

Dynamic Programming Longest Common Subsequence

ANINDITA KUNDU

LECTURER

DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

MILITARY INSTITUTE OF SCIENCE AND TECHNOLOGY

Longest Common Subsequence

A subsequence is a sequence that appears in the same relative order, but not necessarily contiguous.

For example:

X = "abcdefg"

"abc", "abg",

"bdf", "aeg",

"acefg", .. etc are subsequences of X .

Longest Common Subsequence

Given two sequences

$$X = \langle x_1, x_2, \dots, x_m \rangle$$

$$Y = \langle y_1, y_2, \dots, y_n \rangle$$

A subsequence of a given sequence is just the given sequence with zero or more elements left out.

A common subsequence $Z = \langle z_1, z_2, \dots, z_k \rangle$ of X and Y

- Z is a subsequence of both X and Y

Example:

$X =$ A **B** **C** **B** D **A** B

$Y =$ **B** D **C** A **B** **A**

Goal: Find the Longest Common Subsequence (LCS)

Develop a Recursive Solution

- Define $c[i, j]$ to be the length of an LCS of the sequences X_i and Y_j .
 - Goal: Find $c[m, n]$
 - Basis: $c[i, j] = 0$ if either $i = 0$ or $j = 0$
 - Recursion: How to define $c[i, j]$ recursively ?
- Finding an LCS of $X = \langle x_1, x_2, \dots, x_m \rangle$ and $Y = \langle y_1, y_2, \dots, y_n \rangle$
 - If $x_m = y_n$, then we must find an LCS of X_{m-1} and Y_{n-1} .
 - ◆ Appending $x_m = y_n$ to this LCS yields an LCS of X and Y .
 - If $x_m \neq y_n$, then we must solve two subproblems:
 - ◆ Finding an LCS of X_{m-1} and Y
 - ◆ Finding an LCS of X and Y_{n-1}
 - ◆ Whichever of these two LCSs is longer is an LCS of X and Y .

Develop a Recursive Solution

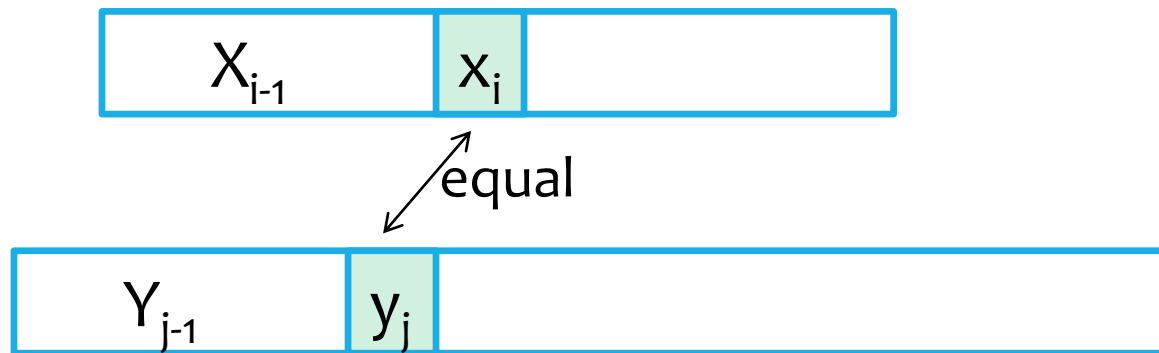
- The recursive formula is

$$c[i, j] = \begin{cases} c[i-1, j-1] + 1 & \text{if } i, j > 0 \text{ and } x[i] = y[j], \\ \max\{c[i, j-1], c[i-1, j]\} & \text{if } i, j > 0 \text{ and } x[i] \neq y[j] \end{cases}$$

