$$\varepsilon \leqslant 8d \sqrt[2d]{\frac{\delta}{2}}$$

$$\Leftrightarrow \varepsilon \leqslant 8d (\sqrt[2d]{\frac{2}{\delta}})^{-1}$$

$$\Leftrightarrow \varepsilon \leqslant 8d2^{-\frac{1}{2d}\log(\frac{2}{\delta})}$$

$$\Leftrightarrow 2^{\frac{1}{2d}\log(\frac{2}{\delta})} \geqslant \frac{8d}{\varepsilon}$$

$$\Leftrightarrow \frac{1}{2d}\log(\frac{2}{\delta}) \geqslant \log(\frac{8d}{\varepsilon})$$

$$\Leftrightarrow \frac{4}{\varepsilon}\log\frac{2}{\delta} \leqslant \frac{8d}{\varepsilon}\log\frac{8d}{\varepsilon}$$

Then  $|A| = \frac{8d}{\varepsilon} \log \frac{8d}{\varepsilon} = max(\frac{4}{\varepsilon} \log \frac{2}{\delta}, \frac{8d}{\varepsilon} \log \frac{8d}{\varepsilon})$ . Therefore, according to theorem ??,  $\mathbb{P}(A \text{ is a } \varepsilon - \text{net}) \ge 1 - \delta$ .

Let's denote the element of A by  $a_1, a_2, ..., a_{|A|}$ . The probability that A is a subset of MIS(G) is  $P_A(\varepsilon)$ .

$$P_A(\varepsilon) = \mathbb{P}(\forall a \in A, \ a \in MIS(G)) = \prod_{i=1}^{|A|} \mathbb{P}(a_i \in MIS(G)) = \prod_{i=1}^{|A|} \frac{|V(G)| \cdot c}{|V(G)|} \geqslant c^{|A|}.$$
 Therefore the probability that  $A$  verify all the condition is  $\mathbb{P}(A \subseteq MIS(G))\mathbb{P}(A \text{ is a } \varepsilon\text{-net}) \geqslant (1 - \varepsilon)$ 

 $\delta c^{f(\frac{1}{\varepsilon})}$ .