

## Day 20

### Assignment 2: Longest Common Subsequence

Implement `int LCS(string text1, string text2)` to find the length of the longest common subsequence between two strings.

A)

*Java implementation of the function LCS that computes the length of the longest common subsequence (LCS) between two strings using dynamic programming.*

**Java:**

```
public class LongestCommonSubsequence {

    public static int LCS(String text1, String text2) {
        int m = text1.length();
        int k = text2.length();
        int[][] dp = new int[m + 1][k + 1];

        // Build dp[][] in bottom-up manner
        for (int i = 1; i <= m; i++) {
            for (int j = 1; j <= k; j++) {
                if (text1.charAt(i - 1) == text2.charAt(j - 1)) {
                    dp[i][j] = dp[i - 1][j - 1] + 1;
                } else {
                    dp[i][j] = Math.max(dp[i - 1][j], dp[i][j - 1]);
                }
            }
        }

        return dp[m][k];
    }

    public static void main(String[] args) {
        String text1 = "abcde"; // Example string 1
        String text2 = "ace"; // Example string 2

        System.out.println("Length of LCS = " + LCS(text1, text2));
    }
}
```

## Explanation:

### 1. Initialization:

- $dp[i][j]$  will hold the length of the LCS of the substrings  $text1[0..i-1]$  and  $text2[0..j-1]$ .
- $dp[0][*] = 0$  and  $dp[*][0] = 0$  because the LCS with any empty string is 0.

### 2. Filling the DP Table:

- Iterate over each character of  $text1$  and  $text2$ .
- If  $text1[i-1] == text2[j-1]$ , then the characters match, and the LCS up to  $i$  and  $j$  is  $dp[i-1][j-1] + 1$ .
- If the characters do not match, then the LCS up to  $i$  and  $j$  is the maximum of the LCS without the current character of  $text1$  ( $dp[i-1][j]$ ) or without the current character of  $text2$  ( $dp[i][j-1]$ ).

### 3. Result:

- The result will be in  $dp[m][k]$ , which represents the length of the LCS of  $text1$  and  $text2$ .

This solution efficiently computes the length of the LCS using dynamic programming with a time complexity of  $O(m \times k)$  and a space complexity of  $O(m \times k)$ , where  $m$  and  $k$  are the lengths of  $text1$  and  $text2$ , respectively.