one other will also change

printf ("%p == value %d\n", p, *p);

printf ("%p == value %d\n", q, *q);

70d8, 70d4, 10 printf ("%p address=%p, p content=%p, value=%d\n",
 $\&p$, p , $*p$)

70e0, 70d8, 70d4 printf ("%p address=%p, p content=%p, value=%d\n",
 $\&q$, q , $*q$)

(**q) (*p) a - value = 10

(*q) (p) a - address = 70d4

a 10 70d4

(q) (&p) p - address = 70d8

p 70d4 70d8

(&q) q - address = 70e0

q 70d8 70e0

* Pointer to a pointer to a pointer to a double

double $***p;$

Arithmetic operations using dereference.

int $a = 10;$

int $*p = \&a;$

int $x = *p + 5;$

$*p = *p + 20$

printf ("%d\n", x); 15

printf ("%d\n", a); 30

$$*p = *p + 20$$

step 1: take value in address in pointer p

step 2: do the addition (30)

step 3: assign 30 to the value in address in pointer p

*∴ you can read and assign values using dereference.

* Dangling pointer is when you got two pointers pointing to the same memory location and you delete one pointer causing the ~~other~~ pointer to point to nothing.

different functions - Q* (Q to 22.660 - Q*)
Q to start no - Q

Functions

Date

No. in file

* Block of functionality which takes inputs & produce an output

* Not all functions take inputs

* Not all functions produce outputs.

* Functions allow the reuse of logic within the same program.

<return-type><function-name> [<params>, *]

{ function-body }

//function body

}

Return type

* Function can have any return type (primitive / user defined)

* If it doesn't return a value "void" type is used

* Void can mean,

1/ No value

2/ No type

3/ No parameter.

* If you are
trying to return
a double it will
only return the
integer part.

ex: int add(int a, int b)

{ return a+b }

int main()

{ printf("%d\n", add(5, 6)); }

Parameters

* Inputs to functions

* Not mandatory to have parameters

* But parameters cannot be optional

* Arguments are the real functional values passed onto the parameter.

int sum (int a, int b)
return a+b

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sum(6,7)

actual parameters

Call by value and call by reference.

Pass by value

- * When parameters are passed using the pass by value methods, the actual parameters are copied to the formal parameters.

void test (int a)

```
{     a → 32 bit  
    printf ("%d\n", a);  
}
```

int main()

```
{ int b=45;  
  test(b);  
  printf ("%d\n", b);  
}
```

- * Here the value of b gets assigned to int a of void test function. By making a copy of b in a.

PASS BY VALUE

- * Actual variables do not change
- * Exists 2 copies of parameters
- * Safer than pass by reference.

Pass by reference

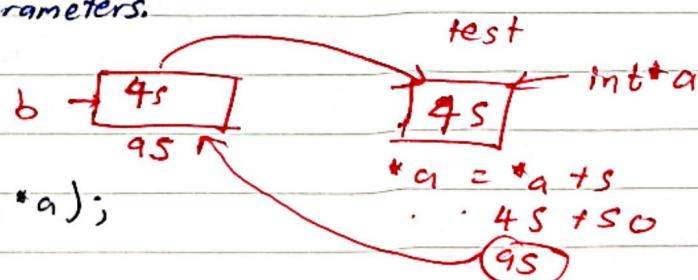
- * When parameters are passed, the address of the actual parameters are passed as formal parameters.

void test (int *a)

```
{ *a = *a + 50  
  printf ("%d %d\n", a, *a);  
}
```

int main()

```
{ int b=45;  
  test(&b);  
  printf ("%d %d\n", b, b);  
}
```



- * Here the address of b is passed as the parameter to the test function.

- * Actual variables are modified within the function.

- * Pointers are used.

- * Preferred when large amounts of data is being passed to the function.

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(for input) and for
the output

scanf()

Date

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Syntax

int scanf (const char *format, variable_address);

Console → in (scanf)
out (printf)

* scanf returns 3 types of values,

1 (>0) values are converted and assigned to the variables

2 (=0) no value has been assigned to the variables

3 (<0) error of stdin has occurred or end of file is reached before the assignment was made.

* Supports scanset specifiers which are represented by %[].

To scan stuff like strings.

* Read can also be used for reading from standard input

notably, there will be nothing in between no returning value & printf

3 files

[25]

printf

scanf

scanf

scanf

* When array element is not defined in modern versions it will give zero, but in old versions it might give garbage values.

* Cannot resize since you cannot guarantee contiguous space remaining in memory to resized.

X	50	14	60	90	87	3	24	X
	32	22	32	32	32	32	32	

Date 3/7/25

Programming in C

Arrays : Group of elements of the same data type that is stored contiguously in memory.

- * AKA homogeneous data structure
- * Fixed size data structures

`<data-type> <array-name> [<size>];`

1, Size either given literally
2, Or determined by the number of elements

`<data-type> <array-name> [] = {<value 1>, ... <value n>};`

- * Has zero based indexing
- * To find the length of the array use `size of C`)
- * Elements of the array can be accessed through pointer arithmetic.

* getting the size of an array: `printf("%d\n", sizeof(array-name)/sizeof(first-element));`

`int a[] = {20, 50, 70, 94, 86};`

`printf("%p\n", a);` // gives the address of a which is the starting point of the array (ie: first element address).

`printf("%p\n", &a[2]);` // address of 3rd element

`printf("%d\n", a[2]);` // value of 3rd element.

a - address of array & `a[X]` - address of an element

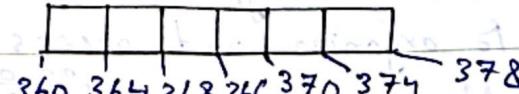
`a[X]` - value of an element.

`for (int i=0; i<5; i++)` Outputs the addresses of the elements.

