Binary Search Tree(BST)

- A binary search tree, also known as ordered or sorted binary tree,
- the nodes are arranged in an order.
- The nodes of the tree store a key and each has two distinguished sub-trees
- binary search property:
 - the key in each node is greater than any key stored in the left sub-tree,
 - and less than or equal to any key stored in the right sub-tree

Binary Search Tree(BST) Operations

- Binary search trees support three main operations:
 - Lookup (checking whether a key is present)
 - Insertion
 - deletion.

New nodes are inserted as leaf nodes in the BST.

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```
case(key < root->key):
     recurse down left subtree
case(key >= root->key):
     recurse down right subtree
case(root == NULL):
     create new node
```

```
struct node *insert(struct node *root, int key){
  if (root==NULL)
    root = createNode(key);
 else if (key < root->key)
    root->lch = insert(root->lch, key);
 else // key >= root->key
    root->rch = insert(root->rch, key);
  return root;
```

root = NULL root = Insert(root, 8)

```
struct node *insert(struct node *root, int key){
  if (root==NULL)
    root = createNode(key);
  else if (key < root->key)
    root->lch = insert(root->lch, key);
  else // key >= root->key
    root->rch = insert(root->rch, key);
  return root;
}
```

```
if(r = NULL) r = createNode(8)
return r
```

```
root = NULL
root = Insert(root, 8)
```

```
0x0 8 0x0
1024
```

```
struct node *insert(struct node *root, int key){
  if (root==NULL)
    root = createNode(key);
  else if (key < root->key)
    root->lch = insert(root->lch, key);
  else // key >= root->key
    root->rch = insert(root->rch, key);
  return root;
}
```

```
if(r = NULL) r = createNode(8)
return r
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```
root = NULL
root = Insert(root, 8)
```

```
0x0 8 0x0
1024
```

```
struct node *insert(struct node *root, int key){
  if (root==NULL)
    root = createNode(key);
  else if (key < root->key)
    root->lch = insert(root->lch, key);
  else // key >= root->key
    root->rch = insert(root->rch, key);
  return root;
}
```

```
If( key < r->key)
r ->lch = Insert(r->lch, 3)
return r
```

```
root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
```

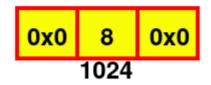
```
0x0 8 0x0
1024
```

```
struct node *insert(struct node *root, int key){
  if (root==NULL)
    root = createNode(key);
  else if (key < root->key)
    root->lch = insert(root->lch, key);
  else // key >= root->key
    root->rch = insert(root->rch, key);
  return root;
}
```

```
if(r = NULL) r = createNode(3)
return r
```

```
If( key < r->key)
r ->lch = Insert(r->lch, 3)
return r
```

```
root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
```

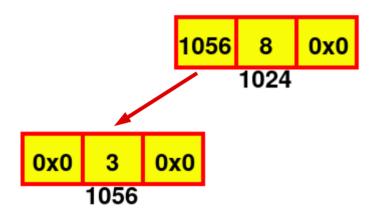


```
0x0 3 0x0
1056
```

```
if(r = NULL) r = createNode(3)
return r
```

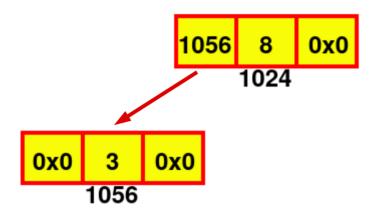
```
If( key < r->key)
r ->lch = Insert(r->lch, 3)
return r
```

```
root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
```



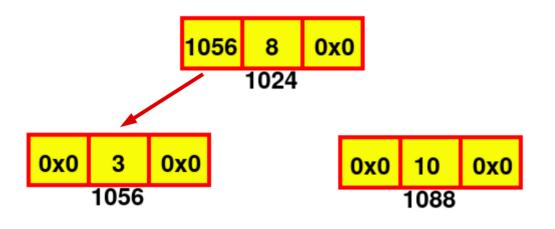
```
If( key < r->key)
r ->lch = Insert(r->lch, 3)
return r
```

