# Morning Worning



# Formal Language And Automata

- > Languages
- > Autamata
- > Computation



# Language

What is a language?

This is a sentence.

This is also a sentence.

So we have

{ sentence 1, sentence 2, sentence 3, .....

the set of sentences  $\iff$  Language

# Sentence and Alphabet

- > sentence/string = sequence of symbols chosen from the alphabet  $\Sigma$  example : Mouse love rice.
- > alphabet = set of symbols

example: ASCII, 中文国标,  $\Sigma$  = {0,1}

 $symbols \Rightarrow sentences \Rightarrow language$ 

### Rules/Grammar

```
> rules = by which sentence is generated
  example: rules for English
      <sentence> → <noun-phrase><predicate>
      <noun-phrase> → <article><noun>
       \langle article \rangle \rightarrow a|an|the
       \langle noun \rangle \rightarrow wolf|sheep
       \langle verb \rangle \rightarrow love | eat
```

# Examples of language

example 1.1

```
L={w|w consists of 0's and 1's, rule and end with 0}
```

```
L={ 0,00,10,000,010,100,110,0000,.....}
```

 $11111100 \in L$ ,  $1 \notin L$ ,  $0001 \notin L$ ,  $20 \notin L$ 

# Examples of language

example 1.2 
$$L=\{ 0^n1^n \mid n \geq 0 \}$$

alphabet =  $\{0,1\}$ 

 $L=\{\varepsilon, 01, 0011, 000111, 00001111, 0000011111, .....\}$ 

# example 1.3 The empty language

$$L = \{ \} = \phi$$

L: contains no string

example 1.4

L={ w | w is a sentence in English }

Bolt won two gold medals in Daegu.

The sheep eats grass.  $\checkmark$ 

The grass eats sheep.

Formal language focus on form, not meaning

# String operations

$$w=a_1a_2....a_n$$

$$v=b_1b_2....b_n$$

### Concatenation

$$wv=a_1a_2....a_nb_1b_2....bn$$

abbabbbaaa

### Reverse

$$w^R = a_n a_{n-1} \dots a_1$$

# Another operation

$$\mathbf{w}^{n} = \mathbf{w} \mathbf{w} \dots \mathbf{w}$$

- $\rightarrow$  w = abb  $\Rightarrow$  w<sup>2</sup>=abbabb, w<sup>3</sup>=abbabbabb
- $\triangleright$  definition:  $w^0 = \varepsilon$

# Operations on languages

> The usual set operations

$$L_1 \cup L_2 = \{ w \mid w \in L_1 \text{ or } w \in L_2 \}$$
 $L_1 \cap L_2 = \{ w \mid w \in L_1 \text{ and } w \in L_2 \}$ 
 $L_1 - L_2 = \{ w \mid w \in L_1 \text{ and } w \notin L_2 \}$ 

> Reverse

$$L^{R} = \{ w^{R} \mid w \in L \}$$

Example 
$$L = \{a^nb^n \mid n \ge 1\} \Rightarrow L^R = \{b^na^n \mid n \ge 1\}$$

# Operations on languages

 $L^2 = ?$ 

Concatenation

$$L_1L_2 = \{ wv \mid w \in L_1 \text{ and } v \in L_2 \}$$

Example 
$$L = \{ a^n b^n \mid n \ge 1 \} , K = \{ b^n a^n \mid n \ge 1 \}$$
 
$$LK = \{ a^n b^n b^n a^n \mid n \ge 1 \}$$
 
$$LK = \{ a^n b^n b^m a^m \mid n \ge 1 , m \ge 1 \}$$

example 1.5

L={  $w \in \{0, 1\}^*$  | all 0's precede all 1's in w } L={  $\epsilon$ , 0, 1, 00, 01, 11, 000, 001, 011, 111, ..... }

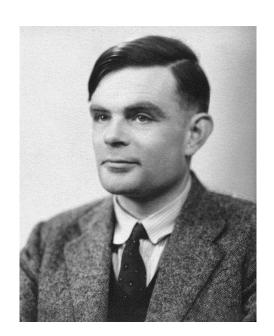
# example 1.6

```
L_1=\{w\in\{0,1\}^*\mid \text{ no prefix of } w \text{ contains } 1\}
   =\{\epsilon, 0, 00, 000, 0000, 00000, 000000, \dots \}
L_2=\{w\in\{0,1\}^*\mid \text{ no prefix of } w \text{ starts with } 1\}
  =\{\epsilon, 0,00,01,000,001,010,011,0000,0001,\ldots\}
  =\{w \in \{0,1\}^* \mid \text{ the first character of } w \text{ is } 0 \} \cup \{\varepsilon\}
 L_3 = \{w \in \{0,1\}^* \mid \text{ every prefix of } w \text{ starts } w \text{ ith } 1\}
   = \phi
```

# Automata

- Alan Marthison Turing
  - On Computable Numbers
     With an Application to
     the Entscheidungs Problem

