

$$\varepsilon \leq 8d \sqrt[2d]{\frac{\delta}{2}} \Leftrightarrow \varepsilon \leq 8d \left( \sqrt[2d]{\frac{2}{\delta}} \right)^{-1} \Leftrightarrow \varepsilon \leq 8d 2^{-\frac{1}{2d} \log(\frac{2}{\delta})} \Leftrightarrow 2^{\frac{1}{2d} \log(\frac{2}{\delta})} \geq \frac{8d}{\varepsilon} \Leftrightarrow \frac{1}{2d} \log(\frac{2}{\delta}) \geq \log(\frac{8d}{\varepsilon})$$

$$\Leftrightarrow \frac{4}{\varepsilon} \log \frac{2}{\delta} \leq \frac{8d}{\varepsilon} \log \frac{8d}{\varepsilon}$$

Then  $|A| = \frac{8d}{\varepsilon} \log \frac{8d}{\varepsilon} = \max(\frac{4}{\varepsilon} \log \frac{2}{\delta}, \frac{8d}{\varepsilon} \log \frac{8d}{\varepsilon})$ . Therefore, according to theorem ??,  $\mathbb{P}(A \text{ is a } \varepsilon\text{-net}) \geq 1 - \delta$ . Let's denote the element of  $A$  by  $a_1, a_2, \dots, a_{|A|}$ . The probability that  $A$  is a subset of  $MIS(G)$  is  $P_A(\varepsilon)$ .

$$P_A(\varepsilon) = \mathbb{P}(\forall a \in A, a \in MIS(G)) = \prod_{i=1}^{|A|} \mathbb{P}(a_i \in MIS(G)) = \prod_{i=1}^{|A|} \frac{|V(G)| \cdot c}{|V(G)|} \geq c^{|A|}.$$

Therefore the probability that  $A$  verify all the condition is  $\mathbb{P}(A \subseteq MIS(G)) \mathbb{P}(A \text{ is a } \varepsilon\text{-net}) \geq (1 - \delta) c^{f(\frac{1}{\varepsilon})}$ .