`printf("%p\n", a+i);`

`}`

`printf("%d\n", *(a+i));`



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Some other terms: member of block for 23 friends person adult *
extreme terms not another blo in the grass out of the fence
memory storage capacities, memory speed up using timer to
decide of promotion of

- * Important: When you change the data type of the array
the gaps in which the pointers of size change to take one
assigned addresses also change according to the
data type.

ex: int 4 byte gaps long 8 byte gaps type of pointers

long
↓

a = \boxed{a}

MATTER,

$$a[0] = a + 1(8) \quad (8) * (a+1)$$

$$a[1] = a + 2(8) \quad (16) * (a+2)$$

$$a[2] = a + 3(8) \quad (24) * (a+3)$$

$$a[i] = \left(\boxed{\text{Base Address} + \text{index} * (\text{Size of data type})} \right)$$

a

j

8

* So if you keep increasing the index it will give random values from memory which are outside of the array. Cuz int pointers doesn't really care about the boundaries of the array.

$$\downarrow \quad a[3]*5 : \{ (a + 3 * (4)) \} * 5 : ((a+3)) * 5$$

- * If its not fixed size (size of data type) variable won't work.
- * If its not contiguous you cannot use indexing since it will point out to random values.

Multidimensional Array

* Arrays can have more than one subscript or an index

< data-type > < name > [size][size]...[size];

* Two ways to organize and access elements. In a multi-dimensional array,

1/ Row Major Order

00	01	10	11
10	11	20	21

2/ Column Major Order

00	10	01	11
10	11	20	21



ProMate

start from left side



int a [4][3][2][2]

Date _____

No. _____

reading :- there is an array of 4 cells

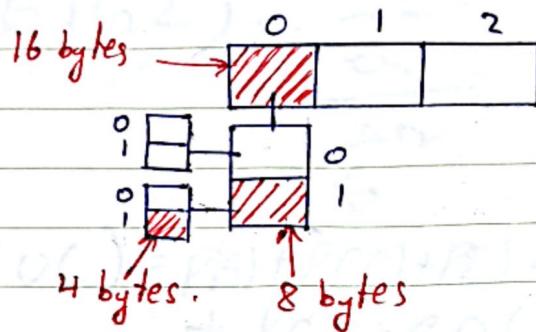
- inside that array there is an array of 3 cells

- inside that array there is an array of 2 cells

- inside that array there is an array of 2 cells.

Total cells = $4 \times 3 \times 2 \times 2 = 48$ cells

ex: $a[3][2][2] \Rightarrow a[0][1][1] = 6$



Row Major Order form Vs Column Major Order

ex: for (int i=0; i<2; i++)
 for (int j=0; j<2; j++)

$a[i][j]$

0 0	c _{i,j}
0 1	
1 0	
1 1	

for (int j=0, i<2; j++, i++)
 for (int i=0; i<2; i++)

$a[j][i]$

0 0	c _{j,i}
1 0	
0 1	
1 1	

* Main advantage of row major order is that it is more efficient than column major order.

* Both ways work but row major is more efficient.

Types of pointers

Void pointers

```
void *Lname;
```

- * This pointer can hold any type of variable address, i.e; primitive, user defined and etc. (Basically can be typecasted to any type).

ex: void *

int a = 10

* p = &a;

float f = 6.0f;

* p = &f;

- * Void pointers cannot be dereferenced

printf("%d\n", *p); ~~gives an error (cannot dereference void pointers)~~

so to dereference a void pointer, you basically need to define the type of variable it ~~points to.~~ points to. (through bottom 2: (1) thought

printf("%d\n", *(Cint)p); ~~works.~~ ~~|| *(<data-type>*) Lpointername~~

- * Standard C doesn't allow for void pointer arithmetic. But gncm3 allows by

- * Enables the implementation of generic functions
- * Mainly used in implementing data structures.
- * Discard type safety which is a disadvantage.

Function Pointers

Date _____

No. _____

return-type (*ptr-name)(param₁, param₂, ...);

* Enables calling a function using a pointer

* It points out to the start of the executable code.

ex: int add(int a, int b) \Rightarrow int (*X)(int, int); x = &add;

int Y() \Rightarrow int (*X)(); x = &Y;

* Useful in implementing call back functions. (can be used to create wrappers around existing functions).

Reading a function pointer: spiral rule is applied clockwise starting from the unknown element.

int *a[20]; : point of array of size twenty with type int

void *f (int, int); : function pointer(f) pointing to two integer parameters and that points to a void pointer.

void * (*f)(int, double (*t)(int, int)); (f) is a function pointer with parameters int and a double (t) which is a function pointer with two integer parameters which points out to a void pointer.

void * (*f (int, double (*t)(int, int))) (int); :

void * (*n) (int)

writing the whole thing as a function pointer.

Constants and Pointers

Date

No

Pointer to a constant

const data-type * <pointer-name>;

* Data pointed by the pointer can only be read, but cannot be changed by the pointer.

* Ensures that you accidentally do not change data pointed by the pointer.

ex: int a = 23;
const int * ptr;

ptr = &a;

int b = 50;

ptr = &b;

printf("%d\n", *ptr);

* But it is possible to change the value by changing the variable the pointer points to, or by directly changing the value of the variable a.

Constant pointers

data-type * const <pointer-name>;

* Should be assigned at initialization

* Constant pointers will always point to (the) same memory location and once assigned, cannot be changed.

Summary

pointer to a constant

Reassign



Change Value initialized



anytime

constant pointer



@ decl.

Constant Pointer to a
constant



@ decl.

(const <data-type> * const <pointer-name>);

ProMate

2 without node begin from point to block position < diagram >

Null Pointers

- * A pointer that doesn't point to any location.
- * Obtained by assigning a pointer to a zero(0) or NULL.
- * It will result in different behaviours across different machines.

ex: `int* a = 0/NULL;`

Null Pointer

- * Initialize pointers without an initial address

- * To check for null before accessing a memory location.

- * To pass to functions when there is

- * ~~not~~^{valid} memory address to provide,

- indicating that the pointer intentionally points to nothing.

Wild Pointers

- * These pointers are uninitialized and may point to an arbitrary memory location.

ex: `int* a;`
`void* x;`

- * These pointers makes the program work unpredictably or cause crashes

Strings

- * In C, string is a sequence of characters terminated with a null character ('\0').

`char<string_name>[array_size]`

ex: Hello

`H E L L O \0`

array size should be n+1 accounting for the null character.

for: `string char string[]` 10 character should be inserted at the end to show the end of the string.

<string.h> - contains inbuilt string manipulation functions.

%s - format specifier

Scansets: format specifiers used with scanf() in C to read strings that match a specific set or range of characters.

* Scansets do not skip leading whitespaces; so leave a space before the scanset in scanf() if you wanna skip whitespaces.

