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# **Recursive Path Ordering on Strings: Two Implementations**

**(Master's Project)**

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## **Abstract**

We have implemented an algorithm to determine whether a given pair of strings can be made RPO-comparable by selecting a partial order on their function symbols. The algorithm is implemented in Python3.

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# 1 Introduction and Basic Theory

Termination orderings play a very important role in automated reasoning, in showing that a sequence of rewrite steps will eventually terminate. One of the most influential is the *recursive path ordering* introduced by Nachum Dershowitz. The key idea is to extend an ordering  $\succ$  on the symbols to terms. Restricted to strings, the ordering is defined as follows:

Let  $\Sigma$  be an alphabet and  $\succ$  be an ordering on  $\Sigma$ . Then  $x >_{rpo} y$  if and only if one of the following conditions hold:

- (1)  $y = \varepsilon$  and  $|x| > 0$ .
- (2)  $x = au$ ,  $y = av$ , and  $u >_{rpo} v$ .
- (3)  $x = au$ ,  $y = bv$ , and either
  - (3a)  $u \geq_{rpo} y$ , or
  - (3b)  $a \succ b$  and  $x >_{rpo} v$ .

One well-known property of the *rpo* is

**Lemma 1.**  $\forall x \forall y \forall z : (xz >_{rpo} yz) \text{ if and only if } (x >_{rpo} y)$ .

Thus before comparing any two strings their common suffixes can be removed.

If the symbol ordering  $\succ$  happens to be total (i.e., the ordering is linear), an alternative characterization was introduced in [3]. Let  $\max(w, \Sigma)$  stand for the *maximal* (highest) symbol of  $\Sigma$  that occurs in  $w$  and let  $\text{mul}(w, \Sigma)$  be the number of times *this* symbol occurs in  $w$ , i.e.,  $\#_{\max(w, \Sigma)}(w)$ . Now  $w >_{rpo} w'$  iff one of the following holds:

- 1.  $\max(w, \Sigma) \succ \max(w', \Sigma)$
- 2.  $\max(w, \Sigma) = \max(w', \Sigma)$  and  $\text{mul}(w, \Sigma) > \text{mul}(w', \Sigma)$
- 3.  $\max(w, \Sigma) = \max(w', \Sigma)$ ,  $\text{mul}(w, \Sigma) = \text{mul}(w', \Sigma)$ ,

$$w = w_0 w_1 w_2 \dots w_k$$

$$w' = u_0 u_1 u_2 \dots u_k$$

and there exists  $0 \leq i \leq k$  such that  $w_i >_{rpo} u_i$  and for all  $j > i$  we have  $w_j = u_j$ .

If the strings do not have any common suffixes, then the last condition can be changed to

$$3'. a = \max(w, \Sigma) = \max(w', \Sigma), \text{mul}(w, \Sigma) = \text{mul}(w', \Sigma),$$

$$w = w_0w_1w_2 \dots w_k$$

$$w' = u_0u_1u_2 \dots u_k$$

and  $w_k >_{rpo} u_k$ .

## 2 Implementation

The program for Basic Theory is implemented as follows:

The program takes the input as the pair of strings and symbol ordering.

1. The program calls `rpo()` function with the input of string1, string2 and precedence ordering where it checks if pair of two strings are equal or not. If pair of two strings are equal then its return false else it calls `rpoeq()` function recursively.
2. `rpoeq()` is function which is having the input of string1, string2 and precedence ordering and is checks for the below conditions:
  - a. Firstly, It removes the common suffix by using function `remove_common_suffix()`.
  - b. Secondly, checks whether the pair two strings are equal, then its return True.
  - c. If Not, it will checks for the `is_subseq()` function of two strings.  
if string1 is sub-sequence of string 2 then return False  
if string2 is sub-sequence of string 1 then return True
  - d. Then If the first symbol of the pair strings are equal then, then It compares the remaining part of the strings by calling the `rpoeq()` function  
else `higher_prec()` function compare the first symbols based on precedence ordering.
  - e. `higher_prec()` function returns True then it will call the `rpoeq()` function with the input of string1, remaining part of the string2 and precedence ordering.  
`higher_prec()` function returns False then It will call the `rpoeq()` function with the input of remaining part of the string1, string2 and precedence ordering.