```
root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
```



```
If( key > r->key)
r ->rch = Insert(r->rch, 10)
return r
```

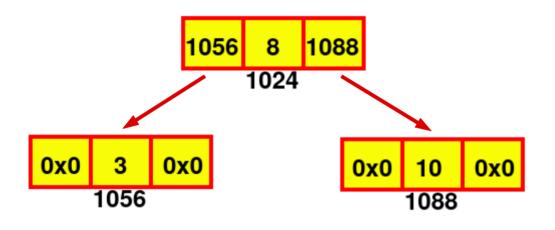
```
root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
root = Insert(root, 10)
```



```
if(r = NULL) r = createNode(10)
return r
```

```
If( key > r->key)
r ->rch = Insert(r->rch, 10)
return r
```

```
root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
root = Insert(root, 10)
```



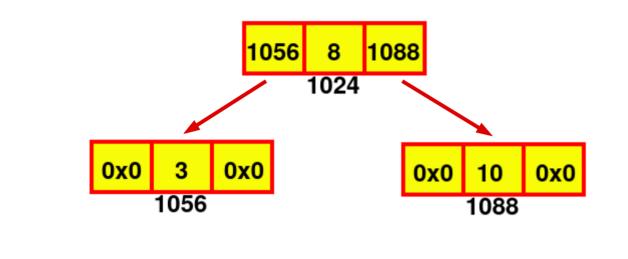
```
If( key > r->key)
r ->rch = Insert(r->rch, 10)
return r
```

```
root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
root = Insert(root, 10)
```

0x0

1120

0x0



```
if(r = NULL) r = createNode(1)
return r

If( key < r->key)
r ->lch = Insert(r->lch, 1)
return r

If( key < r->key)
```

 $r \rightarrow lch = lnsert(r \rightarrow lch, 1)$

return r

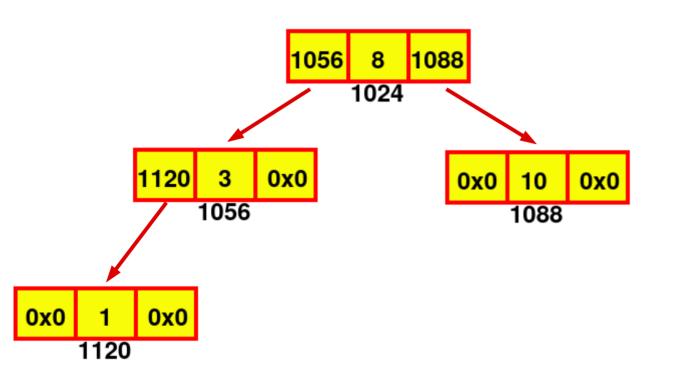
```
root = NULL

root = Insert(root, 8)

root = Insert(root, 3)

root = Insert(root, 10)