Structures (better to create outside main(), but it can also be assigned inside main())

* Structures are a user defined type in C.

struct<name>

{

<data-type 1> <variable-name 1>;

<data-type 2> <variable-name 2>;

...

} [variable name 1, variable name 2, ...];

* Here there is a fixed number of variables for the structure

* All items under a structure are called 'members' of that structure.

struct<name>

{

<data-type 1><variable-name 1>;

<data-type 2><variable-name 2>;

}

* Here you can define any number of variables for

struct<struct-name><variable-name>

the structure.

<variable-name><member-name> <===== access a member. > value

Pointer to a structure

struct<struct-name><variable-name>;

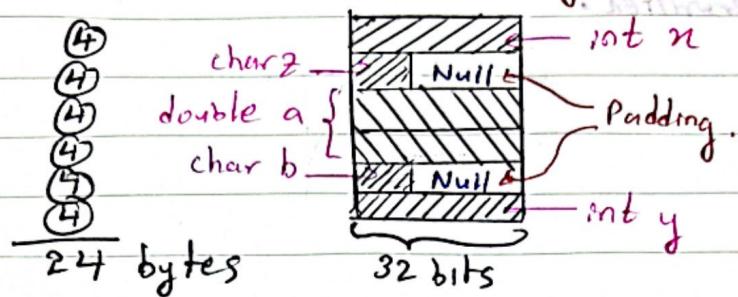
struct<struct-name>*<pointer-name>;

<pointer-name> = &<variable-name>;

<pointer-name> -><member-name> = value;

use arrow operator when accessing members through a pointer to a structure.

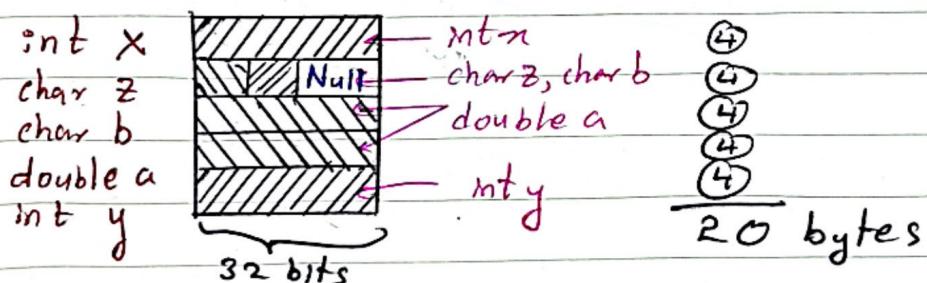
size of structures: $\text{int } x \rightarrow 4$ $\text{int } x \rightarrow 4$ $\text{int } x \rightarrow 4$
~~int y → 4~~ $\frac{8}{8}$ $\text{int } z \rightarrow 1$ $\text{char } z \rightarrow 1$
~~actually = 8~~ $\text{int } y \rightarrow 4$ $\text{double } a \rightarrow 8$
~~9~~ $\text{char } b \rightarrow 1$



$\text{int } y \rightarrow 4$
 $\frac{18}{18}$
~~actually = 24~~

* So basically padding optimizes the memory usage, cus without it the double value will be there in 3 addresses and it will waste computation power cus the machine has to read 3 memory addresses.

* You can optimize the above structure by changing the order of which the variables are declared.



* Members of structures are allocated within a contiguous block of memory.

Unions (size of union = size of largest data member in the union)

* User defined data type that allows data types to be stored in the same memory location.

* Since all members are stored at the same memory location, all members cannot be initialized at the same time.
* Basically it will allocate the size of the largest type.

* Primary application is to save memory.

* Similar syntax to structures when used as a pointer.

* If multiple members are assigned values @ the same time, the data can be ~~overwritten~~ overwritten.

static keyword

• Applied to both variables and functions.

* When applied to variables, it has the property of preserving its value even after they are no longer in scope.

static <data-type> <var-name> = value;

ex: int n()

```
{  
    static int a=0;  
    a++;  
    return a;  
}  
  
main()  
{  
    printf("%d\n", x);  
    printf("%d\n", x);  
}
```

* Without static the output will be 1 and 2.
* But when you use static the value of a will be preserved and next time the function is called the previous call will be incremented. Output is 1 and 2.

* If you don't use static properly you could waste memory, cause if you aren't using the variable again, it will be stored in ram.

* Static variables must not be declared in a structure, but static structures are allowed.

Static functions: static <return-type> <function-name>([params])

```
{  
    . . . . .  
}
```

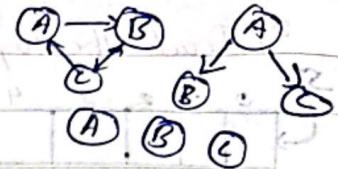
* static functions are restricted to files, they are declared in and mainly used in files.

Data Structures

Stacks

Linear structures : their elements are arranged sequentially, and they are connected to adjacent elements.

Non-Linear structures



* Only Linear structures are done in this course.

* Doesn't necessarily have to be contiguous but arrays are an example for contiguous linear structures.

ex: Arrays, Linked Lists, Stacks, Queues. Linear structures.

ex: Graphs, Trees.

Non-Linear structures.

To get a non-contiguous sequential linear structure, basically the memory blocks are connected by pointers.

Abstract Data Types (ADT)

* It is a class of objects whose behavior is defined by a set of values and operations seen by the user.

* Data structures provide how the behaviour mentioned by an ADT is implemented.

(3) fundamental operations to learn for any data structure,

1/ How to insert data

2/ How to delete data

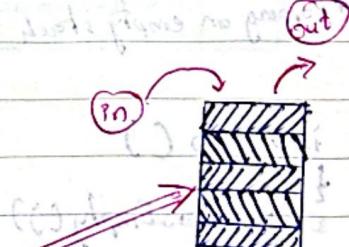
3/ How to search data

* In addition there would be structure specific operations.

Stack (Implemented using arrays or linked lists)

* Follows last in first out policy (LIFO). * LIFO is not only used in stacks.

* Used in backtracking, functional calls, memory management, syntax checking and etc.



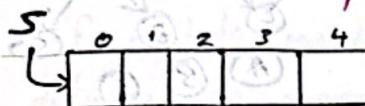
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* Good practice to do is ensure that functions do one thing and one thing only.