### Lemma1 Implementation:



1. The program calls `rpo()` function with the input of string1, string2 and precedence ordering.

2. `rpo()` checks for below conditions:
  - a. Condition1: If pair two strings are equal, then its return false.
  - b. Condition2: If string1 is empty and string2 has one or more symbols, then it will return the false. If it is vice versa, then returns true.
  - c. Condition3: If the pair of strings are not equal then
    - i. Firstly, It removes the common suffix by using function `remove_common_suffix()`.
    - ii. Secondly, If the pair of strings have different max symbols, then max symbol of string1 is higher than max symbol of string2.
    - iii. Next, If the pair of strings are having same max symbols, then `counter_function()` calculates the length of each maximal symbol for both the inputs and next, It will compare the both strings based on the precedence ordering.
    - iv. After comparing both the strings, if strings is having the same number of max symbols, then call `split ()` function at maximal symbols and compare the strings again.  
After calling the `split ()` function, Compare last piece the string that are present in list by calling recursively.

### 3 Requirements and Usage

1. Python 3.x.: You have to install Python 3 and related packages from <https://www.python.org/downloads/>.
2. standard build environment (make, gcc, etc.). To test the API, in command prompt use `$make test`
3. This algorithm is implemented in Python 3. Check for the python versions in your system using commands “`python –version`” or “`python3 –version`”

#### Folder structure:

-  Basic\_Theory
-  Lemma1 - 2Implementations

#### Run the code:

At the command prompt, use command

```
python3 Lemma1.py
python3 Basic.Theory.py
```

## 4 Conclusion

Implemented the two algorithms to determine whether a given pair of strings can be made RPO-comparable by selecting a partial order on their function symbols and compared running time. The Running time of both algorithms is  $O(n)$

## References

- [1] F. Baader and T. Nipkow. *Term rewriting and all that*. Cambridge University Press, 1998.
- [2] N. Dershowitz. Orderings for term-rewriting systems. *Theoretical Computer Science* 17(3): 279–301. 1982.
- [3] P. Narendran and M. Rusinowitch. The theory of total unary RPO is decidable. *Computational Logic – CL 2000, First International Conference, London, UK, 24-28 July, 2000, Proceedings*, volume 1861 of *Lecture Notes in Computer Science*, pages 660–672. Springer, 2000.