root = Insert(root, 1)
```

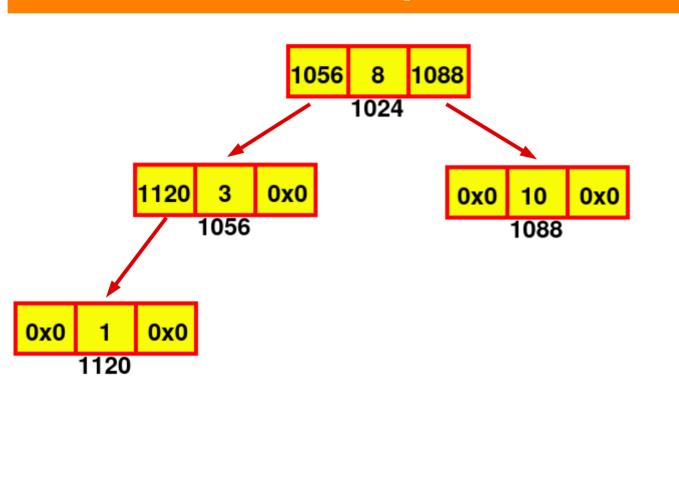


```
If( key < r->key)
r ->lch = Insert(r->lch, 1)
return r

If( key < r->key)
```

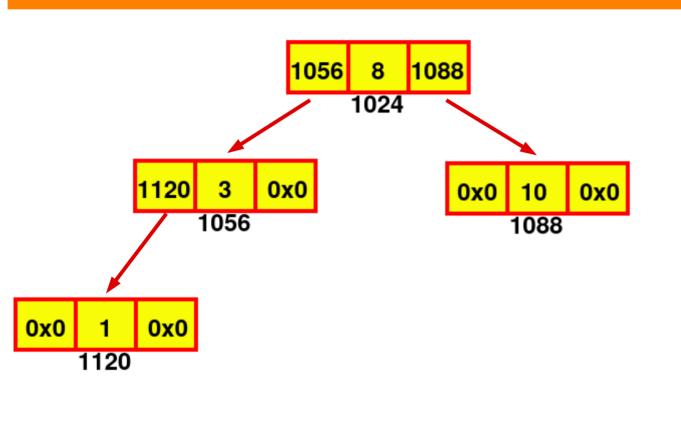
```
If( key < r->key)
r ->lch = Insert(r->lch, 1)
return r
```

```
root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
root = Insert(root, 10)
root = Insert(root, 1)
```



```
If( key < r->key)
I ->rch = Insert(r->lch, 6)
return r
```

```
root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
root = Insert(root, 10)
root = Insert(root, 1)
root = Insert(root, 6)
```



```
If( key > r->key)
r ->rch = Insert(r->rch, 6)
return r

If( key < r->key)
I ->rch = Insert(r->lch, 6)
return r
```

```
root = NULL

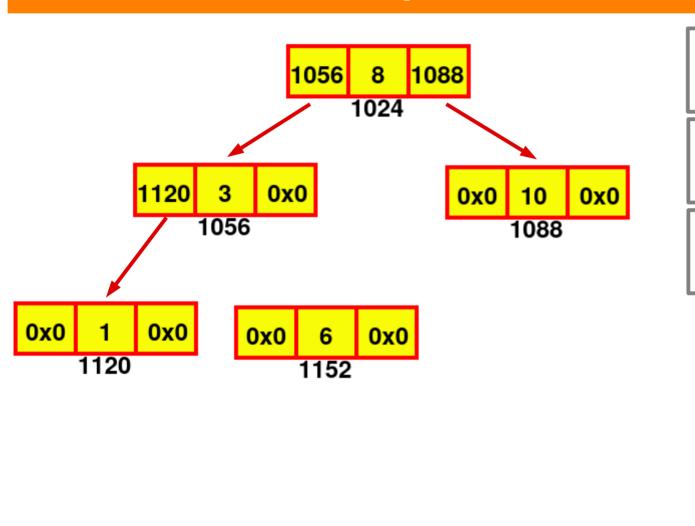
root = Insert(root, 8)

root = Insert(root, 3)

root = Insert(root, 10)

root = Insert(root, 1)

root = Insert(root, 6)
```



```
if(r = NULL) r = createNode(6)
return r

If( key > r->key)
r ->rch = Insert(r->rch, 6)
return r

If( key < r->key)
I ->rch = Insert(r->lch, 6)
return r
```

```
root = NULL

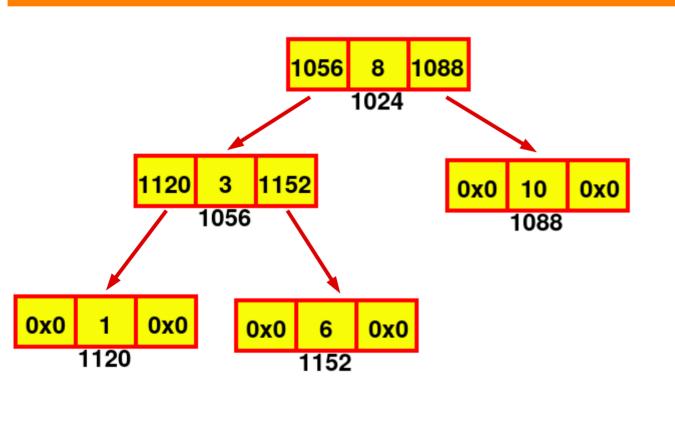
root = Insert(root, 8)

root = Insert(root, 3)

root = Insert(root, 10)

root = Insert(root, 1)