* Things to keep track of when using arrays as stacks,

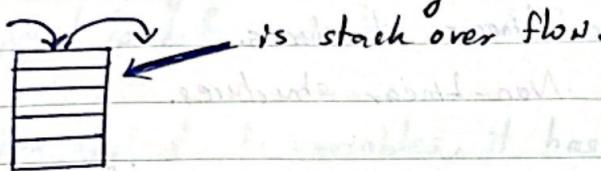
1) Size of array

2) Track the top of stack (TOS) \Rightarrow basically hold the array index of the top element.



Here TOS is -1. exception of in $\text{int isFull}()$ \Leftarrow check if stack is full

insert (push): worst thing that can happen



return $\text{TOS} == \text{size} - 1$; $\boxed{\text{O}(1)}$

} || guards against stack overflow

int isEmpty() \Leftarrow check if stack is empty.

return $\text{TOS} == -1$; $\boxed{\text{O}(1)}$

void push(int v) \Leftarrow To push an element into the stack,

{ if ($\text{!isfull}()$)

{ $\text{S}[++\text{TOS}] = \text{v}$; $\boxed{\text{O}(1)}$

else { printf("Stack is full");

}

}

peek(): retrieves tos without deleting it.

int peek()

{ return TOS ; $\boxed{\text{O}(1)}$

Macros

* It is the name given to a piece of code such that when the preprocessor encounters it in the program, it will be replaced by the code defined.

Object-like Macros.

#define PI 3.146 \nearrow Not terminated by a semicolon

#define add(a,b) (a+b)

* There are also predefined macros

In C like, srand, int, float, etc.

- DATE, - TIME

int pop()

{ if ($\text{!isEmpty}()$)

$\boxed{\text{O}(1)}$

{ return $\text{S}[\text{TOS} - 1]$; \Leftarrow So here you

cannot delete from memory since it is an

array. What is done is, next time an element is pushed, the value will be overwritten.

Memory Layout of a program

(n) Standard Input

(i=0) for i

for loop +

Date

No.

- * C program's memory representation consists of following (5) sections,

Text segment - contains executable instructions, it is shareable so only a single copy is required.

Data segment - Initialized - global and static variables that are initialized

Uninitialized (bss) - for uninitialized global and static variables or ones set to zero.

Heap memory - dynamic memory allocation (basically a collection of memory space where things are stored randomly)

Used by malloc(), free(), realloc() and calloc();

Stack memory - contains program stack

Size mem (in the terminal)

- * Gives the sizes of each section on memory.

* This way, when stack starts from high memory address and

Leaps start from low memory address,

the amount of space allocated is not limited, it can grow until

they touch each other, then program

Low 0

High n

stack

command line parameters.

Text segment

Data segment

empty.

Recursion

- * It is "defining a problem in terms of itself" or in programming terms "a function calling another instance of itself".

- * Recursion has 2 properties: Base case (anchor point)

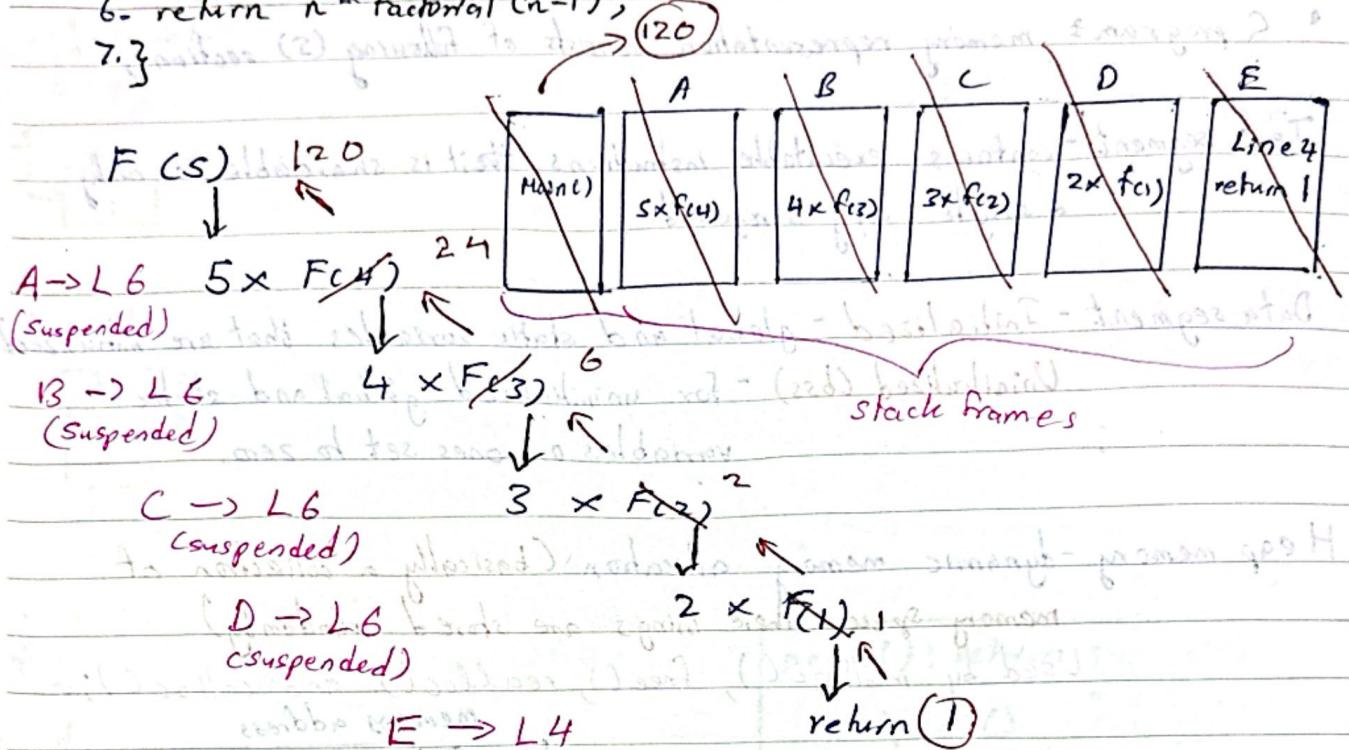
Recursive step

ProMate

eg: 1. int Factorial (int n)

```
2. {  
3.   if (n==1)  
4.     return 1;  
5. else
```

```
6.   return n * Factorial (n-1);  
7. }
```



Excessive Recursion :- when the same problem is solved repeatedly.

ex: int fib (int n)

```
{  
if (n==1 || n==2)  
    return 1;  
else
```

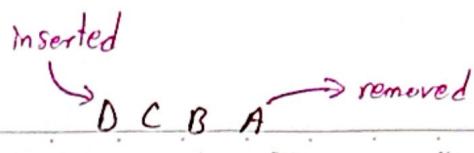
```
    return fib(n-1) + fib(n-2);  
}
```

* Here fib(3), fib(4), fib(2) is

calculated multiple times.