## 5 Appendix

### Python Code

#### 5.1 Basic\_Theory-Implementation.py

```
1
2 import time
3
4 #from prettyprint import *
5
6 import re
7
8 start = time.time()
9
10 def find_common_prefix(s1, s2):
11
12     prefix = ''
13     for i in range(min(len(s1), len(s2))):
14         if s1[i] == s2[i]:
15             prefix += s1[i]
16         else:
17             break
18     return prefix
19
20 """
21 >>> find_common_prefix('asd','aqwq')
22 a
23 >>> find_common_prefix('aabaabaa','aabacaa')
24 aaba
25 >>> find_common_prefix('ab','a')
26 a
27 >>> find_common_prefix('abcd','abef')
28 ab
29
30 """
31
32 def find_common_suffix(s1, s2):
33
34     suffix = ''
35     for i in range(1, min(len(s1), len(s2))+1):
36         if s1[-i] == s2[-i]:
37             suffix = s1[-i] + suffix
38         else:
39             break
40     return suffix
41
42 """
43 >>> find_common_suffix('ababab', 'ccabcb')
44 ab
45 >>> find_common_suffix('ababab', 'ccabcb')
46 b
47 >>> find_common_suffix('ababab', 'ab')
48 ab
```



```

49 """
50
51 def remove_common_prefix(s1, s2):
52
53     prefix = find_common_prefix(s1, s2)
54     return (s1[len(prefix):], s2[len(prefix):])
55
56 """
57 >>> remove_common_prefix('asd','aqwq')
58 ('sd','qwq')
59 >>> remove_common_prefix('aabaabaa','aabacaa')
60 ('abaa','caa')
61 >>> remove_common_prefix('ab','a')
62 ('b','')
63 >>> remove_common_prefix('abcd','abef')
64 ('cd','ef')
65
66 """
67
68 def remove_common_suffix(s1, s2):
69
70     suffix = find_common_suffix(s1, s2)
71     if (suffix == ''):
72         return (s1, s2)
73     else:
74         return (s1[:-len(suffix)], s2[:-len(suffix)])
75
76 """
77 >>> remove_common_suffix('ababab', 'ccabcb')
78 ('abab', 'ccabc')
79 >>> remove_common_suffix('ababab', 'ccabcb')
80 ('ababa', 'ccabc')
81 >>> remove_common_suffix('ababab', 'ab')
82 ('abab', '')
83 """
84
85 """
86 print("Process executed in {0} s".format(round((end-start),1)))
87 """
88
89 rcs = remove_common_suffix
90
91 end = time.time()
92
93 def is_subseq(x, y): # Stephan Pochmann on the internet
94     it = iter(y)
95     return all(c in it for c in x)
96
97 """
98 >>> is_subseq('ac', 'abc')
99 True
100 >>> is_subseq('ac', 'abbc')
101 True
102 >>> is_subseq('ac', 'bbbbabbc')
103 True
104 >>> is_subseq('ac', 'bbbcbab')
105 False

```

```

106 >>> is_subseq('', 'ac')
107 """
108
109 def tail(ls):
110     if (len(ls) == 0):
111         return None
112     else:
113         return ls[1:]
114
115 """
116 >>> tail(['a','d','as'])
117 ['d', 'as']
118 >>> tail(['aa','ad','aass'])
119 ['ad', 'aass']
120
121 """
122
123 def higher_prec(sym1, sym2, prec_list):
124     sym = hd(prec_list)
125     if (sym1 == sym):
126         return True
127     elif (sym2 == sym):
128         return False
129     else:
130         return higher_prec(sym1, sym2, tail(prec_list))
131
132 """
133 >>> higher_prec('b', 'd', pls)
134 True
135 >>> higher_prec('b', 'a', pls)
136 False
137 >>> higher_prec('d', 'a', pls)
138 False
139 """
140
141 def decompose(str, symbol):
142     w = re.split(symbol, str)
143     return w
144
145 """
146 >>> decompose('babbaba', 'a')
147 ['b', 'bb', 'b', '']
148 >>> decompose('babbabccac', 'a')
149 ['b', 'bb', 'bcc', 'c']
150 >>> decompose('babbaba', 'b')
151 ['', 'a', '', 'a', 'a']
152 """
153 def explode(str):
154     return list(str)
155
156 '''
157 >>> explode('ads')
158 ['a', 'd', 's']
159 >>> explode('qwq')
160 ['q', 'w', 'q']
161
162 '''