Develop a Recursive Solution

Case 1: $x_i = y_j$

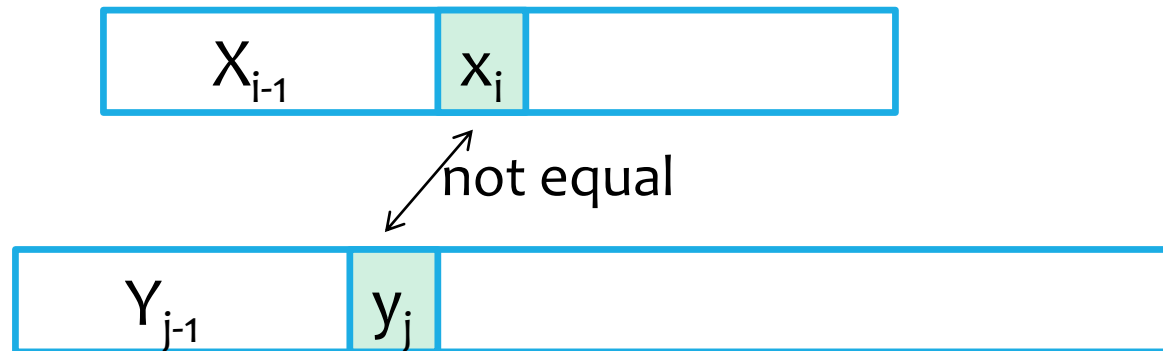
- Recursively find LCS of X_{i-1} and Y_{j-1} and append x_i
- So $c[i, j] = c[i-1, j-1] + 1$ if $i, j > 0$, and $x_i = y_j$