* Recursion should progressively lead to the base case, i.e. it should reduce the problem to smaller and smaller sizes.

* only division and subtraction can lead the recursion to its base case.

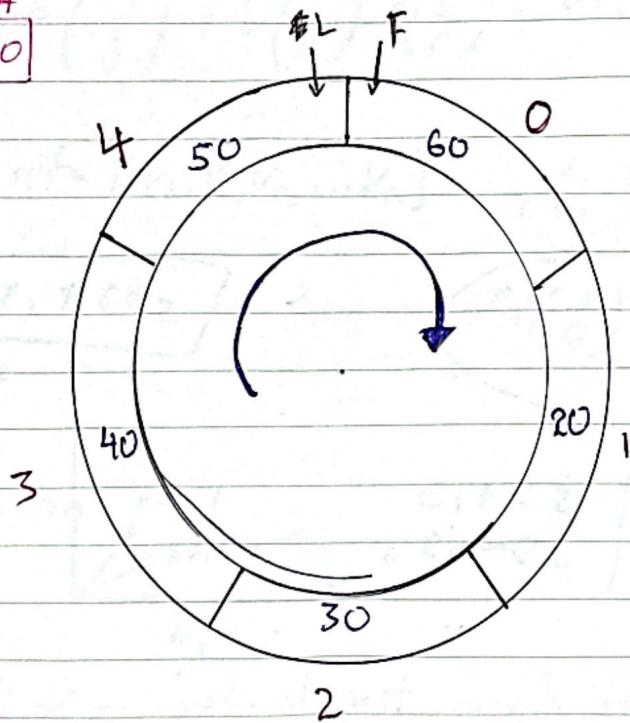


Queues

- * First in first out policy. (FIFO)
- * Implemented using arrays or linked lists.
- * We need two pointers to keep track of first element & last element.

0	1	2	3	4
60	20	30	40	50

1. Enqueue()
2. Dequeue()
3. Front() / Peek()
4. IsFull()
5. isEmpty()



int first, last

↑ keeps track of first element
↑ keeps track of last element

isEmpty

$$F = L = -1$$

OR $F > L$

Is Full

$$L + 1 = F$$

Dynamic Memory Allocation and Linked lists

Dynamic Memory Allocation

- * It is the process of assigning memory to a program @ runtime giving the control to work with data structures of varying sizes.
- * Increases flexibility
- * Increases efficiency: less overhead, no page faults on TLB
- * Occurs in both stacks and the heap memory.

For non-function call based it is assigned from the heap.

`<stdlib.h> => malloc(...), calloc(...)`

`size_t` : guarantees to be big enough to handle the largest possible object in the host system. (Greater negative)

32 bit system \Rightarrow `unsigned int` is most used (.) stack

64 bit system \Rightarrow `unsigned long long`

malloc(..) (Memory allocation)

- * Reserves a block of memory of a specified size from the heap & return a pointer.

`void *malloc(size_t size);`

Creates a void pointer cause it can be type casted to any type

ex: `void *p = malloc(20);`

`int *p = (int *)malloc(20);`

or

`void *a = malloc(20);`

`int *p = (int *)a;`

cast

actually the array size
is 5 ($\frac{20}{4} = 5$)

32 bits
4 bytes
 $\therefore 20 \text{ bytes}$
is 5 integers

* Size parameter expects the size of the memory block required in bytes.

* If no value is not initialized at execution to the function, garbage values present at the time of allocation are the initial values of the memory block.

* If no sufficient space is available to allocate, a null pointer is returned. (∴ check whether it returns a null point)

calloc(...) (contiguous allocation)

Key differences - Initiates each memory block so that each element is zero

- 2 parameters : total elements (n), size of elements
- Returns a void pointer

Void *calloc(current elem, size_t el_size);

malloc(..) faster than calloc(..)

free(..)

* Deallocate memory that was allocated by malloc() and calloc()

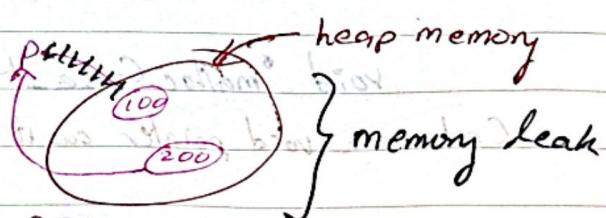
void free(void *pbr)

* Recommend to free memory that has been allocated.

ex: 1. int *p = 0;

2. p = (int*)malloc(100);

3. p = (int*)malloc(200);



* 100 is no longer assigned to anything, therefore no one can use 100 until program ends executing.

• So if the program keeps on leaking memory, it will eventually run slower and slower and finally crash.

automated tools: sudo apt install valgrind

valgrind --leak-check=full <executable> ProMate

realloc(...)

- * Reallocates previously allocated memory when the memory allocated by malloc(..) and calloc() is insufficient.

void *realloc (void *ptr, size_t new-size);

ptr:- pointer previously returned by malloc, calloc or realloc.

- if ptr == NULL, realloc behaves like malloc(new-size)

new-size:- new size in bytes of the memory block.

Register Keyword

register <data-type> <var-name> = <value>;

- * Even though you use register keyword, the compiler ultimately decides whether to put the variable into the register.