Turing Machine

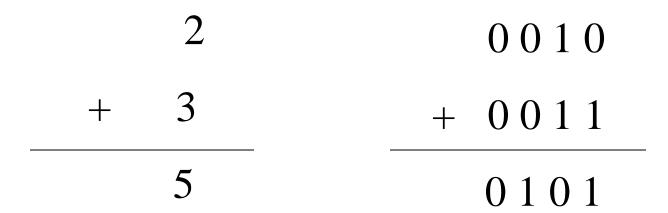


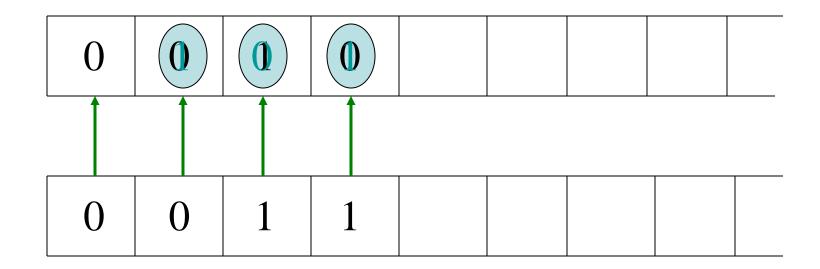


# Automata

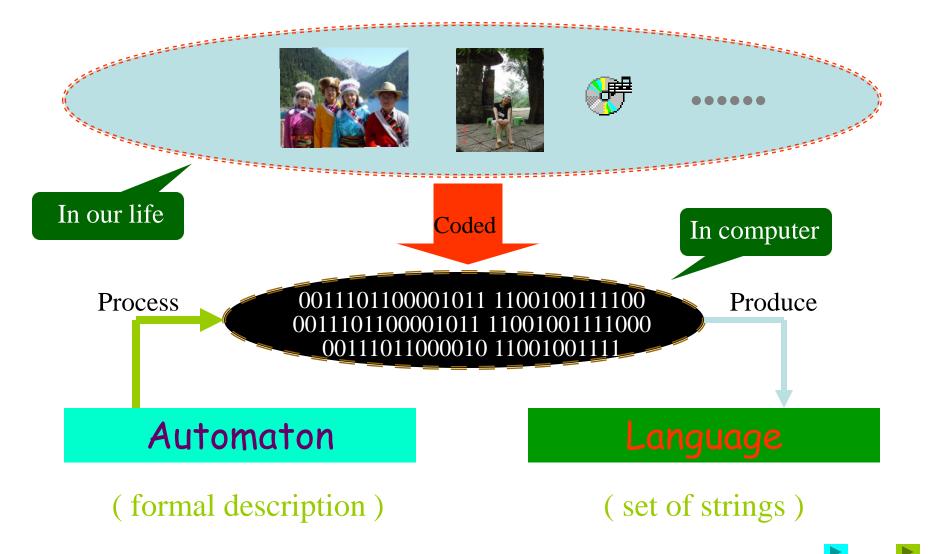
- Finite Automata
  - Deterministic Finite Automata
  - Non-deterministic Finite Automata
- Push Down Automata

# Computation





# Computation for computer



# Computation

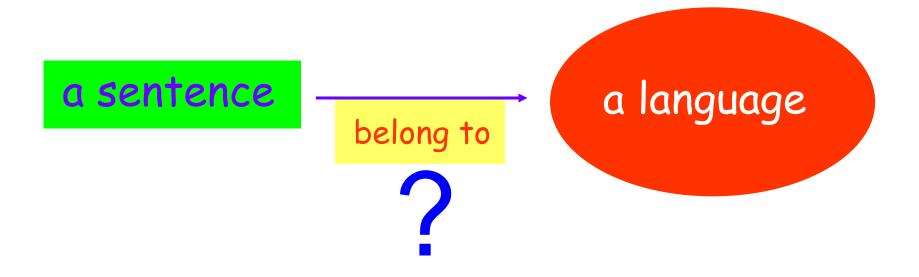
- Computable Problems
  - write a program to solve

- ◆ Intractable Problems
  - find someway to work around

## Undecidable Problem

```
main ()
   Int n, total, x, y, z;
  scanf("%d", &n);
  total=3;
  while(1){
      for(x=1;x<=total-2;x++)
         for(y=1;y<=total-x-1;y++){
            z=total-x-y;
            if(exp(x,n)+exp(y,n)==exp(z,n))
                printf("hello,world\n");
      total++;
```

## Undecidable Problem



automaton

# Content

Grammars Automata Languages Construction Properties Design Finite Regular Regular Expression Automaton Language Generate Recognize Push Down Context Free Context Free Automaton Language Grammar Recursively Turing (Phrase Enumerable Machine Grammar)

# Text book

1. Introduction to Automata Theory, Languages, and Computation (Third Edition)

John E. HopcroftRajeev MotwaniJeffrey D. Ullman

# Text book

2. An Introduction to Formal Languages and Automata (Third Edition )

—— Peter Linz

# Goal

- 1. Understanding "theoretical" concepts
- ----- method of formal description
- 2. Get a sense of how to reason formally
- 3. Improving reading ability with English

# Homework

- All exercises listed on qq-group
- Need not submit
- Check by free talk
- Discussions in group

# Honor and Collaboration

Collaboration is strongly encouraged

Solutions must be written independently

Responsible for Understanding and explaining

# Exam

Only final exam

Open exam

You are allowed to refer to text-book, class handouts, and notes during the exam

Closed exam

Nothing allowed except one pen

# Grading Policy

- ➤ Homework: 40%
- > Final exam: 60%
- > Addition: 10 points
  - Paper
  - Presentation
  - Participation

# Instructor

➤ Name: 孙大烈

➤ Office:综合楼 220

Mobile: 13936169569

> E-mail: sdl@hit.edu.cn

➤ 课程群: 自动机/104804608 (qq)

# Good good study day day day day day who