```

```

163
164 def hd(L):
165     if type(L) == type([]):
166         if len(L) == 0: return None
167         else: return L[0]
168     else: return None
169     """
170 >>> hd(['ad','a'])
171 ad
172 >>> hd(['','a'])
173
174 >>> hd(['a','aasa'])
175 a
176 """
177
178 def rpoeq(str1, str2, preclist):
179     (str1, str2) = remove_common_suffix(str1, str2)
180     if (str1 == str2):
181         return True
182     else:
183         if is_subseq(str1, str2):
184             return False
185         else:
186             if is_subseq(str2, str1):
187                 return True
188             else:
189                 xa = str1[0]
190                 xb = str2[0]
191                 if (xa == xb):
192                     return rpoeq(str1[1:], str2[1:], preclist)
193                 else:
194                     if higher_prec(xa, xb, preclist):
195                         return rpoeq(str1, str2[1:], preclist)
196                     else:
197                         return rpoeq(str1[1:], str2, preclist)
198
199     """
200 >>> rpoeq('abc', 'abb', ['a', 'b', 'c'])
201 False
202 >>> rpoeq('abc', 'bbbbbac', ['a', 'b', 'c'])
203 True
204 >>> rpoeq('abc', 'bbbbbacccc', ['a', 'b', 'c'])
205 True
206 """
207
208 def rpo(str1, str2, preclist):
209     if (str1 == str2):
210         return False
211     else:
212         return rpoeq(str1, str2, preclist)
213
214     """
215 >>> rpo('a', 'bbbb', ['a', 'b', 'c'])
216 True
217 >>> rpo('a', 'a', ['a', 'b', 'c'])
218 False
219 >>> rpo('ab', 'bbbbbbac', ['a', 'b', 'c'])

```

```

220 True
221 >>> rpo('abba', 'baab', ['a', 'b'])
222 True
223 """

```

## 5.2 Lemma1\_2-Implementations.py

```

1
2 import time
3 # from prettyprint import *
4 import re
5 start = time.time()
6
7 # Implmentation 1
8 def pos(x, L):
9     if (len(L) == 0):
10         return None
11     else:
12         if (x == L[0]):
13             return 0
14         else:
15             return (1 + (pos(x, L[1:])))
16
17 """
18 >>> pos('c', ['a','b','c'])
19 2
20 >>> pos('b', ['a','b','d'])
21 1
22 >>> pos('a', ['a','b'])
23 0
24 """
25
26 def rporec(str1, str2, m, n, i, j, flag, symlist):
27     if ((i > m) and (j > n)):
28         return flag
29     else:
30         if (j > n):
31             return True
32         elif (i > m):
33             return False
34         x = pos(str1[i], symlist)
35         y = pos(str2[j], symlist)
36         #print(i, " ", j, " ", x, " ", y);
37         if (x < y):
38             return rporec(str1, str2, m, n, i, j + 1, flag, symlist)
39         elif (x == y):
40             return rporec(str1, str2, m, n, i + 1, j + 1, flag, symlist)
41         )
42     else:
43         if (flag):
44             return rporec(str1, str2, m, n, i + 1, j, flag, symlist)
45         )
46     else:
47         return rporec(str1, str2, m, n, i + 1, j, True, symlist)
48     )

```

```

46
47 '''print(rporec('ac', 'abbc', 0, 0, 0, 0, False, ['a', 'b','c']))
48 False
49 print(rporec('abc', 'abbc', 0, 0, 2, 3, False, ['a', 'b','c']))
50 False
51 print(rporec('ab', 'bbbbbbbac', 3, 2, 0, 0, True, ['a', 'b','c']))
52 True'''
53
54 def rpo_original(str1, str2, symlist):
55     if (str1 == ''):
56         return False
57     elif (str2 == ''):
58         return True
59     m = len(str1) - 1
60     n = len(str2) - 1
61     return rporec(str1, str2, m, n, 0, 0, False, symlist)
62
63 rpoo = rpo_original
64
65 '''
66
67 >>> rpoo("abba","baab",['a', 'b'])
68 False
69 >>> rpoo("abab","baab",['a', 'b'])
70 False
71 >>> rpoo('ab', 'bbbbbbbac', ['a', 'b', 'c'])
72 True
73 >>> rpoo('a', 'a', ['a', 'b', 'c'])
74 False
75 >>> rpoo('a', 'bbbbbb', ['a', 'b', 'c'])
76 True
77
78 '''
79
80 end = time.time()
81
82
83 print("Implementation - Process executed in {0} ms".format((end-
84 start)*10**3))
85
86 #Implementation 2
87 start = time.time()
88 def find_common_suffix(s1, s2):
89
90     suffix = ''
91     for i in range(1, min(len(s1), len(s2))+1):
92         if s1[-i] == s2[-i]:
93             suffix = s1[-i] + suffix
94         else:
95             break
96     return suffix
97
98
99 """
100 >>> find_common_suffix('ababab', 'ccabcab')
101 ab