Auto Keyword

- * In modern C, writing auto is redundant because it is the default behavior for local variables.
- * It basically means variables are allocated storage when the block (function, loop, etc) is entered. And is deallocated when the block is exited.
- * int a; = auto int a;

Command Line Arguments

int main (int argc, char *argv[]) {

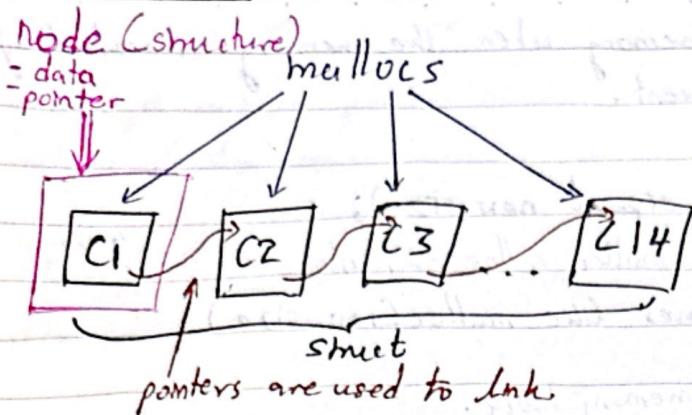
¹ _{ex! .\a.out} ² _x ³ _y ⁴ _z

 argc (argument count) - 4

 argv[] - argument values

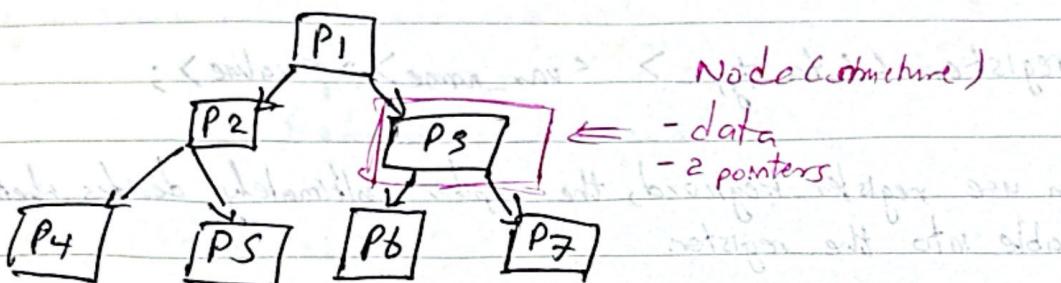
⇒ Basically giving information for a program to run. (NOT USER INPUT)

Linked lists.



```
struct Node {
    int data;
    struct Node* ptr;
}
```

* You can also be more creative,



```
struct Node {
    [<data-type> <member-name>]+; ← + (1 or more)
    [struct Node* <pointer-name>]+; ←
}
```

* The nodes can be created by,

```
void *temp = malloc(sizeof(struct Node)); ←
    struct Node* ←
```

```
struct Node *a = malloc(sizeof(struct Node)); ←
```

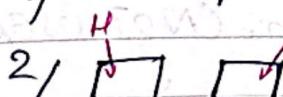
- Add To Head
- Add To Tail
- Delete From Head
- Delete From Tail

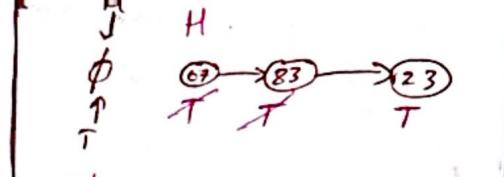
```
a → data = 25
a → ptr = &... ←
```

Singly Linked Lists

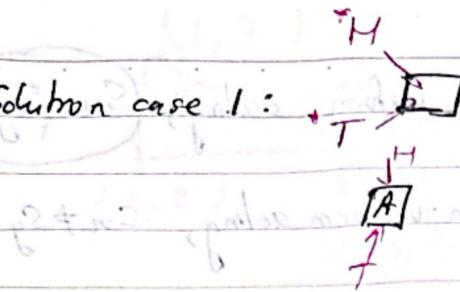
(There is a single link between elements which enables a given node to point to the next node in the list)

$$1/ H = T = \emptyset$$





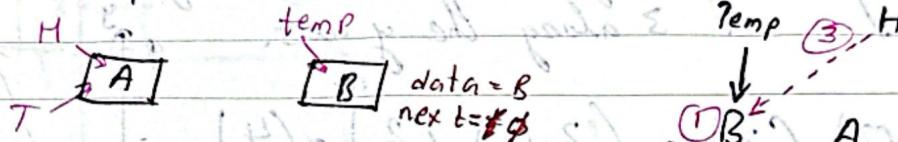
Solution case 1:



Value 67 83 23
1. temp → malloc
2. H = T = temp

addToHead: ⇒ new element becomes the head

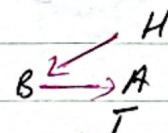
1. temp → malloc (B)



2. temp → next t = H;

3. H = temp

This is wrong!

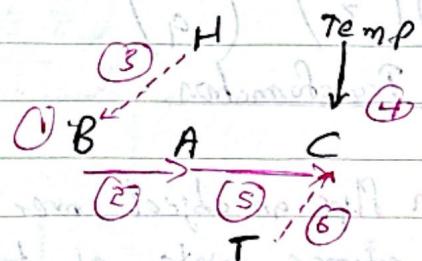


addToTail: ⇒ the element becomes the new tail

3. Temp = malloc

4. T → next = Temp;

5. T = Temp



Inserting node to the middle of the linked list.

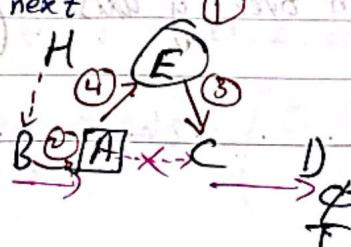
Inserting (E) between A and C

1. temp = malloc (E)

2. find A (start from head and loop until A is found)

3. temp → next = Temp; 2 → next

4. temp2 → next = temp



C programming

Date:

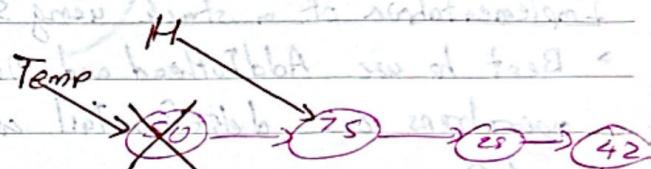
No.

```
#ifndef #YOUR_NAME;    #include "your-name.h" ← include in the file.
#define YOUR_NAME;
xxxxxx
xxxxxx
#endif;
```

- * You can also delete \$ elements from linked lists in 4 ways,
- 1/ delete from head
- 2/ delete from tail
- 3/ delete an arbitrary node
- 4/ delete a single node ($H=T$)

→ ① Delete from head

Case I - if $H=T$ [$H=T=NULL$]