root = Insert(root, 6)
```



```
If( key > r->key)
r ->rch = Insert(r->rch, 6)
return r

If( key < r->key)
I ->rch = Insert(r->lch, 6)
return r
```

```
root = NULL

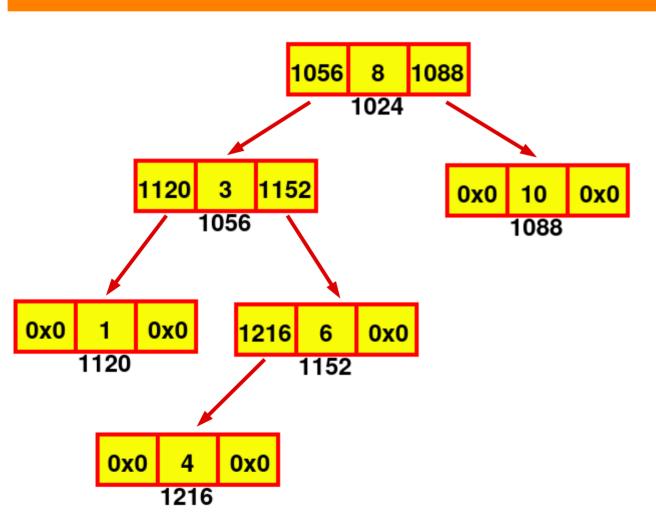
root = Insert(root, 8)

root = Insert(root, 3)

root = Insert(root, 10)

root = Insert(root, 1)