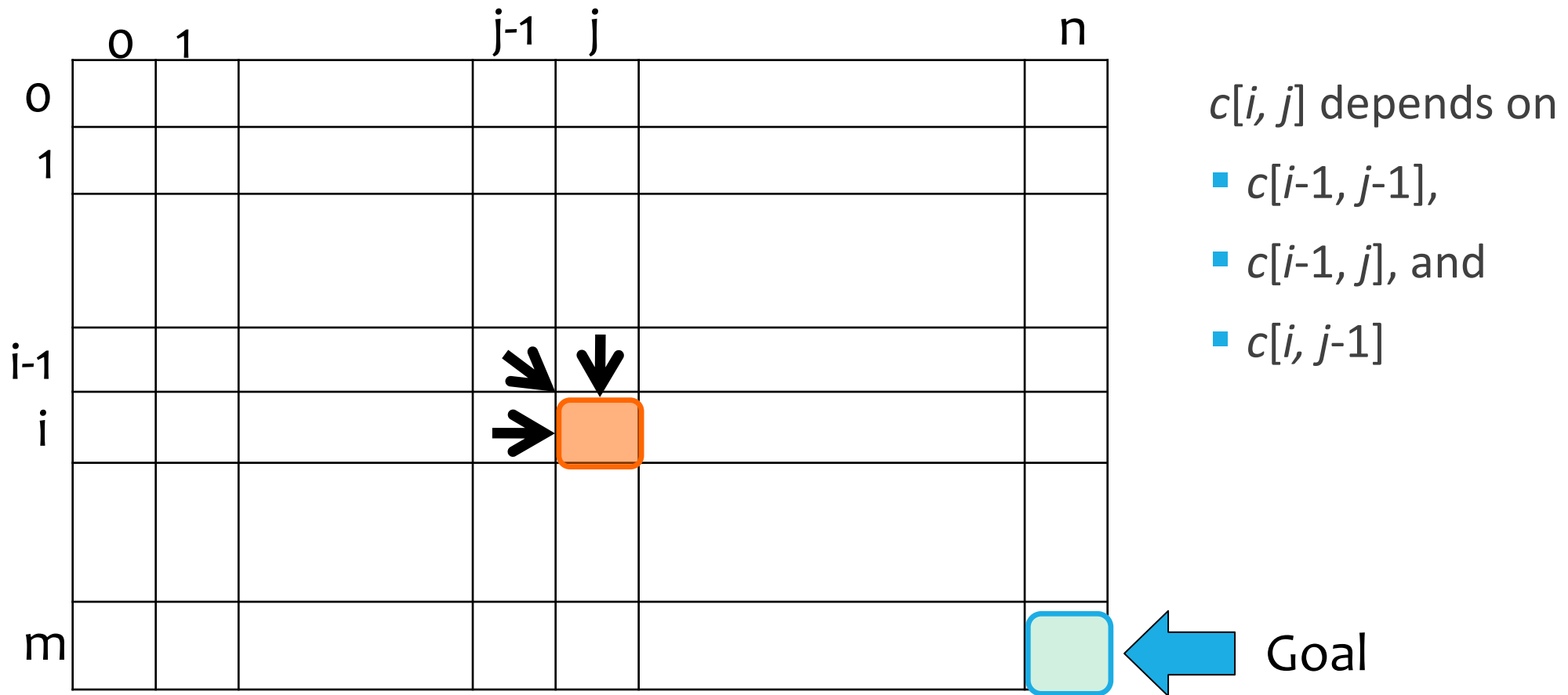
Develop a Recursive Solution

Case 2: $x_i \neq y_j$

- Recursively find LCS of X_{i-1} and Y_j
- Recursively find LCS of X_i and Y_{j-1}
- Take the longer one
- So $c[i, j] = \max\{c[i, j-1], c[i-1, j]\}$ if $i, j > 0$, and $x_i \neq y_j$



Dependencies among Sub-problems



- An order for solving the sub-problems (*i.e.*, filling in the array) that respects the dependencies is row major order:
 - do the rows from top to bottom
 - inside each row, go from left to right

Develop a Recursive Solution

LCS-LENGTH(X, Y)

```
1   $m \leftarrow \text{length}[X]$ 
2   $n \leftarrow \text{length}[Y]$ 
3  for  $i \leftarrow 1$  to  $m$ 
4      do  $c[i, 0] \leftarrow 0$ 
5  for  $j \leftarrow 0$  to  $n$ 
6      do  $c[0, j] \leftarrow 0$ 
7  for  $i \leftarrow 1$  to  $m$ 
8      do for  $j \leftarrow 1$  to  $n$ 
9          do if  $x_i = y_j$ 
10             then  $c[i, j] \leftarrow c[i - 1, j - 1] + 1$ 
11                  $b[i, j] \leftarrow \nwarrow$ 
12             else if  $c[i - 1, j] \geq c[i, j - 1]$ 
13                 then  $c[i, j] \leftarrow c[i - 1, j]$ 
14                      $b[i, j] \leftarrow \uparrow$ 
15                 else  $c[i, j] \leftarrow c[i, j - 1]$ 
16                      $b[i, j] \leftarrow \leftarrow$ 
17  return  $c$  and  $b$ 
```

☞ The algorithm calculates the values of each entry of the array $c[m, n]$.

☞ Each $c[i, j]$ is calculated in constant time, and there are $m \cdot n$ elements in the array.

☞ So the running time is $O(m \cdot n)$.

LCS Example

We'll see how LCS algorithm works on the following example:

$X = \text{ABCG}$

$Y = \text{BDCAG}$

$\text{LCS}(X, Y) = \text{BCG}$

$X = \text{A } \mathbf{B} \quad \mathbf{C} \quad \mathbf{G}$

$Y = \quad \mathbf{B} \mathbf{D} \mathbf{C} \mathbf{A} \mathbf{G}$

LCS Example

ABCG
BDCAG

		j	0	1	2	3	4	5
			y_j	B	D	C	A	G
i								
0	x_i							
1	A							
2	B							
3	C							
4	G							

$X = \text{ABCG};$

$m = |X| = 4$

$Y = \text{BDCAG};$

$n = |Y| = 5$

Allocate array:

$c[5, 4]$

LCS Example

ABCG
BDCAG

i	j	y_j	0	1	2	3	4	5
			x_i	B	D	C	A	G
0				0	0	0	0	0
1	A	0						
2	B	0						
3	C	0						
4	G	0						

for $i = 0$ to m $c[i, 0] = 0$

for $j = 1$ to n $c[0, j] = 0$

LCS Example

ABCG

BDCAG

		j					
		0	1	2	3	4	5
		y_j	B	D	C	A	G
i	x_i						
0		0	0	0	0	0	0
1	A	0	0				
2	B	0					
3	C	0					
4	G	0					

if ($x_i == y_j$)

$c[i, j] = c[i-1, j-1] + 1$

else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABCG

BDCAG

		j	0	1	2	3	4	5
		y_j		B	D	C	A	G
i	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	0		
2	B	0						
3	C	0						
4	G	0						

```

if (  $x_i == y_j$  )
     $c[i, j] = c[i-1, j-1] + 1$ 
else  $c[i, j] = \max( c[i-1, j], c[i, j-1] )$ 
    
```

LCS Example

ABCG
BDCA~~G~~

		j	0	1	2	3	4	5
			y_j	B	D	C	A	G
i	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	0	1	
2	B	0						
3	C	0						
4	G	0						

```

if (  $x_i == y_j$  )
     $c[i, j] = c[i-1, j-1] + 1$ 
else  $c[i, j] = \max( c[i-1, j], c[i, j-1] )$ 
    
```

LCS Example

ABCG

BDCA**G**

i	j	y_j	0	1	2	3	4	5
				B	D	C	A	G
0	x_i		0	0	0	0	0	0
1	A	0	0	0	0	0	1	1
2	B	0						
3	C	0						
4	G	0						

if ($x_i == y_j$)

$c[i, j] = c[i-1, j-1] + 1$

else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABCG

BDCAG

		j					
		0	1	2	3	4	5
		y_j	B	D	C	A	G
i	x_i						
0		0	0	0	0	0	0
1	A	0	0	0	0	1	1
2	B	0	1				
3	C	0					
4	G	0					

if ($x_i == y_j$)

$c[i, j] = c[i-1, j-1] + 1$

else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABCG
BDCAAG

		j	0	1	2	3	4	5
			y_j	B	D	C	A	G
i	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	0	1	1
2	B	0	1	→	1	→	1	↓
3	C	0						
4	G	0						

if ($x_i == y_j$)
 $c[i, j] = c[i-1, j-1] + 1$
 else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABCG

BDCA⁵G

		j	0	1	2	3	4	5
			y_j	B	D	C	A	G
i	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	0	1	1
2	B	0	0	1	1	1	1	1
3	C	0						
4	G	0						

if ($x_i == y_j$)

$c[i, j] = c[i-1, j-1] + 1$

else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABCG

BD CAG

i	j	y_j	0	1	2	3	4	5
				B	D	C	A	G
0	x_i		0	0	0	0	0	0
1	A		0	0	0	0	1	1
2	B		0	1	1	1	1	2
3	C		0	↓	↓			
4	G		0	1	1			

if ($x_i == y_j$)
 $c[i, j] = c[i-1, j-1] + 1$
 else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABC~~G~~

BD~~C~~AG

		j	0	1	2	3	4	5
			y_j	B	D	C	A	G
i	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	0	1	1
2	B	0	1	1	1	1	1	2
3	C	0	1	1	2			
4	G	0						

if ($x_i == y_j$)

$c[i, j] = c[i-1, j-1] + 1$

else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABC~~G~~

BDCAG

		j	0	1	2	3	4	5
			y_j	B	D	C	A	G
i	x_i	0						
1	A		0	0	0	0	1	1
2	B		0	1	1	1	1	1
3	C		0	1	1	2	→ 2	↓ 2
4	G		0					

if ($x_i == y_j$)

$c[i, j] = c[i-1, j-1] + 1$

else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABC**G**

BDCAG

		j					
		0	1	2	3	4	5
		y_j	B	D	C	A	G
i	x_i						
0		0	0	0	0	0	0
1	A	0	0	0	0	1	1
2	B	0	1	1	1	1	1
3	C	0	1	1	2	2	2
4	G	0	1				

if ($x_i == y_j$)