Case II - $Temp = H$

$H = H \rightarrow next$

Afterwards free the temp pointer ($free(Temp)$) *(Big O(1))*

② Delete from tail

Case I - if $H=T$ [$H=T=NULL$]

Case II - loop to find node before tail [$temp \rightarrow next != tail$] $\Rightarrow temp$

$free(Temp \rightarrow next)$

$temp \rightarrow next = NULL$

Big O(n)

for ($temp = head$; $temp \rightarrow next != tail$; $temp = temp \rightarrow next$);

③ Generic Insert *(assuming 2 friends function)*

Insert (X, Location, Pre/Post)

ex: Insert (78, 75, Post) Insert (78, 75, Pre)

$50 \rightarrow (78) \rightarrow 75$

$50 \rightarrow 75 \rightarrow (78)$

if location is head

if Pre \rightarrow addToHead

if location is tail

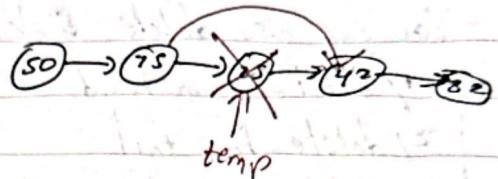
if Post \rightarrow addToTail

* Delete generically: $75 \rightarrow \text{next} = 25 \rightarrow \text{next}$

$\text{temp} = 25$

$\text{temp} \rightarrow \text{next} = \text{null}$

free(temp)



Implementation of a stack using Singly Linked List.

* Best to use AddToHead and DeleteFromHead operations than tail operations cus deleteFromTail costs $O(n)$ while head operations only cost $O(1)$

* push () \rightarrow AddToHead $O(1)$ AddTotal $O(n)$

* pop () \rightarrow DeleteFromHead $O(1)$ DeleteFromTail $O(n)$

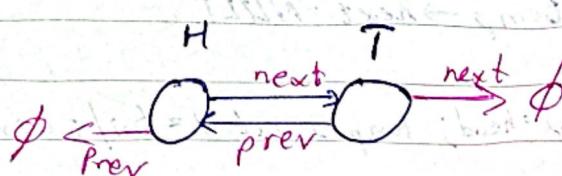
Implementation of a queue using Singly Linked Lists

* Here also DeleteFromTail is best to be avoided.

enqueue () \rightarrow AddToTail $O(1)$ AddTotal $O(n)$

dequeue () \rightarrow DeleteFromHead $O(1)$ DeleteFromTail $O(n)$

Doubly Linked Lists



* More two pointers are used

next and prev. (for the next element & previous element)

struct Node {

[<data_type> member_name]⁺;

struct Node* next;

struct Node* prev;

}

Inserting an element

A ~~→~~ ~~→~~ ~~→~~ B

C \leftarrow C \rightarrow next = B

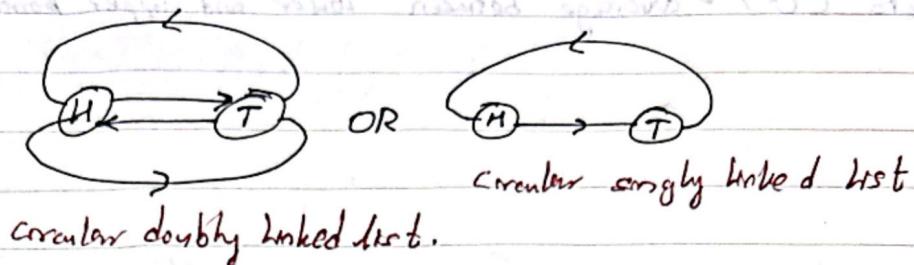
C \rightarrow prev = A

A \rightarrow next = C

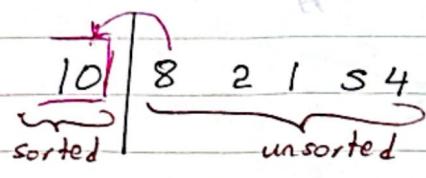
B \leftarrow prev = C

Circular Linked List

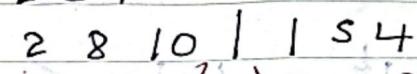
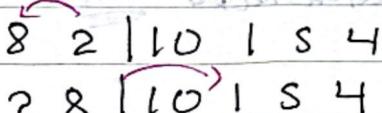
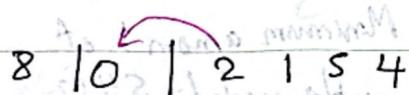
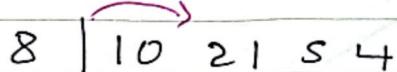
- * Can be created in a singly or doubly linked list manner.
- * Only difference from the previous two is that the head is connected to the tail.



Insertion sort



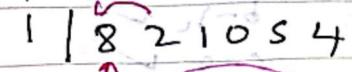
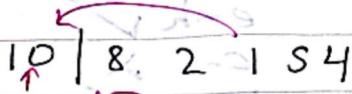
sorting in ascending order



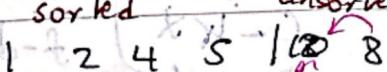
sorted unsorted

Selection Sort

- * Find the minimum element in the array & switch it with the first element in the unsorted part.



sorted unsorted



1 2 4 5 | 8 1 0 (sorted)

Algorithm Analysis

Algorithm sum(A, n)

```

 $s = 0$ 
for ( $i = 0; i < n; i++$ ) —  $n + 1$ 
{
     $s = s + A[i];$  —  $n$ 
}
return  $s;$  —  $1$ 
}

```

Space : $A = n$

$n = 1$

$s = 1$

$i = 1$

$s(n) = n + 3$

$O(n)$

Algorithm Add(A, B, n) * Add two matrices of size $n \times n$.

```

for ( $i = 0; i < n; i++$ ) —  $n + 1$ 
{
    for ( $j = 0; j < n; j++$ ) —  $n \times (n + 1)$ 
        {
             $c[i, j] = A[i, j] + B[i, j]$  —  $n \times n$ 
        }
    }
}

```

Space : $A = n^2$

$B = n^2$

$C = n^2$

$n - 1$

$i - 1$

$j - 1$

$s(n) = 3n^2 + 3$

$O(n^2)$

Algorithm Multiply(A, B, n)

```

for ( $i = 0; i < n; i++$ ) —  $n + 1$ 
{
    for ( $j = 0; j < n; j++$ ) —  $n \times (n + 1)$ 
        {
             $c[i, j] = 0;$ 
            for ( $k = 0; k < n; k++$ ) —  $n \times n \times (n + 1)$ 
                {
                     $c[i, j] = c[i, j] + A[i, k] * B[k, j]$  —  $n \times n \times n$ 
                }
            }
        }
}

