James Lamb, Pawel Manikowski - Coursework 2 - Exercise 3 - Part E

STRONG SCALING

For strong scaling we modified the merge_sort_mpi.py:

$$Nel = int(1e6*1.25)$$

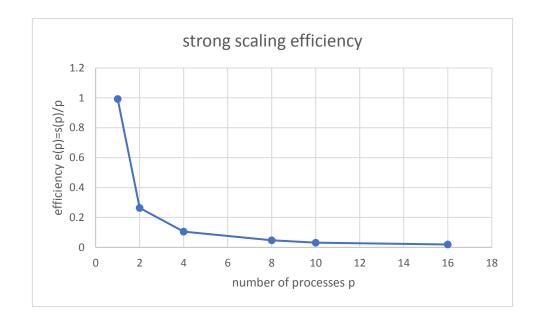
We copied the modified file to room 206 and then run it:

```
> mpirun -N p -f hostfile python myfile.py
```

where p was the number of processes. For each process we calculated the total time three times. We put the obtained values in the table below and calculated the average time, speed up and efficiency:

STRONG SCALING											
	# of processes	total time 1	total time 2	total time 3	average time	speed up	efficiency				
	р	t1	t2	t3	T = (t1+t2+t3)/3	$s(p) = T_s/T_p$	e(p)= s(p)/p				
serial	1	0.238562107	0.236366034	0.240278959	0.238402367	0.992716655	0.992716655				
parallel	2	0.473155022	0.432645083	0.440429926	0.448743343	0.527397238	0.263698619				
parallel	4	0.587161064	0.544538975	0.555813074	0.562504371	0.420736286	0.105184072				
parallel	8	0.621259928	0.637888908	0.621524096	0.626890977	0.377523379	0.047190422				
parallel	10	0.808695078	0.791343927	0.688639879	0.762892962	0.310221764	0.031022176				
parallel	16	0.730319023	0.778112173	0.80500412	0.771145105	0.306902032	0.019181377				

Using the data from the table we obtained the following graph of efficiency for strong scaling:



Clearly, the program does not have a very high strong scaling efficiency, with a 1% efficiency rate at 16 processes. Even at only 2 processes, the runtime is such that using a serial program would be more efficient. I will discuss parallelisation problems further in the next set of results.

WEAK SCALING

For weak scaling to get the local size of 10MB we needed to modify the merge_sort_mpi.py:

```
Nel = size*int(1e6*1.25)
```

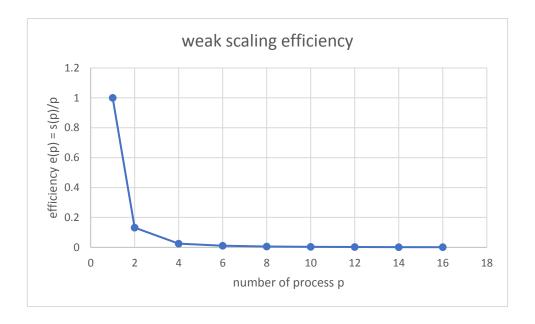
We copied the modified file to room 206 and then run it:

```
> mpirun -N p -f hostfile python myfile.py
```

where p was the number of processes. For each process we calculated the total time three times. We put the obtained values in the table below and calculated the average time, speed up and efficiency:

WEAK SCALING											
	# of processes	total time 1	total time 2	total time 3	average time	speed up	efficiency				
	р	t1	t2	t3	T = (t1+t2+t3)/3	$s(p) = T_s/T_p$	e(p)= s(p)/p				
serial	1	0.24	0.23	0.24	0.236666667	0.999997183	0.999997183				
parallel	2	0.898722887	0.869142056	0.918504	0.895456314	0.264296534	0.132148267				
parallel	4	2.315710783	2.345202923	2.389524937	2.350146214	0.100702671	0.025175668				
parallel	6	3.600500107	3.484977007	3.633229017	3.572902044	0.066239152	0.011039859				
parallel	8	5.287915945	5.462336063	5.492663145	5.414305051	0.043711242	0.005463905				
parallel	10	7.928426027	7.236097097	7.264084101	7.476202408	0.031655911	0.003165591				
parallel	12	8.074198008	7.962728977	8.642838001	8.226588329	0.028768426	0.002397369				
parallel	14	9.587563038	9.552712917	9.159389973	9.433221976	0.025088565	0.00179204				
parallel	16	12.50679994	12.57130003	12.849751	12.64261699	0.0187197	0.001169981				

Using the data from the table we obtained the following graph of efficiency for weak scaling:



Again there is a clear inability of the code to run efficiently on a parallel basis, in increasing the number of elements to be sorted we have reached less than 1% efficiency at 16 processes. Again even at 2 processes it would seem the serial code would be more efficient.

Although, it is true that the sorting of a vector is much faster using pre-sorted subvectors, we must also consider the communication bottlenecks. By breaking down the individual times for each section of the code, we find that as the process number increases, the time taken to presort the individual subvectors decreases. However, in order for this larger number of processes to effectively work together they must communicate these subvectors across the communicator, which then means both processors involved in the communication must take time to do so hence increasing the overall processing time for the program.

In the case of a weak scaling variant, this problem is exemplified further by the fact that as the number of processes increases, the overall workload of each process does not decrease and hence the impact of communication between the processes is no longer mitigated by the sharing of the workload.