```

```

102 >>> find_common_suffix('ababab', 'ccabcb')
103 b
104 >>> find_common_suffix('ababab', 'ab')
105 ab
106 """
107
108 def remove_common_suffix(s1, s2):
109
110     suffix = find_common_suffix(s1, s2)
111     if (suffix == ''):
112         return (s1, s2)
113     else:
114         return (s1[:-len(suffix)], s2[:-len(suffix)])
115
116 """
117 >>> remove_common_suffix('ababab', 'ccabcb')
118 ('abab', 'ccabc')
119 >>> remove_common_suffix('ababab', 'ccabcb')
120 ('ababa', 'ccabc')
121 >>> remove_common_suffix('ababab', 'ab')
122 ('abab', '')
123 """
124
125
126 def ParikhVector(str1, chars):
127     l = []
128     for x in chars:
129         l.append(x)
130     for i,x in enumerate(l):
131         count = len([y for y in str1 if x ==y ])
132         l[i] = count
133     return l
134
135 """
136 >>> ParikhVector('asds',['a','s','d'])
137 [1, 2, 1]
138 >>> ParikhVector('ads',['a','s','d'])
139 [1, 1, 1]
140 >>> ParikhVector('sds',['a','s','d'])
141 [0, 2, 1]
142 """
143
144 def rpo(s1, s2, p_list):
145
146     if (s1 == s2):
147         return False
148     elif len(s1)==0 and len(s2)>0:
149         return False
150     elif len(s1)>0 and len(s2)==0:
151         return True
152     else:
153         s1,s2 = remove_common_suffix(s1,s2)
154         c1 = ParikhVector(s1,p_list)
155         c2 = ParikhVector(s2,p_list)
156         for i,x in enumerate(c1):
157             if x>0 and c2[i]==0:
158                 return True

```

```

159         elif x==0 and c2[i]>0:
160             return False
161         elif x==c2[i]==0:
162             continue
163         elif x>c2[i] and c2[i]>0:
164             return True
165         elif x==c2[i] and c2[i]>0:
166             splitted_str1 = s1.split(p_list[i])
167             splitted_str2 = s2.split(p_list[i])
168             return rpo(splitted_str1[-1], splitted_str2[-1],
p_list)
169         elif x<c2[i] and x>0:
170             return False
171         else:
172             return False
173
174     """
175     >>> rpo("a","bbbb",['a', 'b', 'c'])
176     True
177     >>> rpo("a","a",['a', 'b', 'c'])
178     False
179     >>> rpo("ab","bbbbbbac",['a', 'b', 'c'])
180     True
181     >>> rpo("abba","baab",['a', 'b'])
182     False
183     """
184
185 end = time.time()
186
187
188 print("Implementation2 - Process executed in {0} ms".format((end-
start)*10**3))

```

## 6 Examples:

### 6.1 Basic\_Theory-Implementation.py

Input & Output :

```

>>> rpo('a', 'bbbb', ['a', 'b', 'c'])
True
>>> rpo('a', 'a', ['a', 'b', 'c'])
False
>>> rpo('ab', 'bbbbbbac', ['a', 'b', 'c'])
True
>>> rpo('abba', 'baab', ['a', 'b'])
False

```

## 6.2 Lemma1\_2-Implementations.py

Input & Output :

```
>>> rpo ("a","bbbb", ['a', 'b', 'c'])
True
>>> rpo ("a","a", ['a', 'b', 'c'])
False
>>> rpo ("ab","bbbbbbac", ['a', 'b', 'c'])
True
>>> rpo ("abba","baab", ['a', 'b'])
False
```