```

Space : $A = n^2$

$B = n^2$

$C = n^2$

$n = 1$

$i = 1$

$j = 1$

$k = 1$

$s(n) = 3n^2 + 3$

$O(n^2)$

$f(n) = 3n^3 + 3n^2 + 2n$

$O(n^3)$

$$1) \text{for } (i=0; i < n; i++) - h+1 \\ \{ \text{stmt;} - n \\ \} \frac{n}{O(n)}$$

$$2) \text{for } (i=n; i > 0; i--) - h+1 \\ \{ \text{stmt;} - n \\ \} \frac{n}{O(n)}$$

$$3) \text{for } (i=1; i < n; i=i+20) - h+1 \\ \{ \text{stmt;} - \frac{n}{20} \\ \} \frac{n}{O(n)} \\ f(n) = \frac{n}{20}$$

$$4) \text{for } (i=0; i < n; i++) - h+1 \\ \{ \text{for } (j=0; j < n; j++) - h \times n \\ \{ \text{stmt;} - \frac{h \times n}{n} \\ \} \frac{n}{O(n)}$$

$$5) \text{for } (i=0; i < n; i++) \quad \underline{i \text{ i runtime}} \\ \{ \text{for } (j=0; j < i; j++) \quad \frac{0 \quad 0 \times 0}{1 \quad 0 \times 1} \quad \frac{P}{1 \quad 0+1=1} \\ \{ \text{statement;} \quad \frac{1 \times 0}{2 \quad 0 \quad 2} \quad \frac{2 \times 1}{2 \quad 1+2=3} \\ \} \frac{1+2+3+...+n}{2} \quad \frac{(2 \times n)}{n \quad n} \quad \frac{3 \times 3}{3 \quad 3+3=6} \\ f(n) = \frac{n^2+1}{2} \quad O(n^2) \quad \frac{4 \times 4}{4 \quad 6+4=10} \quad \frac{1+2+3+4+...+k}{k}$$

$$6) P=0; \\ \text{for } (i=1; P \leq n; i++) \quad \therefore P = \frac{k(k+1)}{2}$$

$$\frac{k(k+1)}{2} \geq n \quad O(\sqrt{n})$$

$$7) \text{for } (i=1; i < n; i=i^2)$$

if $n=8$ if $n=10$

$$\{ \text{statement;} \quad \frac{i}{1} \quad \frac{i}{1} \quad \frac{i}{1} \\ \} \quad \frac{i \times 2}{2} = 2 \quad \frac{i}{2} \quad \frac{i}{2} \\ \frac{2 \times 2}{2} = 2^2 \quad \frac{i}{4} \quad \frac{i}{4} \\ \frac{2^2 \times 2}{2} = 2^3 \quad \frac{i}{8} \quad \frac{i}{8} \\ \vdots \quad \vdots \quad \vdots \\ \therefore 2^k > n \quad \log_2 8 = 3 \quad 4$$

Assume $i > 0$

$$\therefore i = 2^k$$

$$\therefore 2^k > n$$

$$\therefore k = \log_2 n. \quad O(\log_2 n)$$

$$2^k$$

$$\log_2 8 = 3$$

$$\log_2 10 = 3.2$$

so we take $\lceil \log_2 n \rceil$

- $i = n$
- 8/ $\{ \text{for } (i=n; i>=1; i=i/2) \}$ $\frac{n}{2^2}$ Assume $i \geq 1$
 { statement; } $\frac{n}{2^3}$ $\therefore \frac{n}{2^k} \leq 1 \text{ for } k \geq 1$
 $\frac{n}{2^k}$ $\frac{n}{2^k} = 1 \text{ for } k = \log_2 n$
 $n = 2^k$ $k = \log_2 n$ $O(\log_2 n)$
- 9/ $\{ \text{for } (i=0; i*i < n; i++) \}$ $\frac{n}{2^k}$ $i^2 = n$
 { statement; } When $i^2 > n$ $i = \sqrt{n}$ $O(\sqrt{n})$
 $i^2 = n$
- 10/ $\{ \text{for } (i=0; i < n; i++) \}$ n
 { statement; }
- 11/ $\{ \text{for } (j=0; j < n; j++) \}$ n
 { statement; }
- 11/ $\{ \text{for } (i=1; i < n; i=i*2) \}$
 { $p++;$ } $P = \log n$
 $j = \log n$
 $\{ \text{for } (j=1; j < p; j=j*2) \}$ $O(\log \log n)$
 { statement; } $\log P$
- 12/ $\{ \text{for } (i=0; i < n; i++) \}$ n
 { $\{ \text{for } (j=1; j < n; j=j*2) \}$ $n \log n$
 { statement; } $\frac{n \times \log n}{2^n \log n + n}$ $O(n \log n)$
- SUMMARY
- $\text{for } (i=0; i < n; i++)$ $- O(n)$
 - $\text{for } (i=0; i < n; i=i+2)$ $- \frac{n}{2} O(n)$
 - $\text{for } (i=n; i>1; i--)$ $- O(n)$
 - $\text{for } (i=1; i < n; i=i*2)$ $- O(\log_2 n)$
 - $\text{for } (i=1; i < n; i=i*3)$ $- O(\log_3 n)$
 - $\text{for } (i=n; i>i; i=i/2)$ $- O(\log_2 n)$
 - $\text{for } (i=0; i < n; i++)$ $- \sqrt{n} O(1)$

File Handling in C

Date

No

1/ File Pointer : File *fp;

2/ Open a file : fopen()

Modes of opening files : 1/ open for reading : "r"

2/ Open for writing (Creates if it doesn't exist) : "w"

3/ Append to end of file : "a"

4/ Read and write : "r+"

5/ Write and read (truncates) : "w+"

6/ Append and read : "a+"

3/ Write to a file : fprintf() or fputs()

4/ Read from a file : fscanf() or fgets()

5/ End of file : feof(fp)

6/ Error checking : ferror(fp)

7/ Closing a file : fclose(fp) (Always close. It prevents memory leaks and corruption.)