root = Insert(root, 6)
```



if(r = NULL) r = createNode(4)
return r

If(key < r->key)
r ->lch = Insert(r->lch, 4)
return r

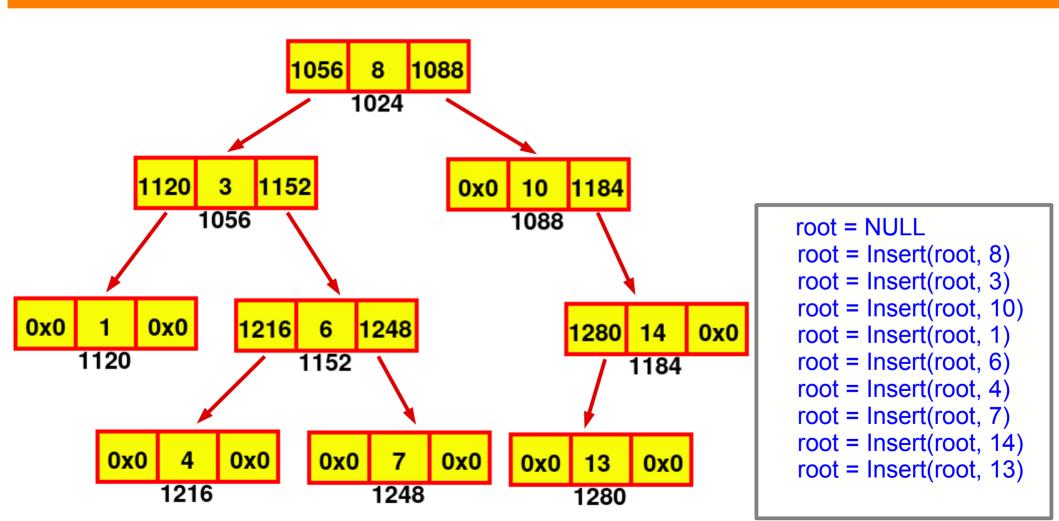
If(key > r->key)

r ->rch = Insert(r->rch, 4)
return r

If(key < r->key)

 $r \rightarrow lch = lnsert(r \rightarrow lch, 4)$

root = NULL
root = Insert(root, 8)
root = Insert(root, 3)
root = Insert(root, 10)
root = Insert(root, 1)
root = Insert(root, 6)
root = Insert(root, 4)



Binary Search Tree(BST)

- Since in insertion we need to traverse to till some leaf node,
- So the running time complexity of BST insertion is O (h) here h is the height of the tree.
- However, the worst case for BST insertion is O (n) here
 n is the total number of nodes in the BST, because an
 unbalanced BST may degenerate to a linked list.
- If the BST is height-balanced the height is O (log n)
- So insertion will take O(log n) time

Binary Search Tree(BST)

- The time complexity of operations on the binary search tree is directly proportional to the height of the tree.
- the nodes in a BST are laid out in such a way that each comparison skips about half of the remaining tree, the lookup performance is proportional to that of binary logarithm.
- The performance of a binary search tree is dependent on the order of insertion of the nodes into the tree;
- The complexity analysis of BST shows that, on average, the insert, delete and search takes O(log n) for n nodes.
- In the worst case, they degrade to that of a singly linked list:
 O(n).

BST Search

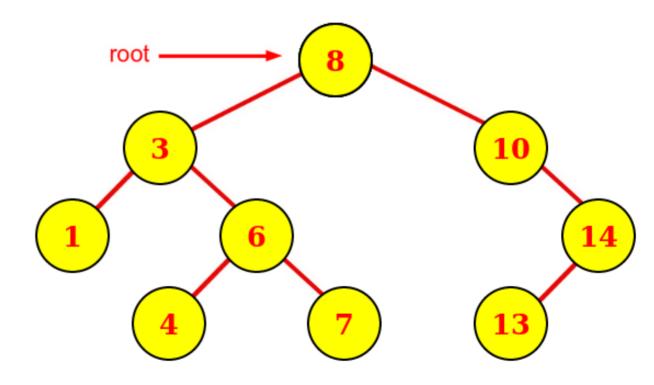
```
struct node* search(struct node *root, int key){
   if(root == NULL || root->key == key)
        return root;
   if(key < root->key)
        return search(root->lch,key);
   if(key > root->key)
        return search(root->rch,key);
}
```

BST search time complexity

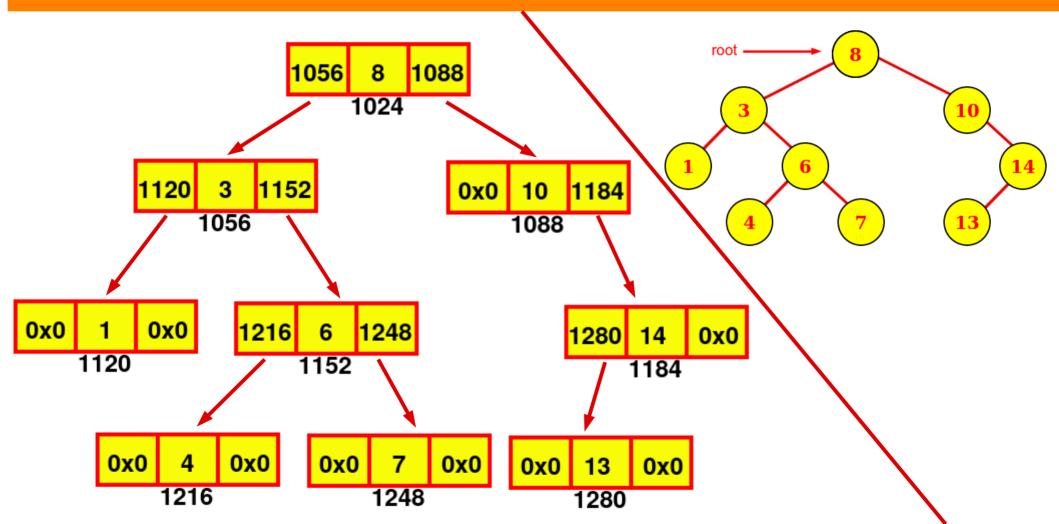
- In BST search we need to traverse to till some leaf node
- So the running time complexity of BST search is O (h) here h
 is the height of the tree.