$c[i, j] = c[i-1, j-1] + 1$

else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABCG

BDCAG

		j					
		0	1	2	3	4	5
i	y_j		B	D	C	A	G
	x_i						
	0	0	0	0	0	0	0
	1	A	0	0	0	1	1
	2	B	0	1	1	1	1
	3	C	0	1	2	2	2
4	G	0	1	1	2	2	

if ($x_i == y_j$)

$c[i, j] = c[i-1, j-1] + 1$

else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

LCS Example

ABC**G**

BDCA**G**

i	j	y_j	0	1	2	3	4	5
				B	D	C	A	G
0	x_i		0	0	0	0	0	0
1	A		0	0	0	0	1	1
2	B		0	1	1	1	1	1
3	C		0	1	1	2	2	2
4	G		0	1	1	2	2	3

if ($x_i == y_j$)

$c[i, j] = c[i-1, j-1] + 1$

else $c[i, j] = \max(c[i-1, j], c[i, j-1])$

Another LCS Example

		j	0	1	2	3	4	5	6
		y_j		B	D	C	A	B	A
i	x_i								
0		0	0	0	0	0	0	0	0
1	A	0	↑	↑	↑	↖	1	←1	↖1
2	B	0	↖	1	←1	←1	↑	↖	2
3	C	0	↑	↑	↑	2	←2	↑	↑
4	B	0	↖	1	↑	↑	↑	↖	3
5	D	0	↑	↖	2	↑	↑	↑	↑
6	A	0	↑	↑	↑	↑	↖	↑	↖
7	B	0	↖	↑	↑	↑	↑	↖	↑

How to Find Actual LCS


So far, we have just found the *length* of LCS, but not LCS itself.

We can modify this algorithm to make it output an LCS of X and Y .

Each $c[i, j]$ depends on $c[i-1, j-1]$, or $c[i-1, j]$ and $c[i, j-1]$.

For each $c[i, j]$ we can say how it was acquired.

2	2
2	3



For example, here

$$c[i, j] = c[i-1, j-1] + 1 = 2 + 1 = 3$$

How to Find Actual LCS

Remember that

$$c[i, j] = \begin{cases} c[i-1, j-1] + 1 & \text{if } x[i] = y[j], \\ \max(c[i, j-1], c[i-1, j]) & \text{otherwise} \end{cases}$$

- So we can start from $c[m, n]$ and go backwards
- Whenever $c[i, j] = c[i-1, j-1] + 1$, remember $x[i]$, because $x[i]$ is a part of LCS
- When $i=0$ or $j=0$ (*i.e.* we reached the beginning), output remembered letters in reverse order

Finding LCS: Example

		j	0	1	2	3	4	5
		y_j		B	D	C	A	G
i	x_i	0	0	0	0	0	0	0
1	A	0	0	0	0	0	1	1
2	B	0	0	1	1	1	1	2
3	C	0	0	1	1	2	2	2
4	G	0	0	1	1	2	2	3

Finding LCS: Example

		j	0	1	2	3	4	5
		y_j		B	D	C	A	G
i	x_i							
0			0	0	0	0	0	0
1	A		0	0	0	0	1	1
2	B		0	1	1	1	1	2
3	C		0	1	1	2	2	2
4	G		0	1	1	2	2	3

LCS (reversed order): **G C B**

LCS (straight order): **B C G**

Finding LCS: Algorithm

PRINT-LCS(b, X, i, j)

1 **if** $i = 0$ or $j = 0$

2 **then return**

3 **if** $b[i, j] = \nwarrow$

4 **then** PRINT-LCS($b, X, i - 1, j - 1$)

5 print x_i

6 **elseif** $b[i, j] = \uparrow$

7 **then** PRINT-LCS($b, X, i - 1, j$)

8 **else** PRINT-LCS($b, X, i, j - 1$)

Trace backwards from $b[m, n]$

Thank You