- However, the worst case for BST search is O (n) here n is the total number of nodes in the BST, because an unbalanced BST may degenerate to a linked list.
- If the BST is height-balanced the height is O (log n)
- So search will take O(log n) time
- The time complexity of operations on the binary search tree is directly proportional to the height of the tree.

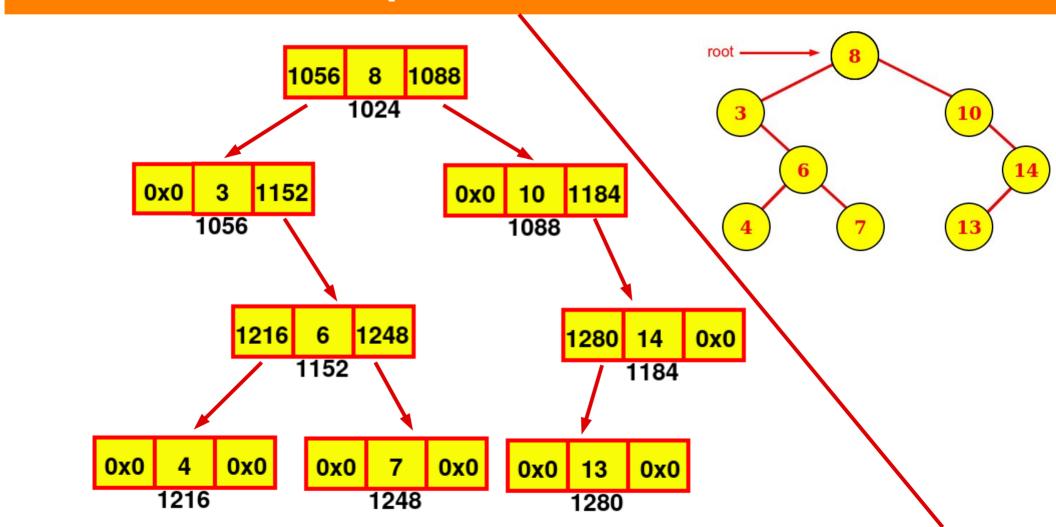
- First we need to find the node to be deleted
- Replace node to be deleted with its successor

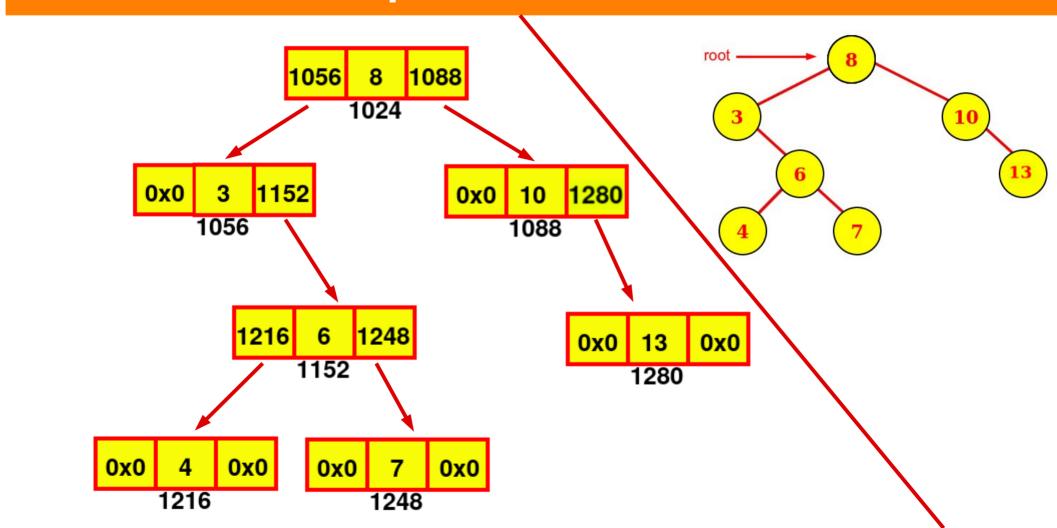
- First we need to find the node to be deleted
- Replace node to be deleted with its successor



- If node to be deleted(z) is a leaf node,
 - the parent node's pointer to the z replaced with null
- If node to be deleted(z) has a single child node
 - the parent node's pointer to the z replaced with z child
- If z has both a left and right child
 - the successor of z (let it be y) takes the position of z in the tree(Replace node with successor).
 - y will be deleted







BST Search

```
struct node* search(struct node *root, int key){
   if(root == NULL || root->key == key)
        return root;
   if(key < root->key)
        return search(root->lch,key);
   if(key > root->key)
        return search(root->rch,key);
}
```

BST Deletion

```
struct node* deletion(struct node* root, int key){
  if (root == NULL) return root;
  else if (key > root->key) root->rch = deletion(root->rch, key);
  else if(key < root->key) root->lch = deletion(root->lch, key);
  else{
     if(root->lch && root->rch){     /*if node has both right and left childs*/
        struct node* sucessor = root->rch;
        while (sucessor->lch != NULL)
           sucessor = sucessor->lch;
        root->key = sucessor->key;
        root->rch = deletion(root->rch, sucessor->key);
     }else {
        if(root->rch) return root->rch; /*if node has right child*/
        else return root->lch; /*if node has left child or no childs*/
        free(root);
   return root;
```